

Memory Allocation II

CSE 351 Winter 2024

Guest Lecturer:

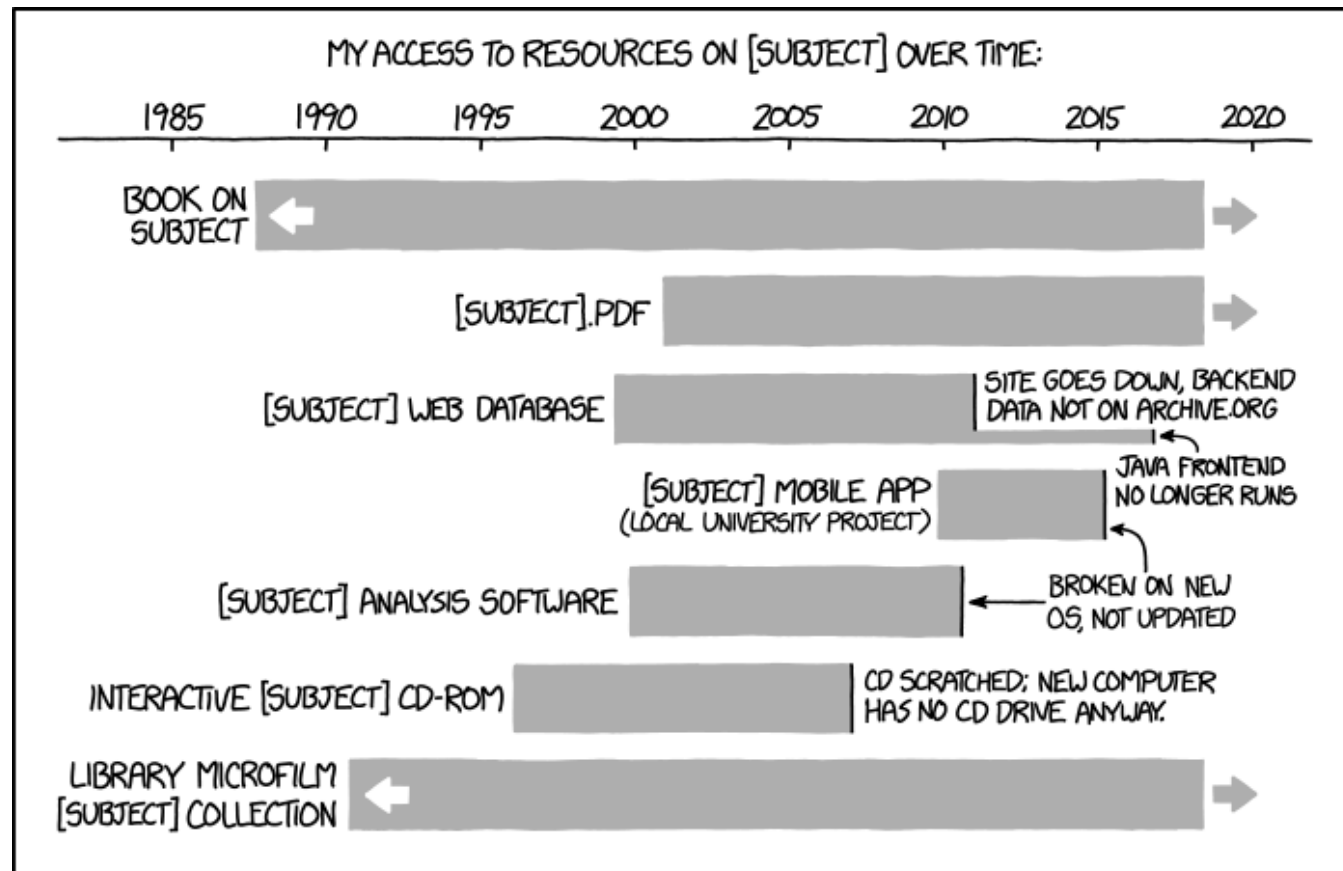
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IT'S UNSETTLING TO REALIZE HOW QUICKLY DIGITAL RESOURCES CAN DISAPPEAR WITHOUT ONGOING WORK TO MAINTAIN THEM.

