x86-64 Programming I

CSE 351 Winter 2022

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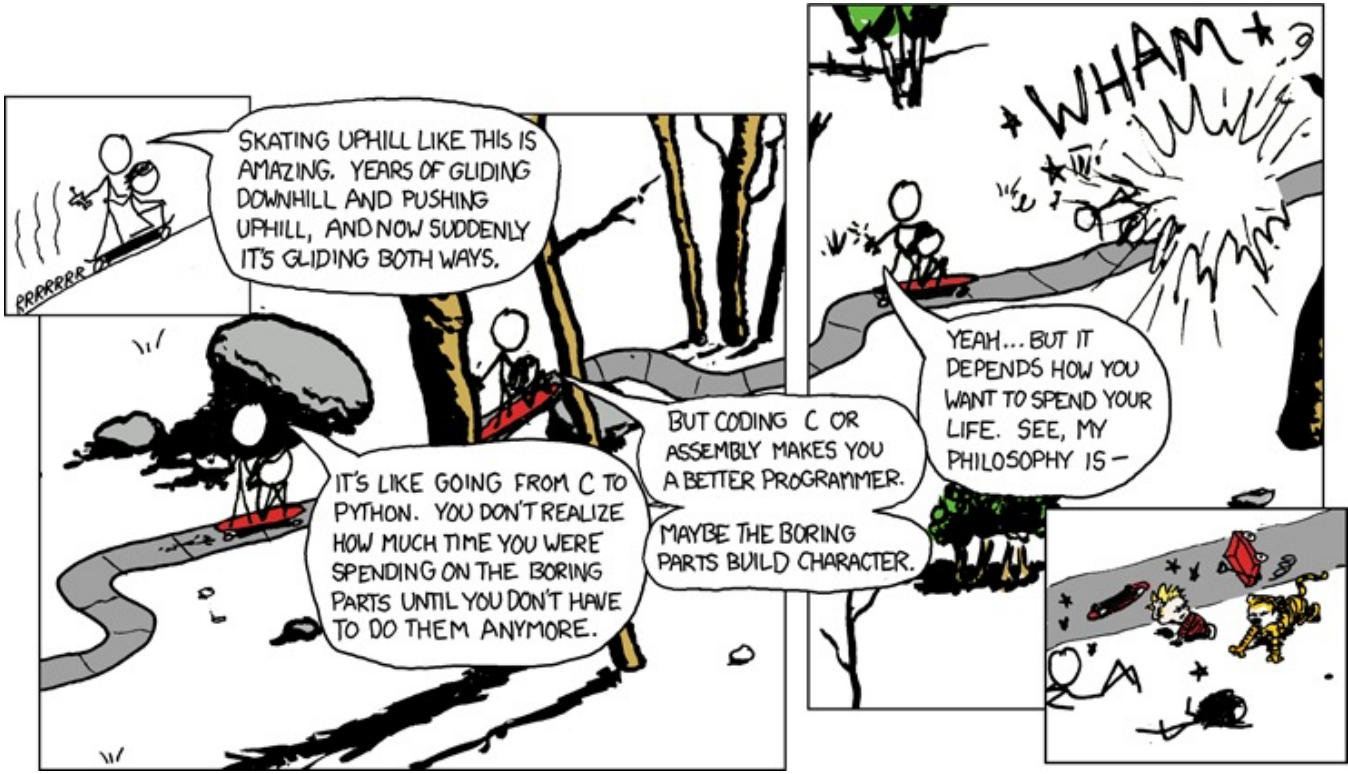


http://xkcd.com/409/

Relevant Course Information

- hw5 due tonight, hw6 due Friday
- Lab 1a due tonight!
  - Submit `pointer.c` and `lab1Asynthesis.txt`
    - Make sure there are no lingering `printf` statements in your code!
  - Make sure you submit *something* to Gradescope before the deadline and that the file names are correct
  - Can use late day tokens to submit up until Friday 11:59 pm

- `lateDays++`
  - Sorry for the inconvenience!
  - This brings you up to **6 late days total**

Relevant Course Information

- Lab 1b due Friday (10/18) at 11:59 pm
  - No major programming restrictions, but should avoid magic numbers by using C macros (`#define`)
  - For debugging, can use provided utility functions `print_binary_short()` and `print_binary_long()`
  - Pay attention to the output of `aisle_test` and `store_test` – failed tests will show you actual vs. expected

Floating Point Conversions in C

!!!

