Procedures and the Stack

The Stack is a region in memory which starts from the highest memory address and grows downwards when necessary as programs execute. We consider the region with the highest address the Stack “bottom” and the region with the lowest address the Stack “top”. %rsp is a dedicated special register which points to the current Stack top.

In x86-64, the Stack can be divided into individual stack frames for each function, which may contain the following sections (in order):

- **Return address**
  - Pushed by callq; address of instruction after callq

- **Callee-saved registers**
  - Only if function modifies/uses them

- **Local variables**
  - Variables that fit in a register may not be allocated on the Stack
  - Unavoidable if variable is too big for a register (e.g., array)
  - Unavoidable if variable needs an address (i.e., uses &var)

- **Caller-saved registers**
  - Only if values are needed across a function call

- **Argument build**
  - Only if function calls a function with more than six arguments

Example: consider the following lines of code and draw out the stack frames for main and foo right before foo returns (i.e., before any deallocation):

```c
int main(int argc, char* argv[]) {
    int x = 351;
    int a[] = {1, 2, 3};
    int y = foo(&x, 2, 3, 4, 5, 6, 7);
    return y + argc;
}

int foo(int* arg1, int arg2, ..., int arg7) {
    return *arg1 + arg7;
}
```
Consider the following x86-64 assembly and C code for the recursive function `rfun`.

```c
// Recursive function rfun
long rfun(char *s) {
    if (*s) {
        long temp = (long)*s;
        s++;
        return temp + rfun(s);
    }
    return 0;
}

// Main Function - program entry
int main(int argc, char **argv) {
    char *s = "CSE351";
    long r = rfun(s);
    printf("r: %ld\n", r);
}
```

00000000004005e6 <rfun>:

```
4005e6: 0f b6 07              movzbl (%rdi),%eax
4005e9: 84 c0                 test %al,%al
4005eb: 74 13                 je 400600 <rfun+0x1a>
4005ed: 53                    push %rbx
4005ee: 48 0f be d8           movsbq %al,%rbx
4005f2: 48 83 c7 01           add $0x1,%rdi
4005f6: e8 eb ff ff ff        callq 4005e6 <rfun>
4005fb: 48 01 d8              add %rbx,%rax
4005fe: eb 06                 jmp 400606 <rfun+0x20>
400600: b8 00 00 00 00 00     mov $0x0,%eax
400605: c3                    retq
400606: 5b                    pop %rbx
400607: c3                    retq
```

a) In terms of the C function, what value is being saved on the stack?

Value: `*s`.

The movsbq instruction at 0x4005ee puts `*s` into `%rbx`, which is pushed onto the stack by the pushq instruction at 0x4005ed.

b) What is the return address to `rfun` that gets stored on the stack during the recursive calls (in hex)?

0x4005fb
c) Assume main calls rfun with char *s = “CSE351” and then prints the result using the printf function, as shown in the C code above. Assume printf does not call any other procedure. Starting with (and including) main, how many total stack frames are created, and what is the maximum depth of the stack?

Total frames: 9      Max depth: 8
main -> rfun(s) -> rfun(s+1) -> rfun(s+2) -> rfun(s+3) -> rfun(s+4) -> rfun(s+5) -> rfun(s+6)-> printf()

The recursive call to rfun(s+6), which handles the null-terminator in the string (base case), still creates a stack frame since we consider the return address pushed to the stack during a procedure call to be part of the callee’s stack frame.

d) Assume main calls rfun with char *s = “CSE351”, as shown in the C code. After main calls rfun, we find that the return address to main is stored on the stack at address 0x7fffffffdb38. On the first call to rfun, the register %rdi holds the address 0x4006d0, which is the address of the input string “CSE 351” (i.e. char *s = 0x4006d0) during the fourth call to rfun.

For each address in the stack diagram below, fill in both the value and a description of the entry.

<table>
<thead>
<tr>
<th>Memory Address</th>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0x7fffffffdb48</td>
<td>Unknown</td>
<td>%rsp when main is entered</td>
</tr>
<tr>
<td>0x7fffffffdb38</td>
<td>0x400616</td>
<td>Return address to main</td>
</tr>
<tr>
<td>0x7fffffffdb30</td>
<td>Unknown</td>
<td>Original %rbx</td>
</tr>
<tr>
<td>0x7fffffffdb28</td>
<td>0x4005fb</td>
<td>Return address</td>
</tr>
<tr>
<td>0x7fffffffdb20</td>
<td>*s, “C”, 0x43</td>
<td>Saved %rbx</td>
</tr>
<tr>
<td>0x7fffffffdb18</td>
<td>0x4005fb</td>
<td>Return address</td>
</tr>
<tr>
<td>0x7fffffffdb10</td>
<td>*s, *(s+1), “S”, 0x53</td>
<td>Saved %rbx</td>
</tr>
<tr>
<td>0x7fffffffdb08</td>
<td>0x4005fb</td>
<td>Return address</td>
</tr>
<tr>
<td>0x7fffffffdb00</td>
<td>*s, *(s+2), “E”, 0x45</td>
<td>Saved %rbx</td>
</tr>
</tbody>
</table>