

Intro to Digital Design

Sequential Logic

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Relevant Course Information

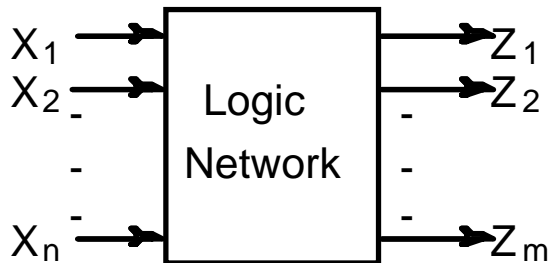
- ❖ Lab 3 Demos due during your assigned demo slots
 - Don't forget to submit your lab materials *before* Wednesday at 2:30 pm, regardless of your demo time

- ❖ Lab 4 – 7-seg displays

- ❖ Quiz 1 is next week in lecture
 - Last 20 minutes, worth 10% of your course grade
 - On Lectures 1-3: CL, K-maps, Waveforms, and Verilog
 - Past Quiz 1 (+ solutions) on website: Course Info → Quizzes

Synchronous Digital Systems (SDS)

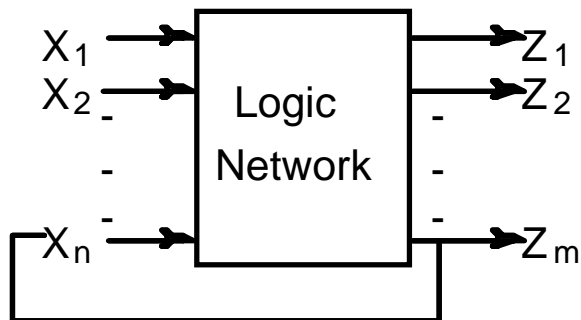
❖ Combinational Logic (CL)



Network of logic gates without feedback.

Outputs are functions only of inputs.

❖ Sequential Logic (SL)



The presence of feedback introduces the notion of "state."

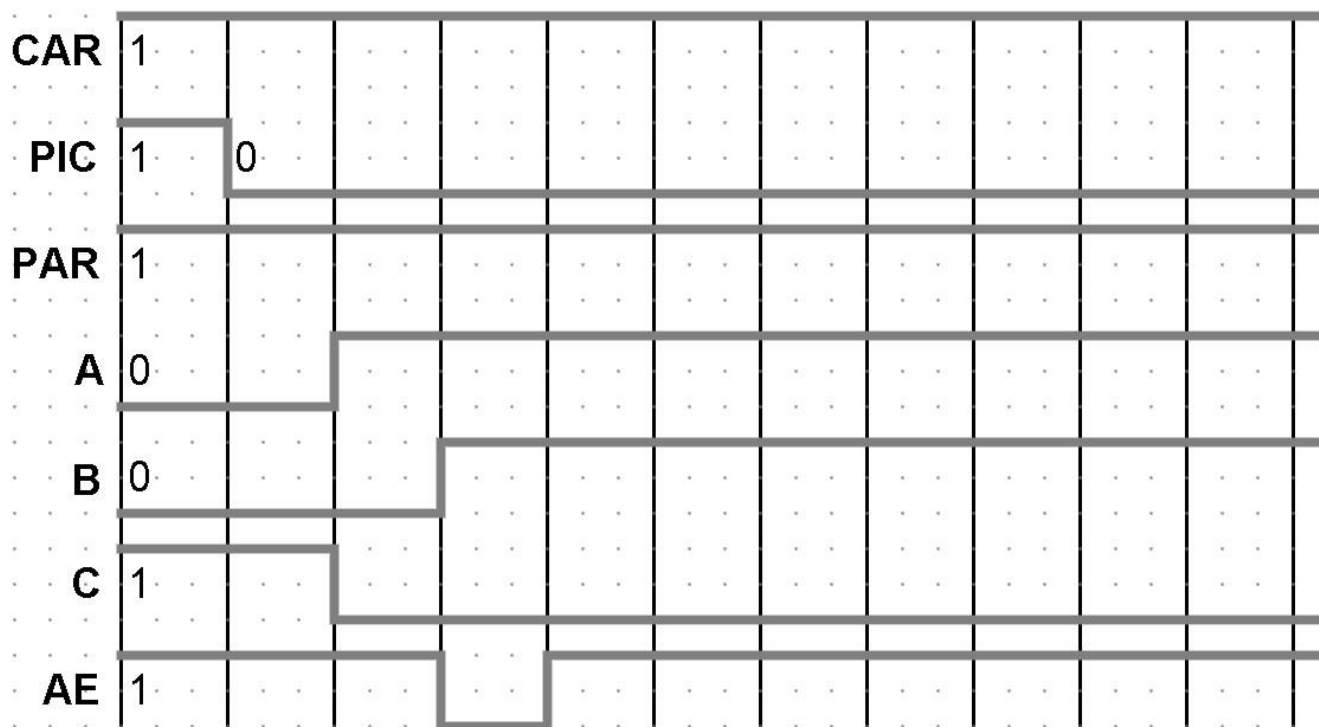
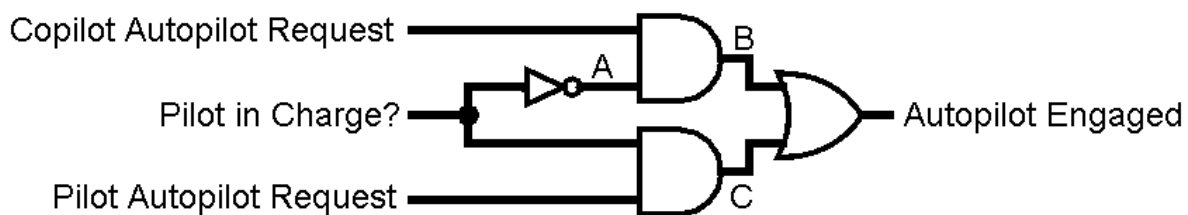
Circuits can "remember" or store information.

Uses for Sequential Logic

- ❖ Place to store values for some amount of time:
 - Registers
 - Memory
- ❖ *Help control flow of information between combinational logic blocks*
 - Hold up the movement of information to allow for orderly passage through CL

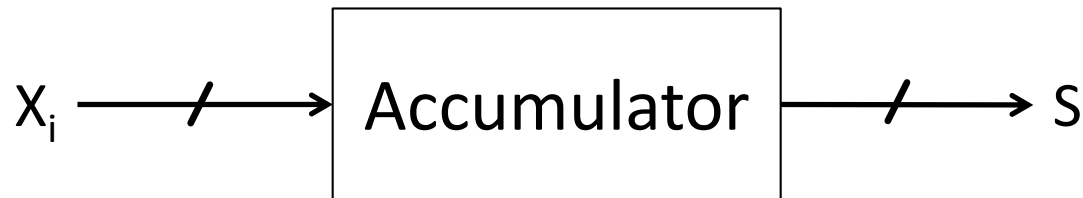
Control Flow of Information?

❖ Circuits can temporarily go to incorrect states!



Accumulator Example

- ❖ An example of why we would need to control the flow of information.



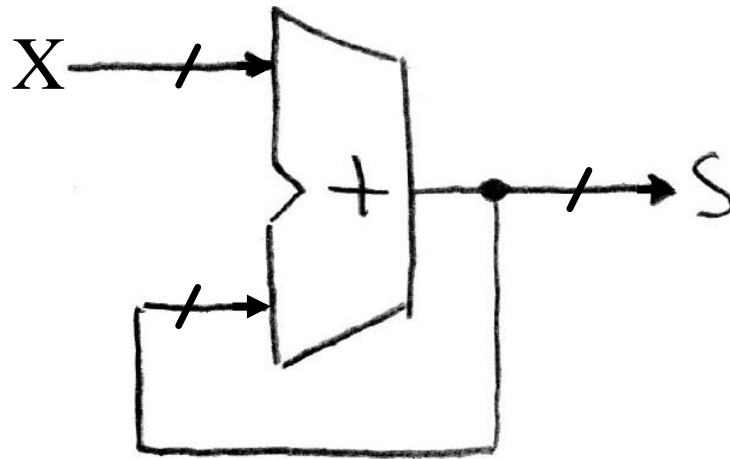
- ❖ Want:

```
S = 0;  
for (i=0; i<n; i++)  
    S = S + X_i;
```
- ❖ Assume:
 - Each X value is applied in succession, one per cycle
 - The sum since cycle 0 is present on S

Accumulator: First Try

Does this work?