- Casting between `int`, `float`, and `double` **changes** the bit representation
  - `int` → `float`
    - May be rounded (not enough bits in mantissa: 23)
    - Overflow impossible
  - `int` or `float` → `double`
    - Exact conversion (all 32-bit `int`s are representable)
  - `long` → `double`
    - Depends on word size (32-bit is exact, 64-bit may be rounded)
  - `double` or `float` → `int`
    - Truncates fractional part (rounded toward zero)
    - "Not defined" when out of range or NaN: generally sets to TMin (even if the value is a very big positive)

Number Representation Really Matters

- **1991:** Patriot missile targeting error
  - clock skew due to conversion from integer to floating point
- **1996:** Ariane 5 rocket exploded ($1 billion)
  - overflow converting 64-bit floating point to 16-bit integer
- **2000:** Y2K problem
  - limited (decimal) representation: overflow, wrap-around
- **2038:** Unix epoch rollover
  - Unix epoch = seconds since 12am, January 1, 1970
  - signed 32-bit integer representation rolls over to TMin in 2038
- **Other related bugs:**
  - 1982: Vancouver Stock Exchange 10% error in less than 2 years
  - 1994: Intel Pentium FDIV (floating point division) HW bug ($475 million)
  - 1997: USS Yorktown "smart" warship stranded: divide by zero
  - 1998: Mars Climate Orbiter crashed: unit mismatch ($193 million)

Reading Review

- Terminology
  - Instruction Set Architecture (ISA): CISC vs. RISC
  - Instructions: data transfer, arithmetic/logical, control flow
    - Size specifiers: `b`, `w`, `l`, `q`
  - Operands: immediates, registers, memory
    - Memory operand: displacement, base register, index register, scale factor

- Questions from the Reading?

Review Questions

- Assume that the register `%rax` currently holds the value `0x 01 02 03 04 05 06 07 08`

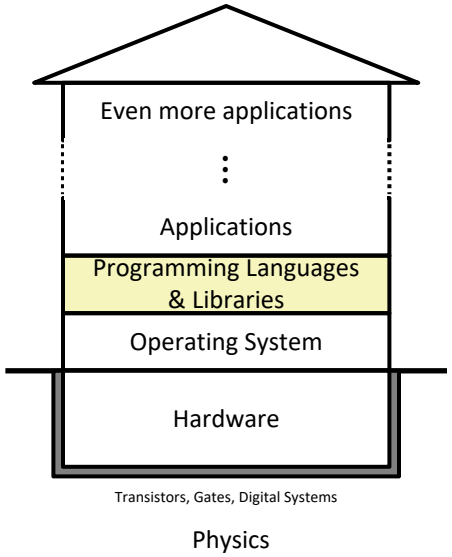
- Answer the questions about the following instruction (<instr> <src> <dst>):

```
xorw $-1, %ax
```

  - Operation type:
  - Operand types:
  - Operation width:
  - (extra) Result in `%rax`:

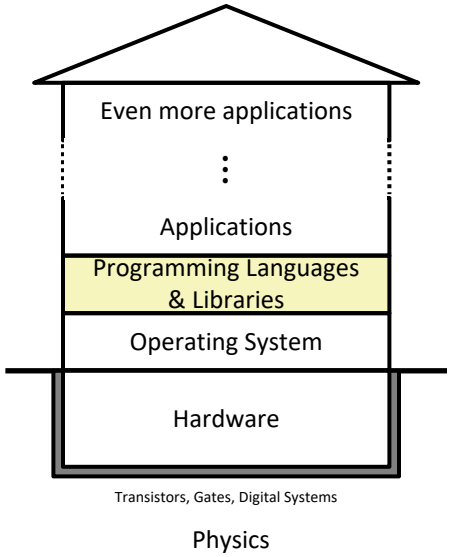
The Hardware/Software Interface

- Topic Group 1: **Data**
  - Memory, Data, Integers, Floating Point, Arrays, Structs
- Topic Group 2: **Programs**
  - **x86-64 Assembly**, Procedures, Stacks, Executables
- Topic Group 3: **Scale & Coherence**
  - Caches, Processes, Virtual Memory, Memory Allocation



The Hardware/Software Interface

- Topic Group 2: **Programs**
  - **x86-64 Assembly**, Procedures, Stacks, Executables



- How are programs created and executed on a CPU?
  - How does your source code become something that your computer understands?
  - How does the CPU organize and manipulate local data?

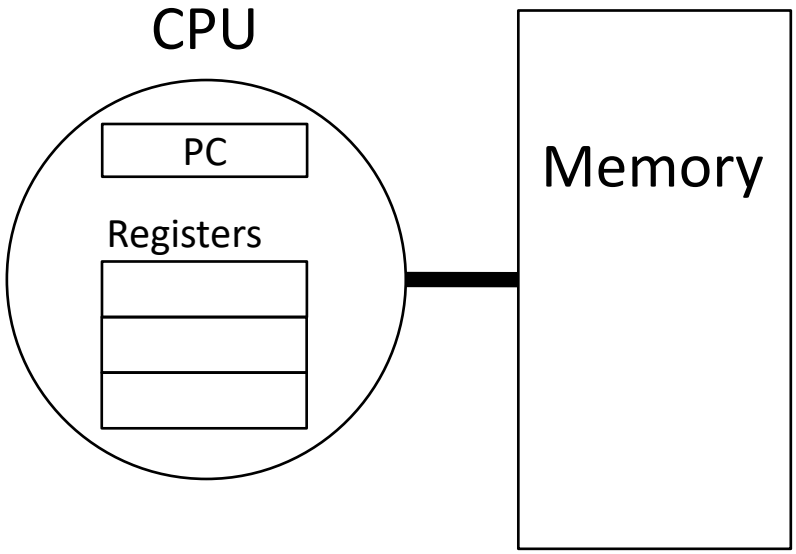
Definitions

- **Architecture (ISA):** The parts of a processor design that one needs to understand to write assembly code
  - What is directly visible to software
  - The "contract" or "blueprint" between hardware and software

