

Intro to Digital Design

L9: Memory, Computer Components, FPGAs

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Administrivia

❖ Lab 8 – Project

- Check-ins this week during your normal lab time
 - **Graded!** Turn in a block diagram and a module implementation (with testbench)!
- Final demo of finished project due before end of next week, reports no later than Friday, March 13 @ 11:59 pm
- Return your lab kit to TAs at your final demo

Administrivia

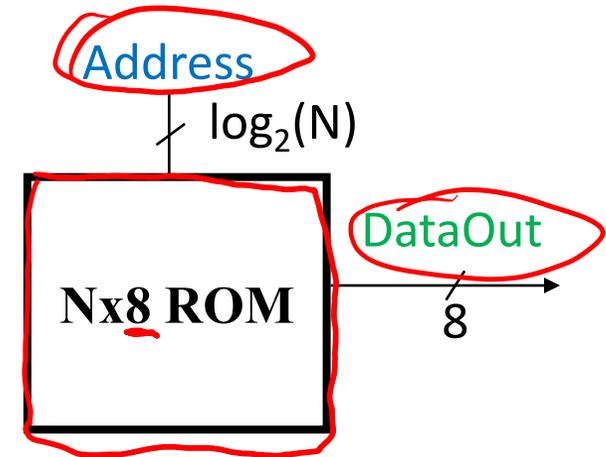
- ❖ **Quiz 3** is next week: Tuesday, March 10
 - We'll start class with a quick wrap-up and then give you time to get settled
 - 60 (+10) minutes, worth 14% of your course grade
 - Topics: Timing, Routing Elements, Computational Building Blocks, Verilog
 - Past Quiz 3 (+ solutions) on website: Course Info → Quizzes
 - **Note**: Your Quiz 3 will be a little different – focus on problem solving

Outline

- ❖ **Memory**
- ❖ CPU teaser
- ❖ Serial I/O
- ❖ FPGAs

Storage Element: Idealized ROM

- ❖ “Read Only Memory”
 - $N \times M$: An array of N M -Bit words
 - Address input selects the word driven on the output
 - Data “burned in” by the manufacturer (or your bitfile) and retained even if power is lost
- ❖ Difference between a ROM and your seven segment decoder is *density*
 - (and thus underlying circuit implementation)
- ❖ What can we do with a truly *read-only* memory?
 - Math look-up tables (trig functions)
 - FSM state transitions
 - Bitmap images and other simple patterns

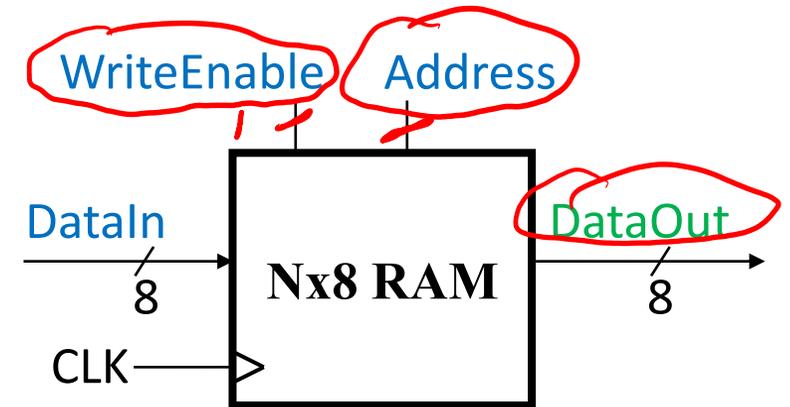


$$\text{DataOut} = \text{ROM}[\text{Address}]$$

Storage Element: Idealized RAM

❖ “Random Access Memory”

- Can now **write** as well as read
 - On clock edge, **DataIn** written to memory cell at **Address** if **WriteEnable** = 1
 - **DataOut** is still completely combinational (changes to **Address** *immediately* change the data output)
 - Data lost on reset or power-off
-
- ❖ Simplest and most robust implementation: “SRAM” – implemented as an array of flip flops with muxed inputs and outputs
 - ❖ Higher-density but more touchy: “DRAM” – implemented with circuit wizardry

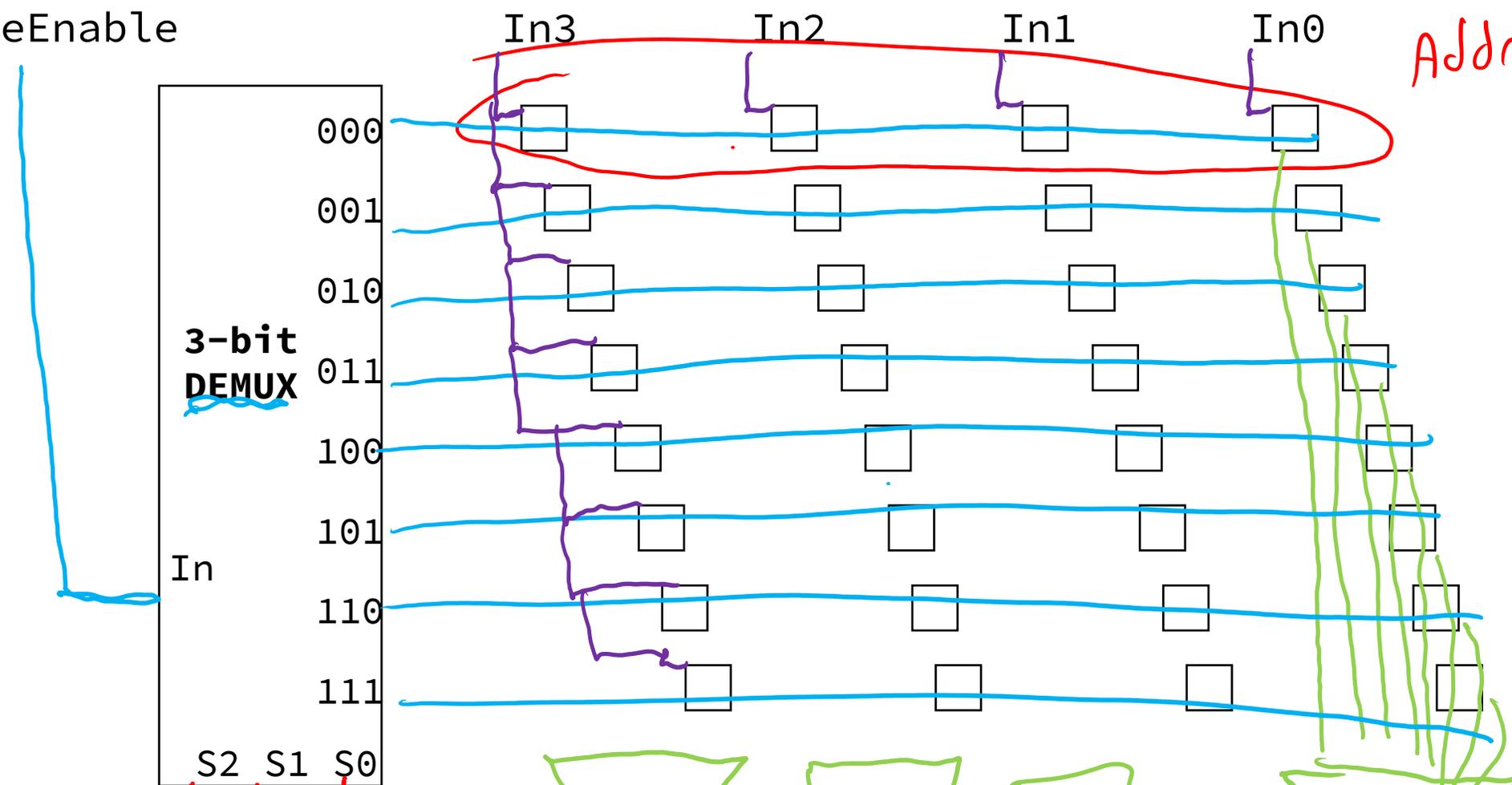


8 4-bit words

Blue = FF En.

8x4 SRAM

WriteEnable

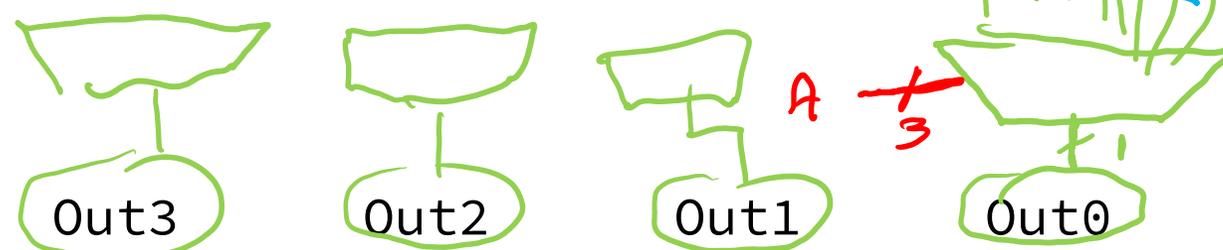


Addr 0

1
2
3
...