<http://xkcd.com/1909/>

Relevant Course Information

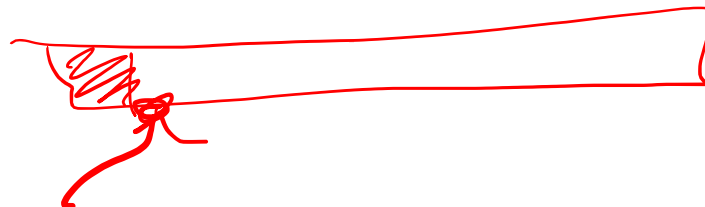
- ❖ HW18 due tonight!
- ❖ HW19 due Monday (2/26)
- ❖ HW20 due Wed (2/28)
 - Mostly a walk through of the heap simulator.
- ❖ Lab 4 due next Friday (3/1)
- ❖ Lab 5 (Mem Alloc) released!
 - Due Friday of finals week so a little over 2 weeks from now.

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions (purple, blue, yellow, green, red) representing different functional blocks and interconnects.

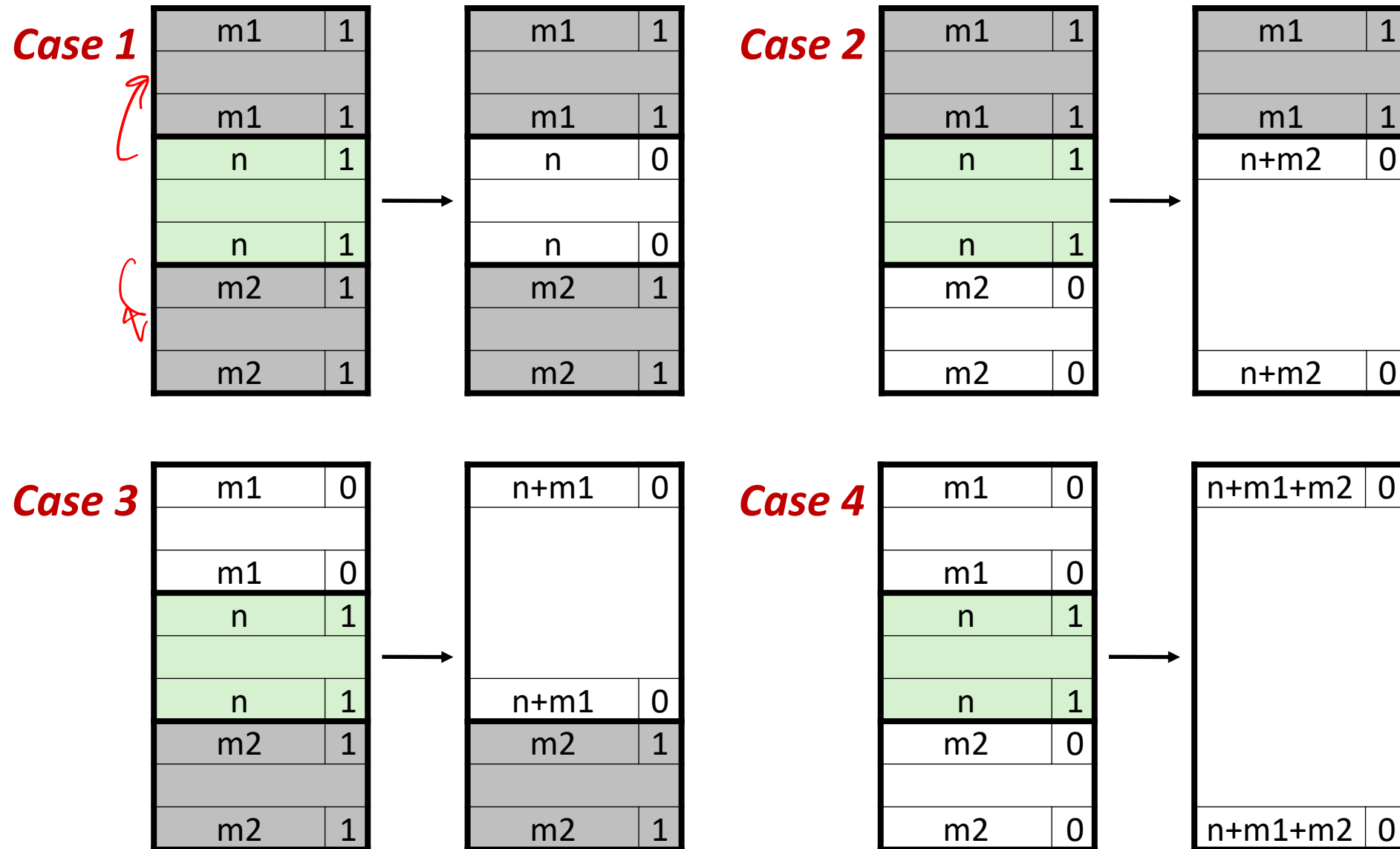
Memory Allocation II

Lesson Summary (1/4) - Fulfilling an Allocation Request

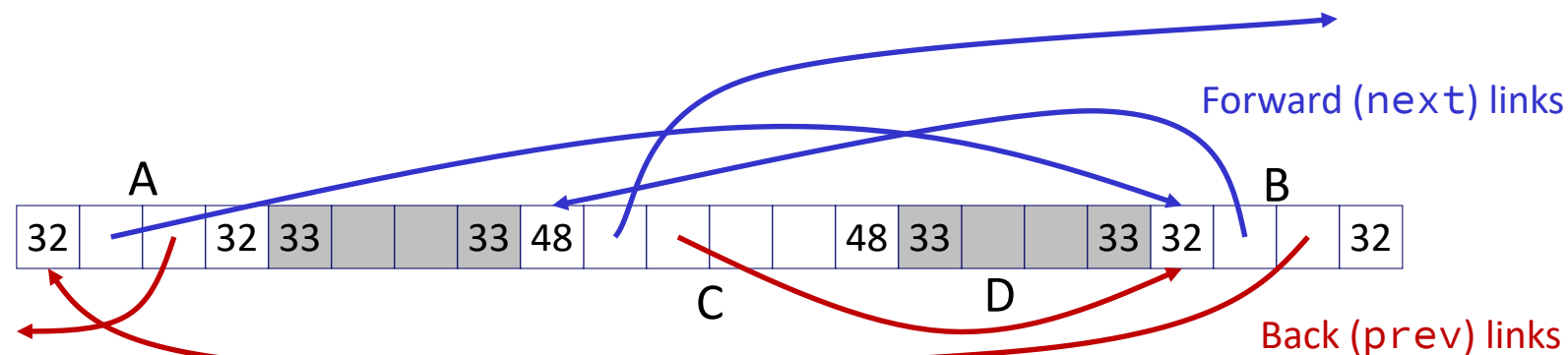
- 1) Compute the necessary block size *payload, padding, metadata*
- 2) Search for a suitable free block using the allocator's allocation strategy
 - If found, continue
 - If not found, return NULL *- first fit*
- next fit
- 3) Compare the necessary block size against the size of the chosen block *- best fit*
 - If equal, allocate the block
 - If not, split off the excess into a new free block before allocating the block
- 4) Return the address of the beginning of the payload



