Memory, Data, & Addressing I

CSE 351 Winter 2021

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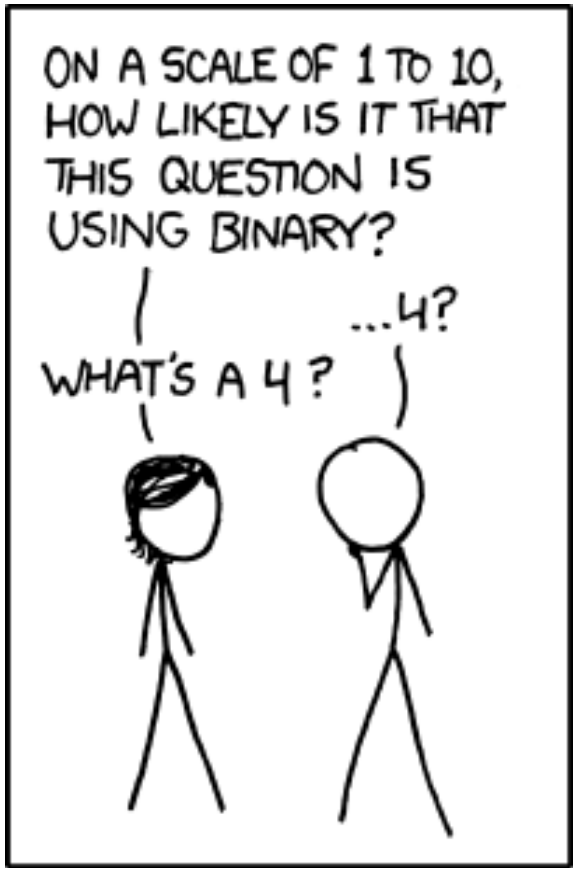


http://xkcd.com/953/

Admin

- Pre-Course Survey and hw0 due tonight @ 11:59 pm
  - Starting Week 2: hw due at 11:00 am (Seattle time)
- hw1 due Friday (1/8) @ 11:59 pm
- hw2 due Monday (1/11) @ 11:00 am
- Lab 0 due Friday (1/8) @ 11:59 pm
  - This lab is *exploratory* and looks like a hw; the other labs will look a lot different

- Ed Discussion etiquette
  - For anything that doesn't involve sensitive information or a solution, post publicly (you can post anonymously!)
  - If you feel like you question has been sufficiently answered, make sure that a response has a checkmark

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

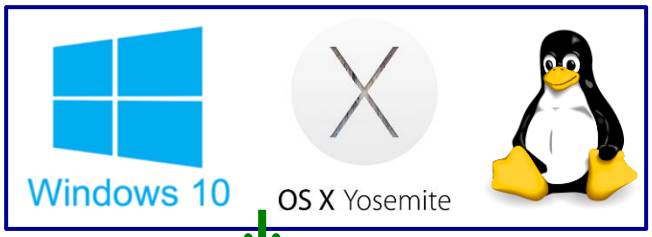
Assembly language:

```
get_mpg:
    pushq   %rbp
    movq    %rsp, %rbp
    ...
