## Lecture 20

- Logistics
- HW6 due Wednesday
- Lab 7 this week (Tuesday exception)
- Midterm 2 Friday (covers material up to simple FSM (today))
- Review on Thursday
- Yoky office hour on Friday moved to Thursday 12-1:20pm online
- Last lecture
- Counter design
- Finite state machine - started vending machine example
- Today
- Continue on the vending machine example
- Moore/Mealy machines


## The "WHY" slide

- Finite State Machine (FSM)
- This is what we have been waiting for in this class. Using combinational and sequential logics, now you can design a lot of clever digital logic circuits for functional products. We will learn different steps you take to go from word problems to logic circuits.
- Moore/Mealy machines
- There are two different ways to express the FSMs with respect to the output. Both have different advantages so it is good to know them.


## FSM design

- Counter-design procedure

1. State diagram
2. State-transition table
3. Next-state logic minimization
4. Implement the design

- FSM-design procedure

1. State diagram
2. state-transition table
3. State minimization
4. State encoding
5. Next-state logic minimization
6. Implement the design

## Example: A vending machine

- 15 cents for a cup of coffee
- Doesn't take pennies or quarters
- Doesn't provide any change
- FSM-design procedure


1. State diagram
2. state-transition table
3. State minimization
4. State encoding
5. Next-state logic minimization
6. Implement the design

## A vending machine: (conceptual) state diagram



## A vending machine: State transition table

|  | present state | inp | N | next state | output open |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | S0 | 0 | 0 | S0 | 0 |
|  |  | 0 | 1 | S1 | 0 |
|  |  | 1 | 0 | S2 | 0 |
|  |  | 1 | 1 | X | X |
|  | S1 | 0 | 0 | S1 | 0 |
|  |  | 0 | 1 | S3 | 0 |
|  |  | 1 | 0 | S4 | 0 |
|  |  | 1 | 1 | X | X |
|  | S2 | 0 | 0 | S2 | 0 |
|  |  | 0 | 1 | S5 | 0 |
|  |  | 1 | 0 | S6 | 0 |
|  |  | 1 | 1 | X | $x$ |
|  | S3 | 0 | 0 | S3 | 0 |
|  |  | 0 | 1 | S7 | 0 |
|  |  | 1 | 0 | S8 | 0 |
|  |  | 1 | 1 | X | X |
|  | S4 | X | X | S4 | 1 |
|  | S5 | X | X | S5 | 1 |
|  | 56 | X | X | 56 | 1 |
|  | S7 | X | X | S7 | 1 |
|  | 58 | X | X | S8 | 1 |
| CSE370, Lecture 20 |  |  |  |  |  |

## A vending machine: State minimization



## A vending machine: State encoding

## -

| present state | inputs | next state | output |
| :---: | :---: | :---: | :---: |
| Q1 Q0 | D N | D1 D0 | open |
| 00 | 0 0 | 00 | 0 |
|  | 01 | 01 | 0 |
|  | 10 | 10 | 0 |
|  | 11 | - - | - |
| 1 | 00 | 01 | 0 |
|  | 01 | 10 | 0 |
|  | 10 | 11 | 0 |
|  | 11 | - - | - |
| 10 | 0 0 | 10 | 0 |
|  | 01 | 11 | 0 |
|  | 10 | 11 | 0 |
|  | 11 | - - | - |
| 1 | - - | 11 | 1 |

A vending machine: Logic minimization


$$
\begin{aligned}
& \mathrm{D} 1=\mathrm{Q} 1+\mathrm{D}+\mathrm{Q} 0 \mathrm{~N} \\
& \mathrm{D} 0=\mathrm{Q} 0^{\prime} \mathrm{N}+\mathrm{Q} 0 \mathrm{~N}^{\prime}+\mathrm{Q} 1 \mathrm{~N}+\mathrm{Q} 1 \mathrm{D} \\
& \mathrm{OPEN}=\mathrm{Q} 1 \mathrm{Q} 0
\end{aligned}
$$

## A vending machine: Implementation

## -



## Generalized FSM model: Moore and Mealy

Combinational logic computes next state and outputs

- Next state is a function of current state and inputs
- Outputs are functions of

KCurrent state (Moore machine)
$\boldsymbol{K}$ Current state and inputs (Mealy machine)


## Moore versus Mealy machines



## Example 10 -> 01: Moore or Mealy?

Circuits recognize $A B=10$ followed by $A B=01$

- What kinds of machines are they?


Moore


## Example 01/10 detector: a Moore machine

- Output is a function of state only
- Specify output in the state bubble



## Example 01/10 detector: a Mealy machine

Output is a function of state and inputs

- Specify outputs on transition arcs


| reset | input | current <br> state | next <br> state | current <br> output |
| :---: | :---: | :---: | :---: | :---: |
| 1 | - | - | A | 0 |
| 0 | 0 | A | B | 0 |
| 0 | 1 | A | C | 0 |
| 0 | 0 | B | B | 0 |
| 0 | 1 | B | C | 1 |
| 0 | 0 | C | B | 1 |
| 0 | 1 | C | C | 0 |

## Comparing Moore and Mealy machines

- Moore machines
+ Safer to use because outputs change at clock edge
- May take additional logic to decode state into outputs
- Mealy machines
+ Typically have fewer states
+ React faster to inputs - don't wait for clock
- Asynchronous outputs can be dangerous
- We often design synchronous Mealy machines
- Design a Mealy machine
- Then register the outputs


## Synchronous (registered) Mealy machine

- Registered state and registered outputs
- No glitches on outputs
- No race conditions between communicating machines

state feedback


## Example 0 -> 1: Moore or Mealy?

- Recognize $A, B=0,1$
- Mealy or Moore?


Registered Mealy
(actually Moore)


