

Boolean algebra

- ◆ Last lecture
 - Binary numbers
 - Base conversion
 - Number systems
 - ↳ Twos-complement
 - A/D and D/A conversion
- ◆ Today's lecture
 - Boolean algebra
 - ↳ Axioms
 - ↳ Useful laws and theorems
 - ↳ Simplifying Boolean expressions

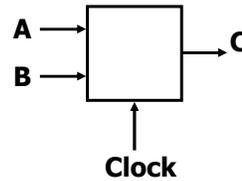
Major topic: Combinational logic

- ◆ Axioms and theorems of Boolean algebra
- ◆ Logic functions and truth tables
 - AND, OR, Buffer, NAND, NOR, NOT, XOR, XNOR
- ◆ Gate logic
 - Networks of Boolean functions
- ◆ Canonical forms
 - Sum of products and product of sums
- ◆ Simplification
 - Boolean cubes and Karnaugh maps
 - Two-level simplification

Combinational versus sequential

◆ Combinational: Memoryless

- Apply fixed inputs A, B
- Wait for clock edge
- Observe C
- Wait for another clock edge
- Observe C again: C will stay the same



◆ Sequential: With Memory

- Apply fixed inputs A, B
- Wait for clock edge
- Observe C
- Wait for another clock edge
- Observe C again: C may be different

Boolean algebra

◆ A Boolean algebra comprises...

- A set of elements B
- Binary operators $\{+, \bullet\}$
- A unary operation $\{ '\}$

◆ ...and the following axioms

- 1. The set B contains at least two elements $\{a, b\}$ with $a \neq b$
- 2. Closure: $a+b$ is in B $a \bullet b$ is in B
- 3. Commutative: $a+b = b+a$ $a \bullet b = b \bullet a$
- 4. Associative: $a+(b+c) = (a+b)+c$ $a \bullet (b \bullet c) = (a \bullet b) \bullet c$
- 5. Identity: $a+0 = a$ $a \bullet 1 = a$
- 6. Distributive: $a+(b \bullet c) = (a+b) \bullet (a+c)$ $a \bullet (b+c) = (a \bullet b) + (a \bullet c)$
- 7. Complementarity: $a+a' = 1$ $a \bullet a' = 0$

Digital (binary) logic is a Boolean algebra

◆ Substitute

- $\{0, 1\}$ for B
- AND for \bullet Boolean Product
- OR for $+$ Boolean Sum
- NOT for $'$

◆ All the axioms hold for binary logic

◆ Definitions

- Boolean function
 - ↳ Maps inputs from the set $\{0,1\}$ to the set $\{0,1\}$
- Boolean expression
 - ↳ An algebraic statement of Boolean variables and operators

AND, OR, Not

◆ AND

$X \bullet Y$

XY



X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1

◆ OR

$X + Y$

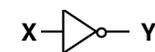


X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1

◆ NOT

\bar{X}

X'



X	Y
0	1
1	0

Logic functions and Boolean algebra

- ◆ Any logic function that is expressible as a truth table can be written in Boolean algebra using +, •, and '.

X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1

$Z = X \cdot Y$

X	Y	X'	Z
0	0	1	0
0	1	1	1
1	0	0	0
1	1	0	0

$Z = X' \cdot Y$

X	Y	X'	Y'	$X \cdot Y$	$X' \cdot Y'$	Z
0	0	1	1	0	1	1
0	1	1	0	0	0	0
1	0	0	1	0	0	0
1	1	0	0	1	0	1

$Z = (X \cdot Y) + (X' \cdot Y')$

Two key concepts

- ◆ Duality (a meta-theorem— *a theorem about theorems*)
 - All Boolean expressions have logical duals
 - Any theorem that can be proved is also proved for its dual
 - Replace: • with +, + with •, 0 with 1, and 1 with 0
 - Leave the variables unchanged
- ◆ de Morgan's Theorem
 - Procedure for complementing Boolean functions
 - Replace: • with +, + with •, 0 with 1, and 1 with 0
 - Replace all variables with their complements

Useful laws and theorems

Identity:	$X + 0 = X$	Dual: $X \cdot 1 = X$
Null:	$X + 1 = 1$	Dual: $X \cdot 0 = 0$
Idempotent:	$X + X = X$	Dual: $X \cdot X = X$
Involution:	$(X')' = X$	
Complementarity:	$X + X' = 1$	Dual: $X \cdot X' = 0$
Commutative:	$X + Y = Y + X$	Dual: $X \cdot Y = Y \cdot X$
Associative:	$(X+Y)+Z=X+(Y+Z)$	Dual: $(X \cdot Y) \cdot Z = X \cdot (Y \cdot Z)$
Distributive:	$X \cdot (Y+Z) = (X \cdot Y) + (X \cdot Z)$	Dual: $X + (Y \cdot Z) = (X+Y) \cdot (X+Z)$
Uniting:	$X \cdot Y + X \cdot Y' = X$	Dual: $(X+Y) \cdot (X+Y') = X$

Useful laws and theorems (con't)

