

## Logistics

- ◆ Always read the Calendar at <http://www.cs.washington.edu/370>
- ◆ HW1 is posted on the web in the calendar --- due 1/14 11:30am
- ◆ Email list: please sign up on the web via the link from the course homepage
- ◆ TA Office Hours:
  - Josh Snyder Mondays 3:30-4:20, CSE 220
  - Aaron Miller Tuesdays 12:30-1:20, CSE 220
  - Sara Rolfe Tuesdays 2:30-3:20, CSE 220
- ◆ My Office Hours:
  - Mondays 12:20-1:00, CSE 668 (grab me after class)
  - TBA

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## Lecture 2: Number Systems

- ◆ Last lecture
  - Class introduction and overview
- ◆ Today
  - Binary numbers
  - Base conversion
  - Number systems
    - ◇ Two's-complement
  - A/D and D/A conversion

There are 10 kinds of people in the world. Those who understand binary... and those who don't.

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## The "WHY" slide

- ◆ Binary numbers
  - All computers work with 0's and 1's so it is like learning alphabets before learning English
- ◆ Base conversion
  - For convenience, people use other bases (like decimal, hexadecimal) and we need to know how to convert from one to another.
- ◆ Number systems
  - There are more than one way to express a number in binary. So 1010 could be 10, -2, -5 or -6 and need to know which one.
- ◆ A/D and D/A conversion
  - Real world signals come in continuous/analog format and it is good to know generally how they become 0's and 1's (and vice versa).

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## Digital

- ◆ Digital = discrete
  - Binary codes (example: BCD)
  - Decimal digits 0-9
- ◆ Binary codes
  - Represent symbols using binary digits (bits)
- ◆ Digital computers:
  - I/O is digital
    - ◇ ASCII, decimal, etc.
  - Internal representation is binary
    - ◇ Process information in bits

Decimal Symbols	BCD Code
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

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## The basics: Binary numbers

- ◆ Bases we will use
  - Binary: Base 2
  - Octal: Base 8
  - Decimal: Base 10
  - Hexadecimal: Base 16 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F
- ◆ Positional number system
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$
  - $63_8 = 6 \times 8^1 + 3 \times 8^0$
  - $A1_{16} = 10 \times 16^1 + 1 \times 16^0$
- ◆ Addition and subtraction

$$\begin{array}{r} 1011 \\ + 1010 \\ \hline 10101 \end{array} \quad \begin{array}{r} 1011 \\ - 0110 \\ \hline 0101 \end{array}$$

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## Binary → hex/decimal/octal conversion

- ◆ Conversion from binary to octal/hex
  - Binary: 10011110001
  - Octal: 10 | 011 | 110 | 001 = 2361<sub>8</sub>
  - Hex: 100 | 1111 | 0001 = 4F1<sub>16</sub>
- ◆ Conversion from binary to decimal
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 5_{10}$
  - $63.4_8 = 6 \times 8^1 + 3 \times 8^0 + 4 \times 8^{-1} = 51.5_{10}$
  - $A1_{16} = 10 \times 16^1 + 1 \times 16^0 = 161_{10}$

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## Decimal → binary/octal/hex conversion

Binary			Octal		
Quotient	Remainder		Quotient	Remainder	
56 ÷ 2 =	28	0	56 ÷ 8 =	7	0
28 ÷ 2 =	14	0	7 ÷ 8 =	0	7
14 ÷ 2 =	7	0			
7 ÷ 2 =	3	1			
3 ÷ 2 =	1	1			
1 ÷ 2 =	0	1			
			56 <sub>10</sub> =	111000 <sub>2</sub>	
			56 <sub>10</sub> =	70 <sub>8</sub>	

- Why does this work?
  - $N = 56_{10} = 111000_2$
  - $Q = N/2 = 56/2 = 111000/2 = 11100$  remainder 0
- Each successive divide liberates an LSB (least significant bit)

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## Number systems

- How do we write negative binary numbers?
- Historically: 3 approaches
  - Sign-and-magnitude
  - Ones-complement
  - Twos-complement
- For all 3, the most-significant bit (MSB) is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative
- twos-complement is the important one
  - Simplifies arithmetic
  - Used almost universally

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## Sign-and-magnitude

- The most-significant bit (MSB) is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative
- The remaining bits are the number's magnitude
- Problem 1: Two representations for zero
  - 0 = 0000 and also -0 = 1000
- Problem 2: Arithmetic is cumbersome

Add		Subtract		Compare and subtract	
4	0100	4	0100	0100	
+ 3	+ 0011	- 3	+ 1011	- 0011	
= 7	= 0111	= 1	≠ 1111	= 0001	
				- 4	1100
				+ 3	+ 0011
				- 1	≠ 1111
					= 1001

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## Ones-complement

- Negative number: Bitwise complement positive number
  - 0111 ≡ 7<sub>10</sub>
  - 1000 ≡ -7<sub>10</sub>
- Solves the arithmetic problem
 

Add	Invert, add, add carry	Invert and add
4    0100	4    0100	- 4    1011
+ 3   + 0011	- 3   + 1100	+ 3   + 0011
= 7   = 0111	= 1   1 0000	- 1   1110
	add carry: +1	
	= 0001	
- Remaining problem: Two representations for zero
  - 0 = 0000 and also -0 = 1111

