

Design of Digital Circuits and Systems

Finite State Machine Review

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Relevant Course Information

- ❖ hw1 due on Monday (4/1)
 - Homework can be completed in groups of up to 4
- ❖ Lab 1 report due Friday (4/5)
 - Labs can be completed in groups of up to 2
- ❖ Lab demos:
 - Lab demo sign up sheet sent out soon (check with partner)
 - 15 minutes for demos, early labs will be quicker
 - Make sure LabsLand is set up and synthesized *beforehand*
- ❖ Quiz 1 is Thursday, April 4 in last 25 min of lecture
 - Draw FSM state diagram & make design decisions

Lecture 1 Review

❖ Useful operators:

- Ternary operator: `<cond> ? <then> : <else>`
- Concatenation: `{sig, ..., sig}`
- Replication: `{n{m}}`

❖ A **parameter** is a named constant

```
parameter N = 8;           // bus width
parameter period = 100;   // timing constant
```

❖ A parameterized module:

- `module <name> #(<parameter list>) (<port list>);`
- Parameters can be given default values
 - e.g., `#(parameter N = 8)`

Review Question

$0b\ 0100$ $\xrightarrow{\text{zero}}$ $0b\ 0000\ 0100$
 $\xrightarrow{\text{sign}}$ $0b\ 0000\ 0100$
 $0b\ 1100$ $\xrightarrow{\text{zero}}$ $0b\ 0000\ 1100$
 $\xrightarrow{\text{sign}}$ $0b\ 1111\ 1100$

- ❖ There are two forms of bit extensions: zero-extension (add 0s) and sign-extension (copy MSB)
- ❖ Write out SystemVerilog pseudocode for a parameterized *extender* module
 - Inputs `sign` (1 bit), `in` (M bits); output `out` (N bits $> M$)
 - `out` should either be the sign-extended version of `in` (`sign` = 1) or the zero-extended version of `in` (`sign` = 0)

Review Question (Possible) Solution

```

module extender #(parameter M = 4, N = 8)
    (output [N-1:0] out,
     input [M-1:0] in,
     input      sign);
    assign out = sign ? {{(N-M){in[M-1]}},in} : {{(N-M){1'b0}},in};
endmodule
    
```

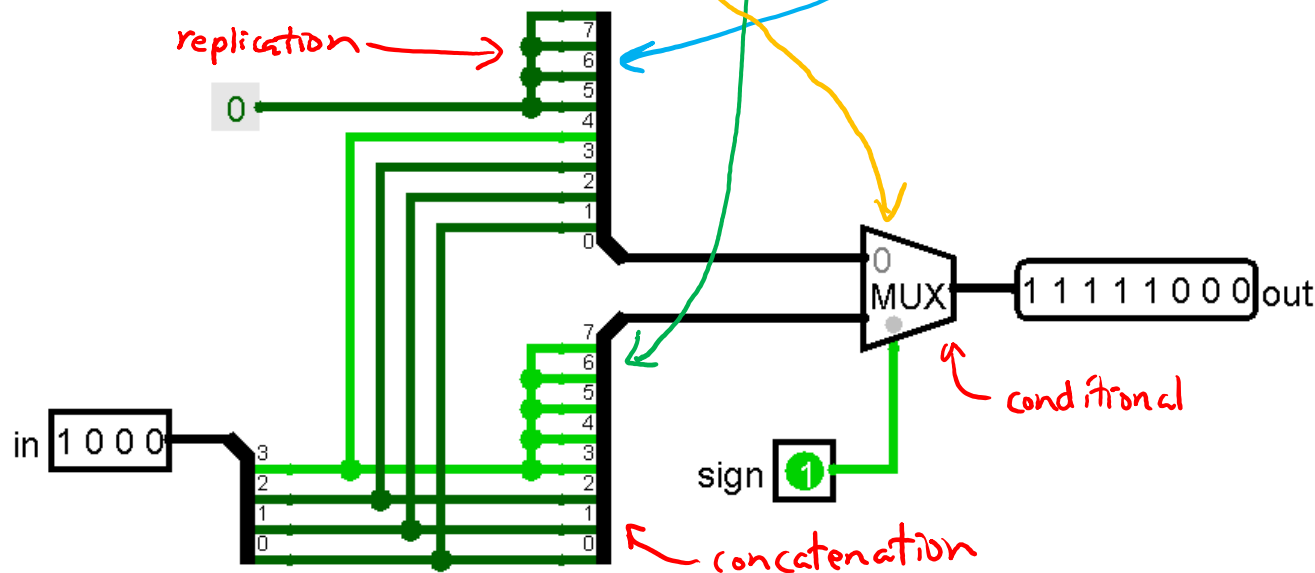
conditional/ternary (pointing to '?')

replication (pointing to `{{(N-M){in[M-1]}}`)

concatenation (pointing to `in` in the first branch)

concatenation (pointing to `in` in the second branch)

❖ Hardware if $M = 4$ and $N = 8$:



Lecture Outline

- ❖ **SystemVerilog Review & Tips (Cont.)**
- ❖ FSMs
- ❖ Test Benches

Structural vs. Behavioral Revisited

- ❖ Not a strict definition of these terms, so exact classification is not that important
- ❖ Structural:
 - Instantiating modules (library and user-defined) and defining port connections
 - **assign**: continuous assignment
 - Used with nets