- No

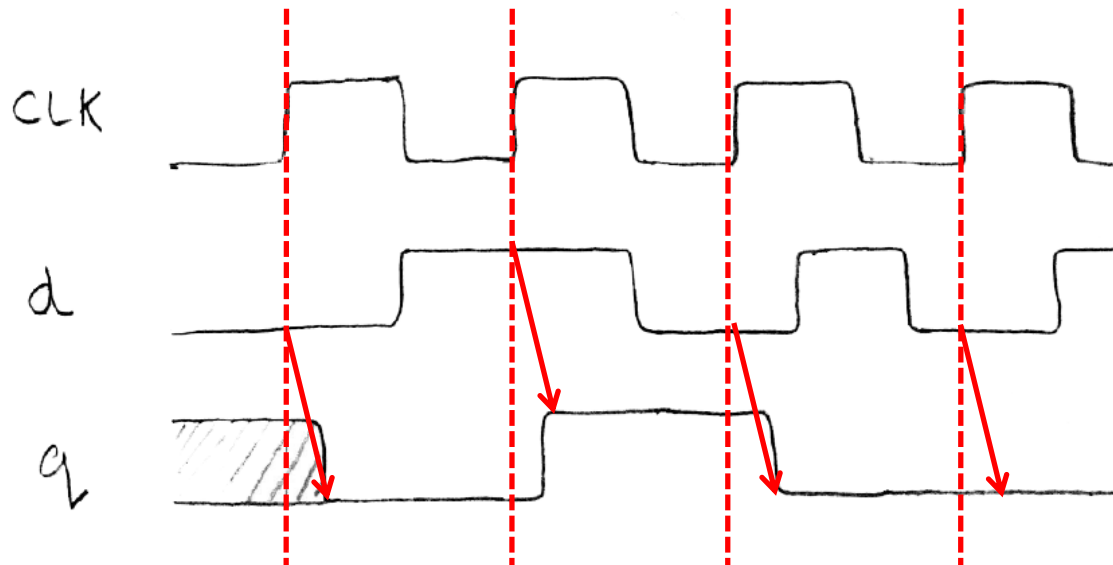
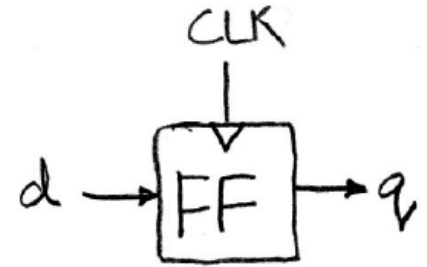


- 1) How to control the next iteration of the 'for' loop?
- 2) How do we say: 'S=0'?

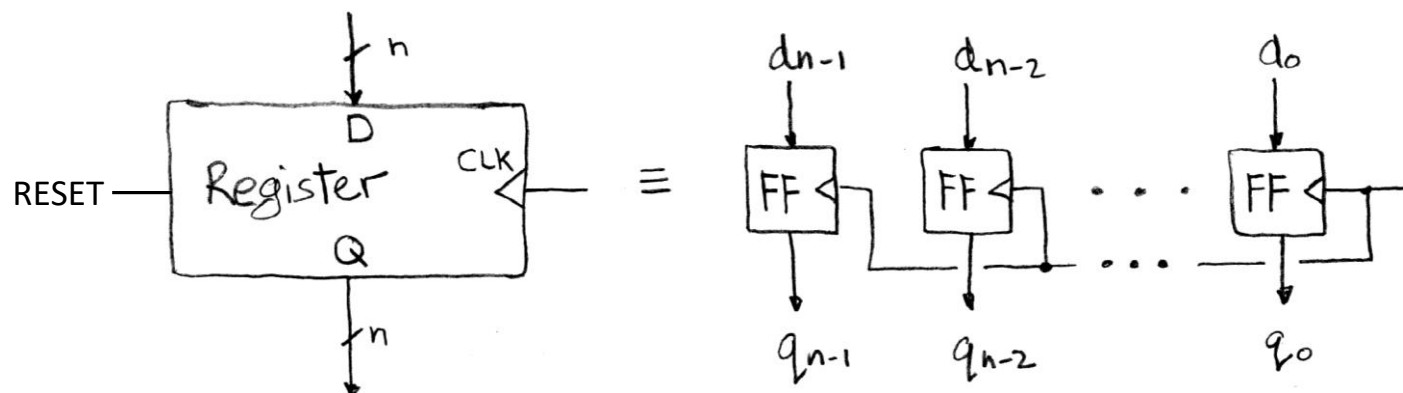
State Element: Flip-Flop

❖ Positive edge-triggered D-type flip flop

- “On the rising edge of the clock (), input d is sampled and transferred to the output q . At other times, the input d is ignored and the previously sampled value is retained.”

