

The Hardware/Software Interface

Integers II

Instructors:

Amber Hu, Justin Hsia

Teaching Assistants:

Anthony Mangus

Grace Zhou

Jiuyang Lyu

Kurt Gu

Mendel Carroll

Naama Amiel

Rose Maresh

Violet Monserate

Divya Ramu

Jessie Sun

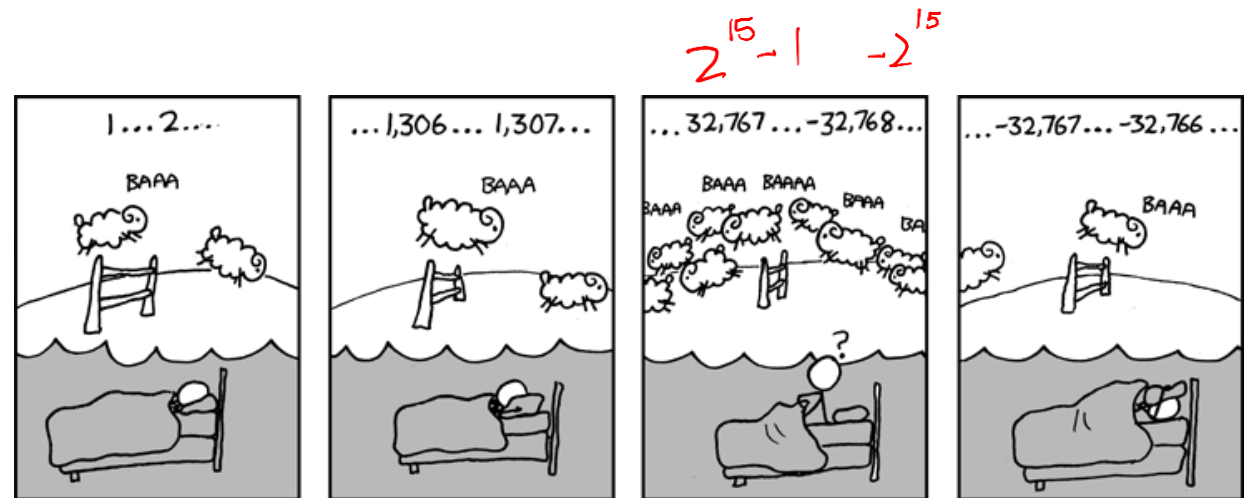
Kanishka Singh

Liander Rainbolt

Ming Yan

Pollux Chen

Soham Bhosale



<http://xkcd.com/571/>

Relevant Course Information

- ❖ HW3 due tonight, HW4 due Monday, HW5 due Wednesday
- ❖ Lab 1a due Monday (10/9)
 - Use `pctest` and `d1c.py` to check your solution for correctness (on the CSE Linux environment)
 - Submit `pointer.c` and `lab1Asynthesis.txt` to Gradescope
 - Make sure you pass the File and Compilation Check – all the correct files were found and there were no compilation or runtime errors
- ❖ Lab 1b released today, due 10/16
 - Bit manipulation on a custom encoding scheme
 - Bonus slides at the end of today's lecture have relevant examples
- ❖ Reading 6 is dense, do it early if you can!

Runnable Code Snippets on Ed

- ❖ Ed allows you to embed runnable code snippets (*e.g.*, readings, homework, discussion)
 - These are *editable* and *rerunnable*!
 - Hides compiler warnings, but will show compiler errors and runtime errors
 - Code must be inside of an `int main()` function
 - To use `printf()`, you must `#include<stdio.h>`
- ❖ Suggested use
 - Good for experimental questions about basic behaviors in C
 - *NOT* entirely consistent with the CSE Linux environment, so should not be used for any lab-related work

Lecture Outline (1/4)

- ❖ **Integer Limitations**
- ❖ Casting in C
- ❖ Bit Shifting
- ❖ Integer Representation Issues in Real Life

Integer Limits for w bits (Review)