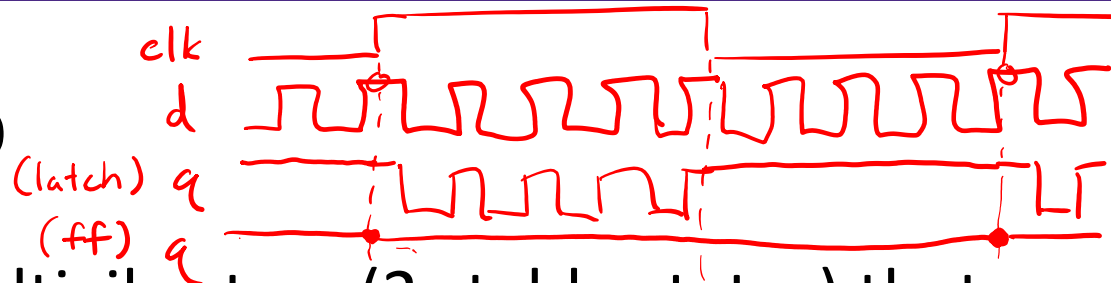
Verilog Procedural Blocks

- ❖ A *procedural block* is made up of behavioral code in the form of procedural statements whose effects are interpreted sequentially
 - The block itself is awakened/triggered in a non-sequential manner
- ❖ **initial**: block triggered once at time zero
 - Non-synthesizable (*i.e.*, for simulation/testbenches only)
- ❖ **always**: block triggered by a *sensitivity list*
 - Any object that is assigned a value in an **always** statement must be declared as a variable (*e.g.*, **logic** or **reg**).

SystemVerilog Procedural Blocks

- ❖ SystemVerilog introduced variants on always that are generally more robust and more specialized
- ❖ ~~@()~~
always_comb: intended for combinational logic
 - Sensitivity list is automatically built
- ❖ **always_latch**: intended for latch-based logic
 - Sensitivity list is automatically built
- ❖ **always_ff**: intended for sequential logic
(*i.e.*, synchronous/clocked)
@(posedge clk)
 - Sensitivity list must be specified

Latch vs. Flip-Flop



- ❖ Both are bistable multivibrators (2 stable states) that can store information
- ❖ A latch is *asynchronous*; a flip-flop is *edge-triggered*

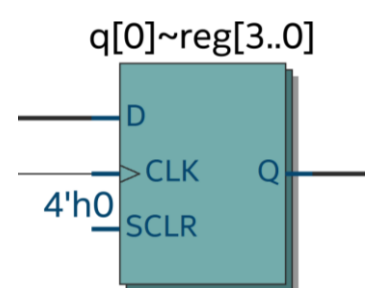
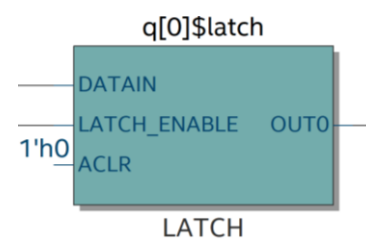
```

module my_latch(input logic clk,
               input logic [3:0] d,
               output logic [3:0] q);

    always_latch
        if (clk) q <= d;

endmodule
    
```

q will update to d as long as clk = 1
← "latched" when clk = 0



```

module my_ff(input logic clk,
            input logic [3:0] d,
            output logic [3:0] q);

    always_ff @(posedge clk)
        q <= d;

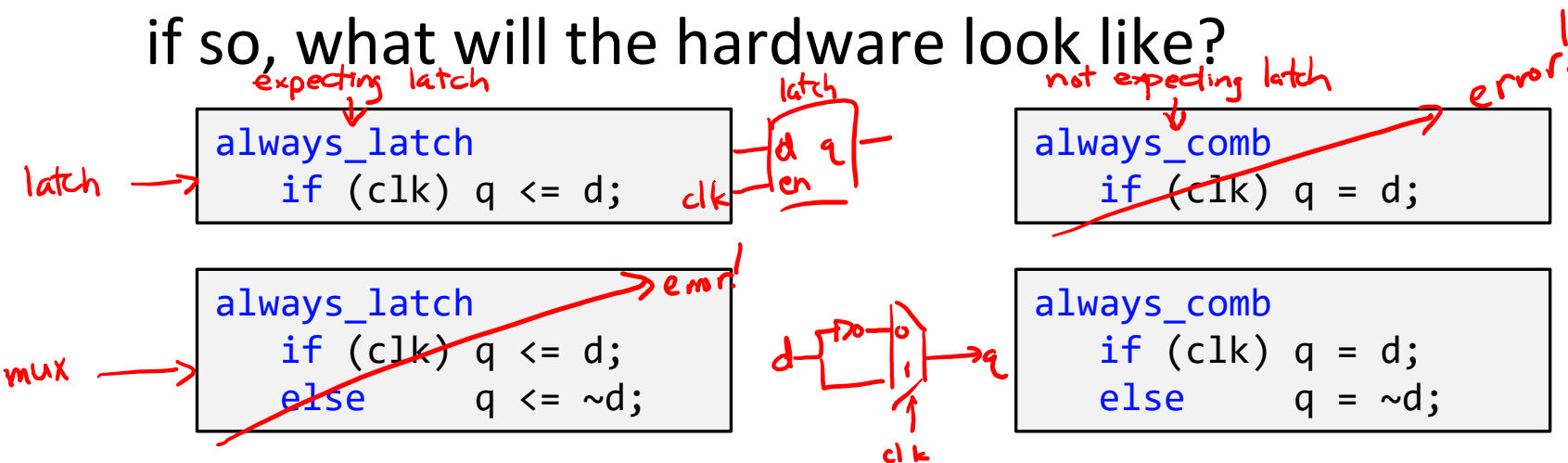
endmodule
    
```

only updates at rising edge of clock

Inferred Latches

- ❖ **Warning:** easy to write code with inadvertent latches
 - Check your synthesis output for “Inferred latch”
 - Usually from incomplete assignments – unspecified branch infers latch behavior

❖ **Question:** which of the following will synthesize and, if so, what will the hardware look like?