Absorption:	$X + X \cdot Y = X$	Dual: $X \cdot (X + Y) = X$
Absorption (#2):	$(X + Y') \cdot Y = X \cdot Y$	Dual: $(X \cdot Y') + Y = X + Y$
de Morgan's:	$(X + Y + \dots)' = X' \cdot Y' \cdot \dots$	Dual: $(X \cdot Y \cdot \dots)' = X' + Y' + \dots$
Duality:	$(X + Y + \dots)^D = X \cdot Y \cdot \dots$	Dual: $(X \cdot Y \cdot \dots)^D = X + Y + \dots$
Multiplying & factoring:	$(X + Y) \cdot (X' + Z) = X \cdot Z + X' \cdot Y$	Dual: $X \cdot Y + X' \cdot Z = (X + Z) \cdot (X' + Y)$
Consensus:	$(X \cdot Y) + (Y \cdot Z) + (X' \cdot Z) = X \cdot Y + X' \cdot Z$	Dual: $(X + Y) \cdot (Y + Z) \cdot (X' + Z) = (X + Y) \cdot (X' + Z)$

Proving theorems

- ◆ Example 1: Prove the uniting theorem-- $X \cdot Y + X \cdot Y' = X$

Distributive	$X \cdot Y + X \cdot Y' = X \cdot (Y + Y')$
Complementarity	$= X \cdot (1)$
Identity	$= X$

- ◆ Example 2: Prove the absorption theorem-- $X + X \cdot Y = X$

Identity	$X + X \cdot Y = (X \cdot 1) + (X \cdot Y)$
Distributive	$= X \cdot (1 + Y)$
Null	$= X \cdot (1)$
Identity	$= X$

Proving theorems

- ◆ Example 3: Prove the consensus theorem--
 $(XY) + (YZ) + (X'Z) = XY + X'Z$

Complementarity	$XY + YZ + X'Z = XY + (X + X')YZ + X'Z$
Distributive	$= XYZ + XY + X'YZ + X'Z$

↳ Use absorption $\{AB + A = A\}$ with $A = XY$ and $B = Z$

$$= XY + X'YZ + X'Z$$

Rearrange terms	$= XY + X'ZY + X'Z$
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↳ Use absorption $\{AB + A = A\}$ with $A = X'Z$ and $B = Y$

$$XY + YZ + X'Z = XY + X'Z$$

de Morgan's Theorem

- ◆ Use de Morgan's Theorem to find complements
- ◆ Example: $F=(A+B) \cdot (A'+C)$, so $F'=(A' \cdot B')+(A \cdot C')$

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

A	B	C	F'
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

Logic simplification

- ◆ Use the axioms to simplify logical expressions
 - Why? To use less hardware
- ◆ Example: A two-level logic expression

$$\begin{aligned}
 Z &= A'BC + AB'C' + AB'C + ABC' + ABC \\
 &= AB'C + AB'C' + A'BC + ABC' + ABC && \text{rearrange} \\
 &= AB'(C + C') + A'BC + AB(C' + C) && \text{distributive} \\
 &= AB' + A'BC + AB && \text{comp.} \\
 &= AB' + AB + A'BC && \text{rearrange} \\
 &= A(B' + B) + A'BC && \text{distributive} \\
 &= A + A'BC && \text{comp.}
 \end{aligned}$$

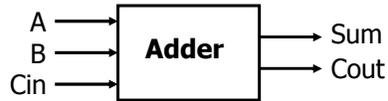
✎ Use absorption #2D $\{(X \cdot Y') + Y = X + Y\}$ with $X=BC$ and $Y=A$

$$Z = A + BC$$

Example: A full adder

◆ 1-bit binary adder

- Inputs: A, B, Carry-in
- Outputs: Sum, Carry-out



A	B	Cin	S	Cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

$$S = A'B'Cin + A'BCin' + AB'Cin' + ABCin$$

$$Cout = A'BCin + AB'Cin + ABCin' + ABCin$$

Simplifying the carry-out function

$$\begin{aligned}
 Cout &= A'BCin + AB'Cin + ABCin' + ABCin \\
 &= A'BCin + AB'Cin + ABCin' + ABCin + ABCin \\
 &= A'BCin + ABCin + AB'Cin + ABCin' + ABCin \\
 &= (A'+A)BCin + AB'Cin + ABCin' + ABCin \\
 &= (1)BCin + AB'Cin + ABCin' + ABCin \\
 &= BCin + AB'Cin + ABCin' + ABCin + ABCin \\
 &= BCin + AB'Cin + ABCin + ABCin' + ABCin \\
 &= BCin + A(B'+B)Cin + ABCin' + ABCin \\
 &= BCin + A(1)Cin + ABCin' + ABCin \\
 &= BCin + ACin + AB(Cin'+Cin) \\
 &= BCin + ACin + AB(1) \\
 &= BCin + ACin + AB
 \end{aligned}$$

associative

idempotent

Some notation

- ◆ Priorities: $\bar{A} \cdot B + C = ((\bar{A}) \cdot B) + C$
- ◆ Variables are sometimes called literals