    popq    %rbp
    ret
```

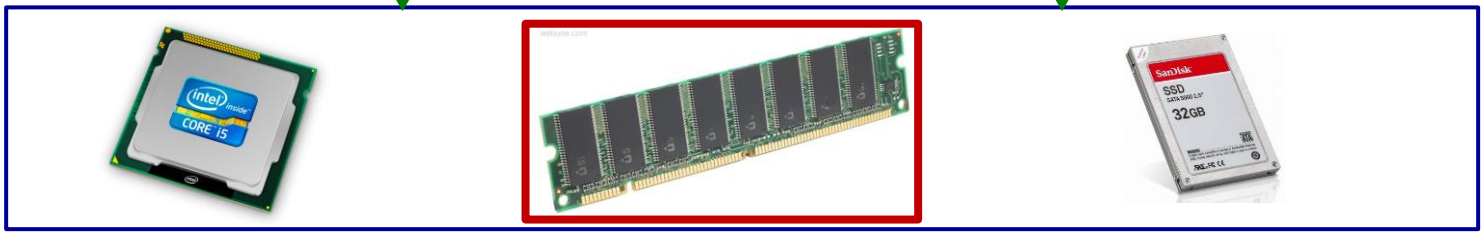
Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
11000001111110100011111
```

OS:



Computer system:



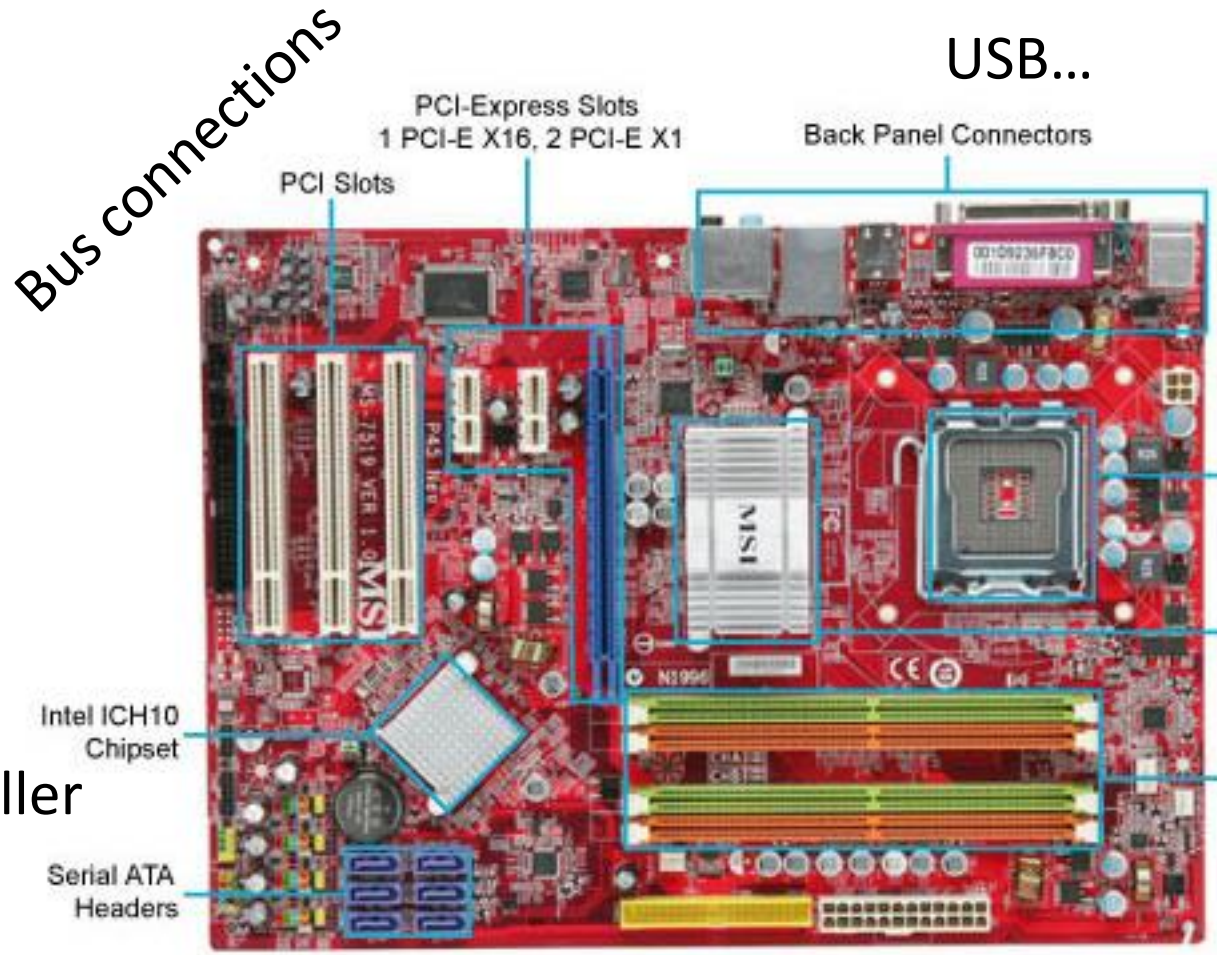
Memory & data
Integers & floats
x86 assembly
Procedures & stacks
Executables
Arrays & structs
Memory & caches