- Demo: Tools → “Netlist Viewers” → “RTL Viewer”

case Statement

- ❖ Create combinational logic and is easier to read than lots of `if/else` statements
 - Must always be inside an `always` block
 - Each case has an implied C-style break

```
module seven_seg(bcd, segs);  
  
    input  logic [3:0] bcd;  
    output logic [6:0] segs;  
  
    always_comb  
        case (bcd)  
            //                abc_defg  
            0: segs = 7'b011_1111;  
            1: segs = 7'b000_0110;  
            2: segs = 7'b101_1011;  
            3: segs = 7'b100_1111;  
            4: segs = 7'b110_0110;  
            5: segs = 7'b110_1101;  
            6: segs = 7'b111_1101;  
            7: segs = 7'b000_0111;  
            8: segs = 7'b111_1111;  
            9: segs = 7'b110_1111;  
  
        endcase  
  
    endmodule
```

case Statement

- ❖ Create combinational logic and is easier to read than lots of `if/else` statements
 - Must always be inside an `always` block
 - Each case has an implied C-style break
 - Remember to use `default` to avoid incomplete assignments!

```
module seven_seg(bcd, segs);  
  
    input  logic [3:0] bcd;  
    output logic [6:0] segs;  
  
    always_comb  
        case (bcd)  
            //                abc_defg  
            0: segs = 7'b011_1111;  
            1: segs = 7'b000_0110;  
            2: segs = 7'b101_1011;  
            3: segs = 7'b100_1111;  
            4: segs = 7'b110_0110;  
            5: segs = 7'b110_1101;  
            6: segs = 7'b111_1101;  
            7: segs = 7'b000_0111;  
            8: segs = 7'b111_1111;  
            9: segs = 7'b110_1111;  
            → default: segs = 7'bX;  
        endcase  
  
endmodule
```

Other SystemVerilog Resources

- ❖ SystemVerilog Language Reference Manual
 - On website, Verilog → Reference Manual
 - 586 pages...
- ❖ SystemVerilog articles
 - <https://www.systemverilog.io/>
 - <http://www.verilogpro.com/>
 - <https://www.chipverify.com/systemverilog/systemverilog-tutorial>
- ❖ One style guide for SystemVerilog
 - <https://www.systemverilog.io/styleguide>
 - We won't enforce, but good guidelines

Technology Break

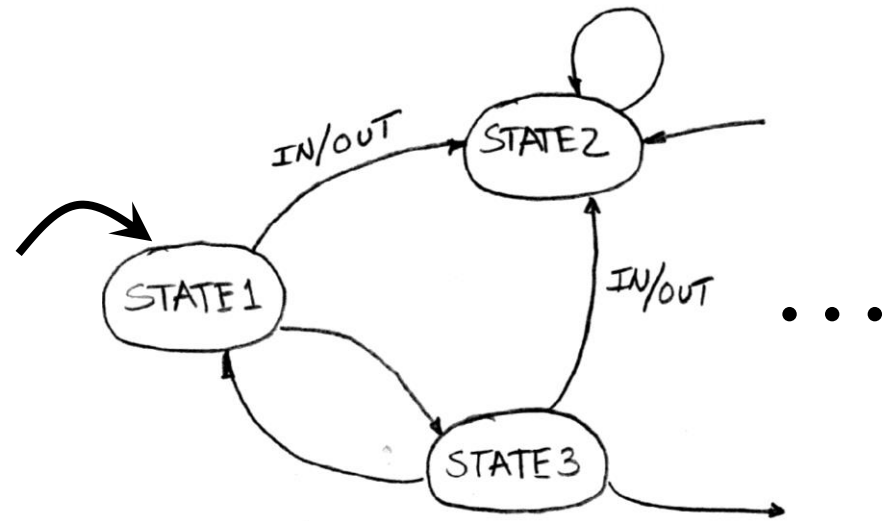
Lecture Outline

- ❖ SystemVerilog Review & Tips (Cont.)
- ❖ **Finite State Machine Design**
- ❖ Test Benches

Finite State Machines (FSMs)

- ❖ A convenient way to conceptualize computation over time using a *state transition diagram*
 - Consists of a *set of states*, an *initial state*, and a *transition function*

- ❖ FSM implementations come in 3 blocks:
 - State register (SL)
 - Next state logic (CL)
 - Output logic (CL)



FSM Implementation Notes

- ❖ States must be assigned a binary encoding
 - More readable by using parameters or an enum
 - Encoding choices can affect logic simplification
- ❖ Reset signal can be synchronous (responds to clk) or asynchronous (responds to reset)
 - Determined by whether or not reset is in sensitivity list
- ❖ State logic (next state logic + state update) can be written as 1 combined block or 2 separate blocks
- ❖ If input is asynchronous, may want to add a two-flip-flop *synchronizer* to deal with metastability



FSM SystemVerilog Design Pattern

alternatives exist!