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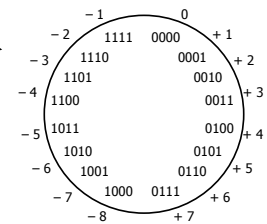
## Why ones-complement works

- The ones-complement of an 8-bit positive  $y$  is  $11111111_2 - y$
- What is  $11111111_2$ ?
  - 1 less than  $1\ 00000000_2 \equiv 2^8 \equiv 256_{10}$
  - So in ones-complement  $-y$  is represented by  $(2^8 - 1) - y$
- Adding representations of  $x$  and  $-y$  where  $x, y$  are positive we get  $(2^8 - 1) + x - y$ 
  - If  $x < y$  then  $x - y < 0$  there is no carry and get -ve number
    - Just add the representations if no carry
  - If  $x > y$  then  $x - y > 0$  there is a carry and get +ve number
    - Need to add 1 and ignore the 2<sup>8</sup> i.e. "add the carry"
  - If  $x = y$  then answer should be 0, get  $2^8 - 1 = 11111111_2$

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## Twos-complement

- Negative number: Bitwise complement plus one
  - 0111 ≡ 7<sub>10</sub>
  - 1001 ≡ -7<sub>10</sub>
- Number wheel
- Only one zero!
- MSB is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative



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## Twos-complement (con't)

- ◆ Complementing a complement  $\rightarrow$  the original number
- ◆ Arithmetic is easy
  - Subtraction = negation and addition
  - ◇ Easy to implement in hardware

Add	Invert and add	Invert and add
$\begin{array}{r} 4 \quad 0100 \\ + 3 \quad + 0011 \\ \hline = 7 \quad = 0111 \end{array}$	$\begin{array}{r} 4 \quad 0100 \\ - 3 \quad + 1101 \\ \hline = 1 \quad 1 \ 0001 \\ \text{drop carry} = 0001 \end{array}$	$\begin{array}{r} -4 \quad 1100 \\ + 3 \quad + 0011 \\ \hline -1 \quad 1111 \end{array}$

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## Why twos-complement works better

- ◆ Recall: The ones-complement of a b-bit positive y is  $(2^b - 1) - y$
- ◆ Adding 1 to get the twos-complement represents  $-y$  by  $2^b - y$ 
  - So  $-y$  and  $2^b - y$  are equal mod  $2^b$  (leave the same remainder when divided by  $2^b$ )
  - Ignoring carries is equivalent to doing arithmetic mod  $2^b$
- ◆ Adding representations of x and  $-y$  yields  $2^b + x - y$ 
  - If there is a carry then that means  $x \geq y$  and dropping the carry yields  $x - y$
  - If there is no carry then  $x < y$  and then we can think of it as  $2^b - (y - x)$

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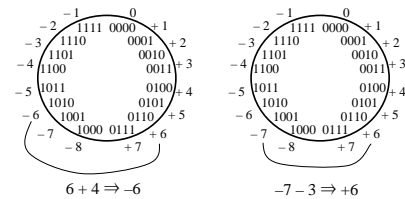
## Miscellaneous

- ◆ Twos-complement of non-integers
  - $1.6875_{10} = 01.1011_2$
  - $-1.6875_{10} = 10.0101_2$
- ◆ Sign extension
  - Write +6 and -6 as twos complement
  - ◇ 0110 and 1010
  - Sign extend to 8-bit bytes
  - ◇ 00000110 and 11111010
- ◆ Can't infer a representation from a number
  - 11001 is 25 (unsigned)
  - 11001 is -9 (sign magnitude)
  - 11001 is -6 (ones complement)
  - 11001 is -7 (twos complement)

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## Twos-complement overflow

- ◆ Answers only correct mod  $2^b$ 
  - Summing two positive numbers can give negative result
  - Summing two negative numbers can give a positive result



- ◆ Make sure to have enough bits to handle overflow

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## BCD (Binary-Coded Decimal) and Gray codes

Decimal Symbols	BCD Code
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

Decimal Symbols	Gray Code
0	0000
1	0001
2	0011
3	0010
4	0110
5	0111
6	0101
7	0100
8	1100
9	1101

Only one bit changes per step

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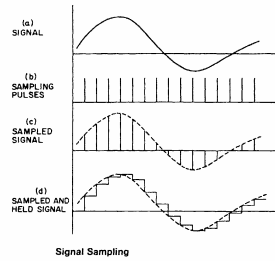
## The physical world is analog

- ◆ Digital systems need to
  - Measure analog quantities
    - ◇ Speech waveforms, etc
  - Control analog systems
    - ◇ Drive motors, etc
- ◆ How do we connect the analog and digital domains?
  - Analog-to-digital converter (A/D)
    - ◇ Example: CD recording
  - Digital-to-analog converter (D/A)
    - ◇ Example: CD playback

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## Sampling

- ◆ **Quantization**
  - Conversion from analog to discrete values
- ◆ Quantizing a signal
  - We sample it



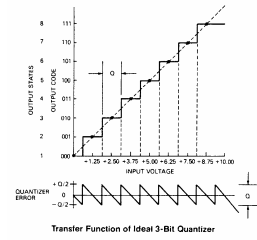
Signal Sampling

Datel Data Acquisition and Conversion Handbook

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## Conversion

- ◆ **Encoding**
  - Assigning a digital word to each discrete value
- ◆ Encoding a quantized signal
  - Encode the samples
  - Typically Gray or binary codes



Transfer Function of Ideal 3-Bit Quantizer

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