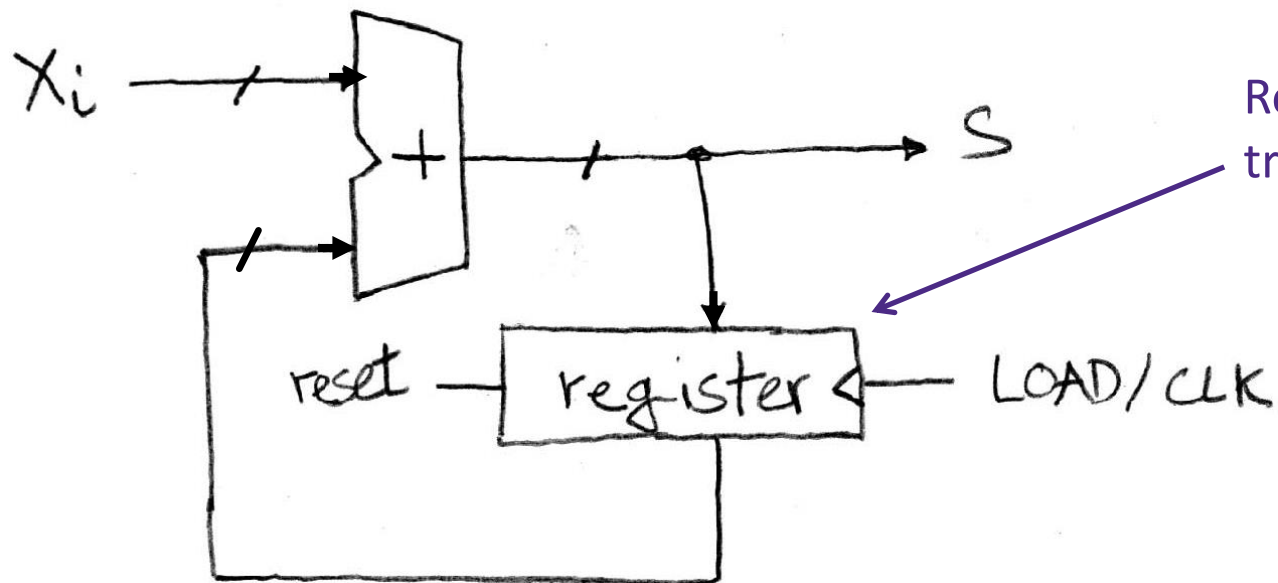


State Element: Register

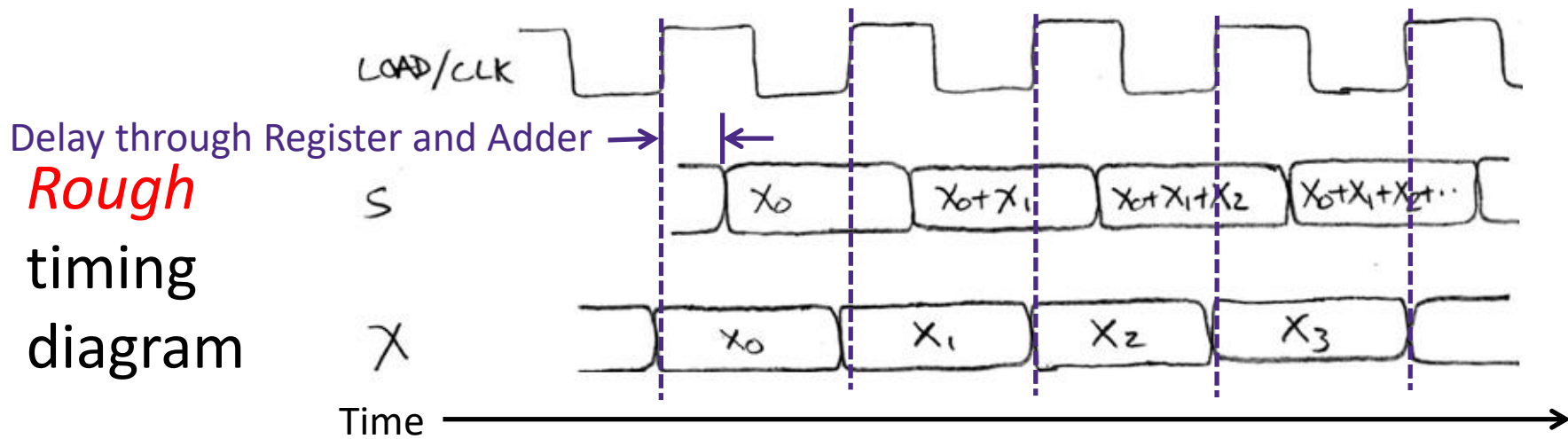


- ❖ n instances of flip-flops together
 - One for every bit in input/output bus width
- ❖ Output Q resets to zero when RESET signal is high *during* clock trigger
 - Some extra circuitry required for this

Accumulator: Second Try

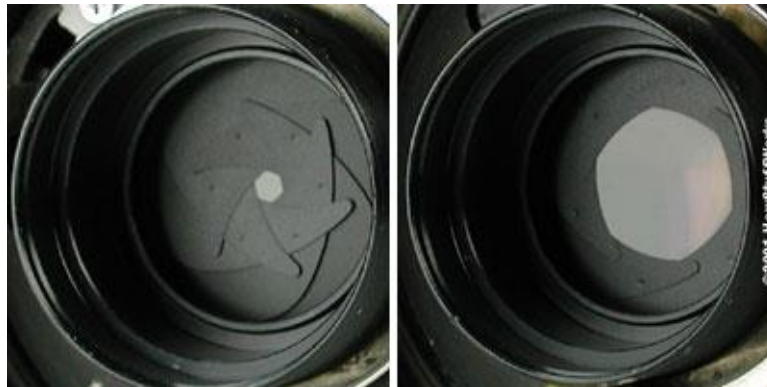


Register holds up the transfer of data to adder



Flip-Flop Timing Terminology (1/2)

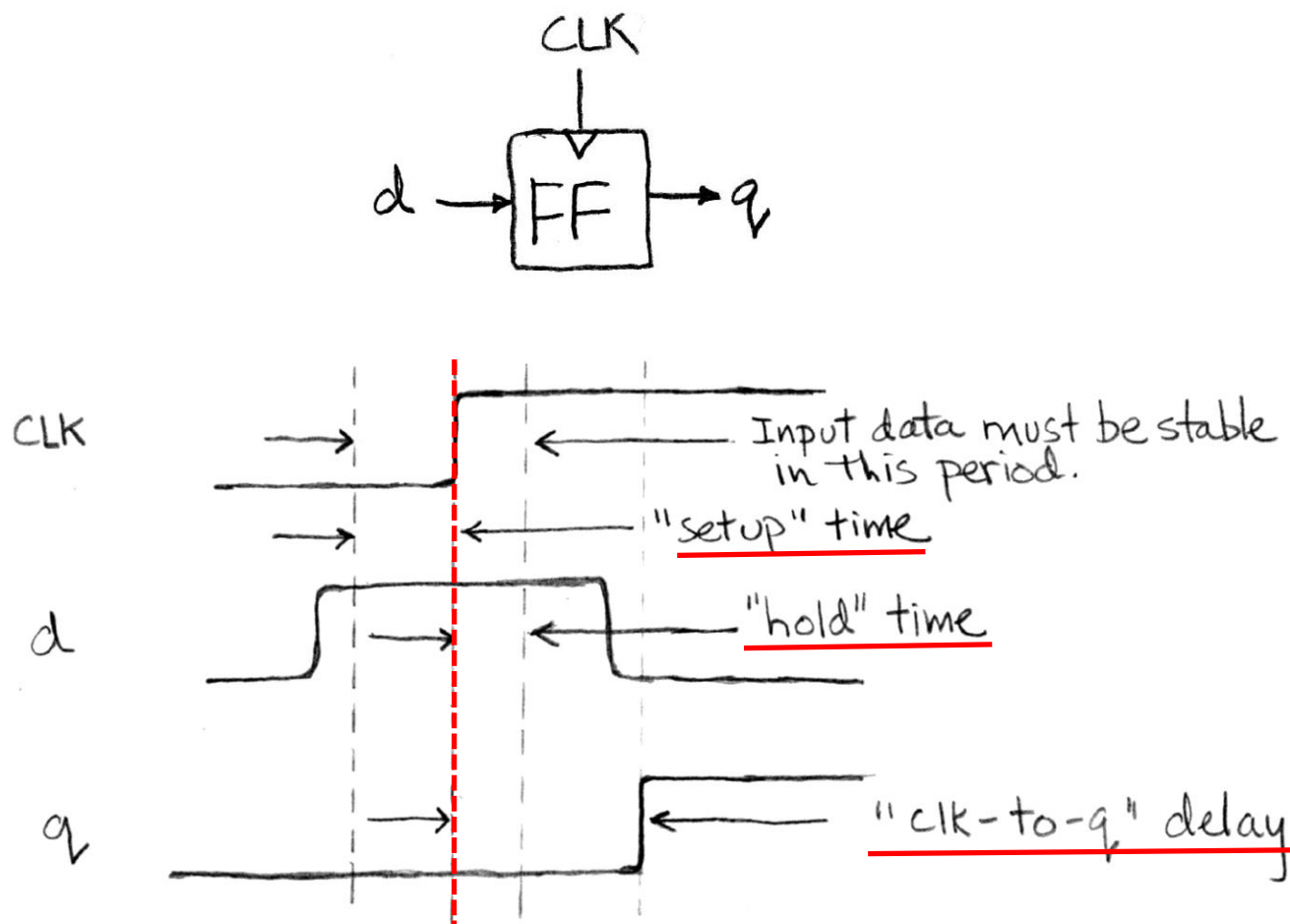
- ❖ Camera Analogy: non-blurry digital photo
 - *Don't move* while camera shutter is opening
 - *Don't move* while camera shutter is closing
 - *Check for blurriness* once image appears on the display



Flip-Flop Timing Terminology (2/2)