Lesson Summary (2/4) - Constant Time Coalescing



Lesson Summary (3/4) - Explicit List Summary

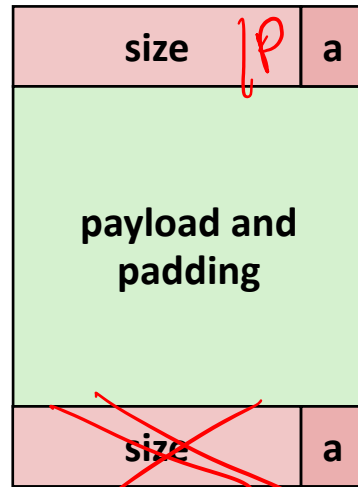


❖ Comparison with implicit list:

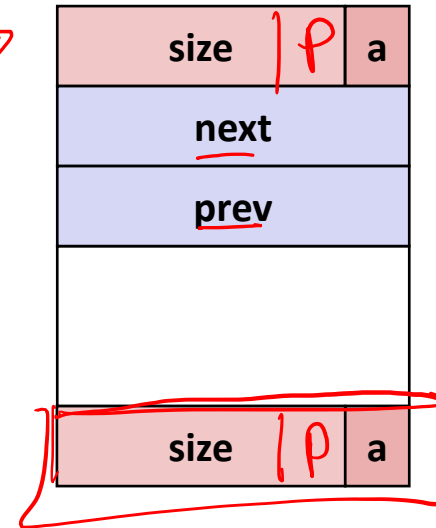
- Block allocation is linear time in number of free blocks instead of all blocks
 - Much faster when most of the memory is full
- Slightly more complicated allocate and free since we need to splice blocks in and out of the list
- Some extra space for the links (2 extra pointers needed for each free block)
 - Increases minimum block size, leading to more internal fragmentation

Lesson Summary (4/4) - Block Anatomy and Minimum Block Size

Allocated block:



Free block:



$$mbs = \text{Max}(\text{min}(\text{alloc}), \text{min}(\text{free}))$$

Lesson Q&A

❖ Learning Objectives:

- ■ Evaluate changes to the state of the heap for a sequence of allocations and deallocations.
 - Explain the tradeoffs between different allocator implementations, policies, and strategies.
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- ❖ What lingering questions do you have from the lesson?