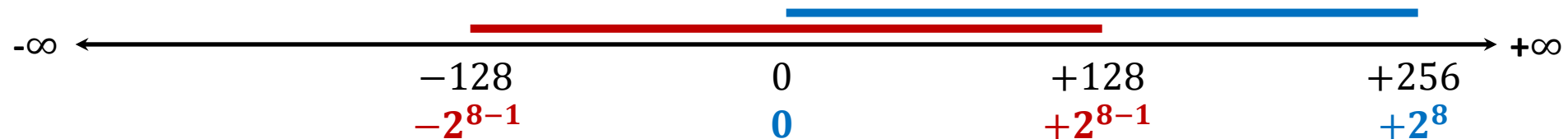
❖ Unsigned range

- $U_{\text{Min}} = 0b00\dots0 = 0$
- $U_{\text{Max}} = 0b11\dots1 = 2^w - 1$

❖ Signed (Two's Complement) values

- $T_{\text{Min}} = 0b10\dots0 = -2^{w-1}$
- $T_{\text{Max}} = 0b01\dots1 = 2^{w-1} - 1$

❖ Example: $w = 8$ (e.g., char)



Integer Arithmetic

❖ The same addition procedure works for both unsigned and signed (Two's Complement) integers

- **Simplifies hardware:** Only one algorithm for addition!
- **Algorithm:** Normal binary addition, **discard the highest carry bit**
 - Called modular addition: result is sum *modulo* 2^w

❖ 4-bit Examples: (HW = hardware, US = unsigned, TC = signed)

HW	US	TC
0100	4	4
+0011	3	3
= 0111	7	7

✓ ✓

HW	US	TC
1100	12	-4
+0011	3	3
= 1111	15	-1

✓ ✓

HW	US	TC
1101	13	-3
+0100	4	4
= 1 0001	1	1

? ✓

Arithmetic Overflow

Bits	Unsigned	Signed
0000	0 <i>U_{Min}</i>	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7 <i>T_{Max}</i>
1000	8	-8 <i>T_{Min}</i>
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15 <i>U_{Max}</i>	-1

- ❖ When a calculation produces a result that can't be represented in the current encoding scheme
 - Integer range limited by fixed width *U_{Min} - U_{Max}*
T_{Min} - T_{Max}
 - Can occur in both the positive and negative directions
- ❖ C and Java ignore overflow exceptions
 - You end up with a bad value in your program and no warning/indication... oops!

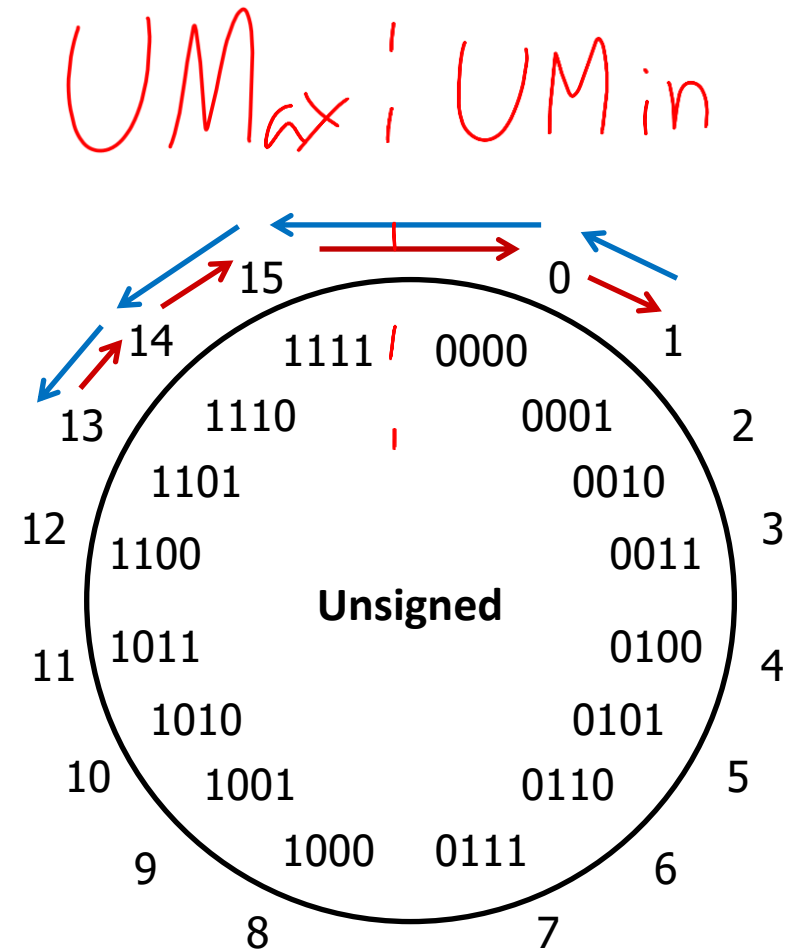
Overflow: Unsigned

- ❖ **Addition:** drop carry bit (-2^w)

