

Java and C (part II)

CSE 351 Autumn 2024

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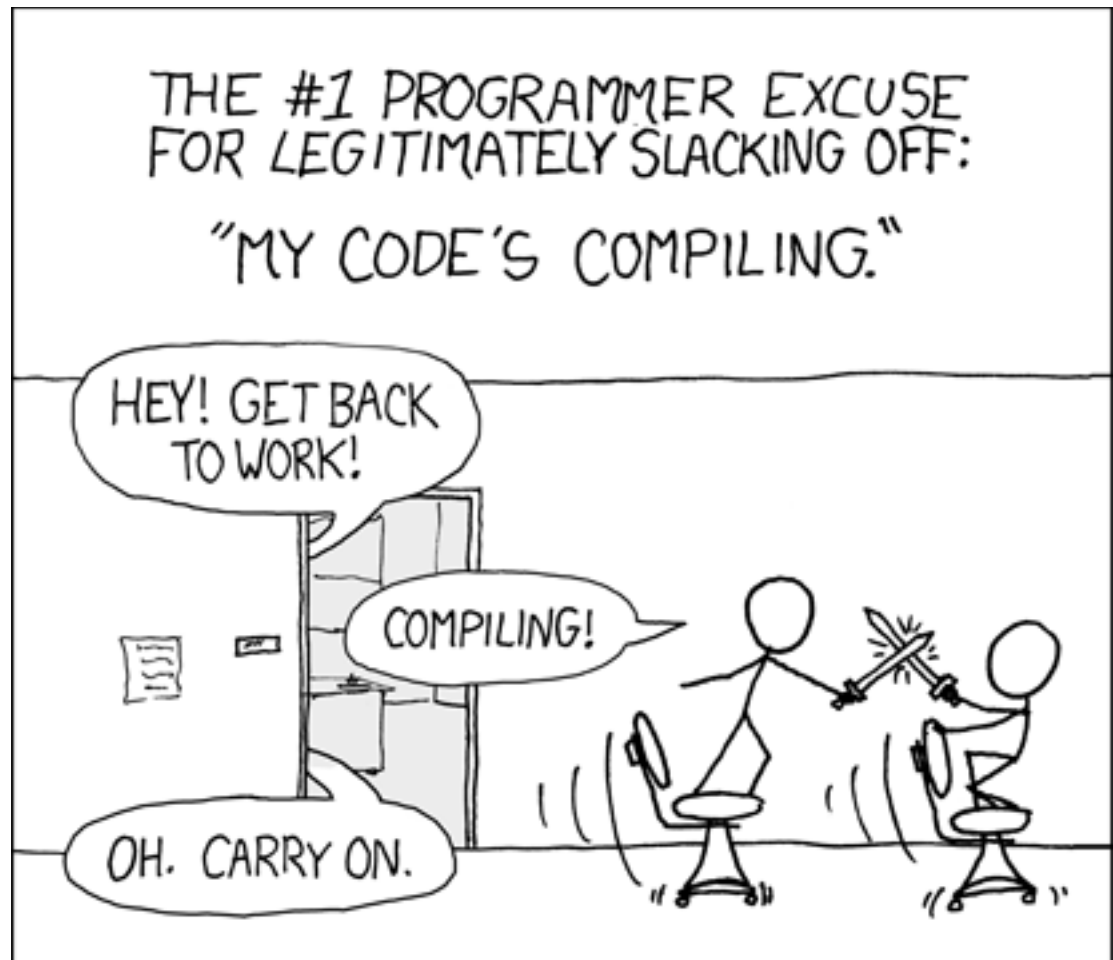
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<https://xkcd.com/303>

Relevant Course Information

- ❖ HW26 due Wednesday (12/04) @ 11:59 pm
- ❖ Final Exam Review in Section tomorrow! (12/05)
- ❖ Lab 5 (on Mem Alloc) due Thurs (12/05) @ 11:59pm
 - Closes Sunday 12/08 @11:59pm
- ❖ OPTIONAL HW on Java posted (for practice only)
- ❖ Final Exam, on Gradescope
 - Released Monday 12/09 at 12:01am
 - Due Wednesday 12/11 at 11:59pm
- ❖ Course evaluations now open - Please fill these out!
 - Separate ones for Lecture and Section

Polling Question

What would you expect to be the order of contents in an instance of the Car class?