Processes
Virtual memory
Memory allocation
Java vs. C

Reading Review

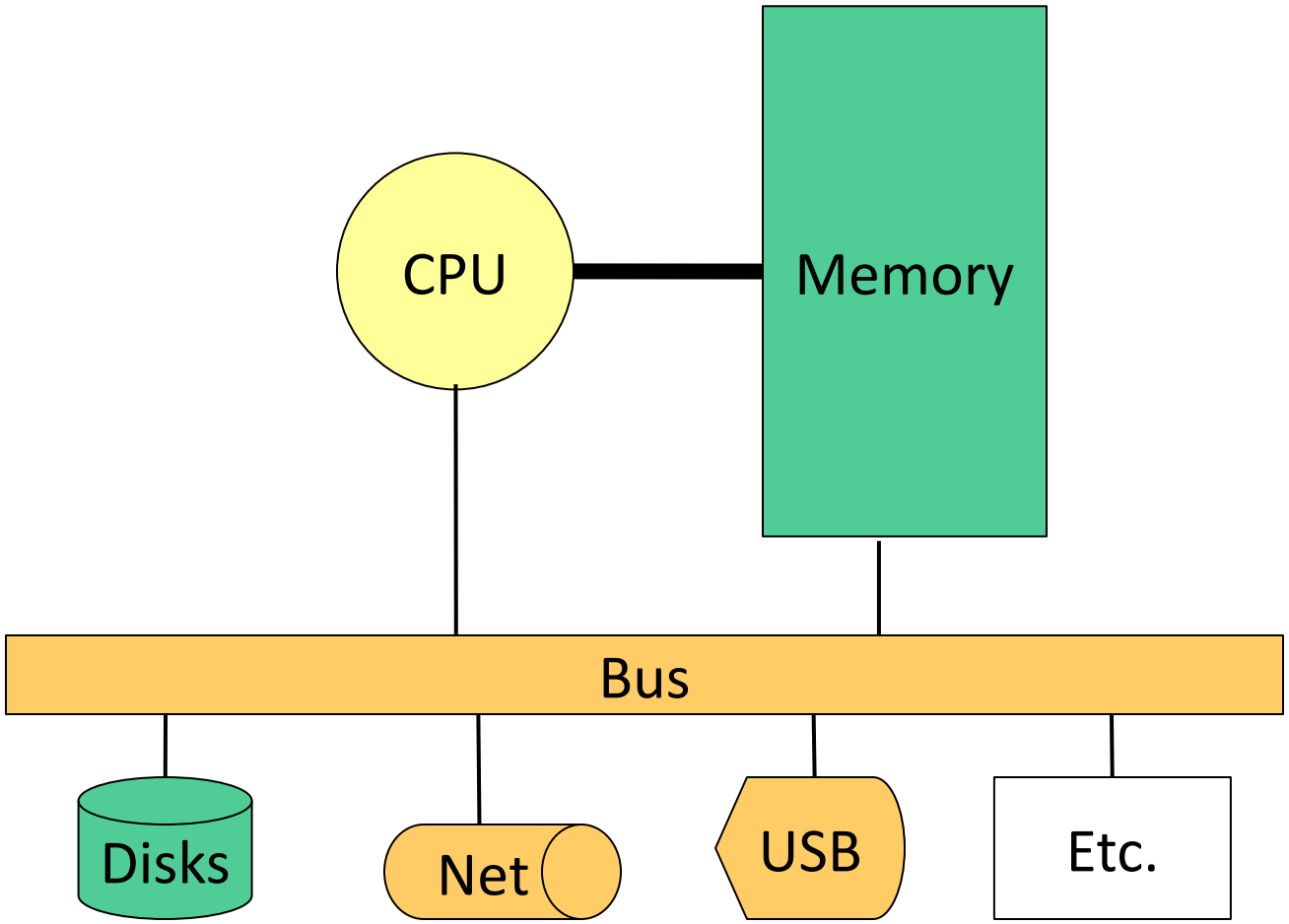
- Terminology:
  - word size, byte-oriented memory
  - address, address space
  - most-significant bit (MSB), least-significant bit (LSB)
  - big-endian, little-endian
  - pointer

- Questions from the Reading?