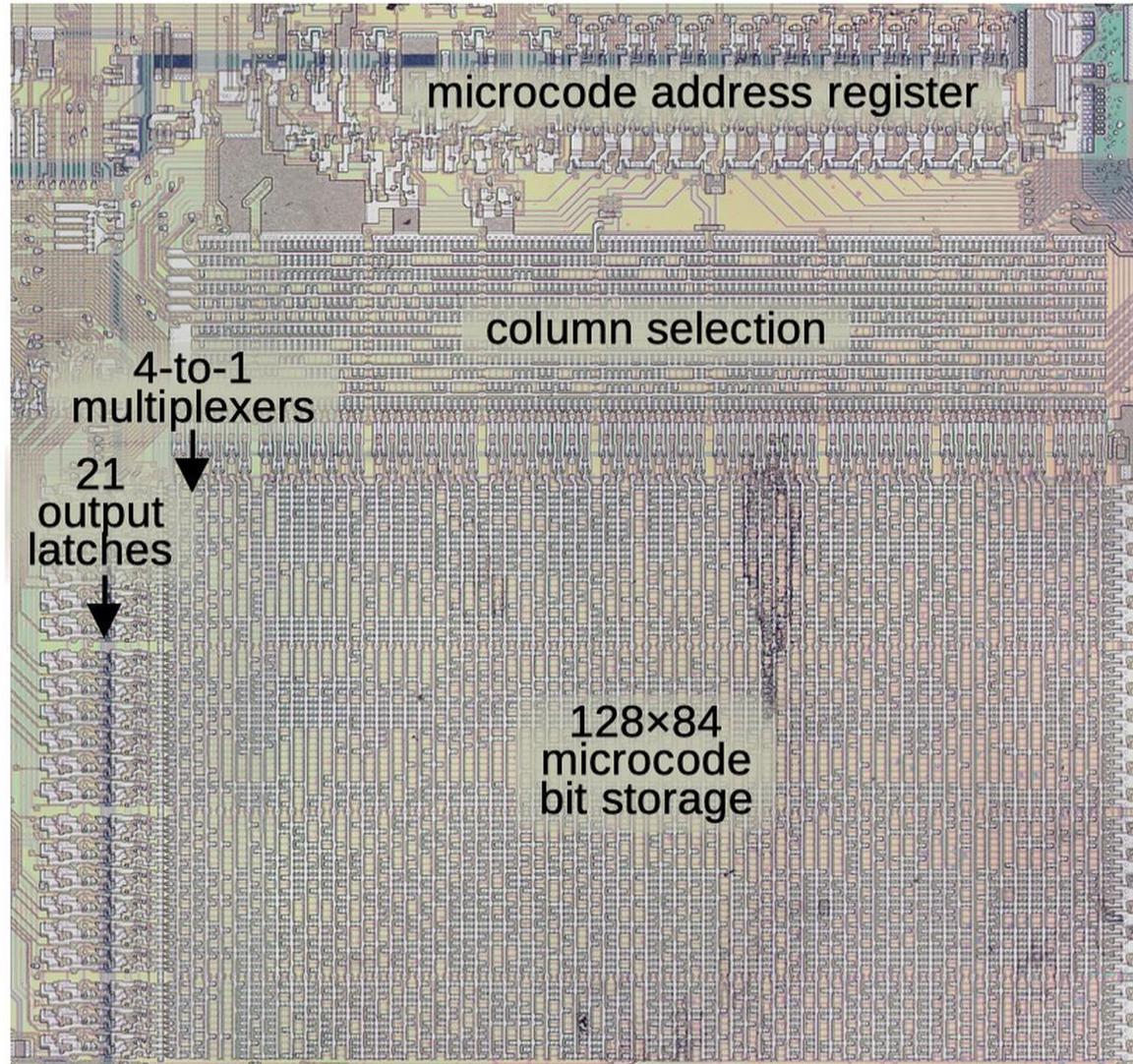
A2
A1
A0

S2 S1 S0



8:1 mux

Memory IRL: Microcode ROM in the original x86



Outline

- ❖ Memory
- ❖ **CPU teaser**
- ❖ Serial I/O
- ❖ FPGAs

What is a computer?

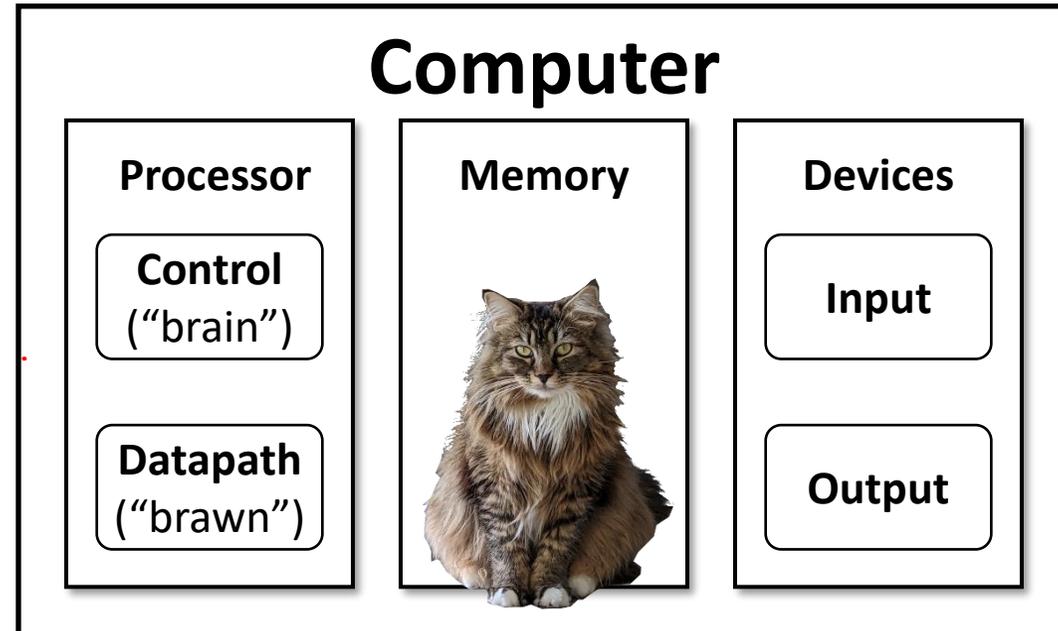
❖ The **central processing unit (CPU)**

- A little machine that takes a list of **instructions** and executes them **sequentially**, one after another
- Most instructions do basic math on local scratch memory called a **register file** (“ALU instructions”)
- Some instructions can set our next position in the instruction list (“**branches**” and “**jumps**”)
- Some instructions save register values to the larger off-chip **main memory** (“stores” and “loads”)

Five Components of a CPU

- ❖ The machine that actually executes those instructions is usually broken down like this:

-  Control
-  Datapath
-  Memory
-  Input
-  Output



Executing an Instruction

*addq (%rdi), %rax
0x48 03 07*

❖ Depends on ISA, but generally:

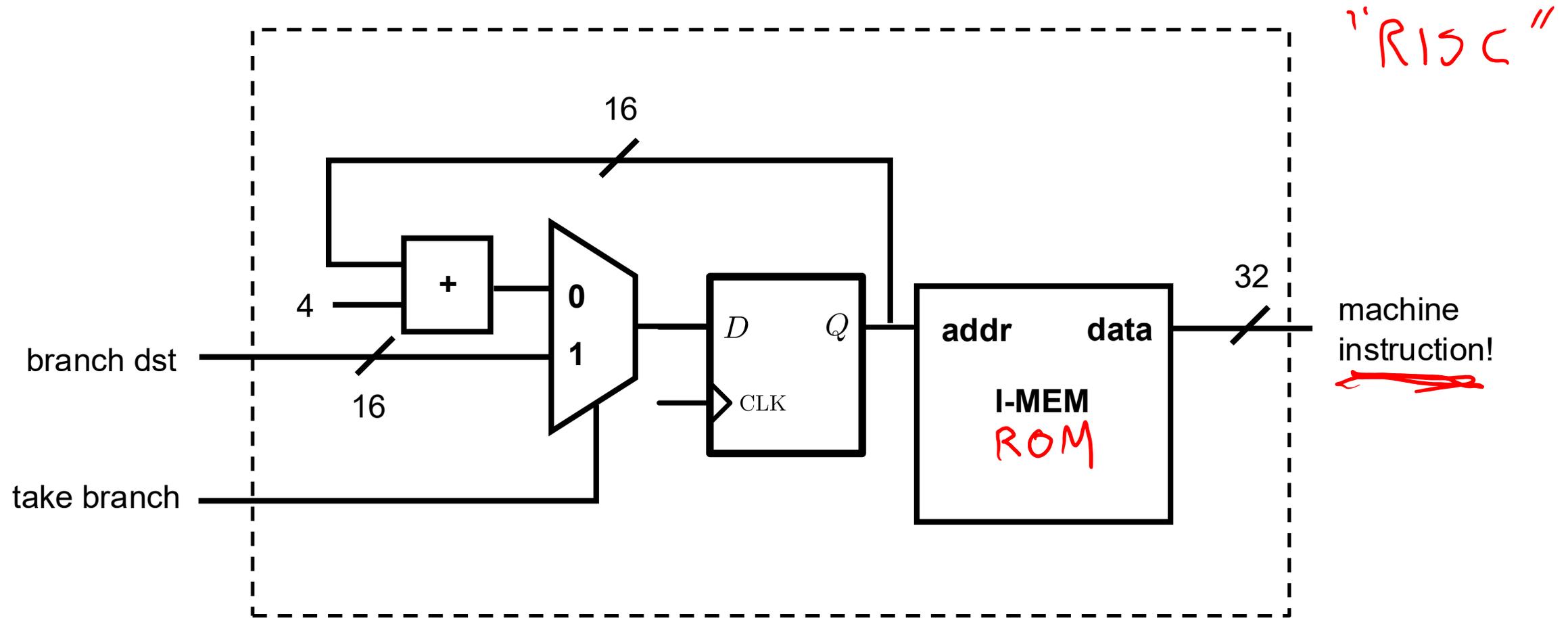
- Instruction Fetch
- Instruction Decode
- Data Fetch
- Computation → ALU
- Store Result



❖ Basic Datapath Components (idealized)