Vote in Ed Lessons

```
class Vehicle {  
    int passengers;  
    // methods not shown  
}  
class Car extends Vehicle {  
    int wheels;  
    // methods not shown  
}
```

- A. header, Vehicle vtable ptr, passengers, Car vtable ptr, wheels
- B. Vehicle vtable ptr, passengers, wheels
- C. header, Vehicle vtable ptr, Car vtable ptr, passengers, wheels
- D. header, Car vtable ptr, passengers, wheels
- E. We're lost...

Roadmap

C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
c.setGals(17);  
float mpg =  
    c.getMPG();
```

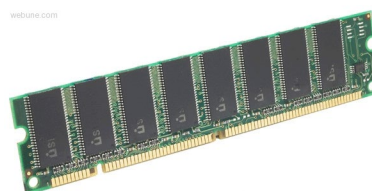
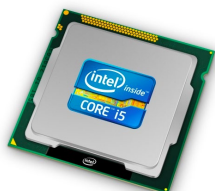
Assembly
language:

```
get_mpg:  
    pushq    %rbp  
    movq     %rsp, %rbp  
    ...  
    popq     %rbp  
    ret
```

Machine
code:

```
0111010000011000  
100011010000010000000010  
1000100111000010  
110000011111101000011111
```

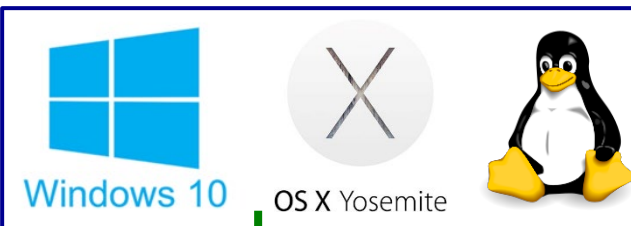
Computer
system:



Memory & data
Integers & floats
x86 assembly
Procedures & stacks
Executables
Arrays & structs
Memory & caches