❖ Which, if any, construct(s) would you expect to use for each of the following basic sections of a module that implements an FSM?

not logic!

enum logic [W-1:0] { A, B, C } ps, ns;

■ *// define states and state variables*
 initial assign always_comb always_ff None

combinational logic

■ *// next state Logic (ns)*
 initial assign always_comb always_ff None

combinational logic

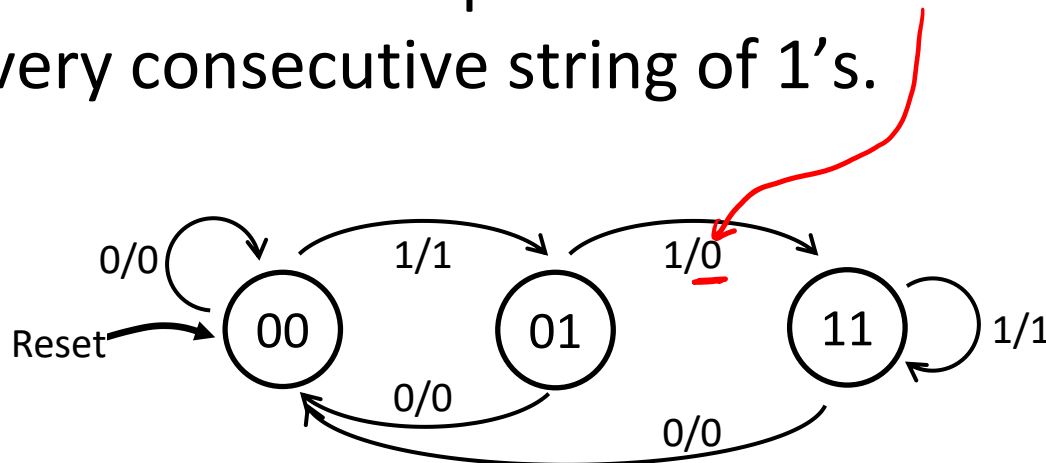
■ *// output Logic*
 initial assign always_comb always_ff None

sequential logic

■ *// state update Logic (ps)*
 initial assign always_comb always_ff None

FSM Example: String Manipulator

- ❖ Takes in a stream of inputs and removes the *second* 1 from every consecutive string of 1's.



- Example inputs: 0 1 0 1 1 0 1 1 1 0 1 1 1 1 ...
 outputs: 0 1 0 1 0 0 1 0 1 0 1 0 1 1 ...

String Manipulator FSM

```
module fsm (input logic clk, reset, in,
            output logic out);
```

```
① // present and next state
enum logic [1:0] {S0, S1, S3} ps, ns;
```

```
② // next state logic
always_comb
```

```
case (ps)
```

```
S0: if (in) ns = S1;
     else ns = S0;
```

```
S1: if (in) ns = S3;
     else ns = S0;
```

```
S3: if (in) ns = S3;
     else ns = S0;
```

```
endcase
```

```
③ // output logic
```

```
assign out = in & (ps[1] | ~ps[0]);
```

```
...
```

```
...
```

```
// sequential logic (DFFs)
```

```
④ // synchronous reset
```

```
always_ff @(posedge clk)
```

```
if (reset)
```

```
ps <= S0; // reset state
```

```
else
```

```
ps <= ns;
```

```
endmodule // fsm
```

could also be: ns = in ? S1 : S0;

Moore vs. Mealy

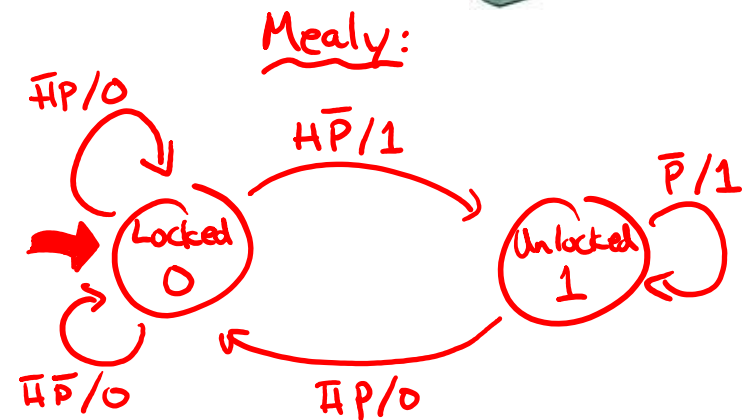
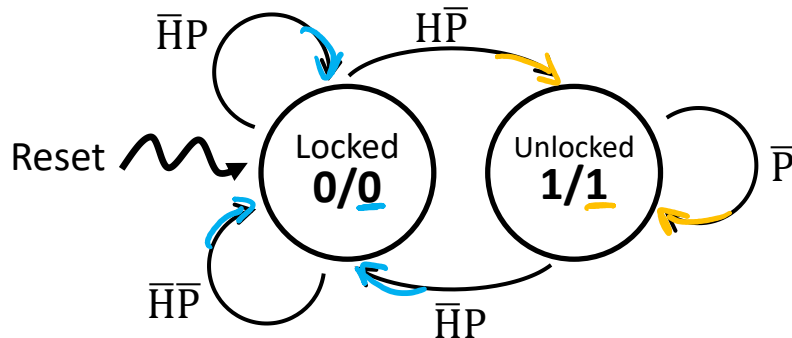
- ❖ **Moore** machines define their outputs based on states ($\textcircled{00/1}$) and **Mealy** machines define outputs based on transitions ($\xrightarrow{0/1}$)
 - Mealy machines are more *flexible*
 - Moore outputs are function of state; Mealy outputs are function of state *and inputs*
 - All FSMs can be expressed in either form, but some systems are more naturally expressed one way versus the other
 - Feel free to use either in this class if not specified
- ★ However, there *are* implementation differences!