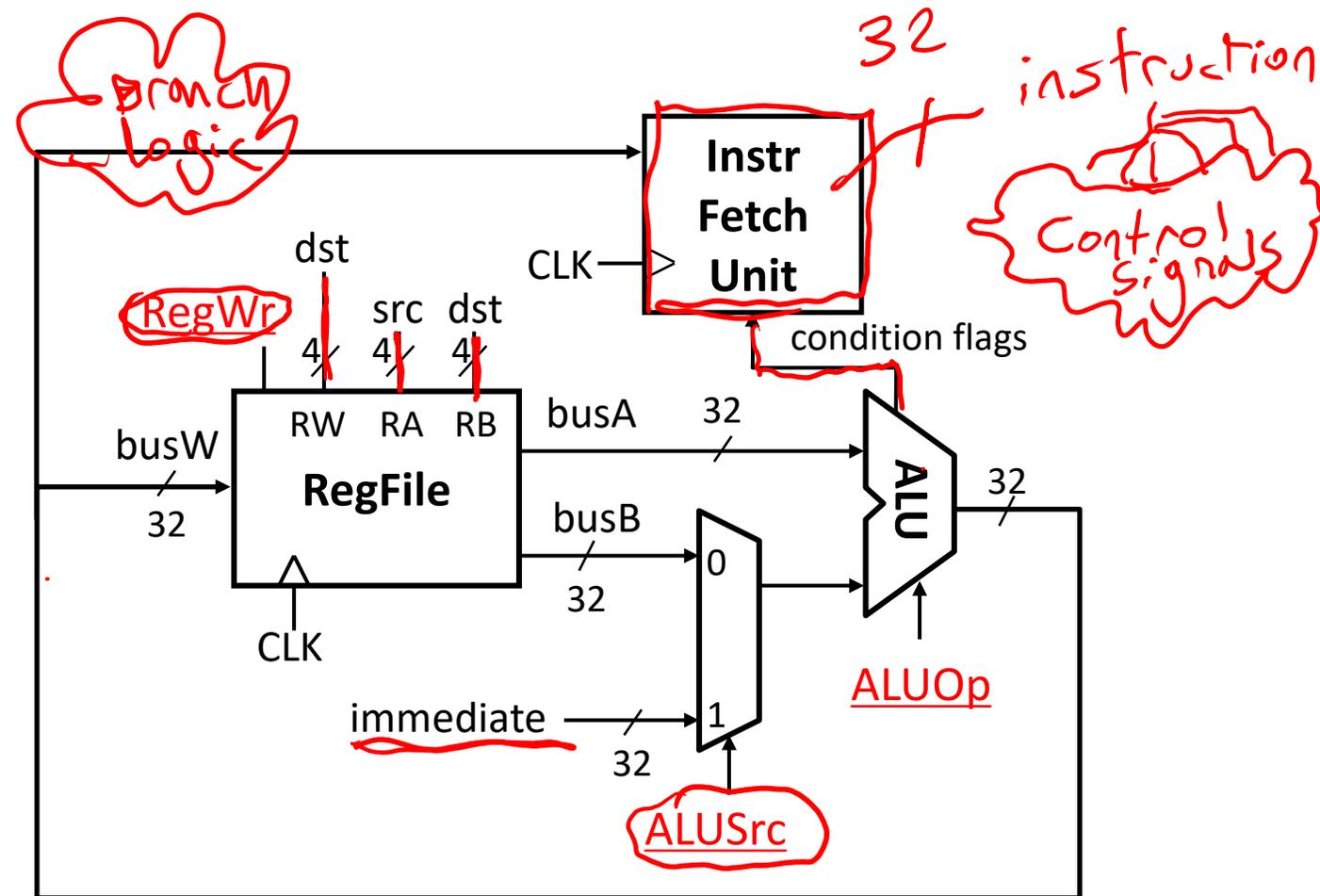
- Register File
 - Memory Management Unit
 - Arithmetic Logic Unit (ALU)
 - Routing Elements
- } ROMs and RAMs
- } Muxes, adders, encoders, etc

Instruction Fetch Unit (16 bit memory, 32 bit instructions)



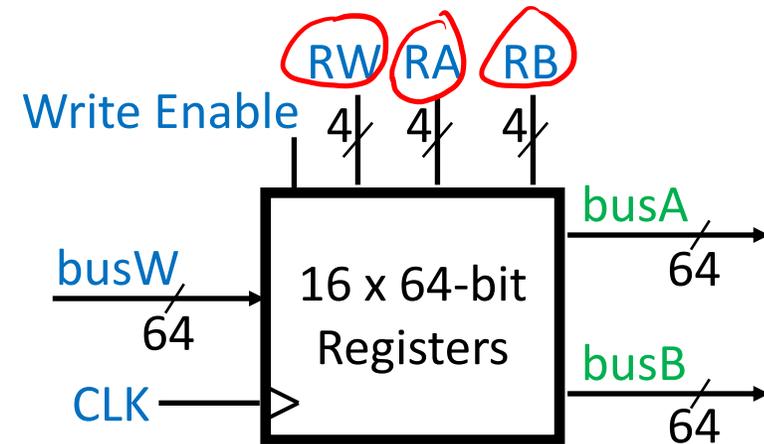
Datapath Teaser

- ❖ Super-simplified and incomplete example of a CPU datapath
 - Assumes instructions of the form: operation src, dst
- ❖ No main memory in this example, just registers
- ❖ Signals in **red** are set by Control based on the instruction's bits

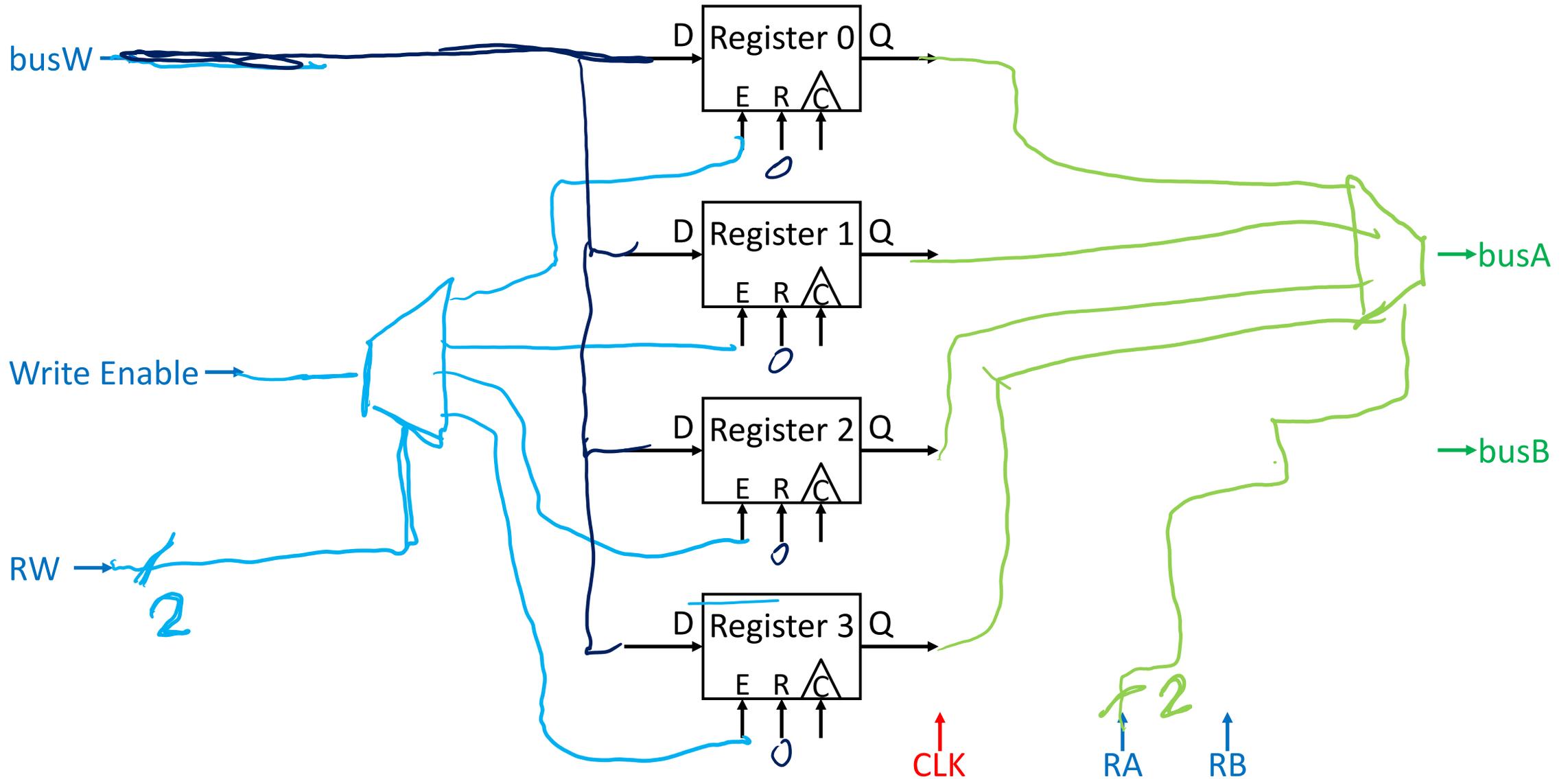


Zoom in on the Register File

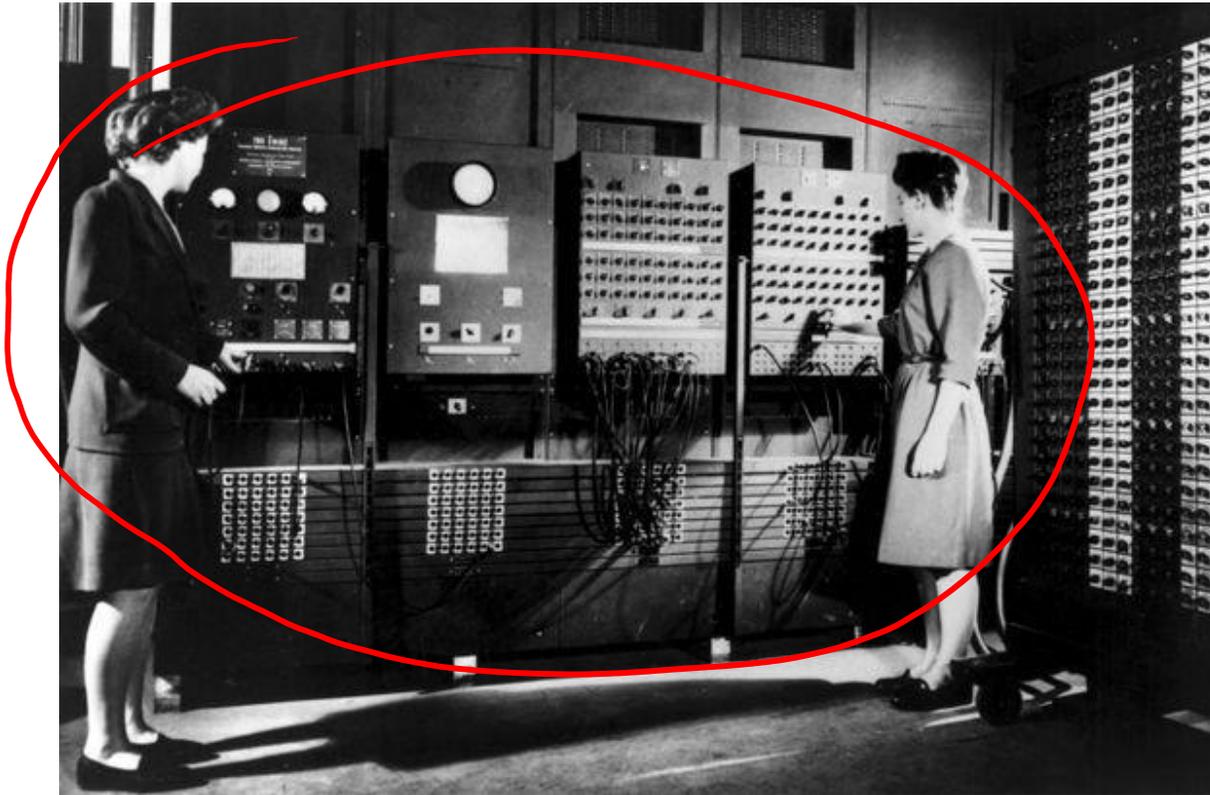
- ❖ Just a regular run o' the mill SRAM
- ❖ Contains all programmer-accessible registers
 - Output buses **busA** and **busB**
 - Input bus **busW**
- ❖ Register selection
 - Place data of registers **RA/RB** (numbers) onto **busA/busB**
 - Store data on **busW** into register **RW** (number) when **Write Enable** is 1