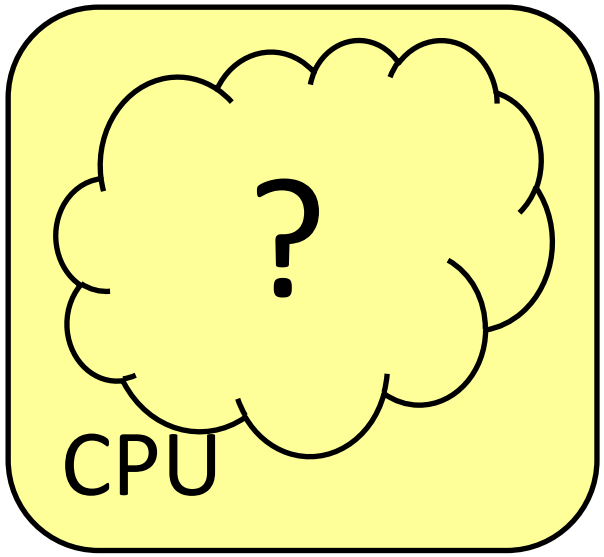
Hardware: Physical View

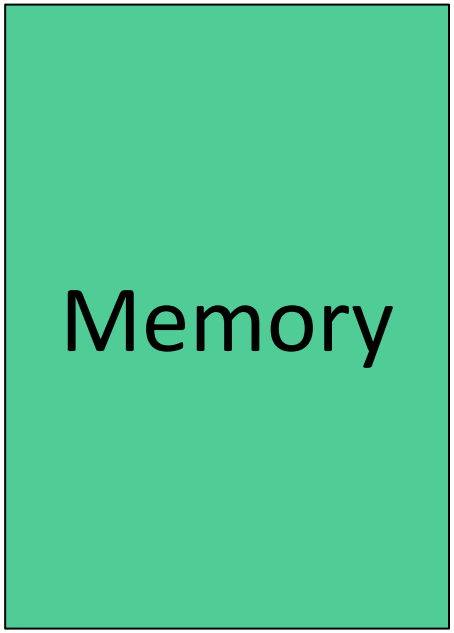


Hardware: Logical View



Hardware: 351 View (version 0)

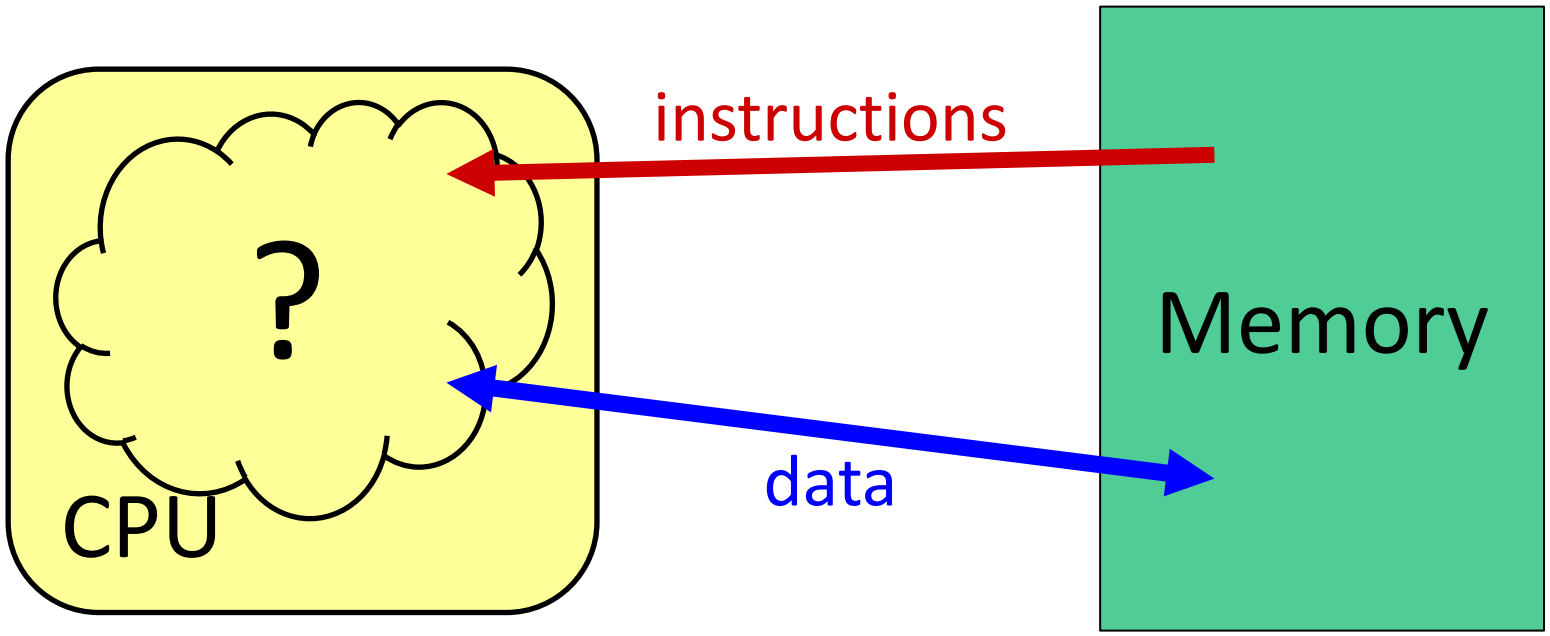




- The CPU executes instructions
- Memory stores data
- Binary encoding!
  - Instructions *are* just data

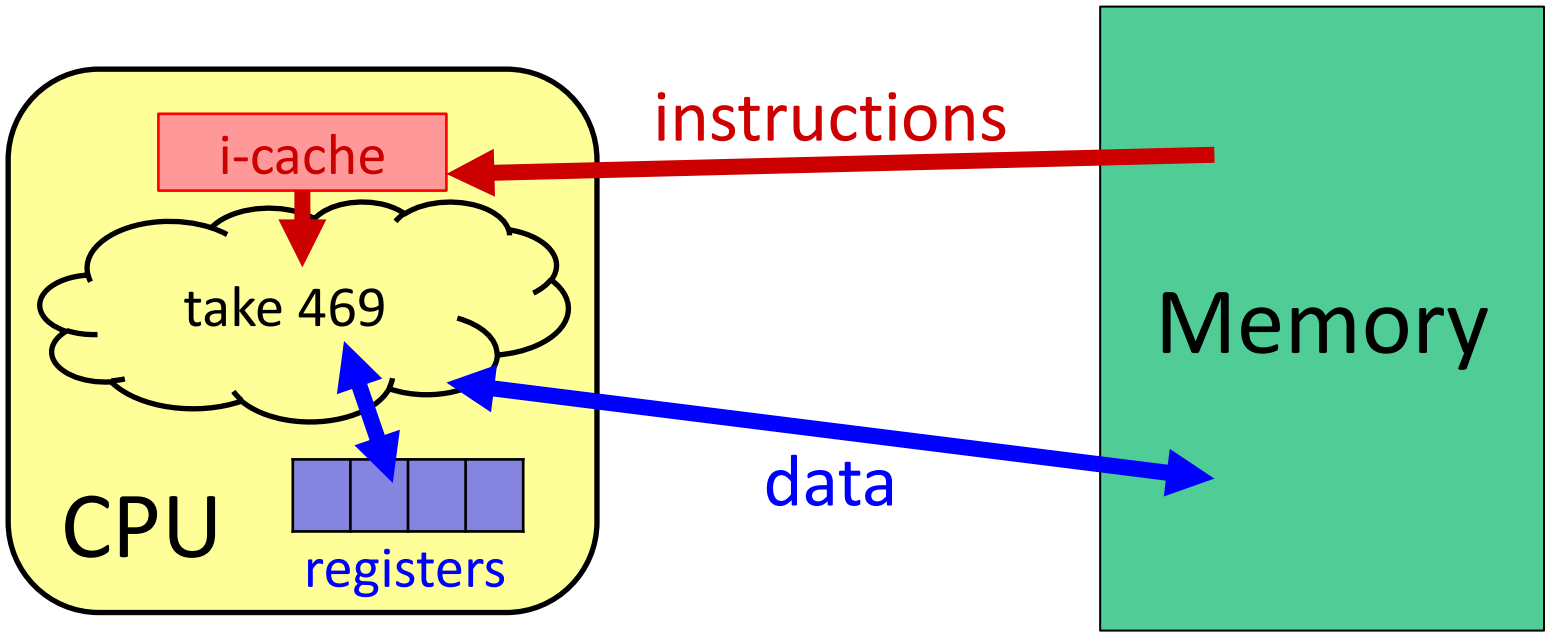