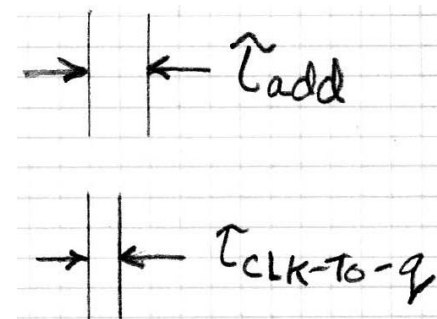
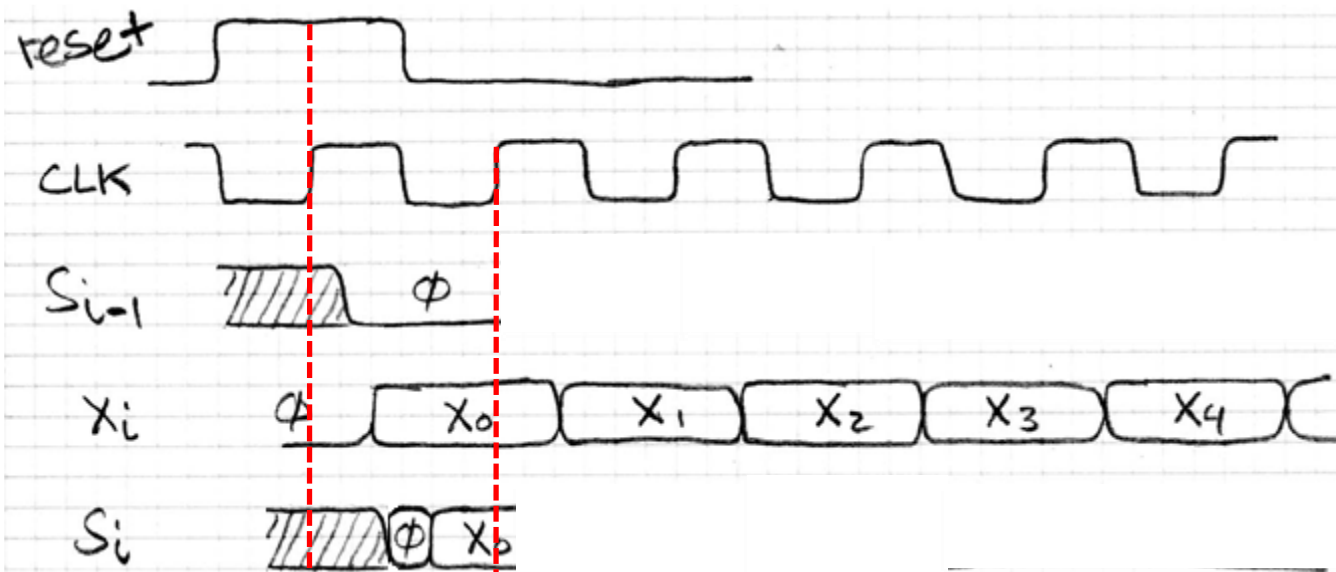
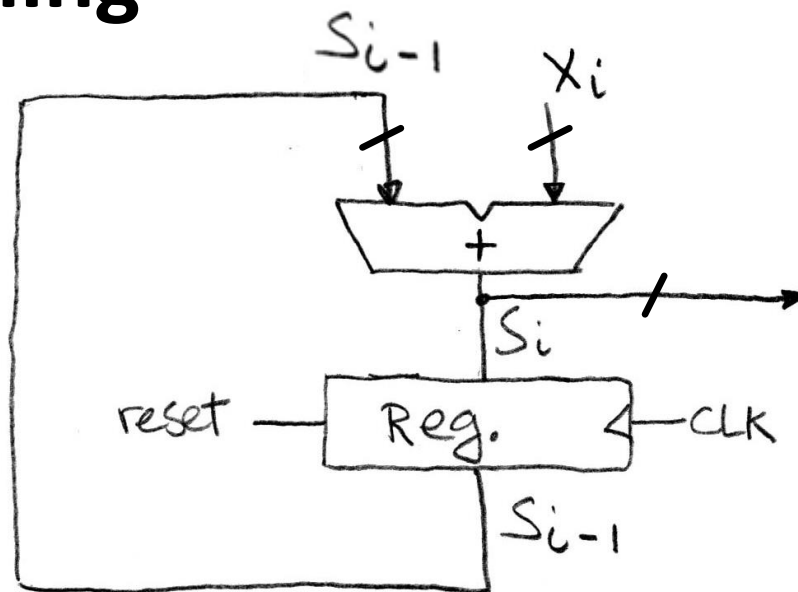
- ❖ Now applied to sequential logic elements:
 - *Setup Time*: how long the input must be stable *before* the CLK trigger for proper input read
 - *Hold Time*: how long the input must be stable *after* the CLK trigger for proper input read
 - *“CLK-to-Q” Delay*: how long it takes the output to change, measured from the CLK trigger

Flip-Flop Timing Behavior



Accumulator: Proper Timing

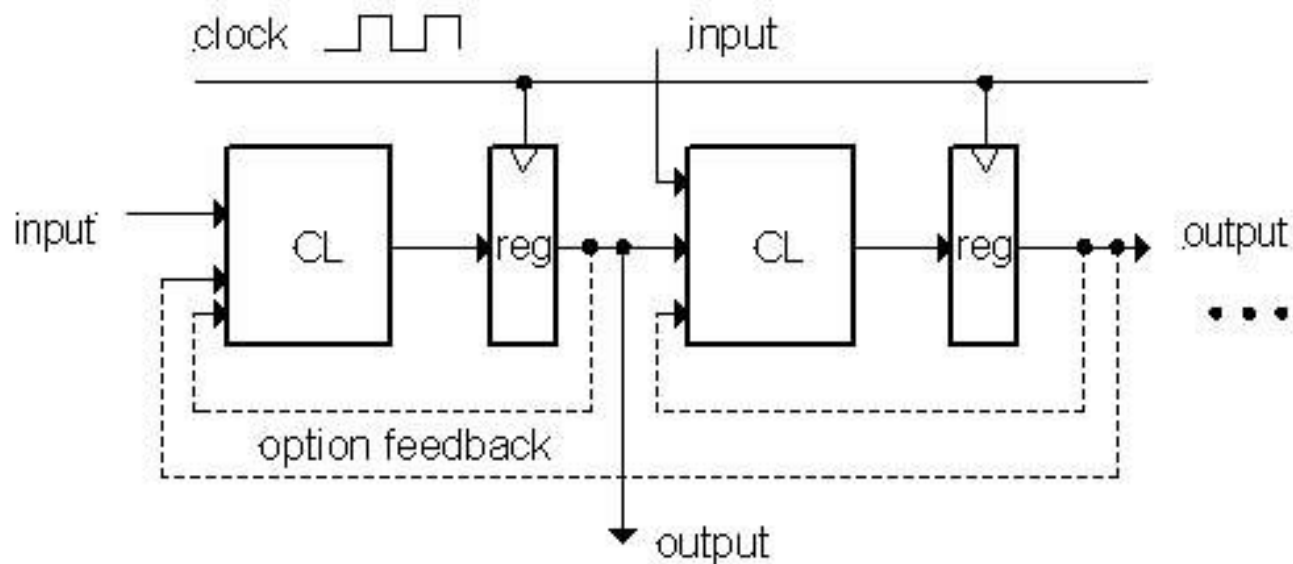
- reset signal shown
- Also, in practice X_i might not arrive at the adder at the same time as S_{i-1}
- S_i temporarily is wrong, but register always captures correct value
- In good circuits, instability never happens around rising edge of CLK



Review Question

- ❖ Which of the following statements is TRUE?
 - (A) The input to a flip-flop must remain stable throughout the CLK-to-Q delay.
 - (B) A flip-flop switches between 0 and 1 on each trigger.
 - (C) In a SDS, we only need to know setup time, hold time, and clk-to-q delay constants to ensure correct behavior.
 - (D) None of the above.