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions (purple, blue, yellow, green, red) representing different functional blocks.

Memory Allocation II – Context

Allocation Policy Tradeoffs

- ❖ Data structure of blocks on lists
 - Implicit (free/allocated), explicit (free), segregated (many free lists) – others possible!
 - Cache implications (how tolerant are we to variable stride access patterns)
 - Alignment (*i.e.*, how many tags can we use in the header/footer)
- ❖ Placement policy: first-fit, next-fit, best-fit
 - Throughput vs. amount of fragmentation
- ❖ When does the allocator free allocated blocks?
 - Deferred coalescing

A detailed, colorful micrograph of a microchip die, showing a complex grid of circuitry and various colored regions (purple, blue, yellow, green, red) representing different functional blocks.

Memory Allocation II – Practice

Practice Question (1/2)

Determine the minimum block sizes (mbs) for the given memory allocators

*allocated blocks must have a payload size of at least 1

*boundary tags (headers and footers) are 8 bytes



Alignment	Allocated blocks	Free blocks	Free list type	<u>mbs alloc.</u>	<u>mbs free</u>	mbs
→ 8	header & footer $8 + 8 + 1 + 7$	header & footer $8 + 8$	<u>implicit</u>	? 24	? 16	? → 24
→ 8	header $8 + 1 \rightarrow 16$	header & footer $8 + 8 + 16$	<u>explicit</u>	? 16	? 32	? → 32
→ 16	header & footer $8 + 8 + 1$ $17 \rightarrow 32$	header & footer $8 + 8 + 16$ 32	explicit	? 32	? 32	? → 32

Practice Question (2/2)

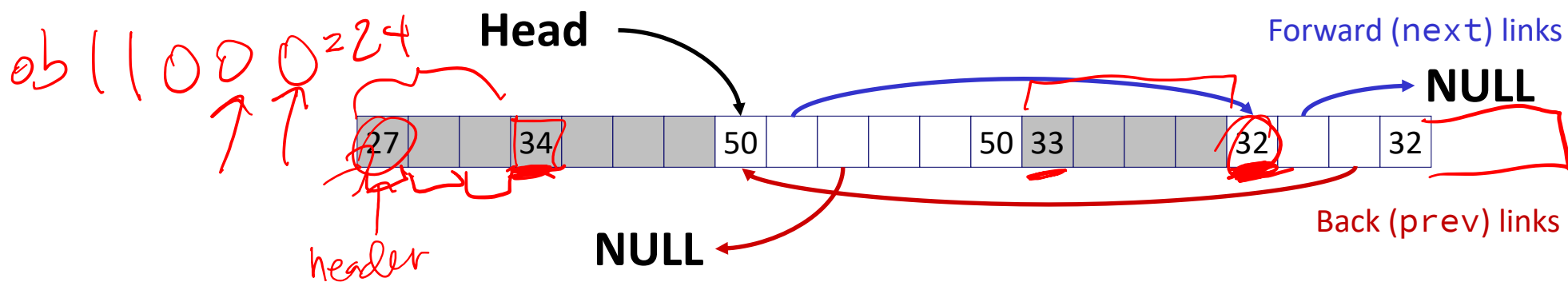
header = size | pre-used | is-allocated

Imagine we take a snapshot of the heap after a series of malloc and free calls. Come up with at least 2 issues/bugs with the below heap (there are 4)

- Explicit free list
 - Alignment: 8
 - Boundary tags: 8 bytes
 - Allocated blocks: header
 - Free blocks: header, footer
- mbs = 32*

□ = 8 bytes

1. $0b100010 \leftarrow$
 $0b100011 = 35$
2. $size(\text{block \#1}) < mbs$
3. $0b100000$
 $0b100010 \rightarrow 34$



4. no term. block at end of heap
- pl1*

Group Work Time

- ❖ During this time, you are encouraged to work on the following:
 - 1) If desired, continue your discussion
 - 2) Work on the homework problems
 - 3) Work on the current lab

- ❖ Resources:
 - You can revisit the lesson material
 - Work together in groups and help each other out
 - Course staff will circle around to provide support