How are data and instructions represented?

Hardware: 351 View (version 0)



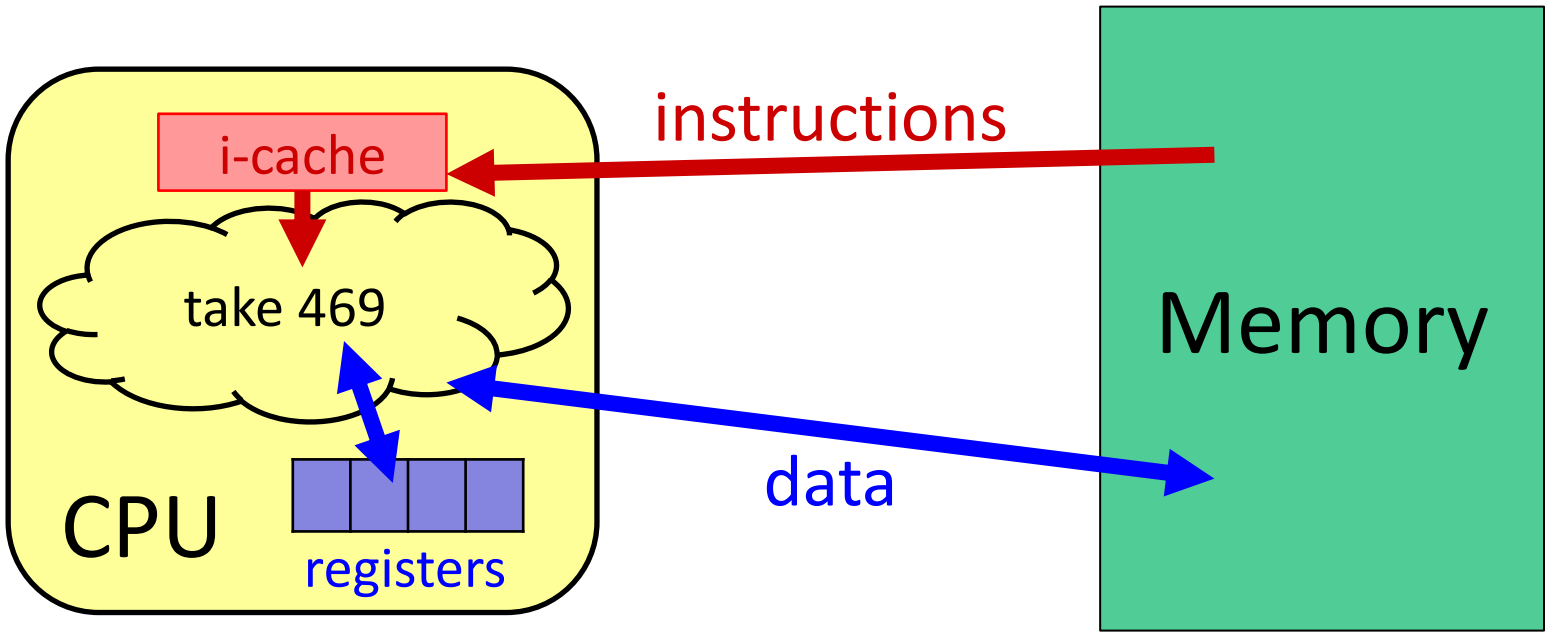
- To execute an instruction, the CPU must:
  1) Fetch the instruction
  2) (if applicable) Fetch data needed by the instruction
  3) Perform the specified computation
  4) (if applicable) Write the result back to memory

Hardware: 351 View (version 1)



- More CPU details:
  - Instructions are held temporarily in the instruction cache
  - Other data are held temporarily in registers
- Instruction fetching is hardware-controlled
- Data movement is programmer-controlled (assembly)

Hardware: 351 View (version 1)



- We will start by learning about Memory

How does a program find its data in memory?