Model for Synchronous Systems



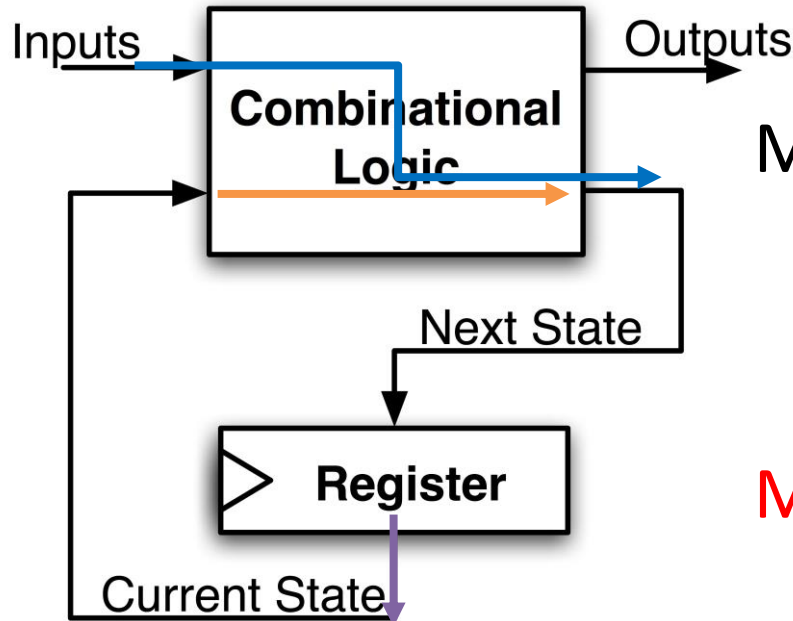
- ❖ Combinational logic blocks separated by registers
 - Clock signal connects only to sequential logic elements
 - Feedback is optional depending on application
- ❖ **How do we ensure proper behavior?**
 - **How fast can we run our clock?**

When Can the Input Change?

- ❖ When a register input changes shouldn't violate hold time (t_{hold}) or setup time (t_{setup}) constraints within a clock period (t_{period})
- ❖ Let $t_{input,i}$ be the time it takes for the input of a register to change for the i -th time in a single clock cycle, measured from the CLK trigger:
 - Then we need $t_{hold} \leq t_{input,i} \leq t_{period} - t_{setup}$ for all i
 - Two separate constraints!

Minimum Delay

- ❖ If shortest path to register input is too short, might violate hold time constraint
 - Input could change before state is “locked in”
 - Particularly problematic with *asynchronous* signals

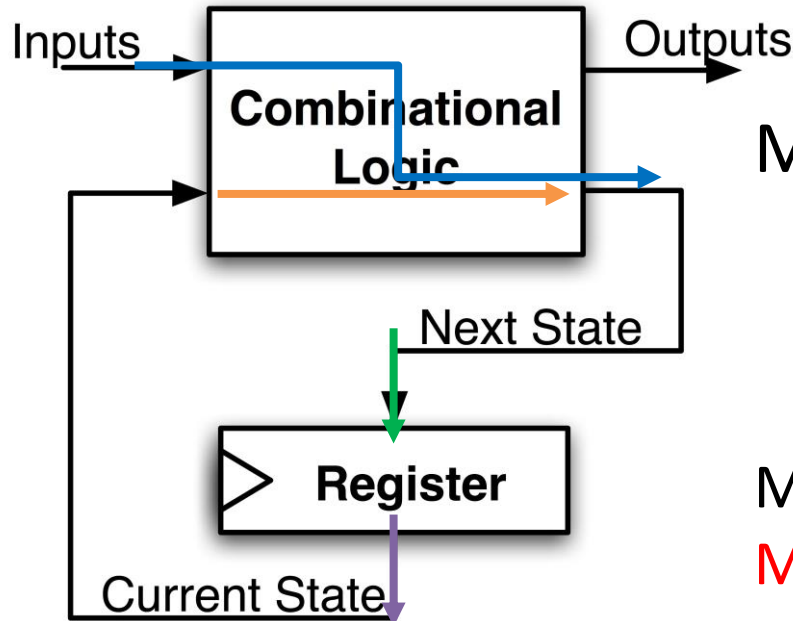


$$\text{Min Delay} = \min(\text{CLK-to-Q Delay} + \text{Min CL Delay}, \text{Min CL Delay})$$

$$\text{Min Delay} \geq \text{Hold Time}$$

Maximum Clock Frequency

- ❖ What is the max frequency of this circuit?
 - Limited by how much time needed to get correct Next State to Register (t_{setup} constraint)

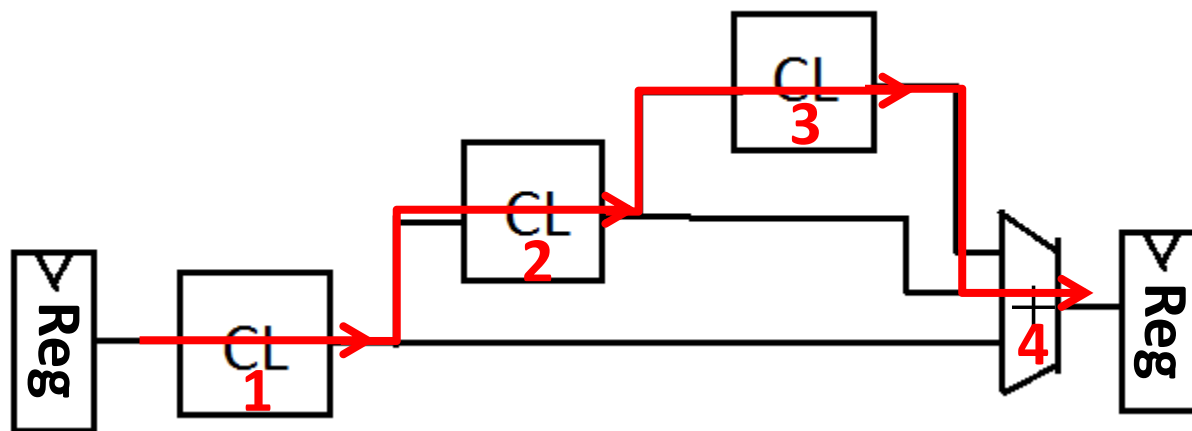


$$\text{Max Delay} = \max(\text{CLK-to-Q Delay} \\ + \text{Max CL Delay}, \\ + \text{Max CL Delay})$$

$$\text{Min Period} = \text{Max Delay} + \text{Setup Time}$$
$$\text{Max Freq} = 1/\text{Min Period}$$

The Critical Path

- ❖ The *critical path* is the longest delay between *any* two registers in a circuit
- ❖ The clock period must be *longer* than this critical path, or the signal will not propagate properly to that next register