13	1101
+ 4	+ 0100
<u>17</u>	<u>10001</u>
17	1 0001
1	

- ❖ **Subtraction:** borrow ($+2^w$)

1	10001
- 4	- 0100
<u>-3</u>	<u>1101</u>
-3	
13	



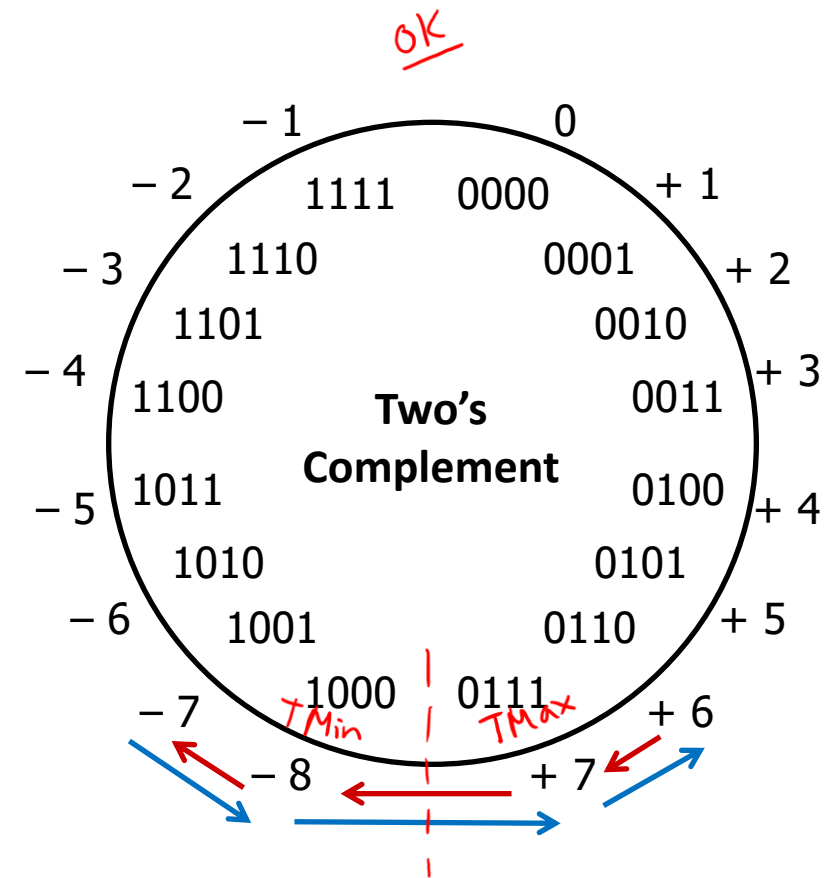
Overflow: Two's Complement

❖ **Addition:** $(+) + (+) = (-)$ result?

$$\begin{array}{r} 6 \qquad 0110 \\ + 3 \qquad + 0011 \\ \hline \cancel{9} \qquad 1001 \\ -7 \end{array}$$

❖ **Subtraction:** $(-) + (-) = (+)$?

$$\begin{array}{r} -7 \qquad 1001 \\ - 3 \qquad - 0011 \\ \hline \cancel{-10} \qquad 0110 \\ 6 \end{array}$$



Arithmetic Overflow Summary

- ❖ Error is always a multiple of $\pm 2^w$ because of modular arithmetic
 - *Unsigned overflow* occurs if result falls outside of [UMin, UMax]
 - There is a carryout from the MSB
 - *Signed overflow* occurs if result falls outside of [TMin, TMax]
 - Signs of both inputs to addition are the same, but the sign of the output is different
- ❖ *Independent* properties of the arithmetic operation
 - All four combinations of signed OF and unsigned OF are possible!

HW	US	TC
1101	13	-3
+0100	+ 4	+ 4
= 1 0001	= 1	= 1

✓ unsigned overflow
 ✗ signed overflow

Polling Questions (1/2)

- ❖ What is the value (and encoding) of TMin for a fictional 6-bit wide integer data type?

$$0b \frac{1}{-2^5} \frac{0}{2^4} \frac{0}{2^3} \frac{0}{2^2} \frac{0}{2^1} \frac{0}{2^0}$$

$$-2^{n-1} = -2^5 = \boxed{-32}$$

signed
most negative
represent $2^6 = 64$ numbers

- ❖ For the following 8-bit integer additions, did signed and/or unsigned overflow occur?

