Alt text: I looked at some of the data dumps from vulnerable sites, and it was ... bad. I saw emails, passwords, password hints. SSL keys and session cookies. Important servers brimming with visitor IPs. Attack ships on fire off the shoulder of Orion, c-beams glittering in the dark near the Tannhäuser Gate. I should probably patch OpenSSL.

http://xkcd.com/1353/
Relevant Course Information

❖ hw13 due Wednesday (11/2)
❖ hw15 due Monday (11/7)

❖ Lab 3 released today, due next Friday (11/11)
  ▪ You will have everything you need by the end of this lecture

❖ Midterm starts Thursday
  ▪ Instructions will be posted on Ed Discussion
  ▪ Gilligan’s Island Rule: discuss high-level concepts and give hints, but not solving the problems together
  ▪ We will be available on Ed Discussion (private posts, please) and office hours to answer clarifying questions
Buffer Overflows

- Address space layout review
- Input buffers on the stack
- Overflowing buffers and injecting code
- Defenses against buffer overflows
Review: General Memory Layout

- **Stack**
  - Local variables (procedure context)

- **Heap**
  - Dynamically allocated as needed
  - `new`, `malloc()`, `calloc()`, ...

- **Statically-allocated Data**
  - Read/write: global variables (Static Data)
  - Read-only: string literals (Literals)

- **Code/Instructions**
  - Executable machine instructions
  - Read-only
Memory Allocation Example

```c
char big_array[1L<<24]; /* 16 MB */

int global = 0;

int useless() { return 0; }

int main() {
  void *p1, *p2;
  int local = 0;
  p1 = malloc(1L << 28); /* 256 MB */
  p2 = malloc(1L << 8); /* 256 B */
  /* Some print statements ... */
}
```

Where does everything go?
Memory Allocation Example

```
char big_array[1L<<24]; /* 16 MB */
int global = 0;
int useless() { return 0; }
int main() {
    void *p1, *p2;
    int local = 0;
    p1 = malloc(1L << 28); /* 256 MB */
    p2 = malloc(1L << 8); /* 256 B */
    /* Some print statements ... */
}
```

Where does everything go?
What Is a Buffer?

❖ A buffer is an array used to temporarily store data

❖ You’ve probably seen “video buffering...”
  ▪ The video is being written into a buffer before being played

❖ Buffers can also store user input
Reminder: x86-64/Linux Stack Frame

- **Caller’s Stack Frame**
  - Arguments (if > 6 args) for this call

- **Current/ Callee Stack Frame**
  - Return address
    - Pushed by `call` instruction
  - Old frame pointer (optional)
  - Caller-saved pushed before setting up arguments for a function call
  - Callee-saved pushed before using long-term registers
  - Local variables
    (if can’t be kept in registers)
  - “Argument build” area
    (Need to call a function with >6 arguments? Put them here)
Buffer Overflow in a Nutshell

- C does not check array bounds
  - Many Unix/Linux/C functions don’t check argument sizes
  - Allows overflowing (writing past the end) of buffers (arrays)

- “Buffer Overflow” = Writing past the end of an array

- Characteristics of the traditional Linux memory layout provide opportunities for malicious programs
  - Stack grows “backwards” in memory
  - Data and instructions both stored in the same memory
Buffer Overflow in a Nutshell

❖ Stack grows down towards lower addresses

❖ Buffer grows up towards higher addresses

❖ If we write past the end of the array, we overwrite data on the stack!

Enter input: hello

No overflow 😊
Buffer Overflow in a Nutshell

- Stack grows down towards lower addresses

- Buffer grows up towards higher addresses

- If we write past the end of the array, we overwrite data on the stack!

Enter input: helloabcdef
Buffer Overflow in a Nutshell

- Stack grows *down* towards lower addresses
- Buffer grows *up* towards higher addresses
- If we write past the end of the array, we overwrite data on the stack!