- **Microarchitecture:** Implementation of the architecture
  - CSE/EE 469

Instruction Set Architectures (Review)

- The ISA defines:
  - The system's **state** (*e.g.*, registers, memory, program counter)
  - The **instructions** the CPU can execute
  - The **effect** that each of these instructions will have on the system state



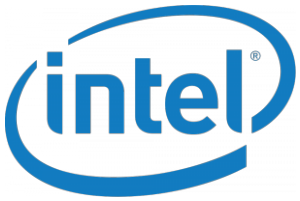
General ISA Design Decisions

- Instructions
  - What instructions are available? What do they do?
  - How are they encoded?
- Registers
  - How many registers are there?
  - How wide are they?
- Memory
  - How do you specify a memory location?

Instruction Set Philosophies (Review)

- *Complex Instruction Set Computing* (CISC):
  Add more and more elaborate and specialized instructions as needed
  - Lots of tools for programmers to use, but hardware must be able to handle all instructions
  - x86-64 is CISC, but only a small subset of instructions encountered with Linux programs
- *Reduced Instruction Set Computing* (RISC):
  Keep instruction set small and regular
  - Easier to build fast, less power-hungry hardware
  - Let software do the complicated operations by composing simpler ones

Mainstream ISAs



x86

Designer	Intel, AMD
Bits	16-bit, 32-bit and 64-bit
Introduced	1978 (16-bit), 1985 (32-bit), 2003 (64-bit)
Design	CISC
Type	Register–memory
Encoding	Variable (1 to 15 bytes)
Branching	Condition code
Endianness	Little

PCs, some Macs
(Core i3, i5, i7, M)
x86-64 Instruction Set

ARM

Designer	Arm Holdings
Bits	32-bit, 64-bit
Introduced	1985
Design	RISC
Type	Register-Register
Encoding	AArch64/A64 and AArch32/A32 use 32-bit instructions, T32 (Thumb-2) uses mixed 16- and 32-bit instructions; ARMv7 user-space compatibility.[1]
Branching	Condition code, compare and branch
Endianness	Bi (little as default)

Smartphone-like devices
(iPhone, iPad, Raspberry Pi)
M1 Macs (new!)
ARM Instruction Set

RISC-V

Designer	University of California, Berkeley
Bits	32 · 64 · 128
Introduced	2010
Design	RISC
Type	Load-store
Encoding	Variable
Endianness	Little[1][3]

Mostly research
(some traction in embedded)
RISC-V Instruction Set

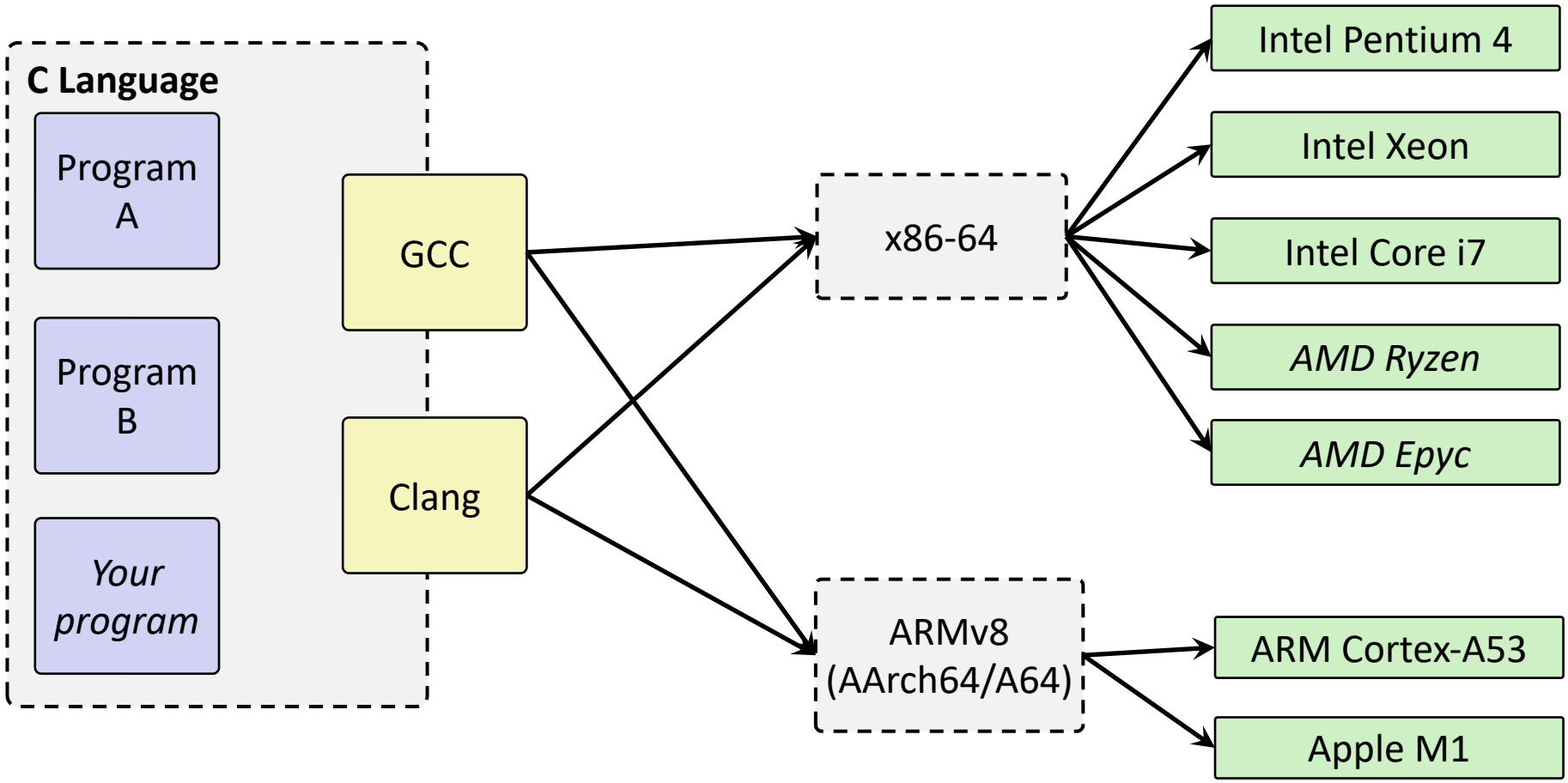
Architecture Sits at the Hardware Interface