- [TMin, TMax] = [-128, 127]
- [UMin, UMax] = [0, 255]

Numeral	Signed	Unsigned
0x27	39	39
0xD9	-39	217
0x7F	127	127
0x81	-127	129

a) **0x27 + 0x81**

signed: $39 + (-127) = -88$
no signed overflow

unsigned: $39 + 129 = 168$
no unsigned overflow

b) **0x7F + 0xD9**

signed: $127 + (-39) = 88$
no signed overflow

unsigned: $127 + 217 = 344$
unsigned overflow

Lecture Outline (2/4)

- ❖ Integer Limitations
- ❖ **Casting in C**
- ❖ Bit Shifting
- ❖ Integer Representation Issues in Real Life

Data Types

- ❖ How does a data type affect a variable?
 - Size of allocation (*e.g.*, `char` vs. `long`)
 - How to interpret the bits (*e.g.*, `int` vs. `unsigned`)
 - Valid operators/operations and their behavior (*e.g.*, can't use subscript notation `[]` on integral types, normal vs. pointer arithmetic)
- ❖ What does it mean or what are the consequences of *changing* your data type?

Literals

- ❖ Constants/literals in your code also have “types”
 - Affect the stored/manipulated data and the behavior of operators
 - In C:
 - By default, literals (decimal or hex) are treated as *signed integers*
 - Use “U” (or “u”) suffix to explicitly force *unsigned* (e.g., `100U`, `4294967259u`)
 - Integer literals generally have an assumed size of 4 bytes unless longer is needed
 - We will learn about floating point literals next lesson
- ❖ Can be confusing if types don't match
 - Example: `signed char c = 255u; printf("%d", c);` *stored as 0xFF* *prints as -1*
 - Example: `int* ip = 0x40210 + 1;` *no pointer arithmetic!*
type mismatch error

Type Casting: Implicit (Review)

- ❖ **Casting** converts data of one data type into a different data type
 - Different programming languages may not allow casting or only in certain cases
- ❖ C is known for having very flexible casts, with different effects:
 - now* {
 - Changes in bit width (e.g., short to int)
 - Changes in interpretations (e.g., int to unsigned int, long int to char*)
 - Full changes in representations (e.g., int to float) *next lesson*
- ❖ An **implicit cast** is done automatically by the compiler to fix type mismatches
 - Needs to be a well-defined conversion between the two types
 - Examples: `int int_var = short_var;`, `printf("%c", short_var);`
↑ print as character

Type Casting: Explicit (Review)

- ❖ **Casting** converts data of one data type into a different data type
 - Different programming languages may not allow casting or only in certain cases
- ❖ C is known for having very flexible casts, with different effects:
 - Changes in bit width (*e.g.*, short to int)
 - Changes in interpretations (*e.g.*, int to unsigned int, long int to char*)
 - Full changes in representations (*e.g.*, int to float)
- ❖ An **explicit cast** can be performed by the programmer by using the syntax: **(data_type)expression**
 - Suppress compiler warnings for implicit casts
 - Forcibly cause changes in interpretation or representation

very powerful, very dangerous!

Casting: Bit Width Change (Mostly Review)

❖ Longer to shorter

- e.g., long → int → short → char
- **Truncation** (i.e., drop upper bytes)

```
short s = 0x0351;  
char c = s;    // 0x51
```