Simple Register File (4 Registers)

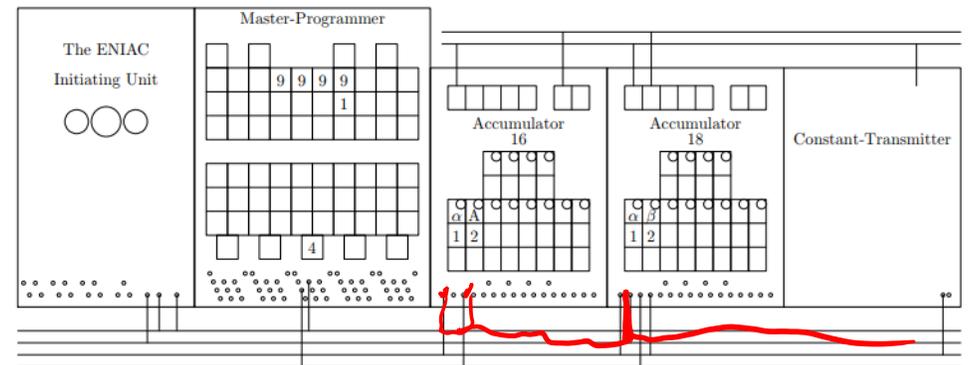
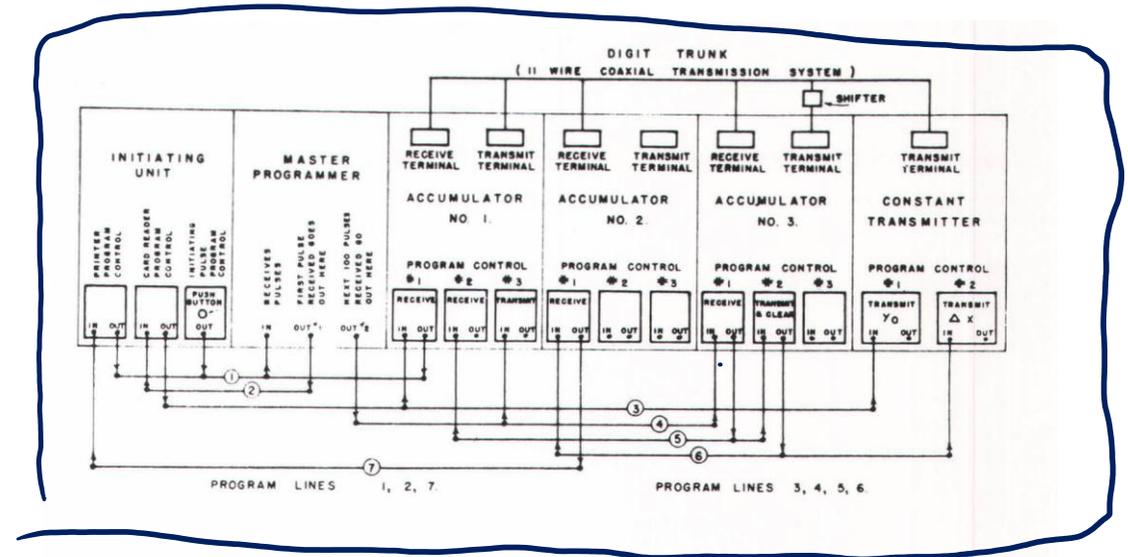


Early Language



Above: Frances Bilas and Betty Jean Jennings in front of the ENIAC

Right: ENIAC program sheets



Miso Moment



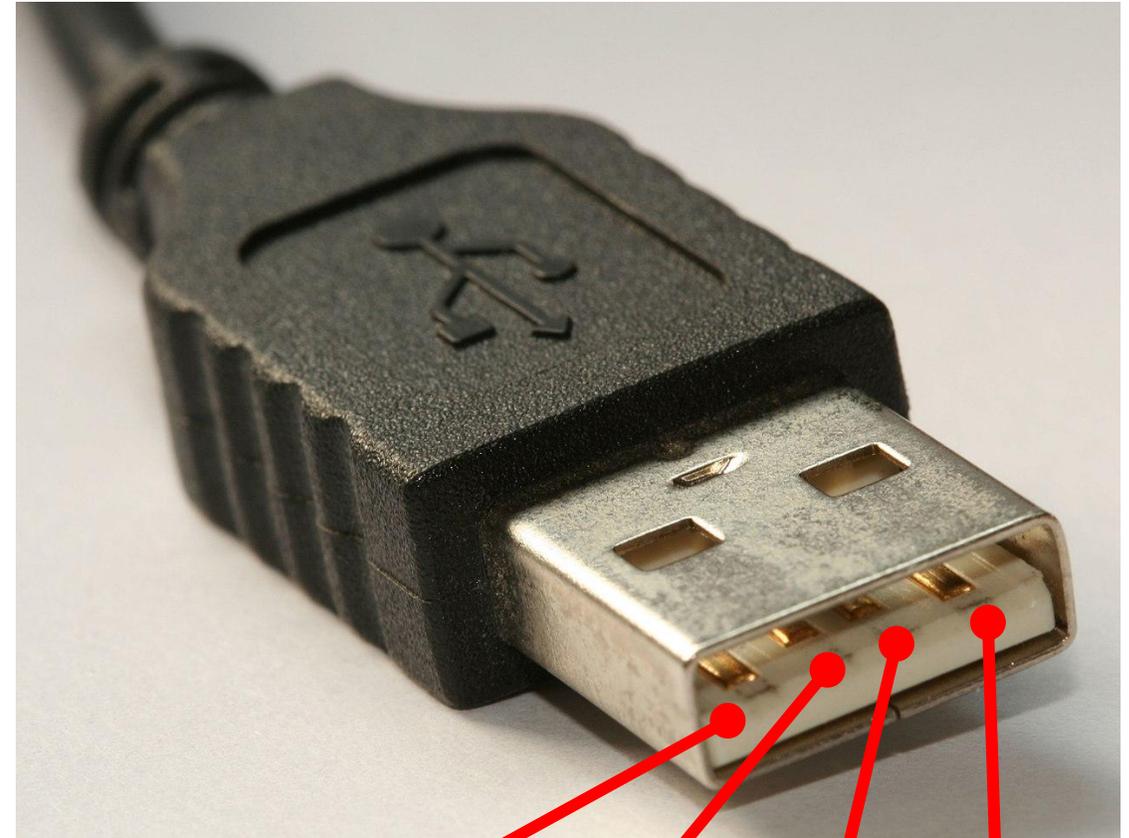
Outline

- ❖ Memory
- ❖ CPU teaser
- ❖ **Serial I/O**
- ❖ FPGAs

https://commons.wikimedia.org/wiki/Category:USB_2_Standard-A_plugs#/media/File:Type_A_USB_Connector_alt.jpg

Data Transfer

- ❖ How do we get stuff **in and out** of a chip?
 - **Parallel:** all bits transferred at once
 - **Serial:** one bit transferred at a time
- ❖ USB: classic example of a **serial** I/O standard
 - Only one (“differential”) data line
- ❖ So how do we get 32-bit data onto the 1-bit USB wire?



Ground

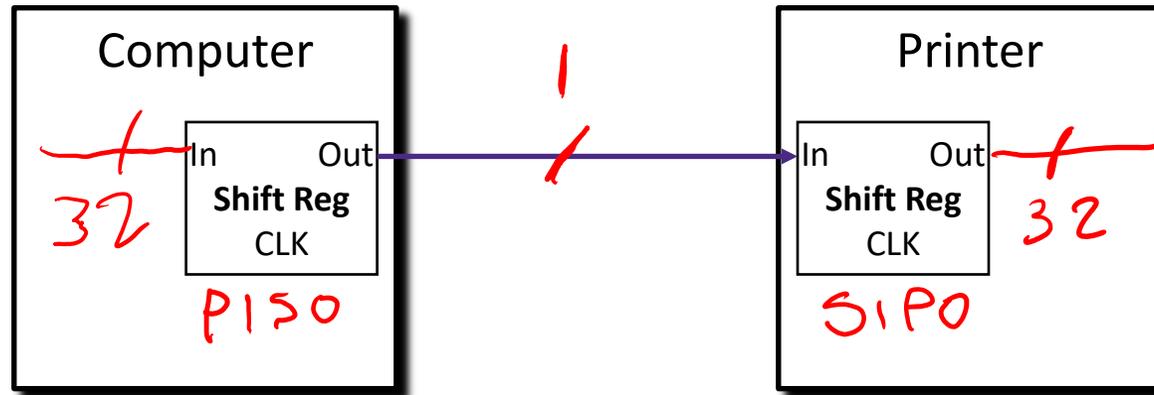
Data +

Data -

V_{bus}

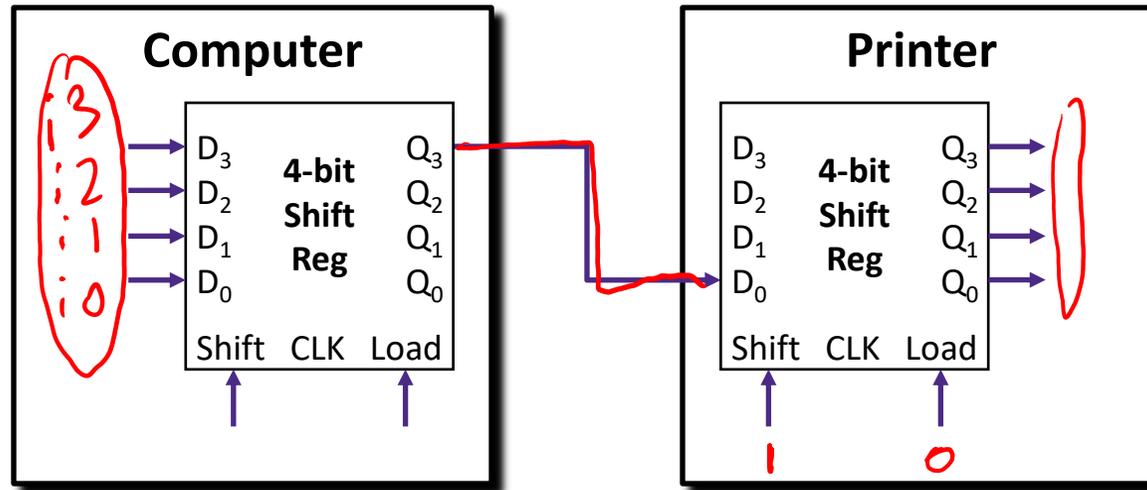
Transfer of Data

- ❖ Shift registers can be used for serial transfer:



- ❖ In general, shift registers can allow for **either** or **both** modes (S for serial, P for parallel) at input and output
 - Examples: SISO, SIPO, PISO

Conversion between Parallel & Serial



Cycle	Shift	Load	Q0	Q1	Q2	Q3
0	0	1	X	X	X	X
1	1	0	0	1	2	3
2	1	0	X	0	1	2
3	1	0	X	X	0	1
4	1	0	X	X	X	0
5	X	X	X	X	X	X
6	X	X				

Cycle	Shift	Load	Q0	Q1	Q2	Q3
0	1	0	X	X	X	X
1	1	0	X	X	X	X
2	1	0	3	X	X	X
3	1	0	2	3	X	X
4	1	0	1	2	3	X
5	1	0	0	1	2	3
6	1	0				

Shifter in Verilog

```
module leftshifter #(parameter WIDTH=8)
  (s_out, p_out, shift, load, d_in, clk);

  output logic s_out;           // serial out
  output logic [WIDTH-1:0] p_out; // parallel out
  input  logic [WIDTH-1:0] d_in; // data in (shift in d0)
  input  logic shift, load, clk;

  always_ff @(posedge clk) begin
    if (load)
      p_out <= d_in;
    else if (shift)
      p_out <= {p_out[WIDTH-2:0], d_in[0]};
  end

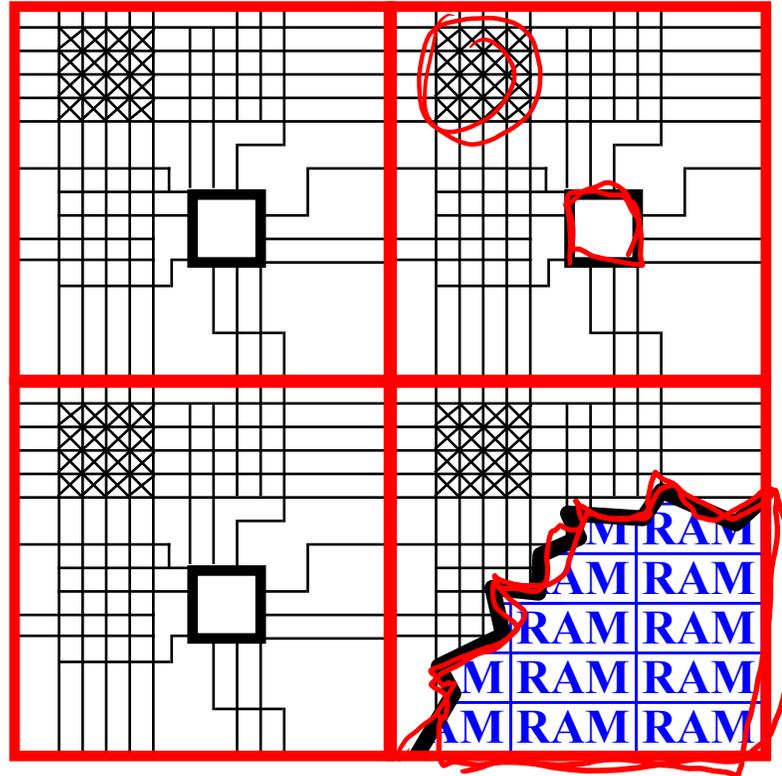
  assign s_out = p_out[WIDTH-1];

endmodule // leftshifter
```

Outline

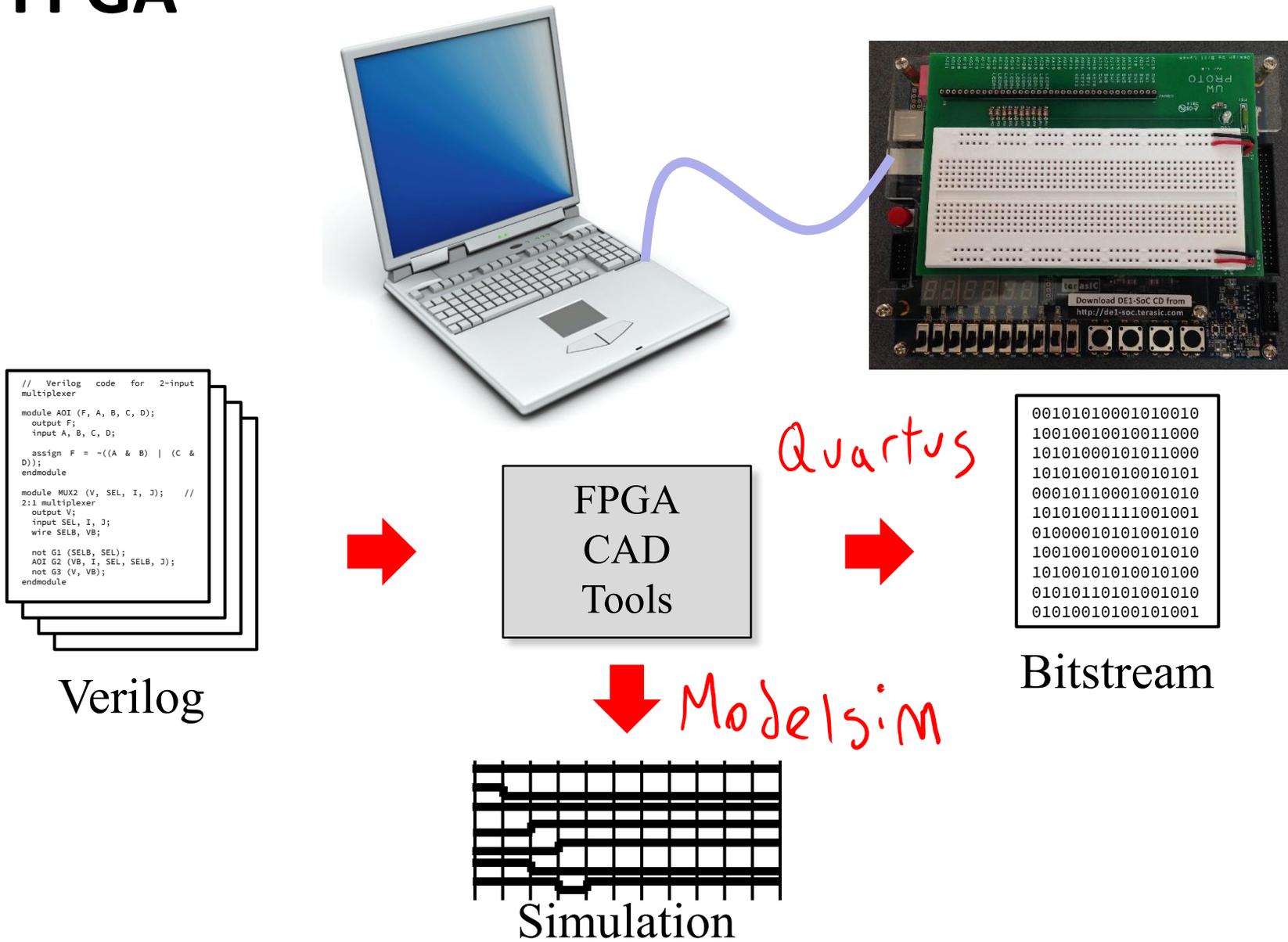
- ❖ Memory
- ❖ CPU teaser
- ❖ Serial I/O
- ❖ **FPGAs**

Field Programmable Gate Arrays (FPGAs)



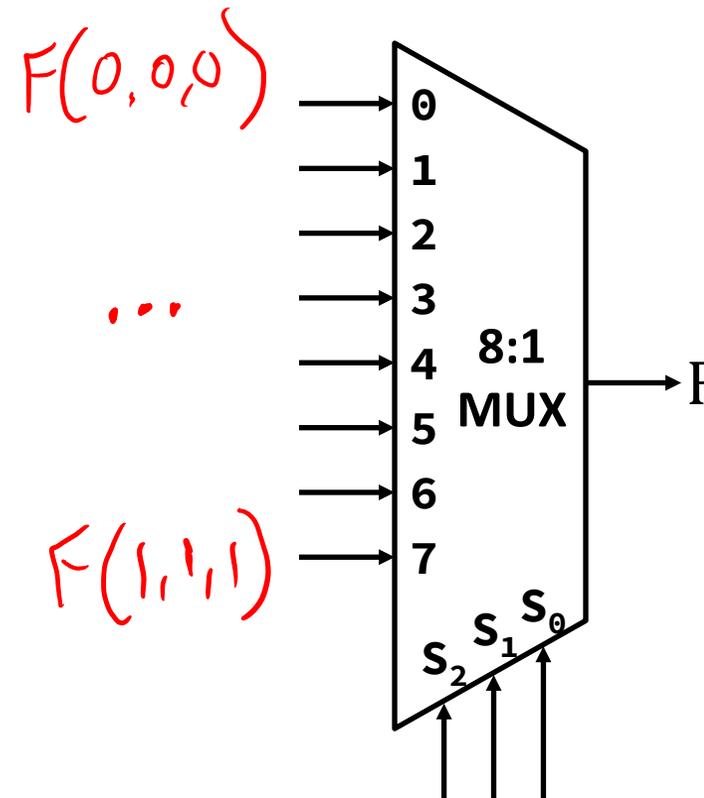
- ❖ Logic cells () embedded in a general routing structure
- ❖ Logic cells usually contain:
 - 6-input Boolean function calculator
 - Flip-flop (1-bit memory)
- ❖ All features are electronically (re)programmable