Mealy \leftrightarrow Moore Conversions

Not testable material

❖ **Moore \rightarrow Mealy:** copy the state output to every transition *entering* the state

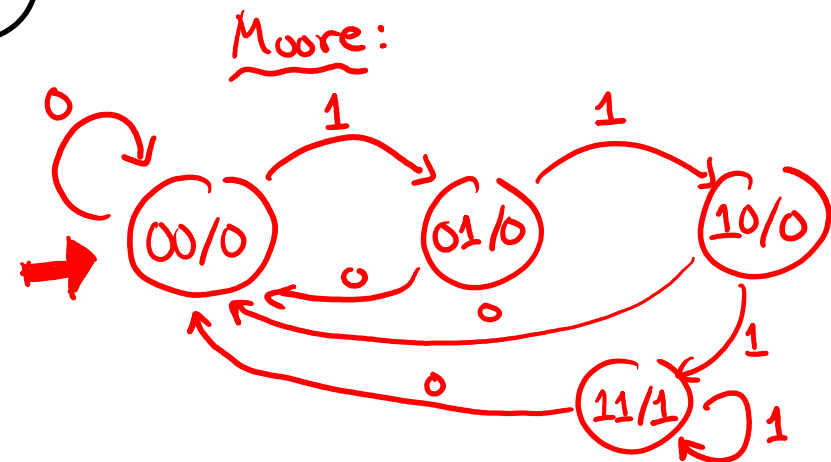
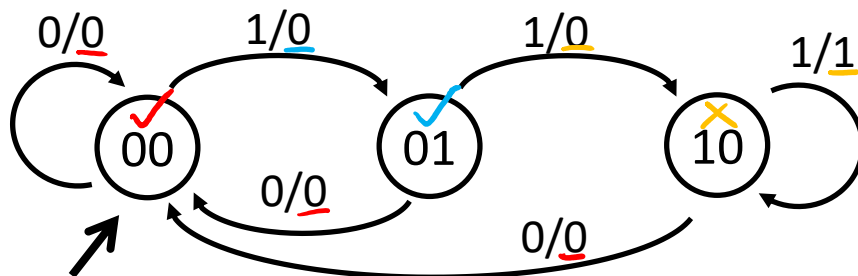
❖ Example: FSM for a *turnstile*, which is locked until someone swipes their Husky ID (input H) and then locks once you push through (input P) the unlocked gate. Outputs a light that glows red (0) or green (1).



Mealy \leftrightarrow Moore Conversions

Not testable
material

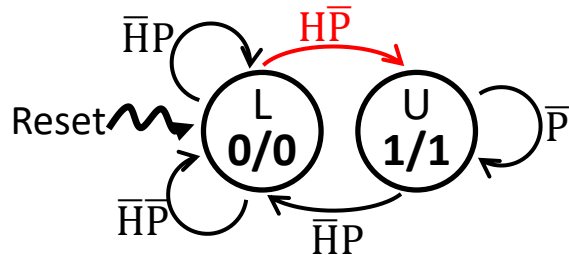
- ❖ **Mealy \rightarrow Moore:** more complicated process; if incoming transitions differ in output, may need to “split” the state
- ❖ Example: the threeOnes FSM from Lecture 1



Moore vs. Mealy Outputs

- ❖ Compare a Moore and Mealy FSM for the turnstile. Complete the statements and waveform below, assuming no delays:

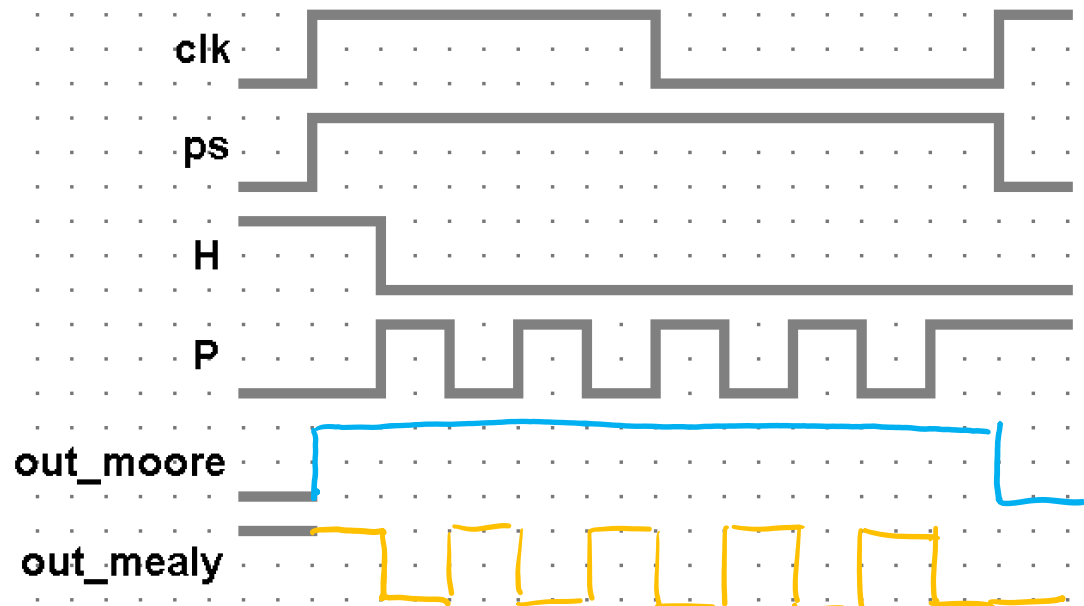
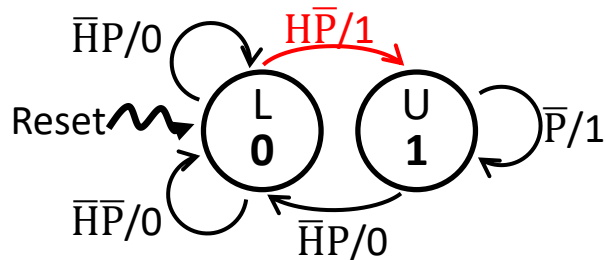
Moore:



```

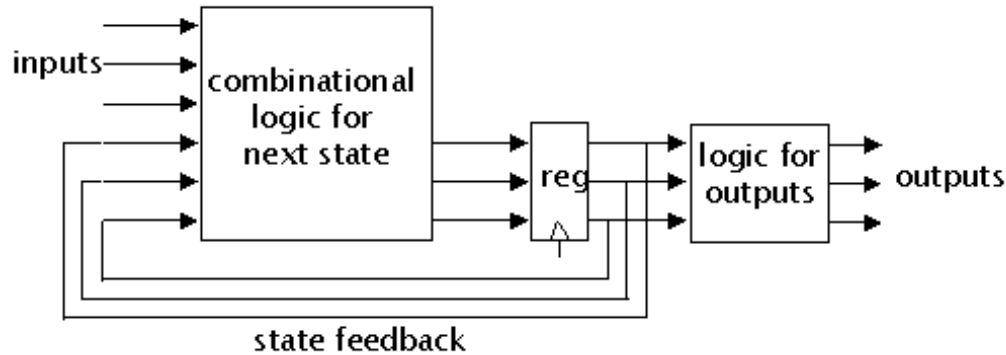
assign out_moore = (ps == U);
assign out_mealy =
    ((ps == U) & ~P) |
    ((ps == L) & H & ~P);
  
```

Mealy:



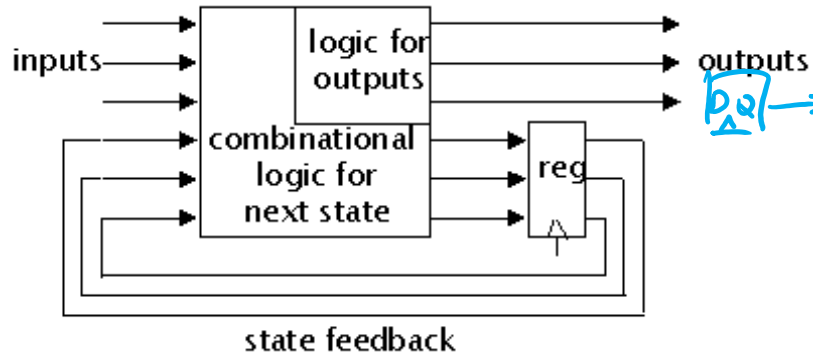
Moore vs. Mealy Outputs

❖ Moore:



- Outputs change synchronously with state changes

❖ Mealy:



can add register to synchronize output changes

- Input changes can cause *immediate* output changes

(often change sooner than Moore outputs)

Lecture Outline

- ❖ SystemVerilog Review & Tips (Cont.)
- ❖ Finite State Machine Design
- ❖ **Test Benches**

Test Benches

- ❖ Special modules *needed for simulation only!*
 - Software constraint to mimic hardware
- ❖ ModelSim runs entirely on your computer
 - Tries to simulate your FPGA environment without actually using hardware – no physical signals available
 - Must create fake inputs for FPGA's physical connections
 - *e.g.*, LEDR, HEX, KEY, SW, CLOCK_50
 - Unnecessary when code is loaded onto FPGA
- ❖ Need to define both input signal combinations as well as their timing

Test Bench Timing Controls

- ❖ Delay: `#<time>` `#10` `#{period/2}`
 - Delays by a specific amount of simulation time
- ❖ Edge-sensitive: `@(<pos/neg>edge <signal>)`

@(posedge clk)

 - Delays next statement until specified transition on signal
- ❖ Level-sensitive Event: `wait(<expression>)`

wait (clk)

 - Waits until `<expression>` evaluates to TRUE
- ❖ Stop simulation: `$stop;`
- ❖ Timescale: ``timescale <time unit> / <precision>`

each arbitrary simulation time unit (#1)