Memory allocation
Processes
Virtual memory

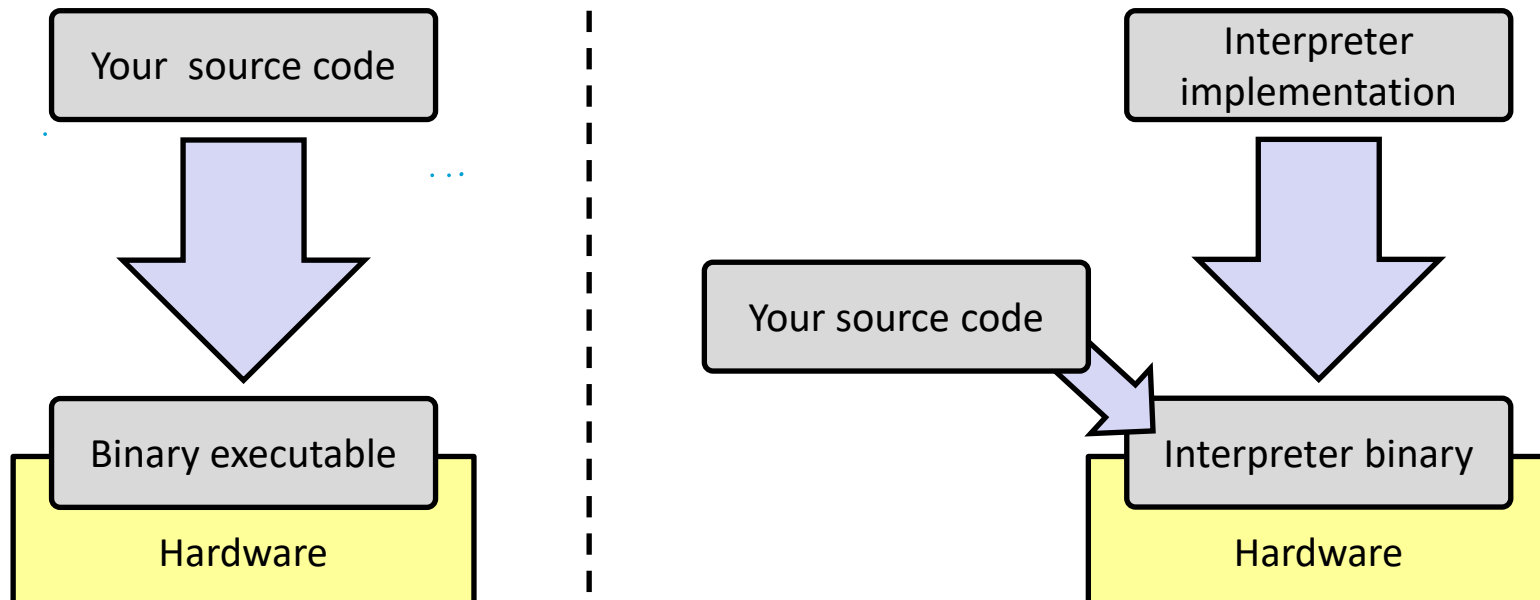
Java vs. C

OS:



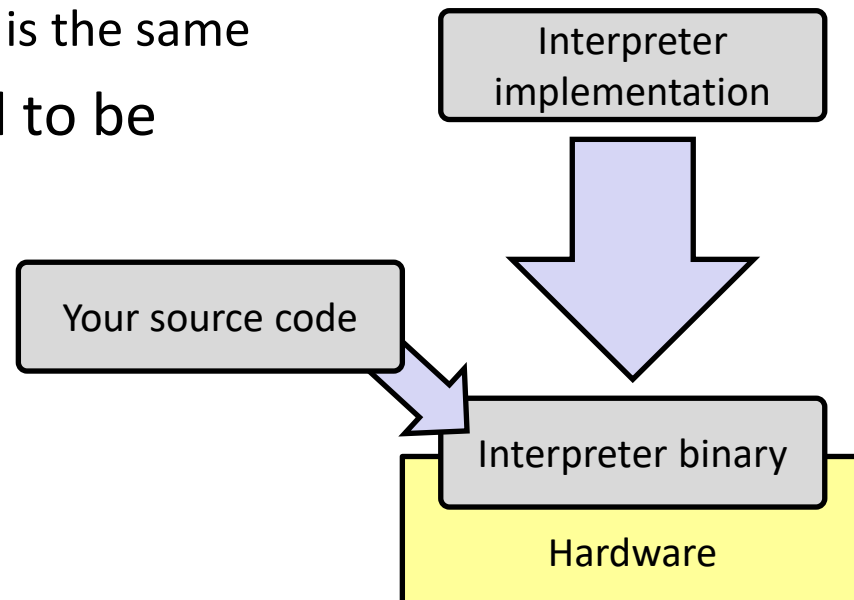
Implementing Programming Languages

- ❖ Many choices in programming model implementation
 - We've previously discussed compilation
 - One can also *interpret*
- ❖ **Interpreters** have a long history and are still in use
 - *e.g.*, Lisp, an early programming language, was interpreted
 - *e.g.*, Python, Javascript, Ruby, Matlab, PHP, Perl, ...



Interpreters

- ❖ Execute (something close to) the *source code* directly, meaning there is less translation required
 - This makes it a simpler program than a compiler and often provides more transparent error messages
- ❖ Easier to run on different architectures – runs in a simulated environment that exists only inside the *interpreter* process
 - Just port the interpreter (program), and then interpreting the source code is the same
- ❖ Interpreted programs tend to be slower to execute and harder to optimize



Interpreters vs. Compilers