Source code
Different applications or algorithms

Compiler
Perform optimizations, generate instructions

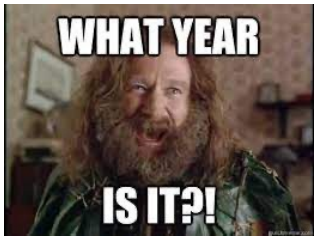
Architecture
Instruction set

Hardware
Different implementations

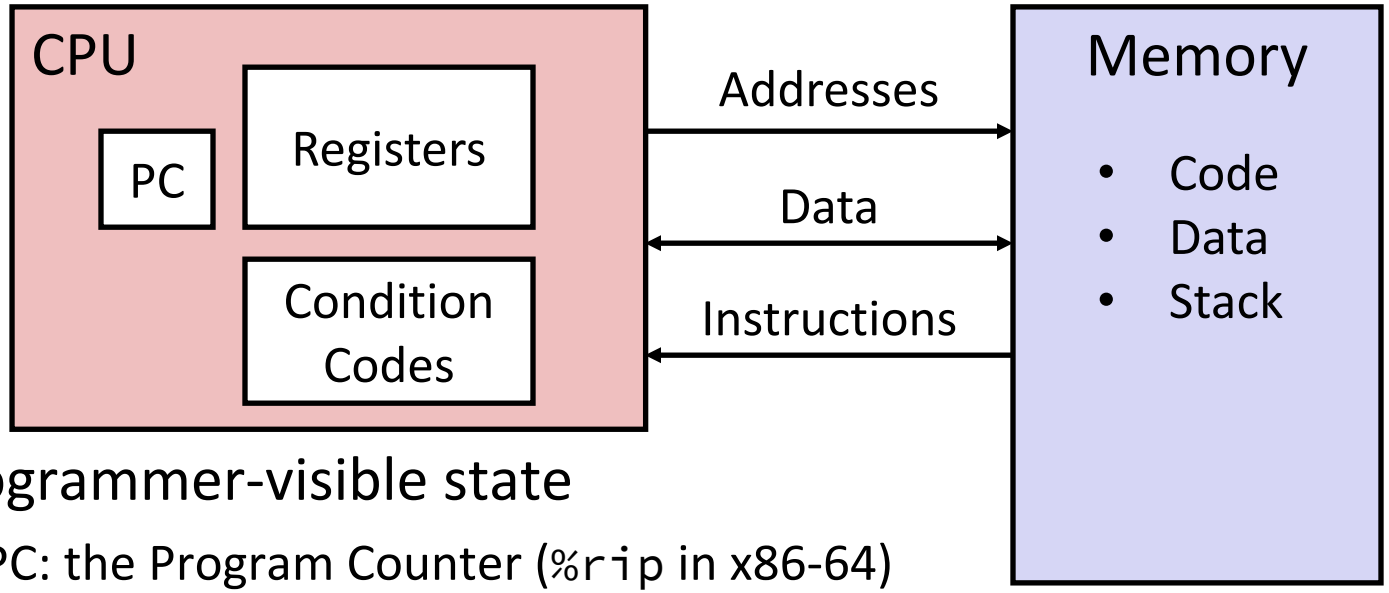


Writing Assembly Code? In $CURRENT_YEAR???