Enter input: helloabcdef

Buffer overflow! 😞
Buffer Overflow in a Nutshell

- Buffer overflows on the stack can overwrite “interesting” data
  - Attackers just choose the right inputs

- Simplest form (sometimes called “stack smashing”)
  - Unchecked length on string input into bounded array causes overwriting of stack data
  - Try to change the return address of the current procedure

- Why is this a big deal?
  - It was the #1 technical cause of security vulnerabilities
    - #1 overall cause is social engineering / user ignorance
String Library Code

- Implementation of Unix function `gets()`

```c
/* Get string from stdin */
char* gets(char* dest) {
    int c = getchar();
    char* p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}
```

- What could go wrong in this code?
String Library Code

❖ Implementation of Unix function `gets()`

```c
/* Get string from stdin */
char* gets(char* dest) {
    int c = getchar();
    char* p = dest;
    while (c != EOF && c != 'n') {
        *p++ = c;
        c = getchar();
    }
    *p = '0';
    return dest;
}
```

❖ No way to specify limit on number of characters to read

❖ Similar problems with other Unix functions:
  - `strcpy`: Copies string of arbitrary length to a dst
  - `scanf`, `fscanf`, `sscanf`, when given `%s` specifier
Vulnerable Buffer Code

```c
/* Echo Line */
void echo() {
    char buf[8]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```c
void call_echo() {
    echo();
}
```

```bash
unix> ./buf-nsp
Enter string: 123456789012345
123456789012345
unix> ./buf-nsp
Enter string: 123456789012345
Segmentation fault (core dumped)
```
Buffer Overflow Disassembly (buf-nsp)

**echo:**

<table>
<thead>
<tr>
<th>Address</th>
<th>Instruction</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>401146</td>
<td><code>48 83 ec 18</code></td>
<td>sub $0x18,%rsp</td>
</tr>
<tr>
<td></td>
<td>...</td>
<td>... calls printf ...</td>
</tr>
<tr>
<td>401159</td>
<td><code>48 8d 7c 24 08</code></td>
<td>lea 0x8(%rsp),%rdi</td>
</tr>
<tr>
<td>40115e</td>
<td><code>b8 00 00 00 00</code></td>
<td>mov $0x0,%eax</td>
</tr>
<tr>
<td>401163</td>
<td><code>e8 e8 fe ff ff</code></td>
<td>callq 401050 <a href="mailto:gets@plt">gets@plt</a></td>
</tr>
<tr>
<td>401168</td>
<td><code>48 8d 7c 24 08</code></td>
<td>lea 0x8(%rsp),%rdi</td>
</tr>
<tr>
<td>40116d</td>
<td><code>e8 be fe ff ff</code></td>
<td>callq 401030 <a href="mailto:puts@plt">puts@plt</a></td>
</tr>
<tr>
<td>401172</td>
<td><code>48 83 c4 18</code></td>
<td>add $0x18,%rsp</td>
</tr>
<tr>
<td>401176</td>
<td><code>c3</code></td>
<td>retq</td>
</tr>
</tbody>
</table>

**call_echo:**

<table>
<thead>
<tr>
<th>Address</th>
<th>Instruction</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>401177</td>
<td><code>48 83 ec 08</code></td>
<td>sub $0x8,%rsp</td>
</tr>
<tr>
<td>40117b</td>
<td><code>b8 00 00 00 00</code></td>
<td>mov $0x0,%eax</td>
</tr>
<tr>
<td>401180</td>
<td><code>e8 c1 ff ff ff</code></td>
<td>callq 401146 &lt;echo&gt;</td>
</tr>
<tr>
<td>401185</td>
<td><code>48 83 c4 08</code></td>
<td>add $0x8,%rsp</td>
</tr>
<tr>
<td>401189</td>
<td><code>c3</code></td>
<td>retq</td>
</tr>
</tbody>
</table>

`return address`
**Buffer Overflow Stack**

*Before call to gets*

```c
/* Echo Line */
void echo()
{
    char buf[8];  /* Way too small! */
    gets(buf);
    puts(buf);
}
```

**Note:** addresses increasing right-to-left, bottom-to-top
Buffer Overflow Example

Before call to gets

Stack frame for call_echo

void echo()
{
    char buf[8];
    gets(buf);
    ...
}

call_echo:

...  
401180:   callq 401146 <echo>  
401185:   add $0x8,%rsp  
...
Buffer Overflow Example #1

After call to gets

Stack frame for call_echo

<p>| | | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>00</td>
<td>00</td>
<td>00</td>
<td>00</td>
<td>00</td>
</tr>
<tr>
<td>00</td>
<td>00</td>
<td>11</td>
<td>85</td>
<td></td>
</tr>
<tr>
<td>00</td>
<td>35</td>
<td>34</td>
<td>33</td>
<td></td>
</tr>
<tr>
<td>32</td>
<td>31</td>
<td>30</td>
<td>39</td>
<td></td>
</tr>
<tr>
<td>38</td>
<td>37</td>
<td>36</td>
<td>35</td>
<td></td>
</tr>
<tr>
<td>34</td>
<td>33</td>
<td>32</td>
<td>31</td>
<td></td>
</tr>
</tbody>
</table>