- ❖ Programs that are designed for use with particular language implementations
 - You can choose to execute code written in a particular language via either a compiler or an interpreter, if they exist
- ❖ “Compiled languages” vs. “interpreted languages” a misuse of terminology
 - But very common to hear this
 - And has *some* validation in the real world (*e.g.*, JavaScript vs. C)
- ❖ Some modern language implementations are a mix
 - *e.g.*, Java compiles to bytecode that is then interpreted
 - Doing just-in-time (JIT) compilation of parts to assembly for performance

Compiling and Running Java

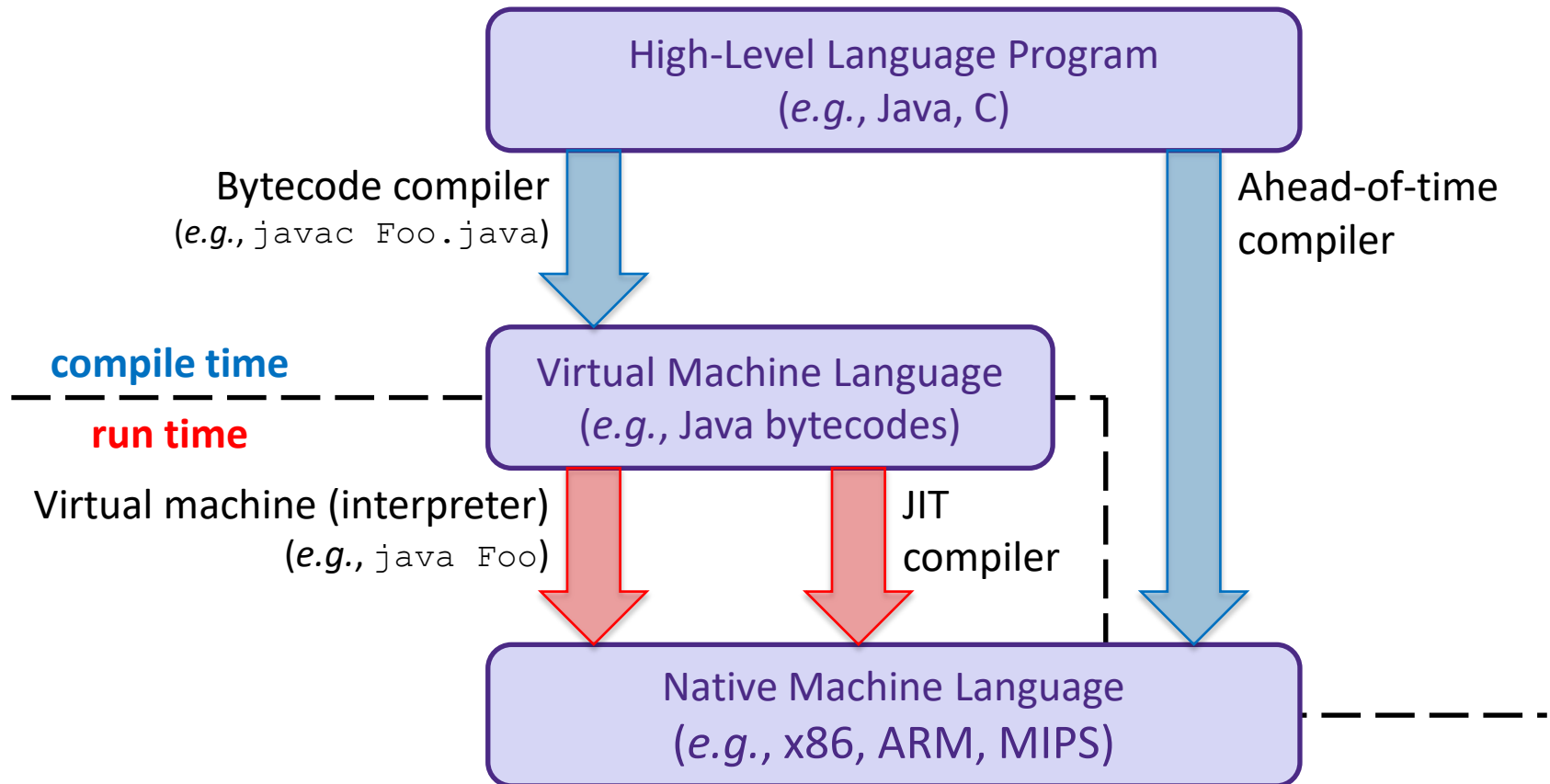
1. Save your Java code in a `.java` file
2. To run the Java compiler:
 - `javac Foo.java`
 - The Java compiler converts Java into *Java bytecodes*
 - Stored in a `.class` file
3. To execute the program stored in the bytecodes, these can be interpreted by the Java Virtual Machine (JVM)
 - Running the virtual machine: `java Foo`
 - Loads `Foo.class` and interprets the bytecodes