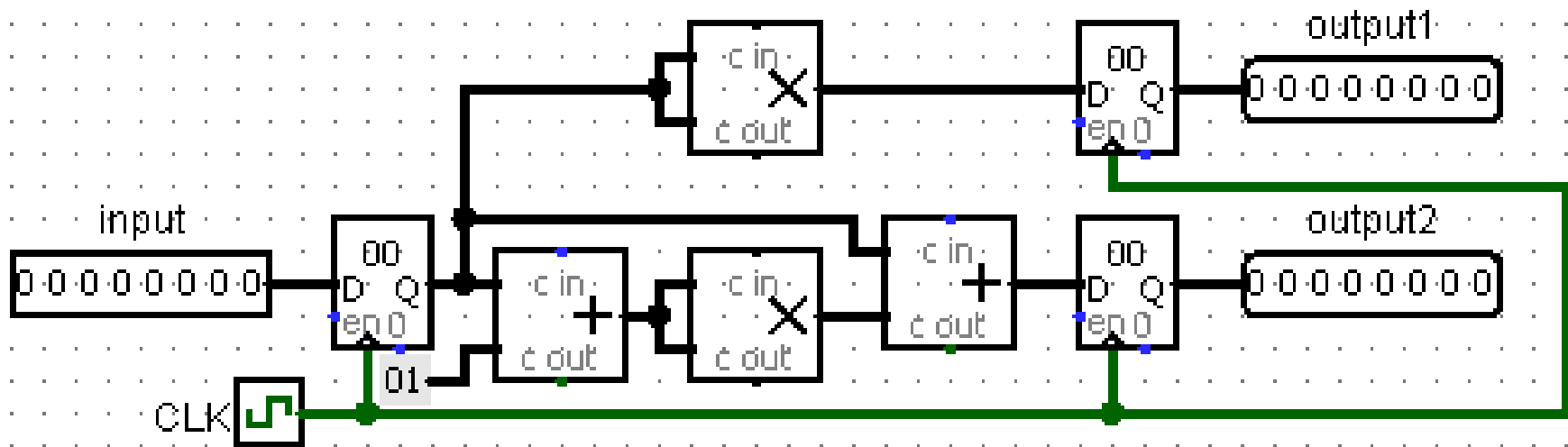


Critical Path =

CLK-to-Q Delay
+ CL Delay 1
+ CL Delay 2
+ CL Delay 3
+ Adder Delay
+ Setup Time

Practice Question

- ❖ We want to run on 1 GHz processor. $t_{add} = 100$ ps. $t_{mult} = 200$ ps. $t_{setup} = t_{hold} = 50$ ps. What is the maximum $t_{clk-to-q}$ we can use?

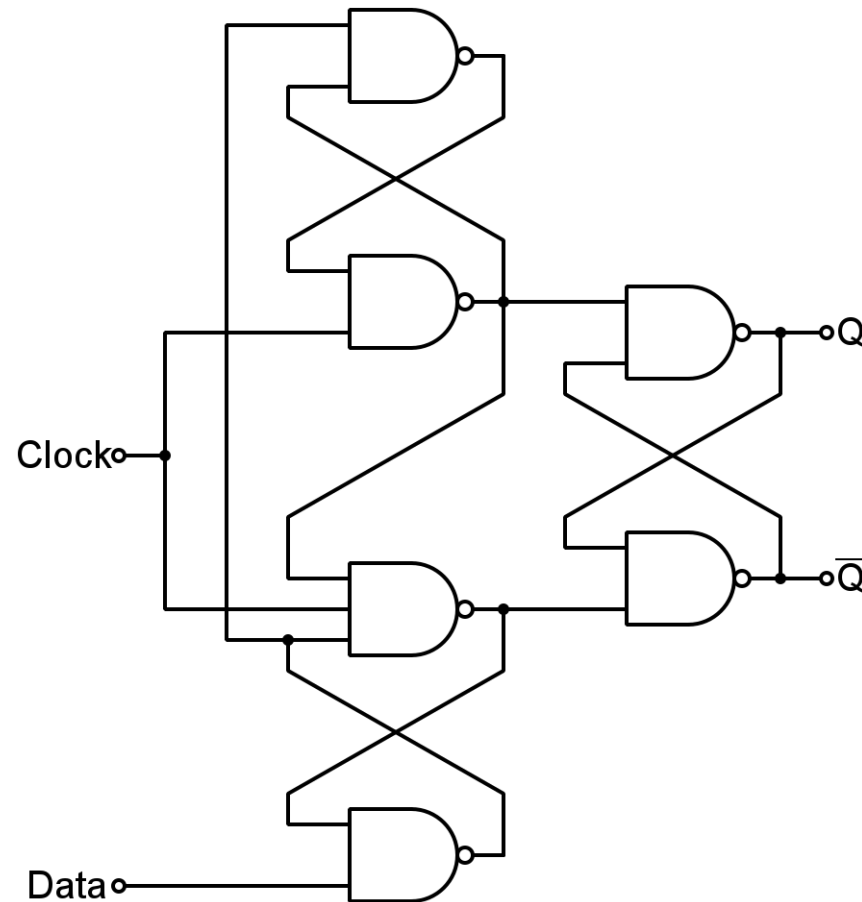


- (A) 550 ps (B) 750 ps (C) 500 ps (D) 700 ps

Technology Break

Where Do Timing Terms Come From?

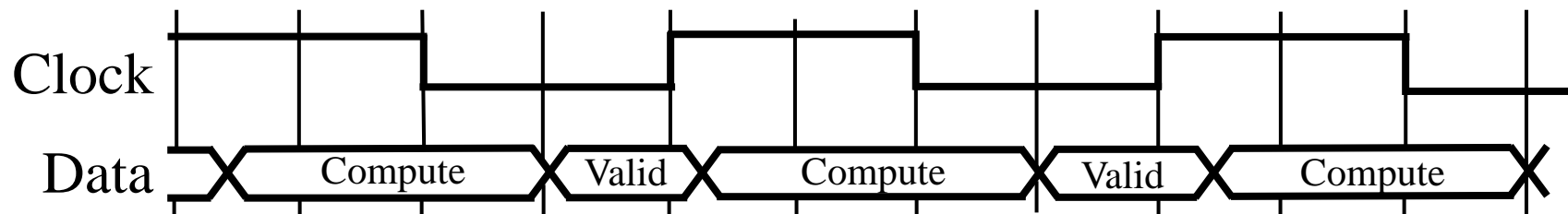
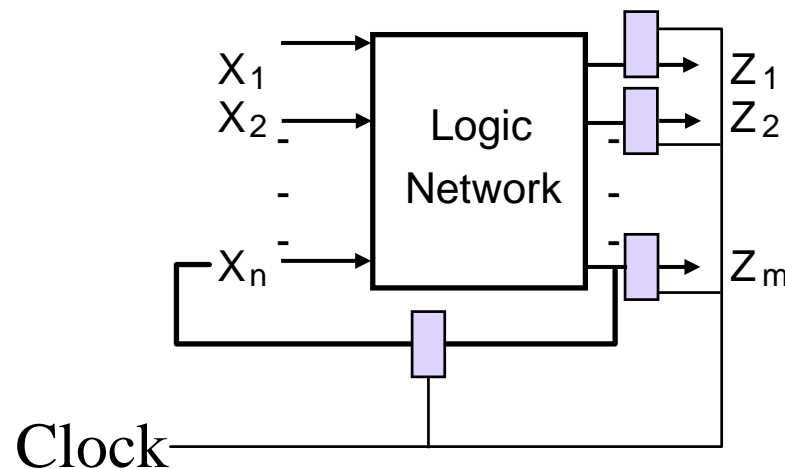
Edge-triggered
D flip-flop:



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<https://commons.wikimedia.org/w/index.php?curid=40852354>