8 bytes unused

Note: Digit “N” is just 0x3N in ASCII!

void echo()
{
    char buf[8];
    gets(buf);
    ...
}

echo:
    subq $24, %rsp
    ...
    leaq 8(%rsp), %rdi
    mov $0x0, %eax
    call gets
    ...

call_echo:

    ...
    401180: callq 401146 <echo>
    401185: add $0x8, %rsp
    ...

unix> ./buf-nsp
Enter string: 1234567890123456789012345
123456789012345
Overflowed buffer, but did not corrupt state
Buffer Overflow Example #2

After call to gets

Stack frame for call_echo

| 00 | 00 | 00 | 00 |
| 00 | 40 | 11 | 00 |
| **36** | 35 | 34 | 33 |
| 32 | 31 | 30 | 39 |
| 38 | 37 | 36 | 35 |
| 34 | 33 | 32 | 31 |

8 bytes unused

void echo()
{
    char buf[8];
    gets(buf);
    ...
}

call_echo:

unix> ./buf-nsp
Enter string: 1234567890123456
Segmentation fault (core dumped)

Overflowed buffer and corrupted return pointer
Buffer Overflow Example #2 Explained

After return from echo

Stack frame for call_echo

Returns” to a valid instruction, but bad indirect jump so program signals SIGSEGV, Segmentation fault
Malicious Use of Buffer Overflow: Code Injection Attacks

- Input string contains byte representation of executable code
- Overwrite return address A with address of buffer B
- When `bar()` executes `ret`, will jump to exploit code

```c
void foo() {
    bar();
    A: ...
}

int bar() {
    char buf[64];
    gets(buf);
    ...
    return ...;
}
```

Stack after call to `gets()`

- High Addresses
- foo stack frame
- A
- B
- data written by `gets()`
- pad
- exploit code
- Low Addresses
- bar stack frame
- buf starts here
- B
Practice Question

- `smash_me` is vulnerable to stack smashing!
- What is the minimum number of characters that `gets` must read in order for us to change the return address to a stack address?
  - For example: (0x00 00 7f ff ca fe f0 0d)

<table>
<thead>
<tr>
<th>Previous stack frame</th>
</tr>
</thead>
<tbody>
<tr>
<td>00 00 00 00 00</td>
</tr>
<tr>
<td>00 40 05 d1</td>
</tr>
<tr>
<td>...</td>
</tr>
<tr>
<td>[0]</td>
</tr>
</tbody>
</table>

smash_me:
- `subq $0x40, %rsp`
- `leaq 16(%rsp), %rdi`
- `call gets`

A. 27
B. 30
C. 51
D. 54
E. We’re lost...
Exploits Based on Buffer Overflows

Buffer overflow bugs can allow attackers to execute arbitrary code on victim machines

- Distressingly common in real programs
  - Programmers keep making the same mistakes 😞
  - Recent measures make these attacks much more difficult

- Examples across the decades
  - Original “Internet worm” (1988)
  - Heartbleed (2014, affected 17% of servers)
    - Similar issue in Cloudbleed (2017)
  - Hacking embedded devices
    - Cars, Smart homes, Planes
Example: the original Internet worm (1988)

❖ Exploited a few vulnerabilities to spread
  ▪ Early versions of the finger server (fingerd) used `gets()` to read the argument sent by the client:
    • `finger droh@cs.cmu.edu`
  ▪ Worm attacked `fingerd` server with phony argument:
    • `finger "exploit-code padding new-return-addr"`
    • Exploit code: executed a root shell on the victim machine with a direct connection to the attacker

❖ Scanned for other machines to attack
  ▪ Invaded ~6000 computers in hours (10% of the Internet)
    • see June 1989 article in *Comm. of the ACM*
  ▪ The author of the worm (Robert Morris*) was prosecuted...
Example: Heartbleed (2014)

HOW THE HEARTBLEED BUG WORKS:

SERVER, ARE YOU STILL THERE? IF SO, REPLY "POTATO" (6 LETTERS).