Using an FPGA

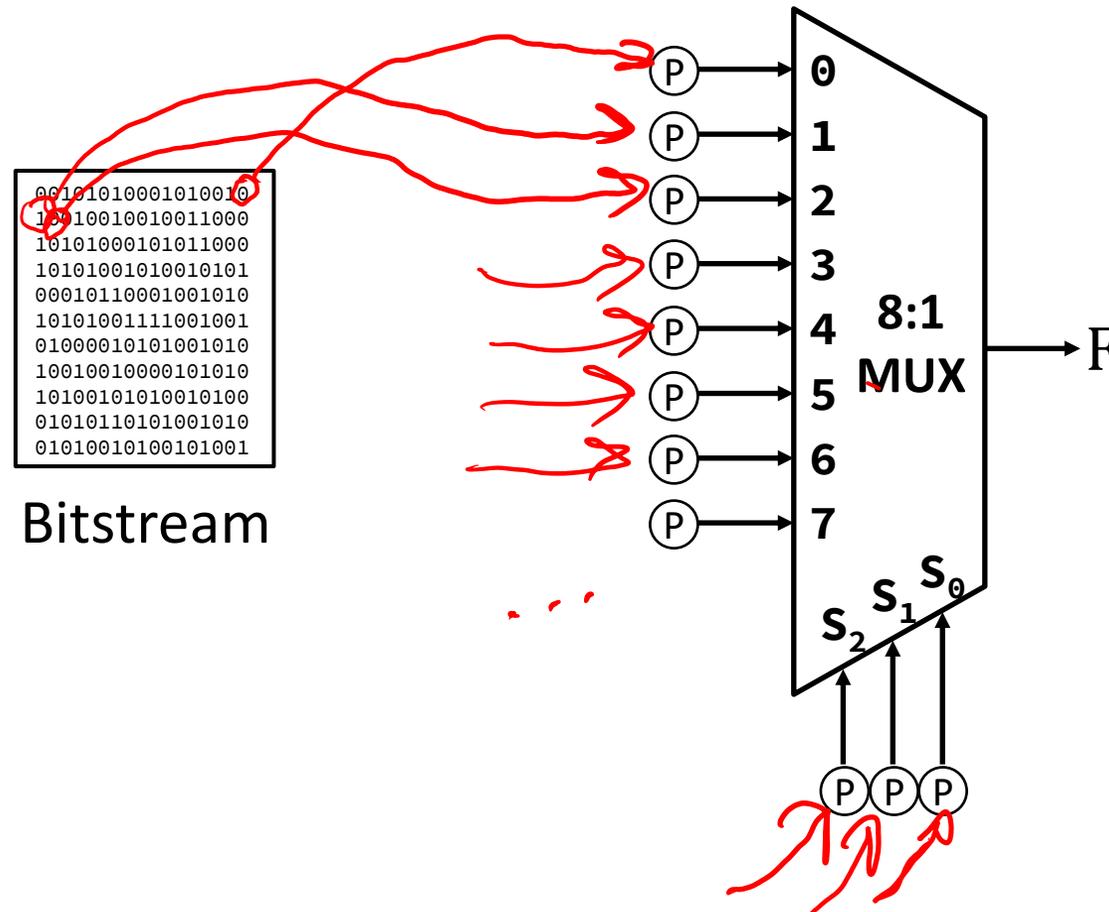


FPGA Combinational Logic

- ❖ Create arbitrary combinational binary function $F(A, B, C)$ using MUXes
 - Creates a **Lookup Table (LUT)**



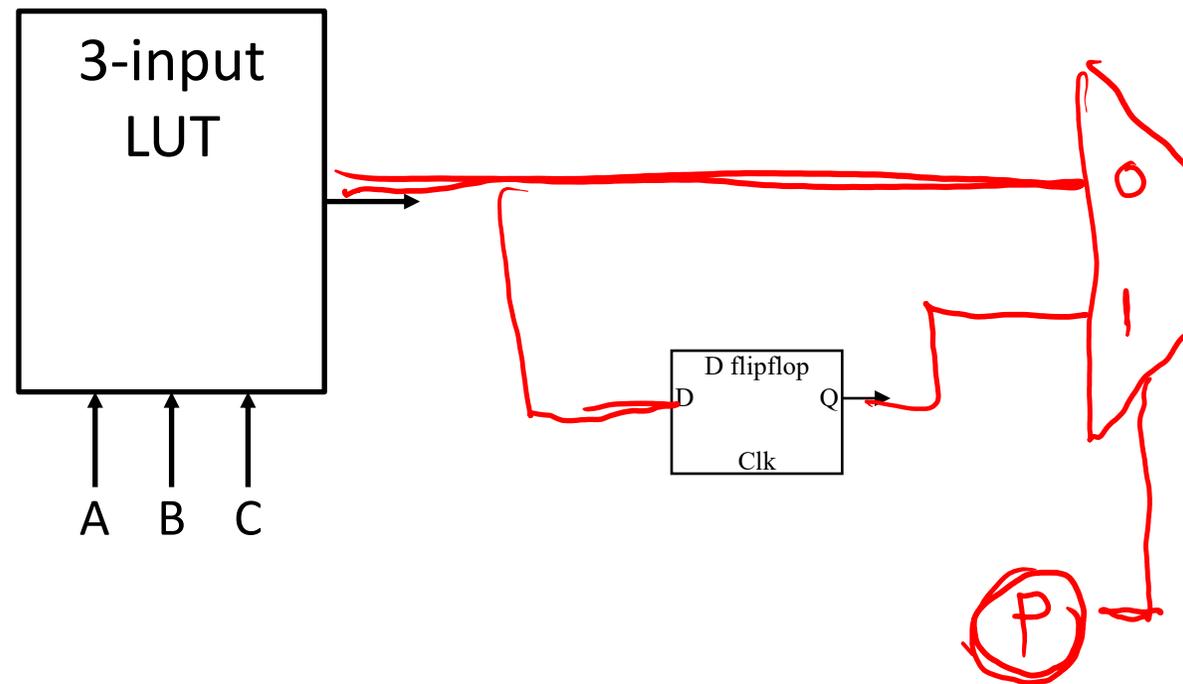
FPGA Programming



Ⓟ = 1 memory cell (stores 1 bit of info)

FPGA Sequential Logic

- ❖ How do we put DFF's onto LUT outputs only when we need them?

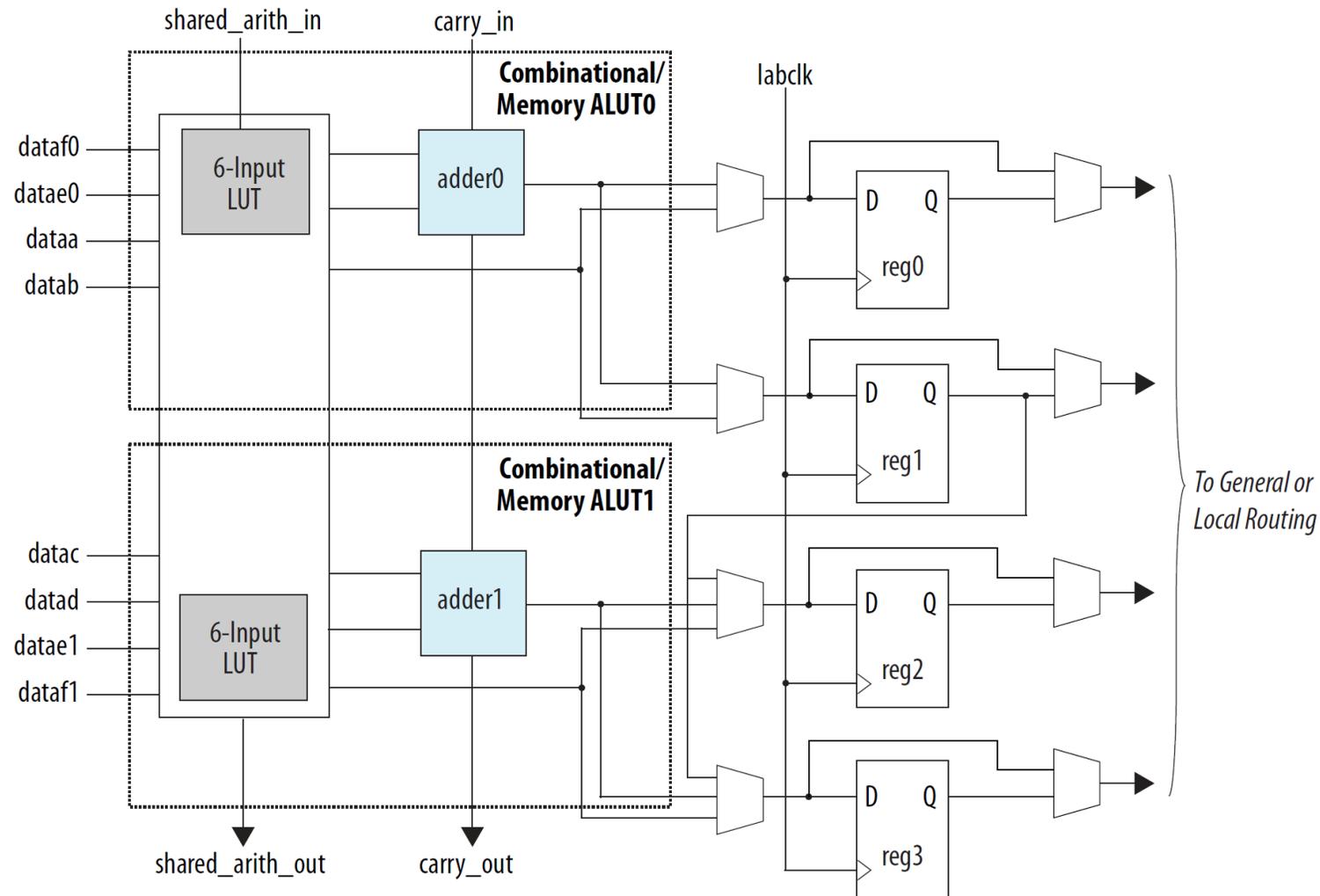


- ❖ Creates a Logic Cell (or Logic Element)