- Chances are, you'll never write a program in assembly, but understanding assembly is the key to the machine-level execution model:
  - Behavior of programs in the presence of bugs
    - When high-level language model breaks down
  - Tuning program performance
    - Understand optimizations done/not done by the compiler
    - Understanding sources of program inefficiency
  - Implementing systems software
    - What are the "states" of processes that the OS must manage
    - Using special units (timers, I/O co-processors, etc.) inside processor!
  - Fighting malicious software
    - Distributed software is in binary form



Assembly Programmer's View



- Programmer-visible state
  - PC: the Program Counter (`%rip` in x86-64)
    - Address of next instruction
  - Named registers
    - Together in "register file"
    - Heavily used program data
  - Condition codes
    - Store status information about most recent arithmetic operation
    - Used for conditional branching
- Memory
  - Byte-addressable array
  - Code and user data
  - Includes *the Stack* (for supporting procedures)

x86-64 Assembly "Data Types"

- Integral data of 1, 2, 4, or 8 bytes
  - Data values
  - Addresses
- Floating point data of 4, 8, 10 or 2x8 or 4x4 or 8x2
  - Different registers for those (*e.g.*, `%xmm1`, `%ymm2`)
  - Come from *extensions to x86* (SSE, AVX, …)

Not covered In 351

- No aggregate types such as arrays or structures
  - Just contiguously allocated bytes in memory
- Two common syntaxes
  - "AT&T": used by our course, slides, textbook, gnu tools, …
  - "Intel": used by Intel documentation, Intel tools, …
  - Must know which you're reading

Instruction Types (Review)

1) Transfer data between memory and register
   - *Load* data from memory into register
     - `%reg` = Mem[address]
   - *Store* register data into memory
     - Mem[address] = `%reg`

Remember: Memory is indexed just like an array of bytes!

2) Perform arithmetic operation on register or memory data
   - `c = a + b;` `z = x << y;` `i = h & g;`

3) Control flow: what instruction to execute next
   - Unconditional jumps to/from procedures
   - Conditional branches

Instruction Sizes and Operands (Review)

- Size specifiers
  - b = 1-byte "byte", w = 2-byte "word",
    l = 4-byte "long word", q = 8-byte "quad word"
  - Note that due to backwards-compatible support for 8086 programs (16-bit machines!), "word" means 16 bits = 2 bytes in x86 instruction names

- Operand types
  - ***Immediate:*** Constant integer data ($)
  - ***Register:*** 1 of 16 general-purpose integer registers (%)
  - ***Memory:*** Consecutive bytes of memory at a computed address (( ))

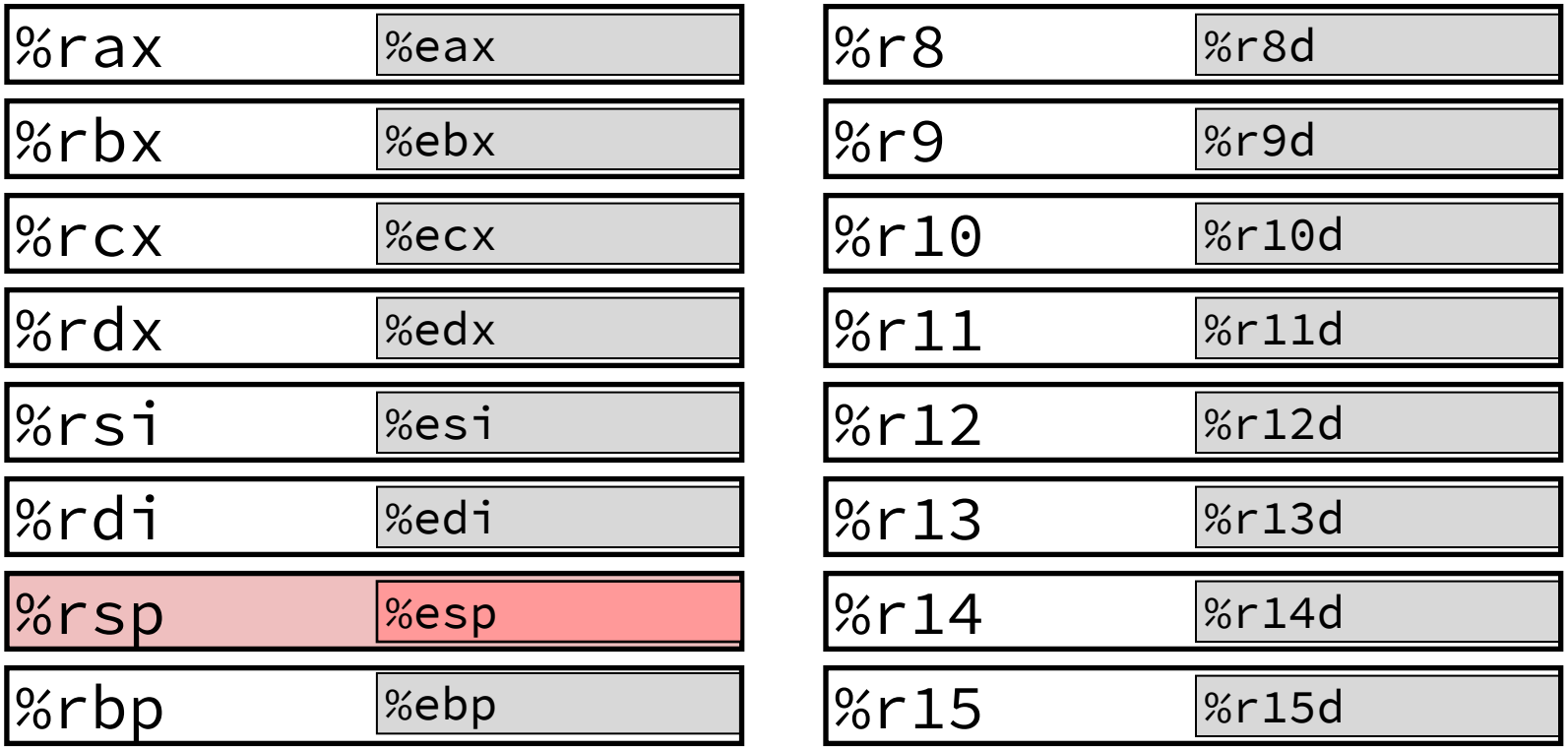
What is a Register? (Review)