User Meg wants these 6 letters: POTATO. User Ada wants pages about "irl games". Unlocking secure records with master key 5130985733435 sends complex encoded data this message: "4538538374224"

POTATO
Example: Heartbleed (2014)

**SERVER, ARE YOU STILL THERE? IF SO, REPLY “BIRD” (4 LETTERS).**

User Olivia from London wants pages about "bees in car why". Note: Files for IP 375.381.83.17 are in /tmp/files-3843. User Meg wants these 4 letters: BIRD. There are currently 34 connections open. User Brendan uploaded the file "cat.jpg" (contents: 334ba962a20cbb96ff90be3b569f890be3b569).

**HMM...**

User Olivia from London wants pages about "bees in car why". Note: Files for IP 375.381.83.17 are in /tmp/files-3843. User Meg wants these 4 letters: BIRD. There are currently 34 connections open. User Brendan uploaded the file "cat.jpg" (contents: 334ba962a20cbb96ff90be3b569f890be3b569).
Example: Heartbleed (2014)

User Meg wants these 500 letters: HAT. Lucas requests the "missed connections" page. Eve (administrator) wants to set server's master key to "14835038534". Isabel wants pages about snakes but not too long". User Karen wants to change account password to "ColdBlood".
Heartbleed Details

- Buffer over-read in OpenSSL
  - Open source security library
  - Bug in a small range of versions
- “Heartbeat” packet
  - Specifies length of message
  - Server echoes it back
  - Library just “trusted” this length
  - Allowed attackers to read contents of memory anywhere they wanted
- Est. 17% of Internet affected
  - “Catastrophic”
  - Github, Yahoo, Stack Overflow, Amazon AWS, ...

![Heartbeat - Normal usage](image1.png)

Server, send me this 4 letter word if you are there: "bird"

Client

![Heartbeat - Malicious usage](image2.png)

Server, send me this 500 letter word if you are there: "bird"

Server has connected. User Bob has connected. User Alice wants 4 letters: bird. Server master key is 31431498531054. User Carol wants to change password to "password 123". By FenixFeather - Own work, CC BY-SA 3.0, https://commons.wikimedia.org/w/index.php?curid=32276981
Hacking Cars (2010)

❖ UW CSE research demonstrated wirelessly hacking a car using buffer overflow

❖ Overwrote the onboard control system’s code
  ▪ Disable brakes, unlock doors, turn engine on/off
Hacking DNA Sequencing Tech (2017)

Computer Security and Privacy in DNA Sequencing
Paul G. Allen School of Computer Science & Engineering, University of Washington

- Potential for malicious code to be encoded in DNA!
- Attacker can gain control of DNA sequencing machine when malicious DNA is read
Dealing with buffer overflow attacks

1) Employ system-level protections
2) Avoid overflow vulnerabilities
3) Have compiler use “stack canaries”
1) System-Level Protections

- **Non-executable code segments**
  - In traditional x86, can mark region of memory as either “read-only” or “writeable”
    - Can execute anything readable
  - x86-64 added explicit “execute” permission
  - Stack marked as non-executable
    - Do *NOT* execute code in Stack, Static Data, or Heap regions
    - Hardware support needed

Any attempt to execute this code will fail
1) System-Level Protections

- Non-executable code segments
  - Wait, doesn’t this fix everything?
  - Works well, but can’t always use it
  - Many embedded devices do not have this protection
    - e.g., cars, smart homes, pacemakers
- Some exploits still work!
  - Return-oriented programming
  - Return to libc attack
  - JIT-spray attack

Stack after call to `gets()`

Any attempt to execute this code will fail
1) System-Level Protections

❖ Randomized stack offsets
   ▪ At start of program, allocate random amount of space on stack
   ▪ Shifts stack addresses for entire program
     • Addresses will vary from one run to another
   ▪ Makes it difficult for hacker to predict beginning of inserted code

❖ Example: Address of variable \texttt{local} for when Slide 5 code executed 3 times:
   • 0x7ffd19d3f8ac
   • 0x7ffe8a462c2c
   • 0x7ffe927c905c

▪ Stack repositioned each time program executes
2) Avoid Overflow Vulnerabilities in Code

/* Echo Line */
void echo()
{
    char buf[8];  /* Way too small! */
    fgets(buf, 8, stdin);
    puts(buf);
}