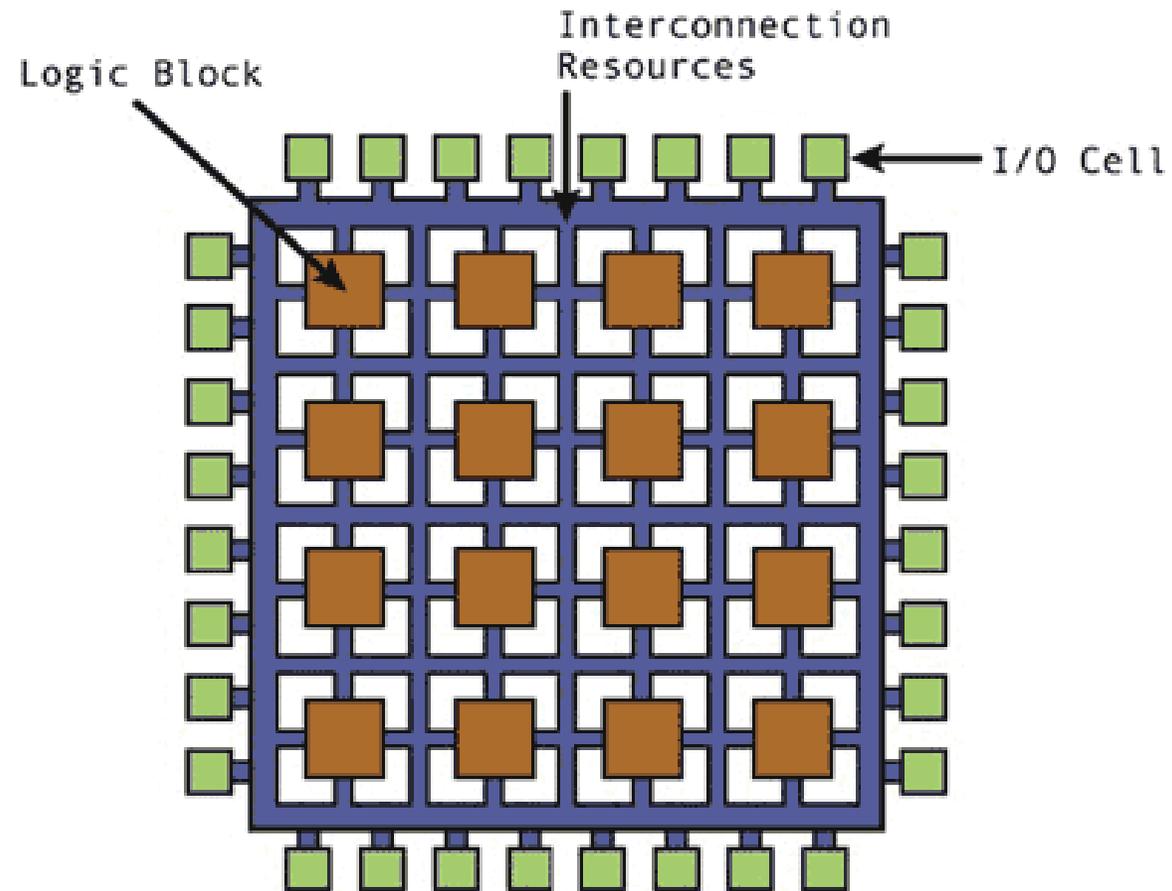
Cyclone V Adaptive Logic Modules

https://www.altera.com/content/dam/altera-www/global/en_US/pdfs/literature/hb/cyclone-v/cv_5v2.pdf

Figure 1-5: ALM High-Level Block Diagram for Cyclone V Devices



Generic FPGA Logic Layout

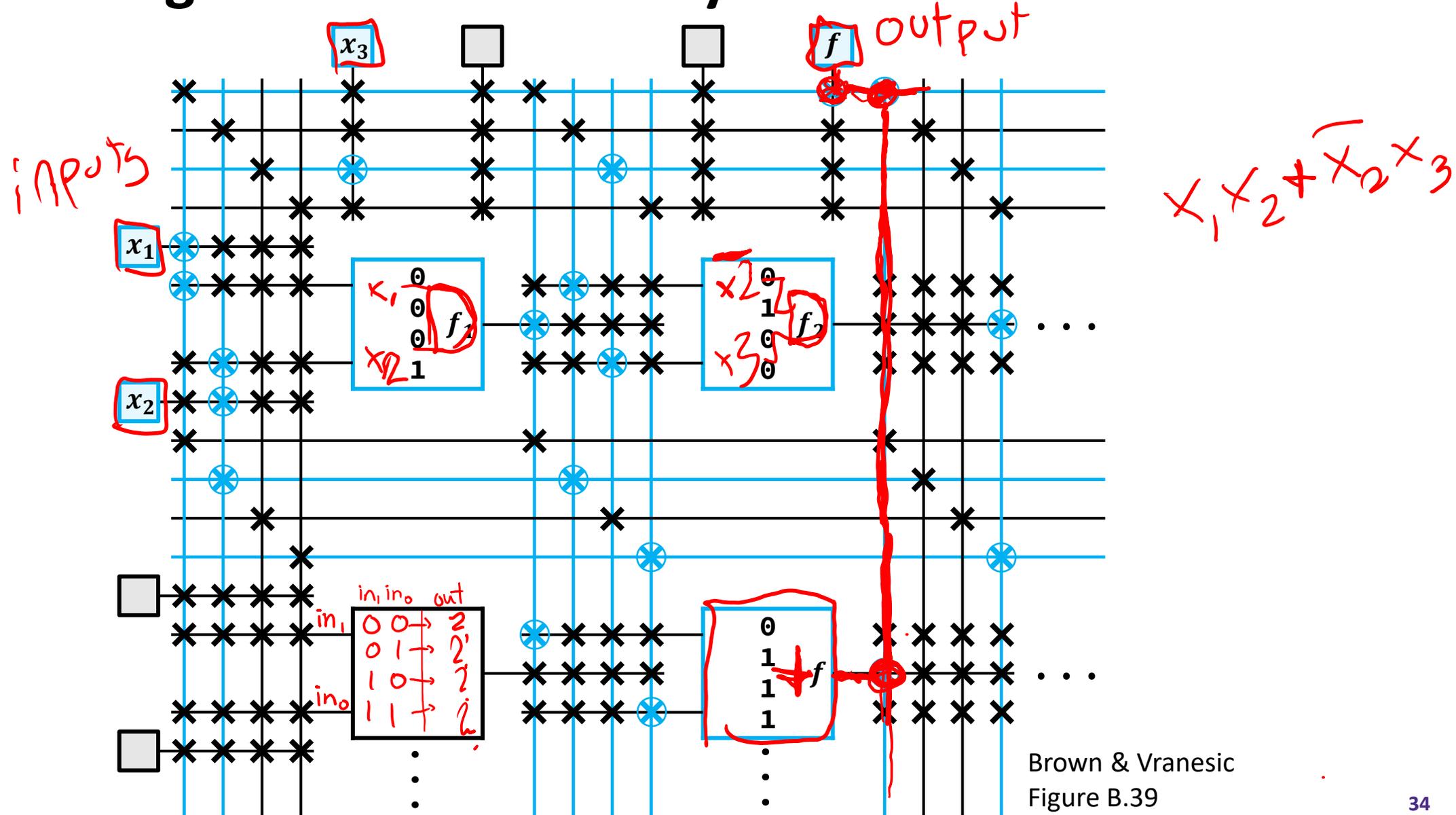


<http://www.chipdesignmag.com/print.php?articleId=434?issueld=16>

Example Programmed Gate Array

⊗ Connected

× Not connected



Brown & Vranesic
Figure B.39

FPGA CAD

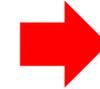
❖ CAD = “Computer-Aided Design”

```
// Verilog code for 2-input
multiplexer
module AOI (F, A, B, C, D);
output F;
input A, B, C, D;
assign F = ~((A & B) | (C &
D));
endmodule
module MUX2 (V, SEL, I, J); //
2:1 multiplexer
output V;
input SEL, I, J;
wire SELB, VB;
not G1 (SELB, SEL);
AOI G2 (VB, I, SEL, SELB, J);
not G3 (V, VB);
endmodule
```

Verilog



FPGA
CAD
Tools



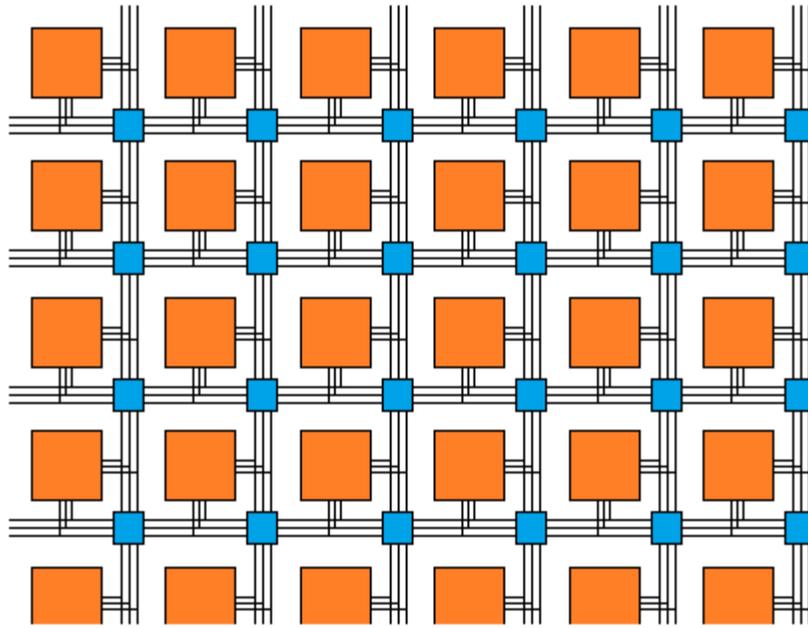
```
00101010001010010
10010010010011000
10101000101011000
10101001010010101
00010110001001010
10101001111001001
01000010101001010
10010010000101010
10100101010010100
01010110101001010
01010010100101001
```