- A location in the CPU that stores a small amount of data, which can be accessed very quickly (once every clock cycle)

- Registers have *names*, not *addresses*
  - In assembly, they start with `%` (*e.g.*, `%rsi`)

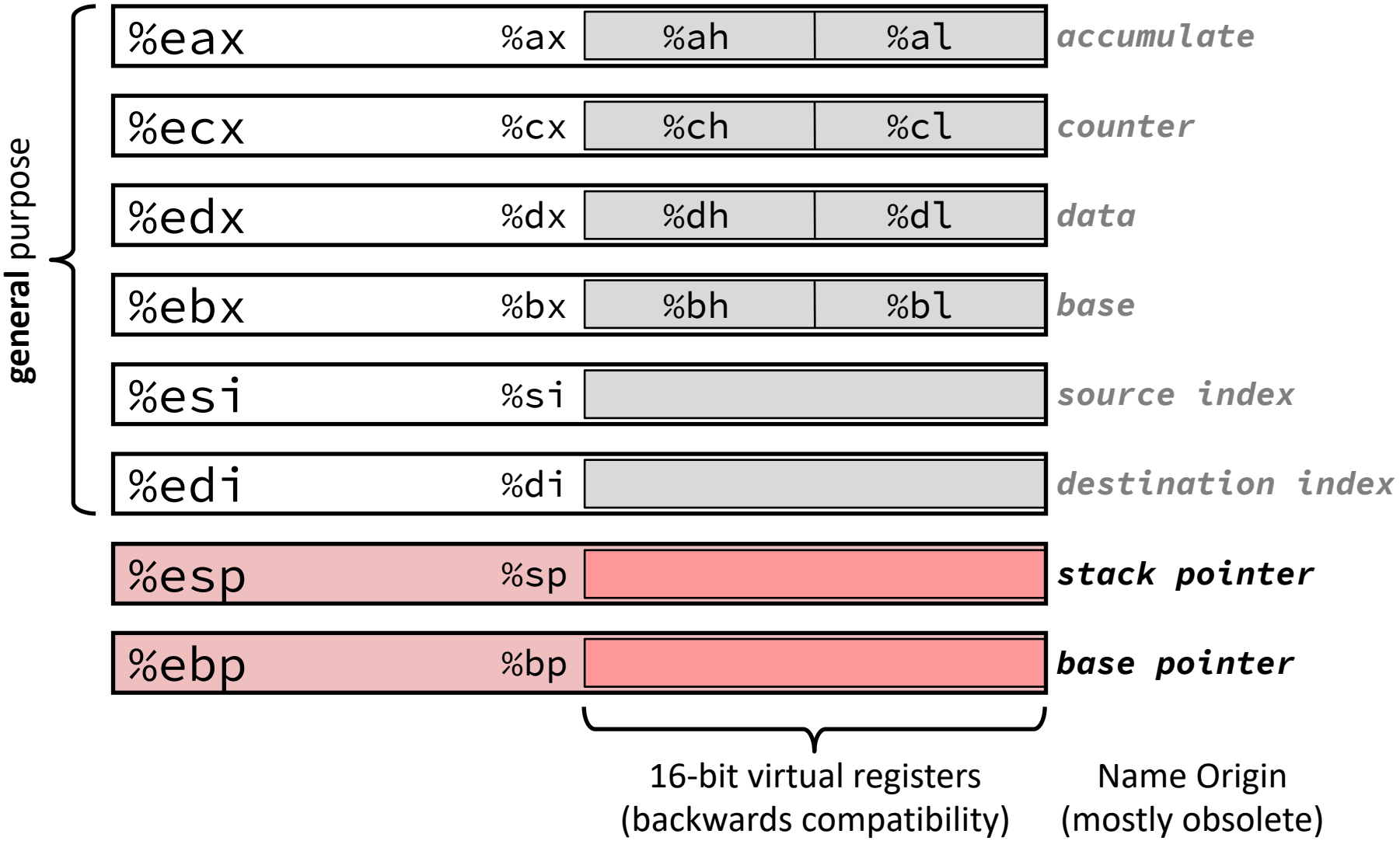
- Registers are at the heart of assembly programming
  - They are a precious commodity in all architectures, but *especially* x86-64