❖ Shorter to longer

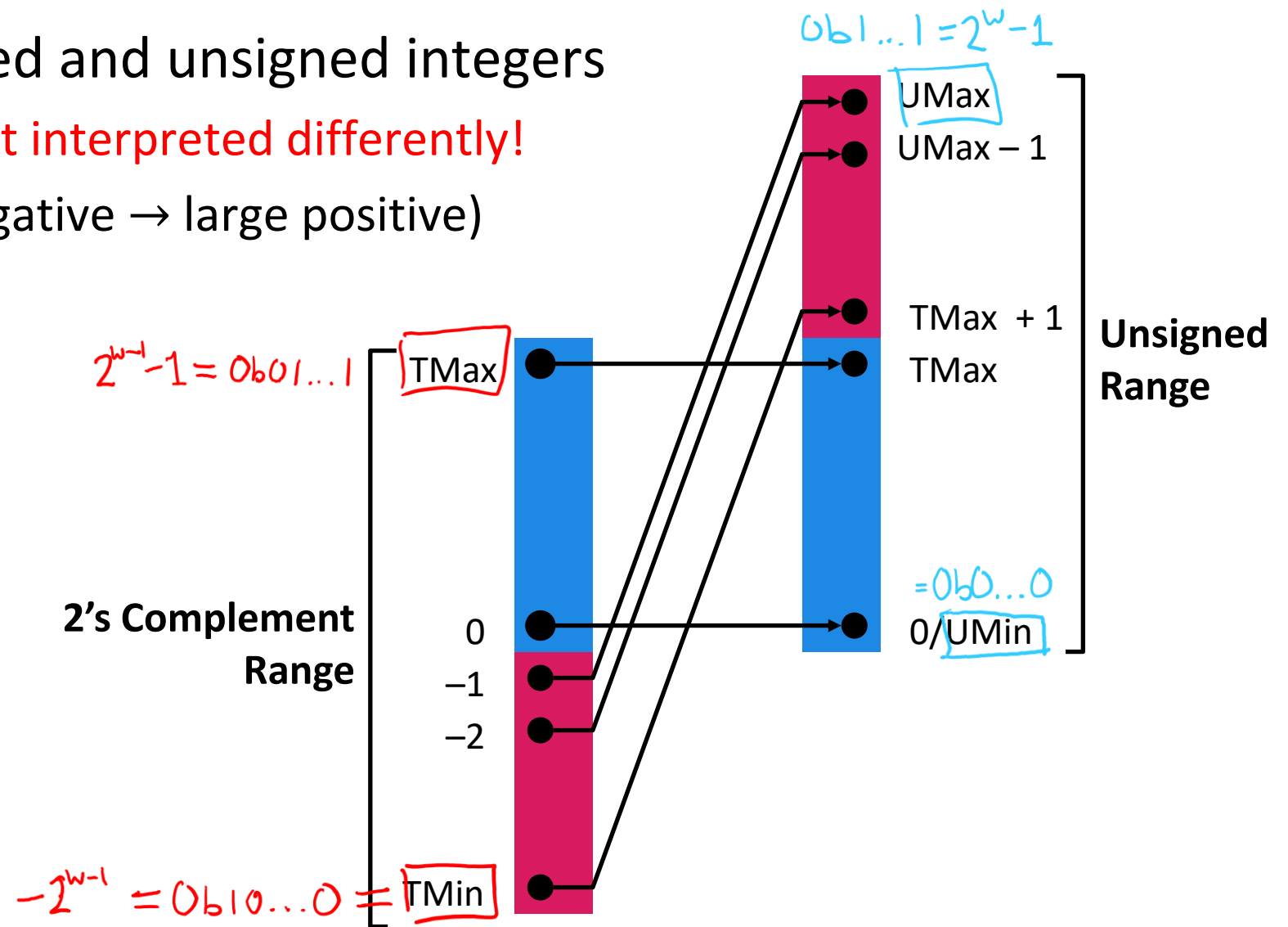
- e.g., char → short → int → long
- **Zero extension:** Add all zeros
 - In C, done for unsigned data
- **Sign extension:** Add all {old sign bit/MSB}
 - In C, done for signed data to preserve value

```
unsigned char uc = 0xFF;  
unsigned short us = uc;    // 0x00FF
```

```
signed char sc = 0xFF;  
short ss = sc;    // 0xFFFF  
sc = 0x10; ss = sc;    // 0x0010
```

Casting: Interpretation Change

- ❖ Casting between signed and unsigned integers
 - Bits are unchanged, just interpreted differently!
 - Ordering Inversion (negative \rightarrow large positive)



Data Types: Operator Behavior (Review)

❖ Expression Evaluation

- When you mixed unsigned and signed in a single expression, then **signed values are implicitly cast to unsigned**
- Including comparison operators $<$, $>$, $==$, $<=$, $>=$

(unsigned
"dominates")

❖ Examples: For 8-bit data, what will the following expressions evaluate to?

signed unsigned
■ $127 < 128u$
 $0b0111\ 1111$ $0b1000\ 0000$

unsigned comparison:

$0b0111\ 1111$ $<$ $0b1000\ 0000$
 $+127$ True $+128$

signed signed
■ $127 < (\text{signed char})\ 128u$
 $0b0111\ 1111$ $0b1000\ 0000$

signed comparison:

$0b0111\ 1111$ $<$ $0b1000\ 0000$
 $+127$ False -128

Lecture Outline (3/4)

- ❖ Integer Limitations
- ❖ Casting in C
- ❖ **Bit Shifting**
- ❖ Integer Representation Issues in Real Life

Shift Operations (Review, 1/2)

- ❖ Throw away (drop) extra bits that “fall off” the end
- ❖ Left shift ($x \ll n$) bit vector x by n positions
 - Fill with 0's on right
- ❖ Right shift ($x \gg n$) bit-vector x by n positions
 - Logical shift (for **unsigned** values)
 - Fill with 0's on left
 - Arithmetic shift (for **signed** values)
 - Replicate most significant bit on left (maintains sign of x)

Shift Operations (Review, 2/2)

digit $d_i \times 2^i$ changes power of 2 by n
because it moved positions

❖ Arithmetic:

- Left shift ($x \ll n$) is equivalent to multiply by 2^n
- Right shift ($x \gg n$) is equivalent to divide by 2^n
- Shifting is faster than general multiply and divide operations!

(compiler will try to optimize for you)

❖ Notes:

- Shifts by $n < 0$ or $n \geq w$ (w is bit width of x) are undefined behavior not guaranteed
- **In C:** behavior of \gg is determined by the compiler arithmetic/logical
 - In gcc / C lang, depends on data type of x (signed/unsigned)
- **In Java:** logical shift is \ggg and arithmetic shift is \gg

Left Shifting 8-bit Example

- ❖ No difference in left shift operation for unsigned and signed numbers (just manipulates bits)
 - Difference comes during interpretation: $x * 2^n$?

		Signed	Unsigned
$x = 25;$	00011001 =	25	25
$L1 = x \ll 2;$	00 01100100 =	100	100
$L2 = x \ll 3;$	000 11001000 =	-56	200
$L3 = x \ll 4;$	0001 10010000 =	-112	144

signed overflow
 unsigned overflow

200
 -256
 → 2^8
 400
 -256
 → 2^8

Logical Right Shifting 8-bit Example

- ❖ **Reminder:** C operator `>>` does *logical* shift on **unsigned** values and *arithmetic* shift on **signed** values

- **Logical** Shift: $x / 2^n$?

$$\begin{array}{llll}xu = 240u; & 11110000 & = & 240 \\ & \swarrow \swarrow \swarrow \swarrow \searrow \searrow & & /8 = 30 \\ R1u = xu >> 3; & 00011110 & = & 30 \\ & \text{(last two bits crossed out)} & & /4 = 7.5 \\ R2u = xu >> 5; & 00000111 & = & 7 \\ & \text{(last three bits crossed out)} & & \text{rounding (down)}\end{array}$$

Arithmetic Right Shifting 8-bit Negative Example

- ❖ **Reminder:** C operator `>>` does *logical* shift on **unsigned** values and *arithmetic* shift on **signed** values

- **Arithmetic** Shift: $x / 2^n$?

$x_s = -16;$ 11110000 $= -16$

$R1s = x_s >> 3;$ 11111110 $= -2$ $\frac{1}{4} = -0.5$

$R2s = x_s >> 5;$ 11111111 $= -1$

rounding (down)

Arithmetic Right Shifting 8-bit Positive Example

- ❖ **Reminder:** C operator `>>` does *logical* shift on **unsigned** values and *arithmetic* shift on **signed** values

- **Arithmetic** Shift: $x / 2^n$?

$x_s = 112;$ 01110000 $= 112$

$R3s = x_s >> 3;$ 00001110 $= 14$ $\frac{14}{4} = 3.5$

$R4s = x_s >> 5;$ 00000011 $= 3$

rounding (down)

Polling Questions (2/2)

- ❖ For unsigned char `uc = 0xA1;`, what are the produced data for the cast **(unsigned short)uc**?

2 bytes

unsigned \rightarrow zero extension

0x0DA1

- ❖ What is the result of the following expressions?

