## Lecture 6

## - Logistics

- HW2 out, due 4/16 Wednesday
- Lab 2 ongoing
- Last lecture
- Canonical forms
- NAND and NOR
- Today's lecture
- One more pushing bubble example
- Logic simplification

K Boolean cubes
$\boldsymbol{K}$ Karnaugh maps


## Example of bubble pushing: NOR/NOR



## Goal: Minimize two-level logic expression

- Algebraic simplification
- not an systematic procedure
- hard to know when we reached the minimum
- Computer-aided design tools
- require very long computation times (NP hard)

■ heuristic methods employed - "educated guesses"

- Visualization methods are useful
- our brain is good at figuring things out over computers
- many real-world problems are solvable by hand


## Key tool: The Uniting Theorem

- The uniting theorem $\rightarrow A\left(B^{\prime}+B\right)=A$
- The approach:
- Find some variables don't change (the A's above) and others do (the B's above)
- Eliminate the changing variables (the B's)

| A B | F | A has the same value in both "on-set" rows $\Rightarrow$ keep A |
| :---: | :---: | :---: |
| $0{ }^{0} 0$ | 1 |  |
| $0{ }^{1}$ | 1 |  |
| 10 | 0 | $B$ has a different value in the two rows $\Rightarrow$ eliminate B |
| 1 | 0 |  |
|  |  | $F=A^{\prime} B^{\prime}+A^{\prime} B=A^{\prime}\left(B+B^{\prime}\right)=A^{\prime}$ |

## Boolean cubes

- Visualization tool for the uniting theorem
- $n$ input variables $=n$-dimensional "cube"



## Mapping truth tables onto Boolean cubes

- ON set = solid nodes
- OFF set = empty nodes



## Example using Boolean cube

- Binary full-adder carry-out logic
- On-set is covered by the OR of three 2-D subcubes

| A | B | Cin | Cout |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |



Cout $=B C i n+A B+A C i n$

## M-dimensional cubes in n-dimensional space

- In a 3-cube (three variables):
- A 0 -cube (a single node) yields a term in 3 literals
- A 1-cube (a line of two nodes) yields a term in 2 literals
- A 2-cube (a plane of four nodes) yields a term in 1 literal
- A 3-cube (a cube of eight nodes) yields a constant term "1"



## Karnaugh maps (K-map)

- Flat representation of Boolean cubes
- Easy to use for 2- 4 dimensions
- Hard for 4-6 dimensions
- Virtually impossible for 6+ dimensions $\boldsymbol{K}$ Use CAD tools
- Help visualize adjacencies

|  | $A$ | $B$ | $F$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 2 | 1 | 0 | 1 |
| 3 | 1 | 1 | 0 |

- On-set elements that have one variable changing are adjacent

|  | ${ }^{\text {A }} 1$ |  |
| :---: | :---: | :---: |
| 0 | ${ }_{0} 1$ | 2 |
| 1 | 10 | ${ }_{3} 0$ |

## 2, 3, and 4 dimensional K-maps

- Uses Gray code: Only one bit changes between cells
- Example: $00 \rightarrow 01 \rightarrow 11 \rightarrow 10$




## Adjacencies

- Wrap-around at edges
- First column to last column
- Top row to bottom row



## K-map minimization example: 2 variables

|  | A | B | F | $B^{A} 0 \quad 1$ |  |  | $F=B^{\prime}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 |  |  |  |  |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 |  |
| 2 | 1 | 0 | 1 | 1 |  | 2 |  |
| 3 | 1 | 1 | 0 | 1 | 10 | 30 |  |

## K-map minimization example: 3 variables

| A | B | Cin | Cout |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |



K-map minimization example: minterms

$$
\begin{aligned}
F(A, B, C) & =\Sigma m(0,4,5,7) \\
& =B^{\prime} C^{\prime}+A C
\end{aligned}
$$

## K-map minimization example: complement



$$
\begin{aligned}
\mathrm{F}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}) & =\Sigma \mathrm{m}(0,4,5,7) \\
& =\mathrm{B}^{\prime} \mathrm{C}^{\prime}+\mathrm{AC} \\
\mathrm{~F}^{\prime}(\mathrm{A}, \mathrm{~B}, \mathrm{C}) & =\Sigma \mathrm{m}(1,2,3,6) \\
& =\mathrm{A}^{\prime} \mathrm{C}+\mathrm{BC}^{\prime}
\end{aligned}
$$

## K-map minimization example: 4 variables

- Minimize $F(A, B, C, D)=\Sigma m(0,2,3,5,6,7,8,10,11,14,15)$
- Find the least number of subcubes, each as large as possible, that cover the ON-set



## K-map minimization example: 4 variables

- Minimize $F(A, B, C, D)=\Sigma m(0,2,3,5,6,7,8,10,11,14,15)$
- Answer: F = C+A'BD+B'D'



## K-map minimization examples: do it yourself

$$
\begin{array}{ll}
\mathrm{F}(\mathrm{~A}, \mathrm{~B}, \mathrm{C})=\Sigma \mathrm{m}(0,3,6,7) & \mathrm{F}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}, \mathrm{D})=\Sigma \mathrm{m}(0,3,7,8,11,15) \\
\mathrm{F}(\mathrm{~A}, \mathrm{~B}, \mathrm{C})= & \mathrm{F}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}, \mathrm{D})= \\
\mathrm{F}^{\prime}(\mathrm{A}, \mathrm{~B}, \mathrm{C})= & \mathrm{F}^{\prime}(\mathrm{A}, \mathrm{~B}, \mathrm{C}, \mathrm{D})=
\end{array}
$$