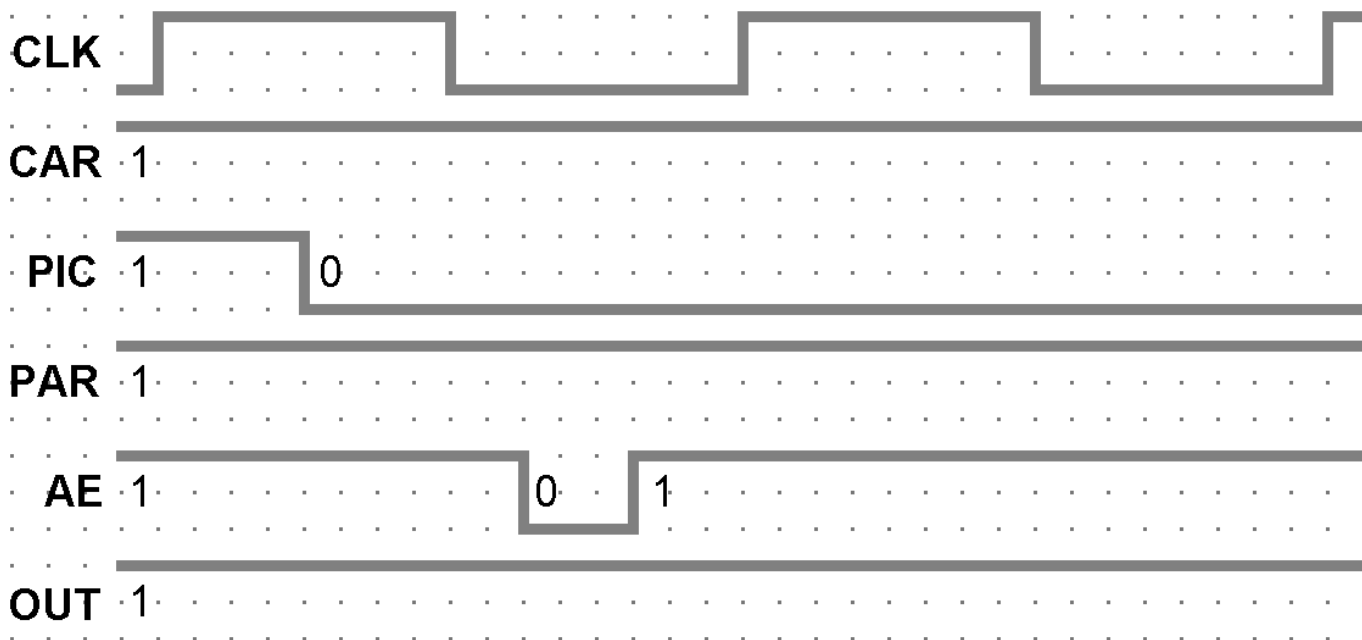
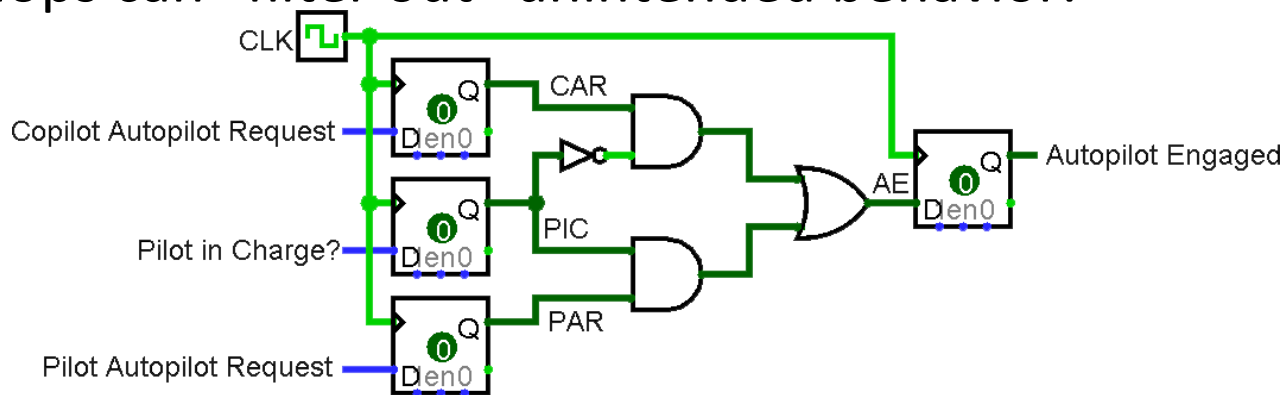
Safe Sequential Circuits

- ❖ Clocked elements on feedback, perhaps outputs
 - Clock signal synchronizes operation
 - Clocked elements hide glitches/hazards



Autopilot Revisited

- ❖ Flip-flops can “filter out” unintended behavior:

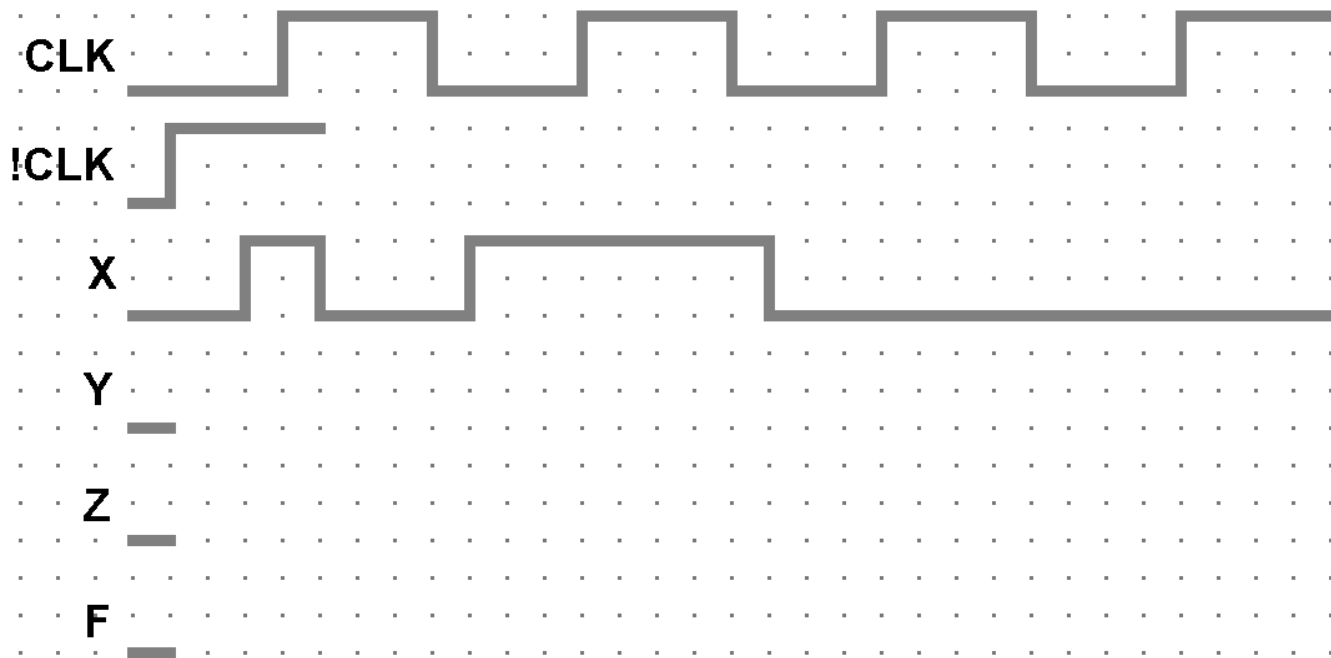
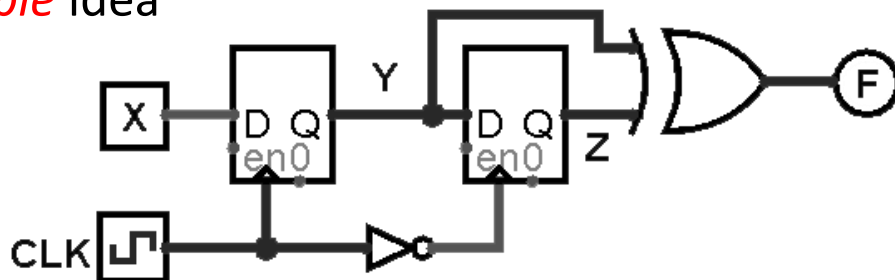


Waveform Diagrams Revisited

- ❖ Easiest to start with CLK on top
 - Solve signal by signal, from inputs to outputs
 - Can only draw the waveform for a signal if *all* of its input waveforms are drawn
- ❖ When does a signal update?
 - A *state element* updates based on CLK triggers
 - A *combinational element* updates ANY time ANY of its inputs changes

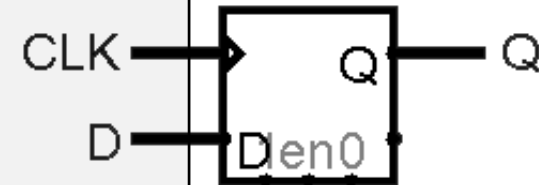
Example: SDS Waveform Diagram

- ❖ Assume: $t_{C2Q} = 3$ ticks, $t_{XOR} = 2$ ticks, $t_{NOT} = 1$ tick; $t_s = t_h = 0$
 - Note: clocking the gate is a *terrible* idea