- **(signed char)uc >> 2**

- **(unsigned char)uc >> 3**

signed: 0b 1010 0001 $\xrightarrow{\text{arithmetic}}$ 0b 1110 1000 = 0xE8

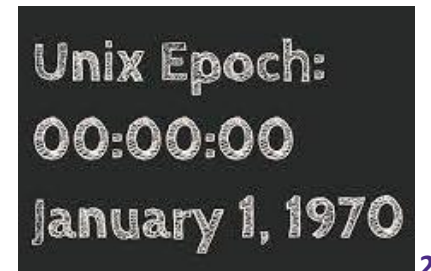
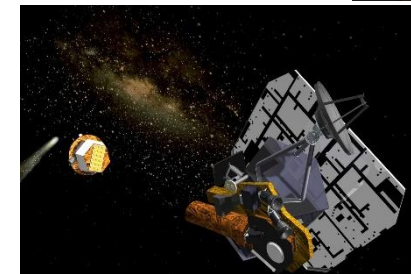
unsigned: 0b 1010 0001 $\xrightarrow{\text{logical}}$ 0b 0001 0100 = 0x14

Lecture Outline (4/4)

- ❖ Integer Limitations
- ❖ Casting in C
- ❖ Bit Shifting
- ❖ **Integer Representation Issues in Real Life**

Integer Representation Issues in Real Life

- ❖ **1985:** Therac-25 radiation therapy machine
 - Overdoses of radiation due to arithmetic overflow of incrementing a 1-byte safety flag variable
- ❖ **2000:** Y2K problem
 - Limited representation (two-digit decimal year)
- ❖ **2013:** Deep Impact spacecraft lost
 - Suspected integer overflow from storing time as tenth-seconds in unsigned int: 8/11/2013, 00:38:49.6
- ❖ **2038:** Unix epoch time rollover (seconds since 1/1/1970)
 - Signed 32-bit integer representation rolls over to TMin in 2038



Discussion Question

- ❖ Discuss the following question(s) in groups of 3-4 students
 - I will call on a few groups afterwards so please be prepared to share out
 - Be respectful of others' opinions and experiences

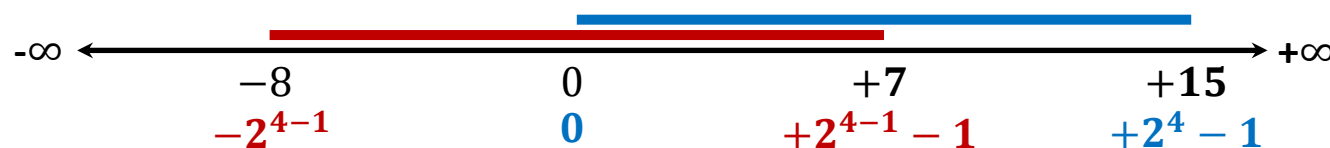
- ❖ Given that **arithmetic overflow** is a well-known property of integers in computing, what do you think are some of the *causes* and *pressures* that perpetuate these issues?
 - Think broadly! Ideas could be technical, economic, societal, etc.

Summary (1/3)

❖ We can only represent a limited range of numbers in w bits (2^w things)

■ Unsigned: [UMin, UMax]

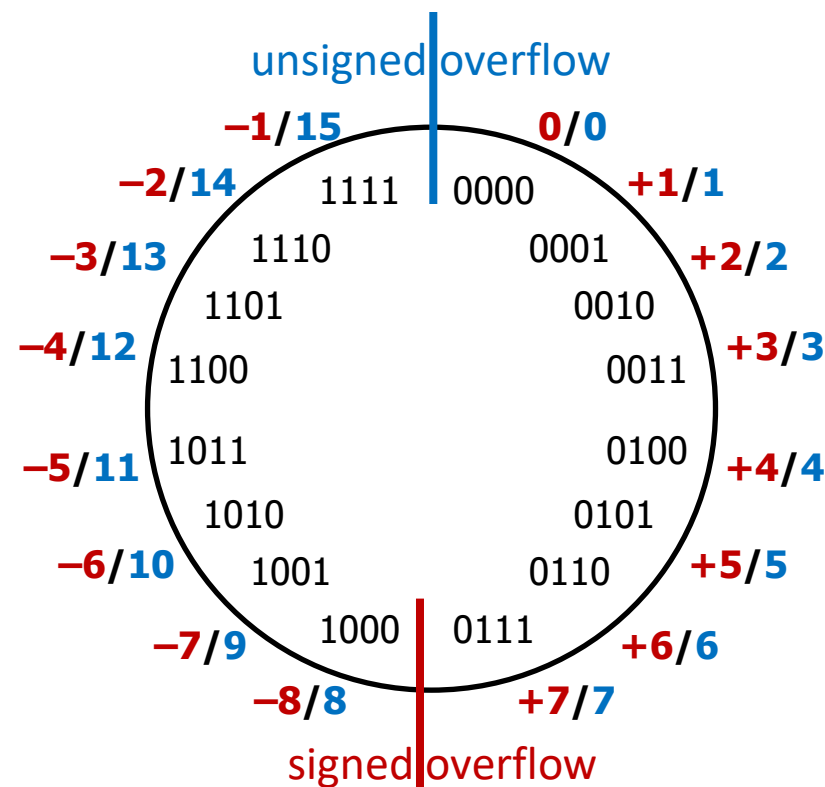
■ Signed: [TMin, TMax]