 - e.g., ``timescale 1 ns / 1 ps`

rounding

Extender Test Bench

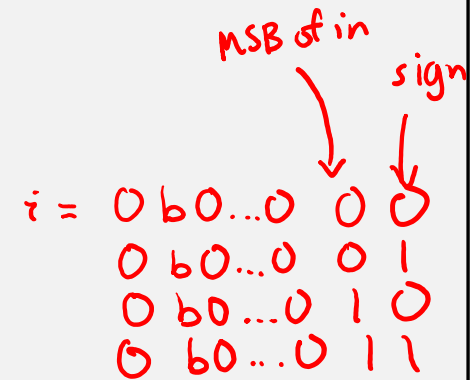
```

`timescale 1 ns / 1 ns
module extender_tb();

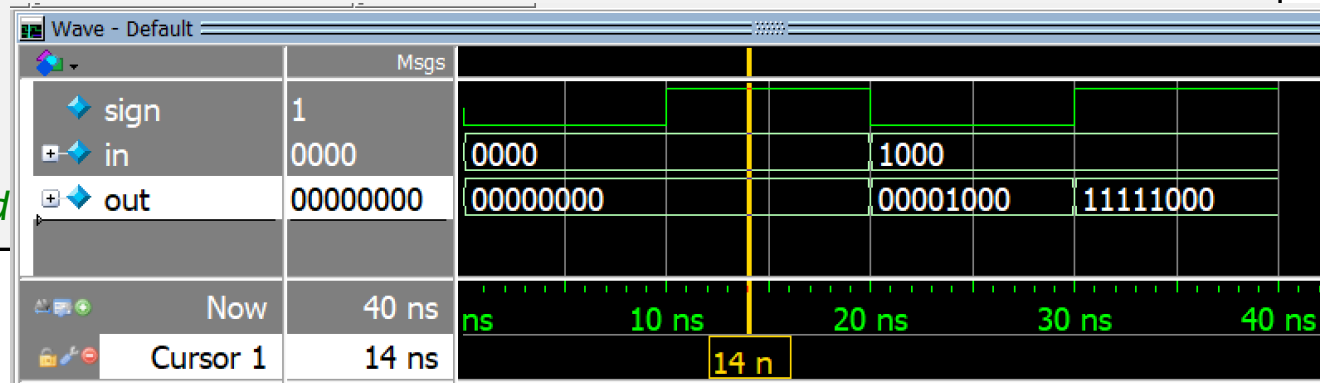
    parameter M = 4, N = 8;
    logic [M-1:0] in;
    logic [N-1:0] out;
    logic sign;

    extender #(M, N) dut (.*);

    int i;
    initial begin
        for (i = 0; i < 2**2; i++) begin
            sign = i[0]; in = {i[1], {(M-1){1'b0}}}; #10;
        end // for
        $stop;
    end // initial
endmodule // extend
    
```



exponentiation (2^2)



FSM Test Bench Notes

- ❖ Your main goal is to test every transition that we care about – may take extra clock cycles
- ❖ For simulation, you need to generate a clock signal
 - Assume we have `parameter clock_period;`

Explicit

Edges:

```
initial
  clk = 0;

always_comb begin
  #(clock_period/2) clk <= 1;
  #(clock_period/2) clk <= 0;
end
```

Toggle:

```
initial begin
  clk <= 0;
  forever #(clock_period/2) clk <= ~clk;
end
```

String Manipulator Test Bench

```

module fsm_tb();

  logic clk, reset, in, out;

  fsm dut (.*);

  // simulated clock
  parameter period = 100;
  initial begin
    clk <= 0;
    forever
      #(period/2)
        clk <= ~clk;
  end // initial clock

  ...