x86-64 Integer Registers – 64 bits wide



- Can reference low-order 4 bytes (also low-order 2 & 1 bytes)

Some History: IA32 Registers – 32 bits wide



Memory vs. Registers

- Addresses **vs.** Names
  - `0x7FFFD024C3DC` `%rdi`
- Big **vs.** Small
  - ~ 8 GiB (16 x 8 B) = 128 B
- Slow **vs.** Fast
  - ~50-100 ns sub-nanosecond timescale
- Dynamic **vs.** Static
  - Can "grow" as needed while program runs fixed number in hardware

Moving Data

- General form: `mov_ source, destination`
  - Really more of a "copy" than a "move"
  - Like all instructions, missing letter (_) is the size specifier
  - Lots of these in typical code

Operand Combinations

	Source	Dest	Src, Dest	C Analog
movq	Imm	Reg	`movq $0x4, %rax`	`var_a = 0x4;`
		Mem	`movq $-147, (%rax)`	`*p_a = -147;`
	Reg	Reg	`movq %rax, %rdx`	`var_d = var_a;`
		Mem	`movq %rax, (%rdx)`	`*p_d = var_a;`
	Mem	Reg	`movq (%rax), %rdx`	`var_d = *p_a;`

- *Cannot do memory-memory transfer with a single instruction*
  - How would you do it?

Some Arithmetic Operations

- Binary (two-operand) Instructions:
  - **Maximum of one memory operand**
  - Beware argument order!
  - No distinction between signed and unsigned
    - Only arithmetic vs. logical shifts

Format	Computation	
addq *src, dst*	*dst = dst + src*	(*dst += src*)
subq *src, dst*	*dst = dst – src*	
imulq *src, dst*	*dst = dst * src*	signed mult
sarq *src, dst*	*dst = dst >> src*	**A**rithmetic
shrq *src, dst*	*dst = dst >> src*	**L**ogical
shlq *src, dst*	*dst = dst << src*	(same as `salq`)
xorq *src, dst*	*dst = dst ^ src*	
andq *src, dst*	*dst = dst & src*	
orq *src, dst*	*dst = dst | src*	

operand size specifier

Practice Question

- Which of the following are valid implementations of `rcx = rax + rbx`?

  - ```
    addq %rax, %rcx
    addq %rbx, %rcx
    ```

  - ```
    movq %rax, %rcx
    addq %rbx, %rcx
    ```

  - ```
    movq $0, %rcx
    addq %rbx, %rcx
    addq %rax, %rcx
    ```

  - ```
    xorq %rax, %rax
    addq %rax, %rcx
    addq %rbx, %rcx
    ```

Arithmetic Example

Register	Use(s)
%rdi	1st argument (x)
%rsi	2nd argument (y)
%rax	return value

```
long simple_arith(long x, long y)
{
  long t1 = x + y;
  long t2 = t1 * 3;
  return t2;
}
```

```
y += x;
y *= 3;
long r = y;
return r;
```

```
simple_arith:
  addq    %rdi, %rsi
  imulq     $3, %rsi
  movq    %rsi, %rax
  ret
```

Example of Basic Addressing Modes

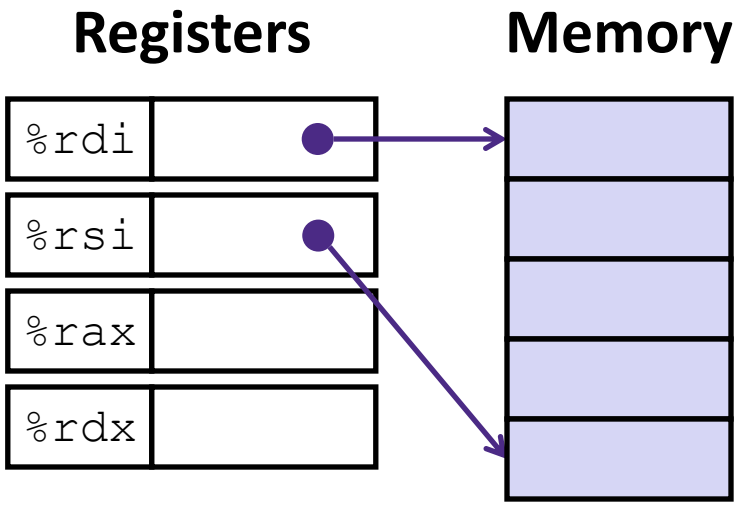
```
void swap(long* xp, long* yp)
{
  long t0 = *xp;
  long t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
    movq   (%rdi), %rax
    movq   (%rsi), %rdx
    movq   %rdx, (%rdi)
    movq   %rax, (%rsi)
    ret
```