Bitstream

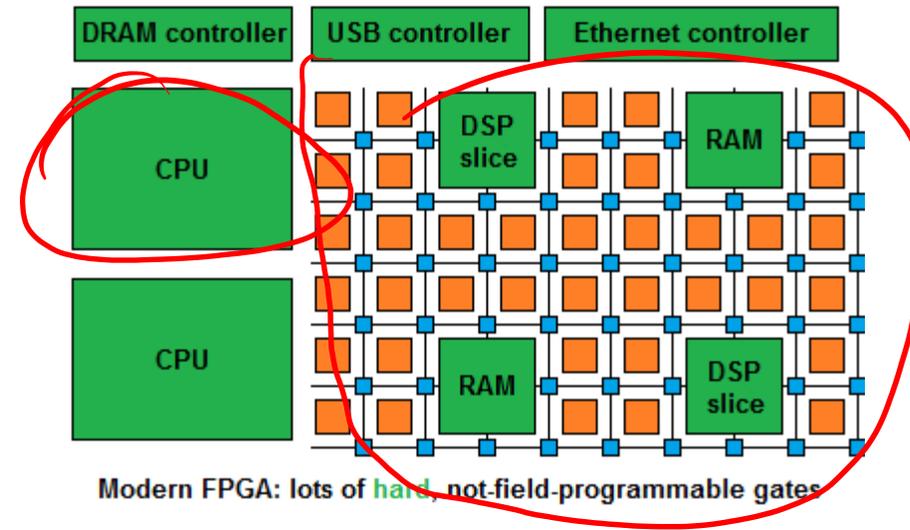
- 1) **Tech Mapping:** Convert Verilog to LUTs
- 2) **Placement:** Assign LUTs to specific locations
- 3) **Routing:** Wire inputs to outputs
- 4) **Bitstream Generation:** Convert mapping to bits

Gate Array vs. SoC

❖ SoC = “System on a Chip”

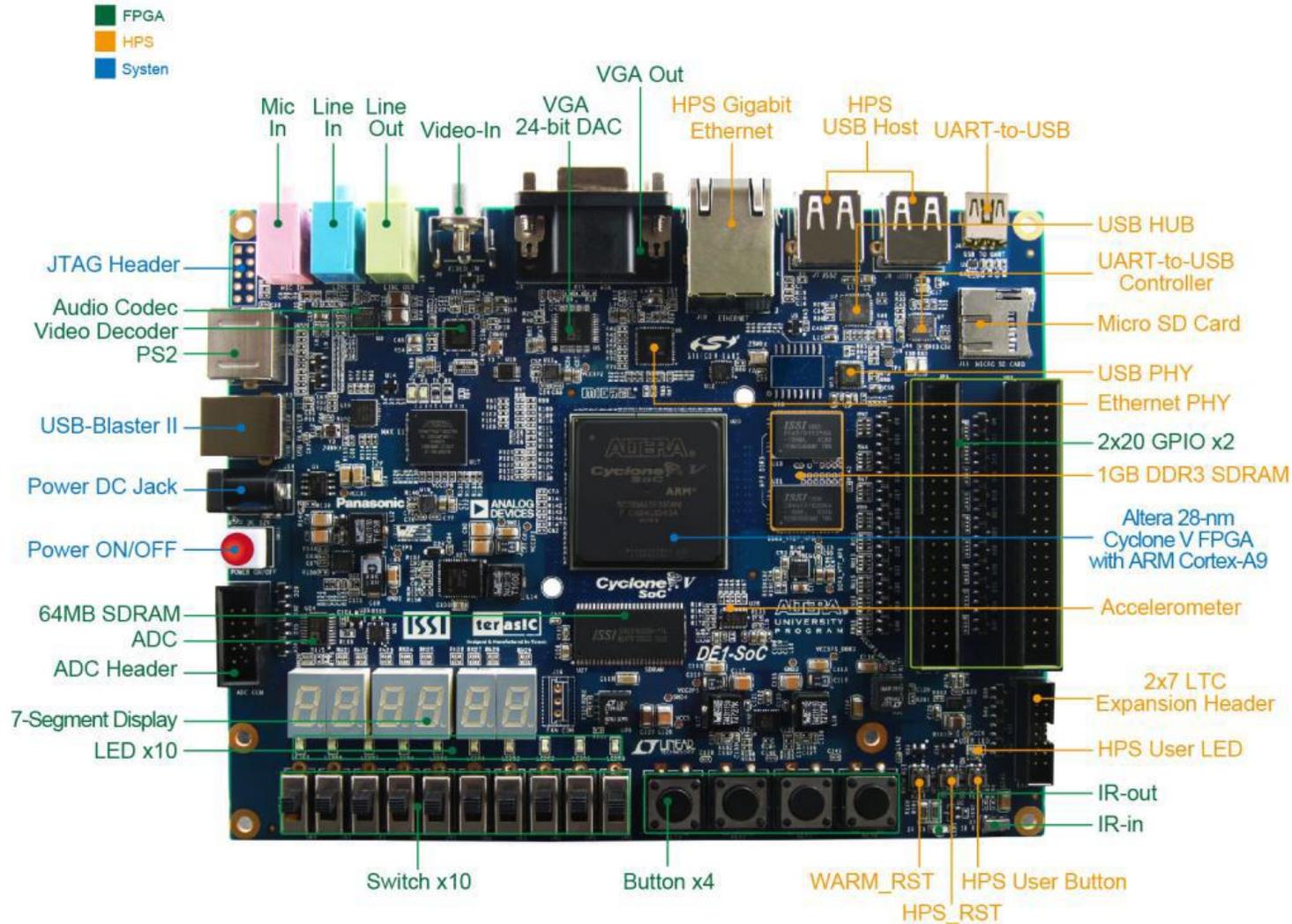


"Clean slate" FPGA: programmable gates and routers

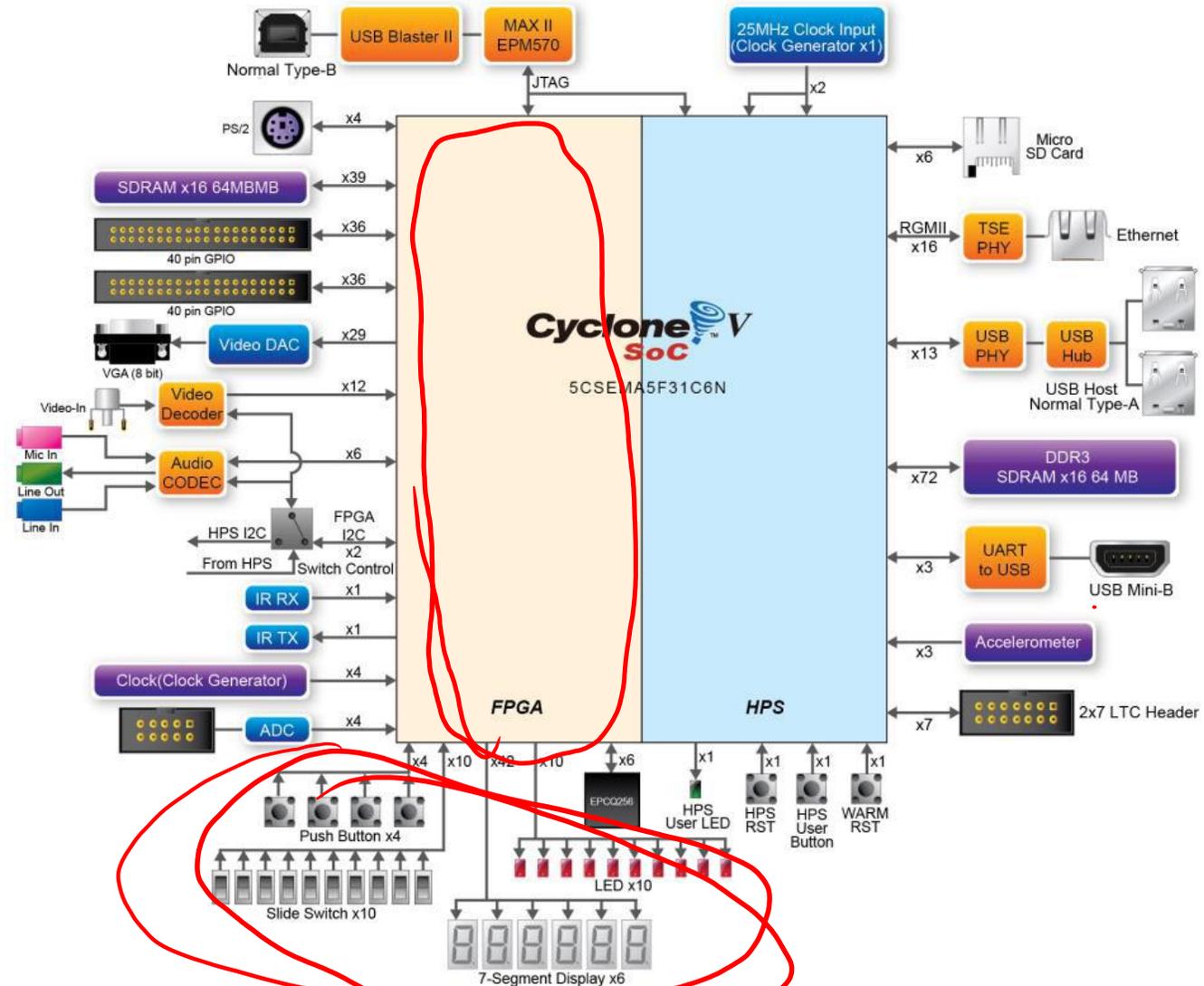


Modern FPGA: lots of hardwired, not-field-programmable gates

DE1-SoC Board



DE1-SoC Board



FPGAs vs. CPUs

- ❖ Individual CPUs designed to handle sequential instruction execution
 - Design and layout determined by architecture
 - Computations “limited” to instruction set
 - Powerful, but also requires things like OS

- ❖ FPGAs
 - Programmable, so specially-designed logic for application
 - But can be difficult to specify certain applications
 - Good for parallelizing computations
 - Can perform separate computations if separated via routing