```

```

...
initial begin
  reset <= 1; in <= 0; @(posedge clk);
  reset <= 0; in <= 0; @(posedge clk);
  in <= 0; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 0; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 0; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 1; @(posedge clk);
  in <= 1; @(posedge clk);
  $stop; // end simulation
end // initial signals

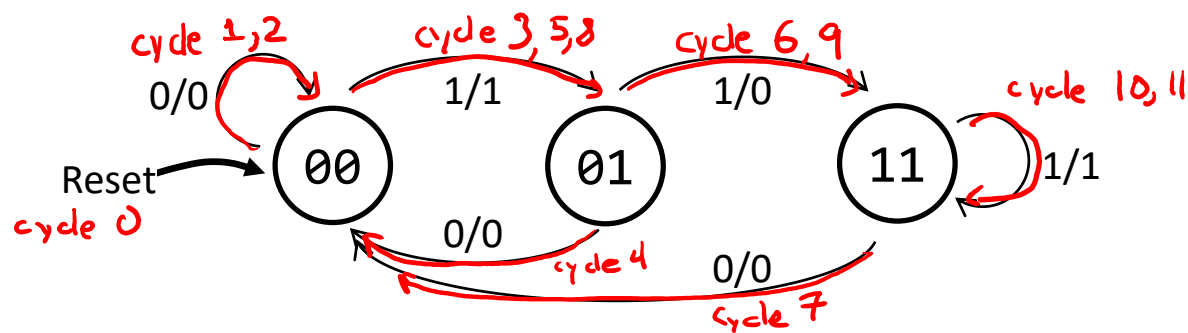
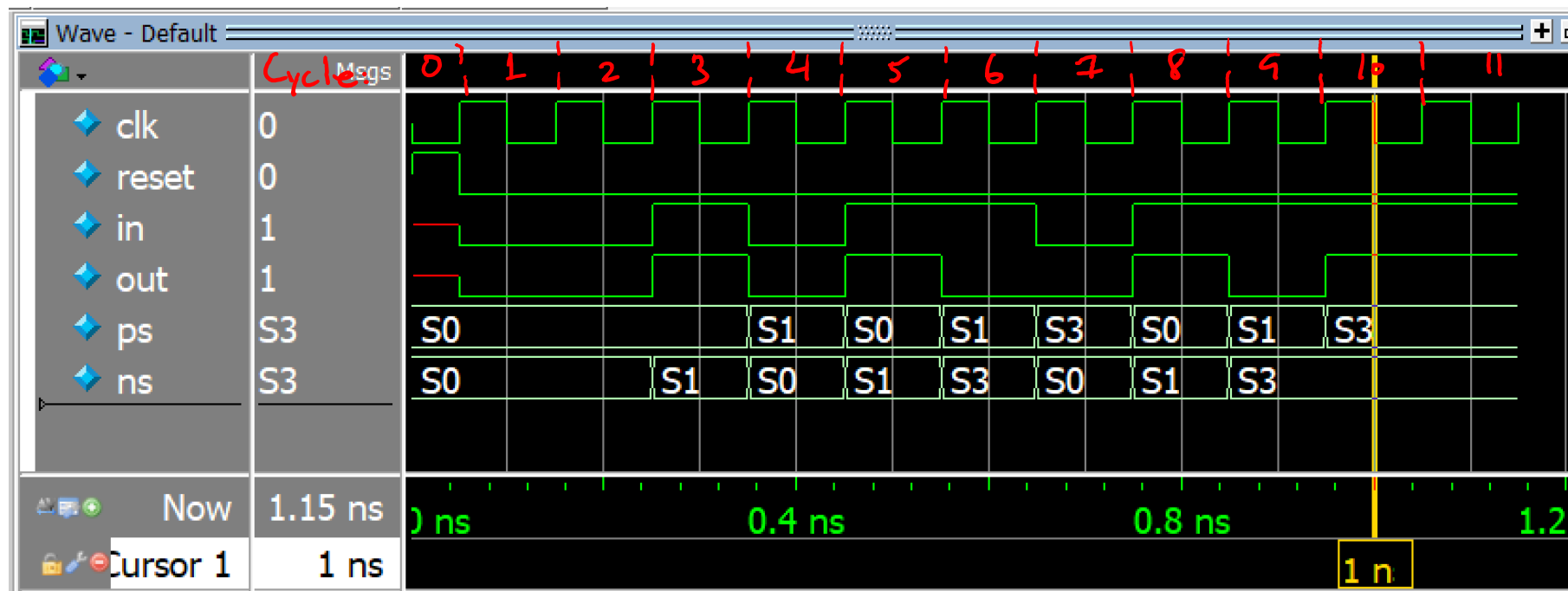
endmodule // fsm_tb

```


① signal changes

② the clock edge that reads the new signal values

String Manipulator Waveforms



Checking Responses

- ❖ Visually checking simulated waveforms quickly becomes impractical for large designs simulated over thousands of clock cycles
 - Displaying and explaining your waveforms for labs can be tedious
- ❖ There are simulator-independent system tasks to write messages to the user/tester!
 - Look similar to `printf()` in C
 - `$<system_task>(<format_string>, <sig_1>, <sig_2>, ...)`The diagram shows three red arrows. One arrow points from the `printf()` function in the previous bullet point to the `<format_string>` argument in the system task syntax. A second arrow points from the `printf()` function to the `<sig_1>` argument. A third arrow points from the `printf()` function to the `<sig_2>` argument.
 - Will look at `$display` today and others later on

Checking Responses: \$display

- ❖ Triggers once when encountered, prints the given format string and adds a new line:

```
// define test inputs
int i;
initial begin
  for (i = 0; i < 2**2; i++) begin
    sign = i[0]; in = {i[1], {(M-1){1'b0}}}; #10;
    $display("t = %0t, %b %s %b",
             (current simulation time)$time, in, sign ? "-+->" : "-0->", out);
  end // for
  $stop;
end // initial
```

```
Transcript
VSI...
# t = 10, 0000 -+-> 00000000
# t = 20, 0000 -+-> 00000000
# t = 30, 1000 -0-> 00001000
# t = 40, 1000 -+-> 11111000
```

Format Specifiers

Table 5.7: Format Specifiers.

| Specifier | Meaning |
|-----------|---------------------------------------|
| %h | Hexadecimal format |
| %d | Decimal format (signed) |
| %o | Octal format |
| %b | Binary format |
| %c | ASCII character format |
| %v | Net signalstrength |
| %m | Hierarchical name of current scope |
| %s | String |
| %t | Time |
| %e | Real in exponential format |
| %f | Real in decimal format |
| %g | Real in exponential or decimal format |

escape character
Table 5.8: Special characters.

| Symbol | Meaning |
|--------|---|
| \n | New line |
| \t | Tab |
| \\ | \ character |
| \" | " character |
| \xyz | Where xyz is are octal digits - the character given by that octal code |
| %% | % character |

- **Warning:** these differ from the specifiers for printf
- The *minimum* field width is specified by numbers between the '%' and specifier letter
 - e.g., %3d will pad out to 3 digits if necessary,
%0d will show just the minimum number of digits needed