#### Overview

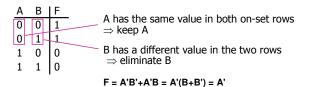
- Last lecture
  - deMorgan's theorem
  - NAND and NOR
  - Canonical forms

    - ✔ Product-of-sums (maxterms)
- ◆ Today's lecture
  - Logic simplification
    - **∠** Boolean cubes
    - ★ Karnaugh maps

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### Logic-function simplification

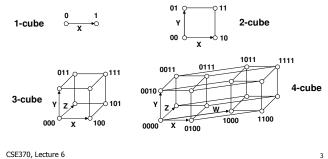
- Key tool: The uniting theorem  $\rightarrow A(B'+B) = A$
- The approach:
  - Find subsets of the ON-set where some variables don't change (the A's above) and others do (the B's above)
  - Eliminate the changing variables (the B's)



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#### Boolean cubes

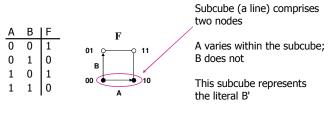
- ◆ *Visualize* when we can apply the uniting theorem
  - n input variables = n-dimensional "cube"



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## Mapping truth tables onto Boolean cubes

- ◆ ON set = solid nodes
- ◆ OFF set = empty nodes

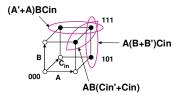


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#### Logic minimization using Boolean cubes

- Uniting theorem = find reduced-dimensionality subcubes
- ◆ Example: Binary full-adder carry-out logic
  - On-set is covered by the OR of three 2-D subcubes

В	Cin	Cout
0	0	0
0	1	0
1	0	0
1	1	1
0	0	0
0	1	1
1	0	1
1	1	1
	0 0 1 1 0 0	0 0 0 1 1 0 1 1 0 0 0 1 1 1



Cout = BCin+AB+ACin

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# Karnaugh maps

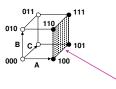
- ◆ Flat representation of Boolean cubes
  - Easy to use for 2– 4 dimensions
  - Hard for 4 6 dimensions
  - Virtually impossible for 6+ dimensions ✓ Use CAD tools
- Help visualize adjacencies
  - On-set elements that have one variable changing are adjacent
    Unlike a truth-table
  - Visual way to apply the uniting theorem

	Α	В	F
0	0	0	1
1	0	1	0
2	1	0	1
3	1	1	0

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#### M-dimensional cubes in n-dimensional space

- ◆ In a 3-cube (three variables):
  - A 0-cube (a single node) yields a term in 3 literals
  - A 1-cube (a line of two nodes) yields a term in 2 literals
  - A 2-cube (a plane of four nodes) yields a term in 1 literal
  - A 3-cube (a cube of eight nodes) yields a constant term "1"



 $F(A,B,C) = \sum m(4,5,6,7)$ 

On-set forms a square (a 2-D cube)

A is asserted (true) and unchanging B and C vary

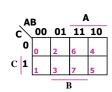
This subcube represents the literal A

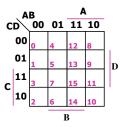
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## K-map cell numbering

- ◆ Gray-code: Only one bit changes between cells
  - Example:  $00 \rightarrow 01 \rightarrow 11 \rightarrow 10$
- ◆ Layout for 2 4 dimension K-maps:







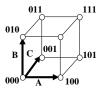
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## Adjacencies

- Wrap—around at edges
  - First column to last column
  - Top row to bottom row

AB			Α	
c/	00	01	11	10
0←	000-	→010 2	110 6	100 4
C 1	001	011	111 7	101 5
R				



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# K-map minimization (con't)

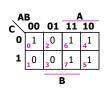
◆ Obtain the complement by covering 0s with subcubes

$$\begin{matrix} \textbf{AB} & \textbf{OO} & \textbf{OI} & \frac{\textbf{A}}{\textbf{11} \ \textbf{10}} \\ \textbf{O} & 0 & 0 & 0 & 0 \\ \textbf{O} & 0 & 0 & 0 & 0 \\ \textbf{I} & 0 & 0 & 0 & 0 \\ \textbf{I} & 0 & 0 & 0 & 0 \\ \hline \textbf{B} \end{matrix}$$

$$F(A,B,C) = \Sigma m(0,4,5,7)$$
  
= B'C'+AC

$$F'(A,B,C) = \Sigma m(1,2,3,6)$$
  
= A'C + BC'

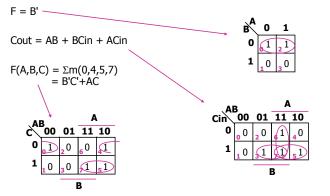
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$$F(A,B,C) = ???$$

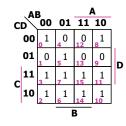
$$F'(A,B,C) = ???$$

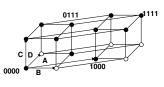
K-map minimization: 2 and 3 variables



#### K-map minimization: 4 variables

- Minimize  $F(A,B,C,D) = \Sigma m(0,2,3,5,6,7,8,10,11,14,15)$ 
  - Find the least number of subcubes, each as large as possible, that cover the ON-set





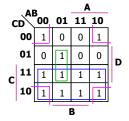
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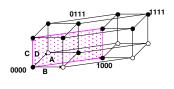
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## Karnaugh map: 4-variable example (con't)

- Minimize  $F(A,B,C,D) = \Sigma m(0,2,3,5,6,7,8,10,11,14,15)$
- ◆ Answer: F = C+A'BD+B'D'





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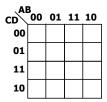
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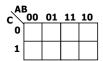
## K-map class examples

 $F(A,B,C,D) = \Sigma m(0,3,7,8,11,15)$ 

 $F(A,B,C) = \Sigma m(0,3,6,7)$ 

F(A,B,C) = ??? F'(A,B,C) = ???





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