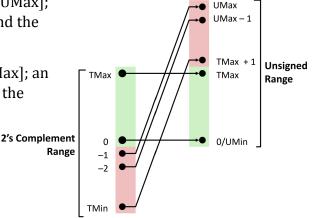
# CSE 351 Section 3 – Integers and Floating Point

Welcome back to section, we're happy that you're here  $\ensuremath{\textcircled{}}$ 

## **Integers and Arithmetic Overflow**

Arithmetic overflow occurs when the result of a calculation can't be represented in the current encoding scheme (*i.e.*, it lies outside of the representable range of values), resulting in an incorrect value.

- <u>Unsigned overflow</u>: the result lies outside of [UMin, UMax]; an indicator of this is when you add two numbers and the result is smaller than either number.
- <u>Signed overflow</u>: the result lies outside of [TMin, TMax]; an indicator of this is when you add two numbers with the same sign and the result has the opposite sign.



### Exercises:

1) [Spring 2016 Midterm 1C] Assuming these are all signed two's complement 6-bit integers, compute the result of each of the following additions. For each, indicate if it resulted in overflow.

001001	110001	011001	101111
<u>+ 110110</u>	<u>+ 111011</u>	<u>+ 001100</u>	<u>+ 011111</u>

2) [Autumn 2019 Midterm 1C] Find the largest 8-bit unsigned numeral (answer in hex) such that c + 0x80 causes NEITHER signed nor unsigned overflow in 8 bits.

# **IEEE 754 Floating Point Standard**

#### Goals

- ★ Represent a large range of values (both very small and very large numbers),
- $\star$  Include a high amount of precision, and
- ★ Allow for real arithmetic results (*e.g.*,  $\infty$  and NaN).

### Encoding

The <u>value</u> of a real number can be represented in normalized scientific binary notation as:

 $(-1)^{\text{sign}} \times \text{Mantissa}_2 \times 2^{\text{Exponent}} = (-1)^{\text{S}} \times 1.\text{M}_2 \times 2^{\text{E-bias}}$ 

The <u>binary representation</u> for floating point encodes these three components into separate fields:

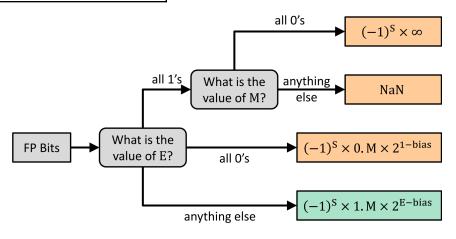
	S	E	М
float:	1	8 bits	23 bits
double:	1	11 bits	52 bits

- S: the sign of the number (0 for positive, 1 for negative)
- E: the exponent in biased notation (unsigned with a bias of 2<sup>w-1</sup>-1)
- M: the mantissa (also called the significand or fraction) without the implicit leading 1

#### **Special Cases and Interpretations**

The interpretation depends on the values in the exponent and mantissa fields:

E	М	Meaning
0b00	anything	denormalized number (denorm)
anything else	anything	normalized number
0b11	zero	infinity (∞)
0b11	nonzero	not-a-number (NaN)



#### **Mathematical Properties**

- Not <u>associative</u>:  $(2 + 2^{50}) 2^{50} \neq 2 + (2^{50} 2^{50})$
- Not <u>distributive</u>:  $100 \times (0.1 + 0.2) \neq 100 \times 0.1 + 100 \times 0.2$
- Not <u>cumulative</u>:  $2^{25} + 1 + 1 + 1 + 1 \neq 2^{25} + 4$

#### Exercises:

- 3) Let's say that we want to represent the number 3145728.125 (broken down as  $2^{21} + 2^{20} + 2^{-3}$ )
  - a) Convert this number to into single precision floating point representation:

- b) Which limitation of floating point representation does this result highlight?
- 4) [Summer 2018 Midterm 1E-G] We are working with a new floating point datatype (flo) that follows the same conventions as IEEE 754 except using 8 bits split into the following fields:

Sign (1) Exponent (3)	Mantissa (4)
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- a) What is the encoding of the most negative real number that we can represent (∞ is not a real number) in this floating point scheme (binary)?
- b) If we have signed char x = 0b10101000 = -88, what will occur if we cast flo f = (flo) x (*i.e.*, try to represent the value stored in x as a flo)?

Rounding Underflow 0	Overflow	None of these
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- 5) Based on the floating point representation, explain why each of the three mathematical property examples shown on the previous page occurs.
  - a) Not associative:
  - b) Not distributive:
  - c) Not cumulative:
- 6) If we have float x, y;, give two *different* reasons why (x+2\*y)-y == x+y might evaluate to false.