Compiler Explorer:
https://godbolt.org/z/zc4Pcq

Understanding swap()

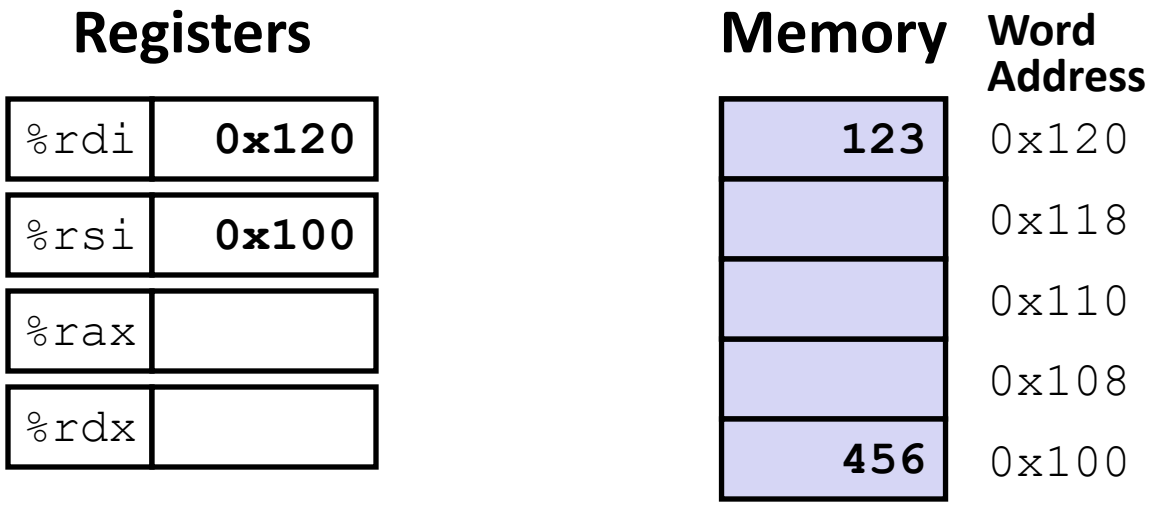
```
void swap(long* xp, long* yp)
{
  long t0 = *xp;
  long t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```




```
swap:
    movq   (%rdi), %rax
    movq   (%rsi), %rdx
    movq   %rdx, (%rdi)
    movq   %rax, (%rsi)
    ret
```

Register		Variable
%rdi	⇔	xp
%rsi	⇔	yp
%rax	⇔	t0
%rdx	⇔	t1

Understanding `swap()`




```
swap:
    movq   (%rdi), %rax   #  t0 = *xp
    movq   (%rsi), %rdx   #  t1 = *yp
    movq   %rdx, (%rdi)   # *xp =  t1
    movq   %rax, (%rsi)   # *yp =  t0
    ret
```

Complete Memory Addressing Modes

- **General:**
  - `D(Rb,Ri,S)` Mem[Reg[`Rb`]+Reg[`Ri`]*`S`+`D`]
    - `Rb`: Base register (any register)
    - `Ri`: Index register (any register except `%rsp`)
    - `S`: Scale factor (1, 2, 4, 8) – *why these numbers?*
    - `D`: Constant displacement value (a.k.a. immediate)

- **Special cases** (see CSPP Figure 3.3 on p.181)
  - `D(Rb,Ri)` Mem[Reg[`Rb`]+Reg[`Ri`]+`D`] `(S=1)`
  - `(Rb,Ri,S)` Mem[Reg[`Rb`]+Reg[`Ri`]*`S`] `(D=0)`
  - `(Rb,Ri)` Mem[Reg[`Rb`]+Reg[`Ri`]] `(S=1,D=0)`
  - `(,Ri,S)` Mem[Reg[`Ri`]*`S`] `(Rb=0,D=0)`

Address Computation Examples

%rdx	0xf000
%rcx	0x0100

D(Rb,Ri,S) → Mem[Reg[Rb]+Reg[Ri]*S+D]

ignore the memory access for now

Expression	Address Computation	Address (8 bytes wide)
0x8(%rdx)		
(%rdx,%rcx)		
(%rdx,%rcx,4)		
0x80(,%rdx,2)		

Summary

- x86-64 is a complex instruction set computing (CISC) architecture
  - There are 3 types of operands in x86-64
    - Immediate, Register, Memory
  - There are 3 types of instructions in x86-64
    - Data transfer, Arithmetic, Control Flow