“The JVM”

Note: The JVM is different than the CSE VM running on VMWare. Yet *another* use of the word “virtual”!

- ❖ Java programs are usually run by a
Java *virtual machine* (JVM)
 - JVMs interpret an intermediate language called *Java bytecode*
 - Many JVMs compile bytecode to native machine code
 - **Just-in-time (JIT) compilation**
 - http://en.wikipedia.org/wiki/Just-in-time_compilation
 - Java is sometimes compiled ahead of time (AOT) like C

Virtual Machine Model



Polling Question – Answer in Ed Lessons

- ❖ You type **javac** and **java** at the command line.
You provide an argument to both commands.

javac:

- ❖ A) Is a: java source file/bytecode file/executable
- ❖ B) Its argument should refer to: _____
- ❖ C) It does this: _____


java:

- ❖ A) Is a java source file/bytecode file/executable
- ❖ B) Its argument should refer to _____
- ❖ C) It does this: _____

Java Bytecodes

- ❖ “Assembly code” for the Java Virtual Machine (JVM)
 - works on *all* JVMs
 - Hardware-independent! The JVM is just a program that has been compiled to run on this particular hardware
- ❖ Bytecodes are typed (unlike x86 assembly)

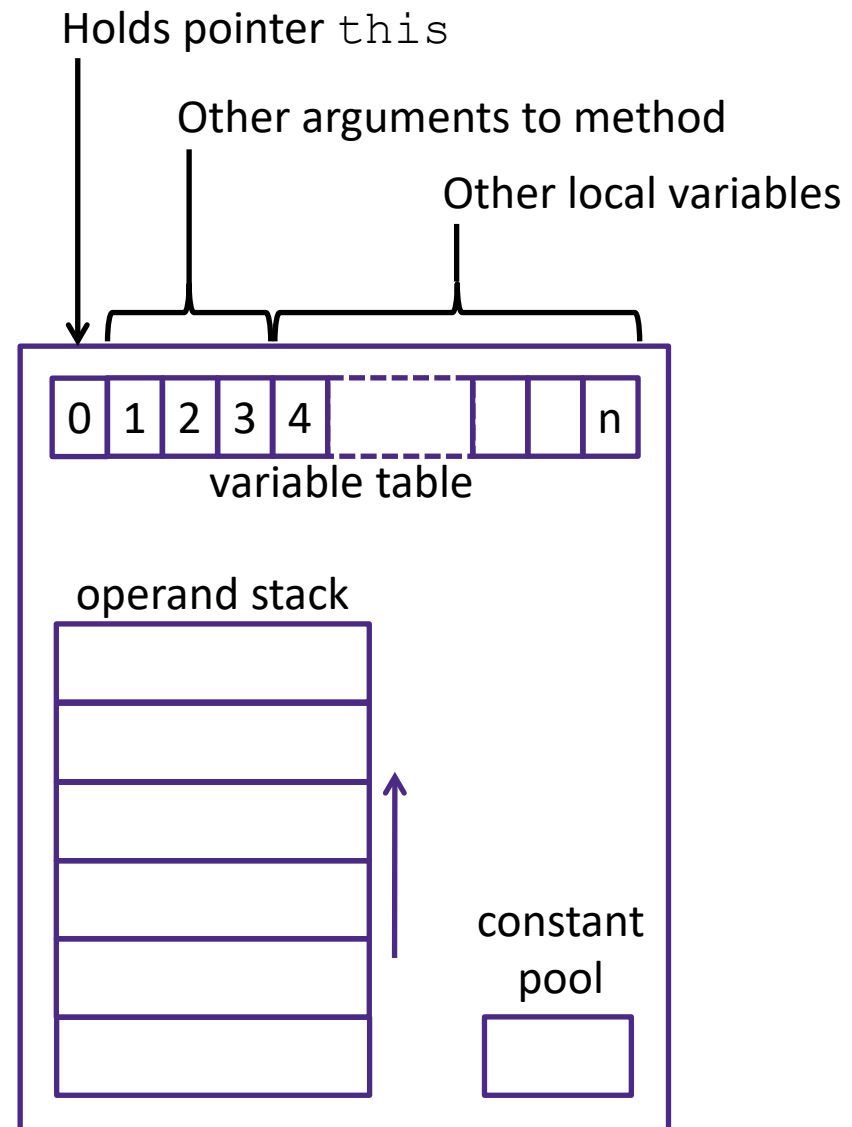
*'i' = integer,
'a' = reference,
'b' for byte,
'c' for char,
'd' for double, ...*



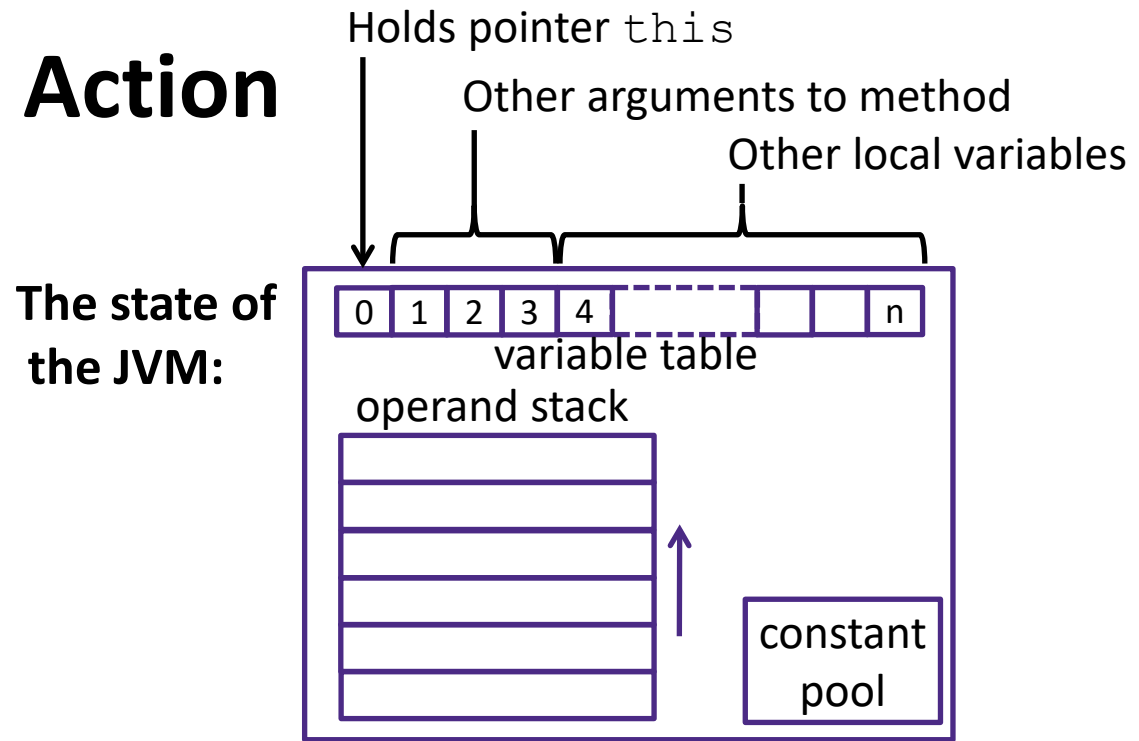
```
iload 1    // push 1st argument from variable table onto operand stack
iload 2    // push 2nd argument from variable table onto operand stack
iadd      // pop top 2 elements from operand stack, add together, and
            // push result back onto operand stack
istore 3   // pop result and put it into third slot in variable table
```

The Java Virtual Machine (JVM)

- ❖ Similar to how we described the state that x86 assembly instructions could modify: registers, memory, condition codes
- ❖ Java Bytecodes modify the state of the JVM: operand stack, variable table
- ❖ The state that x86 assembly modifies is **actual hardware**!
- ❖ The state that Java bytecodes modify is the **state of a program**!



Java Bytecode in Action



Java
Bytecode:

```

iload 1    // push 1st argument from table onto stack
iload 2    // push 2nd argument from table onto stack
iadd      // pop top 2 elements from stack, add together, and
            // push result back onto stack
istore 3   // pop result and put it into third slot in table
  
```

No registers or stack locations!
All operations use operand stack

Compiled
to (IA32) x86:

```

mov 8(%ebp), %eax
mov 12(%ebp), %edx
add %edx, %eax
mov %eax, -8(%ebp)
  
```

A Simple Java Method

```
Method java.lang.String getEmployeeName()
```

```
0 aload 0          // "this" object is stored at 0 in the var table

1 getfield #5 <Field java.lang.String name>
    // getfield instruction has a 3-byte encoding
    // Pop an element from top of stack, retrieve its
    //   specified instance field and push it onto stack
    // "name" field is the fifth field of the object

4 areturn          // Returns object at top of stack
```

Byte number: 0

1

4

aload_0	getfield	00	05	areturn
---------	----------	----	----	---------

As stored in the .class file:

2A	B4	00	05	B0
----	----	----	----	----

http://en.wikipedia.org/wiki/Java_bytecode_instruction_listings

Class File Format

- ❖ Every class in Java source code is compiled to its own class file
- ❖ 10 sections in the Java class file structure:
 - **Magic number:** 0xCAFEBAE (legible hex from James Gosling – Java’s inventor)
 - **Version of class file format:** The minor and major versions of the class file
 - **Constant pool:** Set of constant values for the class
 - **Access flags:** For example whether the class is abstract, static, final, etc.
 - **This class:** The name of the current class
 - **Super class:** The name of the super class
 - **Interfaces:** Any interfaces in the class
 - **Fields:** Any fields in the class
 - **Methods:** Any methods in the class
 - **Attributes:** Any attributes of the class (for example, name of source file, etc.)
- ❖ A `.jar` file collects together all of the class files needed for the program, plus any additional resources (e.g. images)

Disassembled Java Bytecode

```
> javac Employee.java
> javap -c Employee
```

http://en.wikipedia.org/wiki/Java_bytecode_instruction_listing

```
Compiled from Employee.java
class Employee extends java.lang.Object {
    public Employee(java.lang.String,int);
    public java.lang.String getEmployeeName();
    public int getEmployeeNumber();
}

Method Employee(java.lang.String,int)
0  aload_0
1  invokespecial #3 <Method java.lang.Object()>
4  aload_0
5  aload_1
6  putfield #5 <Field java.lang.String name>
9  aload_0
10 iload_2
11 putfield #4 <Field int idNumber>
14 aload_0
15 aload_1
16 iload_2
17 invokespecial #6 <Method void
    storeData(java.lang.String, int)>
20 return

Method java.lang.String getEmployeeName()
0  aload_0
1  getfield #5 <Field java.lang.String name>
4  areturn

Method int getEmployeeNumber()
0  aload_0
1  getfield #4 <Field int idNumber>
4  ireturn