❖ Use library routines that limit string lengths
   - fgets instead of gets (2\textsuperscript{nd} argument to fgets sets limit)
   - strncpy instead of strcpy
   - Don’t use scanf with \texttt{\%s} conversion specification
     - Use fgets to read the string
     - Or use \texttt{\%ns} where \texttt{n} is a suitable integer
2) Avoid Overflow Vulnerabilities in Code

- Alternatively, don’t use C - use a language that does array index bounds check
  - Buffer overflow is impossible in Java
    - ArrayIndexOutOfBoundsException
  - Rust language was designed with security in mind
    - Panics on index out of bounds, plus more protections
3) Stack Canaries

❖ Basic Idea: place special value ("canary") on stack just beyond buffer
  ▪ Secret value that is randomized before main()
  ▪ Placed between buffer and return address
  ▪ Check for corruption before exiting function

❖ GCC implementation
  ▪ \texttt{-fstack-protector}

\texttt{unix}> ./buf
Enter string: \texttt{12345678}
12345678

\texttt{unix}> ./buf
Enter string: \texttt{123456789}
*** stack smashing detected ***
Protected Buffer Disassembly (buf)

**echo:**

```
401156:  push   %rbx
401157:  sub    $0x10,%rsp
40115b:  mov    $0x28,%ebx
401160:  mov    %fs:(%rbx),%rax
401164:  mov    %rax,0x8(%rsp)
401169:  xor    %eax,%eax
        ...     ... call printf ...
40117d:  callq  401060 <gets@plt>
401182:  mov    %rsp,%rdi
401185:  callq  401030 <puts@plt>
40118a:  mov    0x8(%rsp),%rax
40118f:  xor    %fs:(%rbx),%rax
401193:  jne   40119b <echo+0x45>
401195:  add    $0x10,%rsp
401199:  pop    %rbx
40119a:  retq
40119b:  callq  401040 <__stack_chk_fail@plt>
```

This is extra (non-testable) material
Setting Up Canary

Before call to `gets`

```
/* Echo Line */
void echo()
{
    char buf[8]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

echo:

```
...  
    movq %fs:40, %rax  # Get canary
    movq %rax, 8(%rsp) # Place on stack
    xorl %eax, %eax   # Erase canary
    ...  
```

buf ← %rsp

This is extra (non-testable) material
## Checking Canary

### After call to `gets`

### Stack frame for `call_echo`

<table>
<thead>
<tr>
<th>Return address (8 bytes)</th>
<th>Canary (8 bytes)</th>
</tr>
</thead>
<tbody>
<tr>
<td>00 37 36 35</td>
<td>34 33 32 31</td>
</tr>
</tbody>
</table>

### Code Snippet

```c
/* Echo Line */
void echo()
{
    char buf[8]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```assembly
...movq 8(%rsp), %rax       # retrieve from Stack
xorq %fs:40, %rax          # compare to canary
jne .L4                   # if not same, FAIL
...
.L4: call __stack_chk_fail
```

Input: 1234567

This is extra (non-testable) material
Summary of Prevention Measures

1) Employ system-level protections
   ▪ Code on the Stack is not executable
   ▪ Randomized Stack offsets

2) Avoid overflow vulnerabilities
   ▪ Use library routines that limit string lengths
   ▪ Use a language that makes them impossible

3) Have compiler use “stack canaries”
Think this is cool?

- You’ll love Lab 3 😊
  - Released Wednesday, due next Friday (11/13)
  - Some parts *must* be run through GDB to disable certain security features

- Take CSE 484 (Security)
  - Several different kinds of buffer overflow exploits
  - Many ways to counter them

- Nintendo fun!
  - Using glitches to rewrite code: [https://www.youtube.com/watch?v=TqK-2jUQBUY](https://www.youtube.com/watch?v=TqK-2jUQBUY)
  - Flappy Bird in Mario: [https://www.youtube.com/watch?v=hB6eY73sLV0](https://www.youtube.com/watch?v=hB6eY73sLV0)
Discussion Questions

❖ In Lab 3, you will run a buffer overflow code injection attack; students love this lab because it “makes you feel like a hacker”
  - What connotations (i.e., ideas or feelings evoked) does this statement carry for you and where do those come from?
  
  - While it is easy to say that you should not exploit security vulnerabilities, does the target of an attack change how you feel about it? Why?