❖ Integer arithmetic is the same in hardware regardless of interpretation

■ When we exceed the limits, *arithmetic overflow* occurs following the rules of modular arithmetic

- Signed vs. unsigned overflow depends on interpretation of numbers:



Summary (2/3)

- ❖ Data types determine size, interpretations, and operator behaviors
- ❖ Casting (implicit or explicit) can convert values between different data types
 - Be careful of the possible consequences of casting (truncation, zero/sign extension, change in interpreted value, change in operator behaviors like comparisons and shifting)

```
int i = -1;  
long c = i;           // changed size (sign extension)  
unsigned int ui = i;  // changed interpretation  
  
// i < 1 (True) is different than ui < 1 (False)
```


Summary (3/3)

- ❖ Shifting is a useful bitwise operator
 - Throw away (drop) extra bits that “fall off” the end
 - Left shifting always fills with 0’s
 - Right shifting can be **arithmetic** (fill with copies of sign bit) or **logical** (fill with 0’s)
 - Shifts by $n < 0$ or $n \geq w$ (w is bit width) are *undefined*
- ❖ Common use cases: constant multiplication, bit masking
 - `x = x << 3; // equivalent to 8*x`
 - `x = (x >> 8) << 8; // zeros out lowest byte of x`

BONUS SLIDES

Some examples of using shift operators in combination with bitmasks, which you may find helpful for Lab 1b.

- ❖ Extract the 2nd most significant byte of an `int`
- ❖ Extract the sign bit of a signed `int`
- ❖ Conditionals as Boolean expressions

Using Shifts and Masks

❖ Extract the 2nd most significant *byte* of an `int`:

- First shift, then mask: $(x \gg 16) \& 0xFF$

x	000000001 000000010 000000011 000000100
x >> 16	000000000 000000000 000000001 000000010
0xFF	000000000 000000000 000000000 111111111
(x >> 16) & 0xFF	000000000 000000000 000000000 000000010

- Or first mask, then shift: $(x \& 0xFF0000) \gg 16$

x	000000001 000000010 000000011 000000100
0xFF0000	000000000 111111111 000000000 000000000
x & 0xFF0000	000000000 000000010 000000000 000000000
(x & 0xFF0000) >> 16	000000000 000000000 000000000 000000010

Using Shifts and Masks

- ❖ Extract the sign bit of a signed int:
 - First shift, then mask: $(x \gg 31) \& 0x1$
 - Assuming arithmetic shift here, but this works in either case
 - Need mask to clear 1s possibly shifted in

x	00000001 00000010 00000011 00000100
$x \gg 31$	00000000 00000000 00000000 00000000
0x1	00000000 00000000 00000000 00000001
$(x \gg 31) \& 0x1$	00000000 00000000 00000000 00000000

x	10000001 00000010 00000011 00000100
$x \gg 31$	11111111 11111111 11111111 11111111
0x1	00000000 00000000 00000000 00000001
$(x \gg 31) \& 0x1$	00000000 00000000 00000000 00000001

Using Shifts and Masks

❖ Conditionals as Boolean expressions

- For **int** x , what does $(x \ll 31) \gg 31$ do?

$x = !123$	00000000 00000000 00000000 00000000 1
$x \ll 31$	1 00000000 00000000 00000000 00000000
$(x \ll 31) \gg 31$	11111111 11111111 11111111 11111111
$!x$	00000000 00000000 00000000 00000000 0
$!x \ll 31$	0 00000000 00000000 00000000 00000000
$(!x \ll 31) \gg 31$	00000000 00000000 00000000 00000000

- Can use in place of conditional:
 - In C: `if(x) {a=y;} else {a=z;}` equivalent to `a=x?y:z;`
 - `a=(((!x<<31)>>31)&y) | (((!x<<31)>>31)&z);`