Method void storeData(java.lang.String, int)
...
```

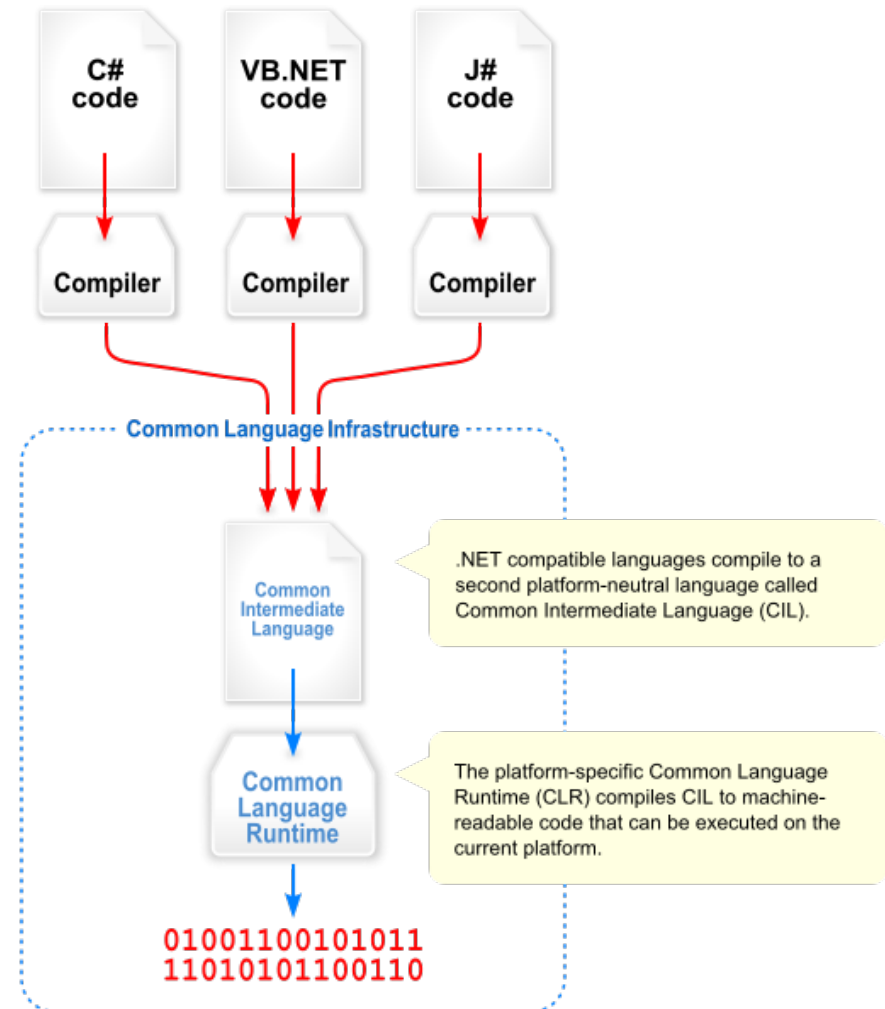
Other languages for JVMs

- ❖ JVMs run on so many computers that compilers have been built to translate many other languages to Java bytecode:
 - **AspectJ**, an aspect-oriented extension of Java
 - **ColdFusion**, a scripting language compiled to Java
 - **Clojure**, a functional Lisp dialect
 - **Groovy**, a scripting language
 - **JavaFX Script**, a scripting language for web apps
 - **JRuby**, an implementation of Ruby
 - **Jython**, an implementation of Python
 - **Rhino**, an implementation of JavaScript
 - **Scala**, an object-oriented and functional programming language
 - And many others, even including C!
- ❖ Originally, JVMs were designed and built for Java (still the major use) but JVMs are also viewed as a safe, GC'ed platform

Microsoft's C# and .NET Framework

❖ C# has similar motivations as Java

- Virtual machine is called the *Common Language Runtime*
- *Common Intermediate Language* is the bytecode for C# and other languages in the .NET framework



We made it! 🤪 😎 🤪

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c->miles = 100;
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float mpg = get_mpg(c);
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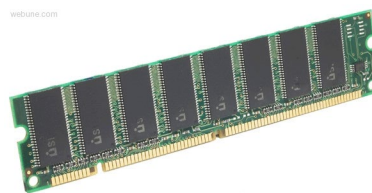
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