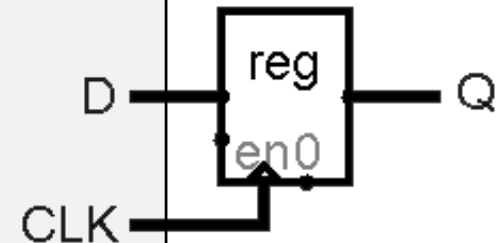


Verilog: Basic D Flip-Flop, Register

```
module basic_D_FF (q, d, clk);  
    output logic q; // q is state-holding  
    input  logic d, clk;  
  
    always_ff @(posedge clk)  
        q <= d; // use <= for clocked elements  
endmodule
```



```
module basic_reg (q, d, clk);  
    output logic [7:0] q;  
    input  logic [7:0] d;  
    input  logic      clk;  
  
    always_ff @(posedge clk)  
        q <= d;  
endmodule
```



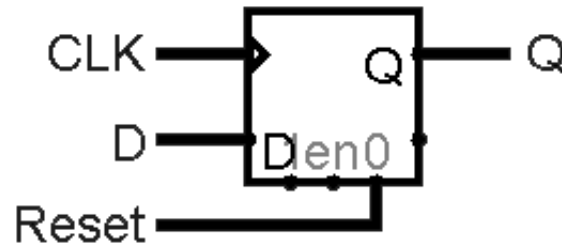
Procedural Blocks

- ❖ `always`: loop to execute over and over again
 - Block gets triggered by a *sensitivity list*
 - Any object that is assigned a value in an `always` statement must be declared as a variable (`logic` or `reg`).
 - Example:
 - `always @ (posedge clk)`
- ❖ `always_ff`: special SystemVerilog for SL
 - *Only for use with sequential logic – signal intent that you want flip-flops*
 - Example:
 - `always_ff @ (posedge clk)`

SystemVerilog Coding Guidelines

- 1) When modeling sequential logic, use *nonblocking* assignments
- 2) When modeling combinational logic with an `always_comb` block, use *blocking* assignments
- 3) When modeling both sequential and combinational logic within the same `always_ff` block, use *nonblocking* assignments
- 4) Do not mix *blocking* and *nonblocking* assignments in the same `always_*` block
- 5) Do not make assignments to the same variable from more than one `always_*` block

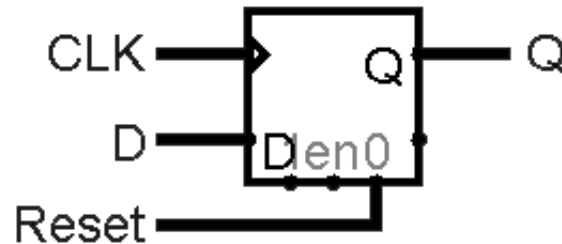
Verilog: Reset Functionality



❖ Option 1: synchronous reset

```
module D_FF1 (q, d, reset, clk);  
    output logic q; // q is state-holding  
    input  logic d, reset, clk;  
  
    always_ff @(posedge clk)  
        if (reset)  
            q <= 0; // on reset, set to 0  
        else  
            q <= d; // otherwise pass d to q  
  
endmodule
```


Verilog: Reset Functionality



❖ Option 2: asynchronous reset

```
module D_FF1 (q, d, reset, clk);  
    output logic q; // q is state-holding  
    input  logic d, reset, clk;  
  
    always_ff @(posedge clk or posedge reset)  
        if (reset)  
            q <= 0; // on reset, set to 0  
        else  
            q <= d; // otherwise pass d to q  
  
endmodule
```

Verilog: Simulated Clock

- ❖ For simulation, you need to generate a clock signal:
 - For entirety of simulation/program, so use `always` block

Explicit Edges:

```
initial
    clk = 0;
always begin
    #50  clk <= 1;
    #50  clk <= 0;
end
```

Toggle:

```
initial
    clk = 0;
always
    #50  clk <= ~clk;
```

- ❖ Define clock period:
 - Define **parameter**

```
parameter period = 100;
initial
    clk = 0;
always
    #(period/2)  clk <= ~clk;
```

Verilog Testbench with Clock

```
module D_FF_testbench;
  logic CLK, reset, d;
  logic q;

  parameter PERIOD = 100;

  D_FF dut (.q, .d, .reset, .CLK); // Instantiate the D_FF

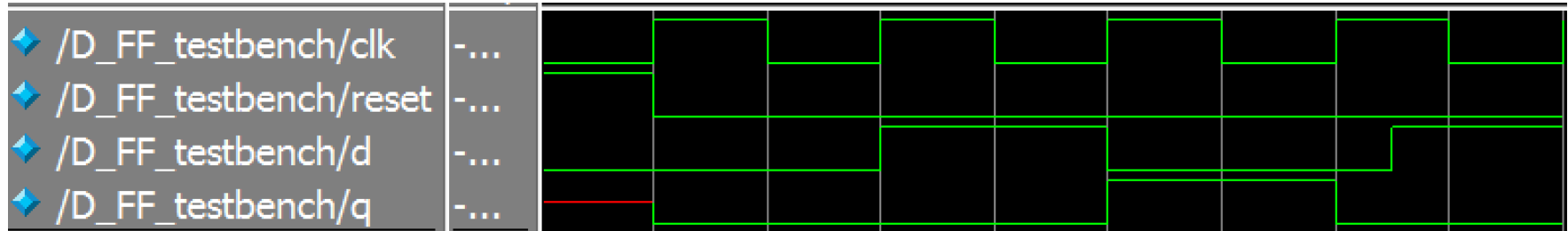
  initial CLK <= 0; // Set up clock
  always #(PERIOD/2) CLK<= ~CLK;

  initial begin // Set up signals
    d <= 0; reset <= 1;
    @(posedge CLK); reset <= 0;
    @(posedge CLK); d <= 1;
    @(posedge CLK); d <= 0;
    @(posedge CLK); #(PERIOD/4) d <= 1;
    @(posedge CLK);
    $stop(); // end the simulation
  end
endmodule
```

Timing Controls

- ❖ Delay: `#<time>`
 - Delays by a specific amount of simulation time
 - Can do calculations in `<time>`
 - Examples: `# (PERIOD/4)`, `#50`
- ❖ Edge-sensitive: `@ (<pos/negedge> signal)`
 - Delays next statement until specified transition on signal
 - Example: `@ (posedge CLK)`
- ❖ Level-sensitive Event: `wait (<expression>)`
 - Delays next statement until `<expression>` evaluates to TRUE
 - Example: `wait (enable == 1)`

ModelSim Waveforms



```
initial begin  
    d <= 0; reset <= 1;  
    @ (posedge CLK) ;      reset <= 0;  
    @ (posedge CLK) ; d <= 1;  
    @ (posedge CLK) ; d <= 0;  
    @ (posedge CLK) ; # (PERIOD/4) d <= 1;  
    @ (posedge CLK) ;  
    $stop() ;  
end
```

Summary (1/2)

- ❖ State elements controlled by clock
 - Store information
 - Control the flow of information between other state elements and combinational logic
- ❖ Registers implemented from flip-flops
 - Triggered by CLK, pass input to output, can reset
- ❖ Critical path constrains clock rate
 - Timing constants: setup time, hold time, clk-to-q delay, propagation delays

Summary (2/2)

- ❖ Generating a clock
 - Manually create using `always` block
 - Need to decide on period
- ❖ Blocking vs. Non-blocking
 - Blocking: Statements executed “in series”
 - Non-blocking: Statements executed “in parallel”
 - Always use non-blocking for clocked elements
- ❖ Synchronous vs. Asynchronous
 - Whether signals are controlled by clock or not