- Addresses!
  - Can be stored in *pointers*

Review Questions – Ed Lessons (1.5)

- By looking at the bits stored in memory, I can tell what a particular 4 bytes is being used to represent.

  A. True **B. False**

- We can fetch a piece of data from memory as long as we have its address.

  A. True **B. False**

- Which of the following bytes have a most-significant bit (MSB) of 1?

  A. 0x63 **B. 0x90** **C. 0xCA** **D. 0xF**

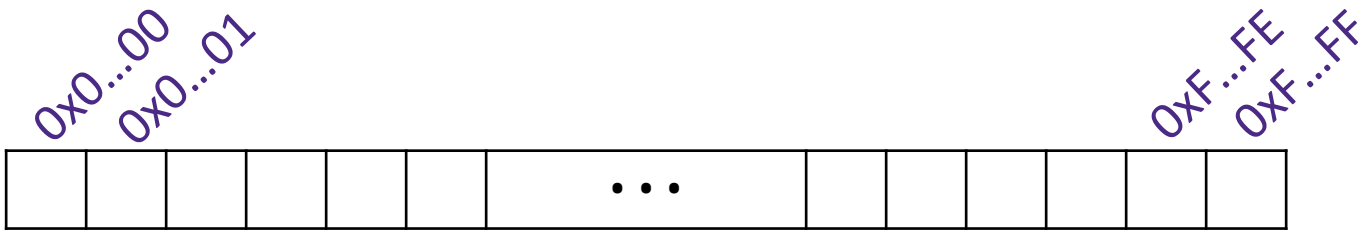
Binary Encoding Additional Details

- Because storage is finite in reality, everything is stored as "fixed" length
  - Data is moved and manipulated in fixed-length chunks
  - Multiple fixed lengths (*e.g.*, 1 byte, 4 bytes, 8 bytes)
  - Leading zeros now *must* be included up to "fill out" the fixed length

- Example: the "eight-bit" representation of the number 4 is 0b00000100

Bits and Bytes and Things

- 1 byte = 8 bits
- n bits can represent up to 2^n things
  - Sometimes (oftentimes?) those "things" are bytes!
- If addresses are a-bits wide, how many distinct addresses are there?
- What does each address refer to?



Machine "Words"

- Instructions encoded into machine code (0's and 1's)
  - Historically (still true in some assembly languages), all instructions were exactly the size of a word
- We have *chosen* to tie word size to address size/width
  - word size = address size = register size
  - word size = w bits $\rightarrow 2^w$ addresses
- Current x86 systems use **64-bit (8-byte) words**
  - Potential address space: $\mathbf{2^{64}}$ addresses
    2^{64} bytes $\approx$ **1.8 x 10^{19} bytes**
    = 18 billion billion bytes = 18 EB (exabytes)
  - Actual physical address space: **48 bits**

Data Representations

- Sizes of data types (in bytes)

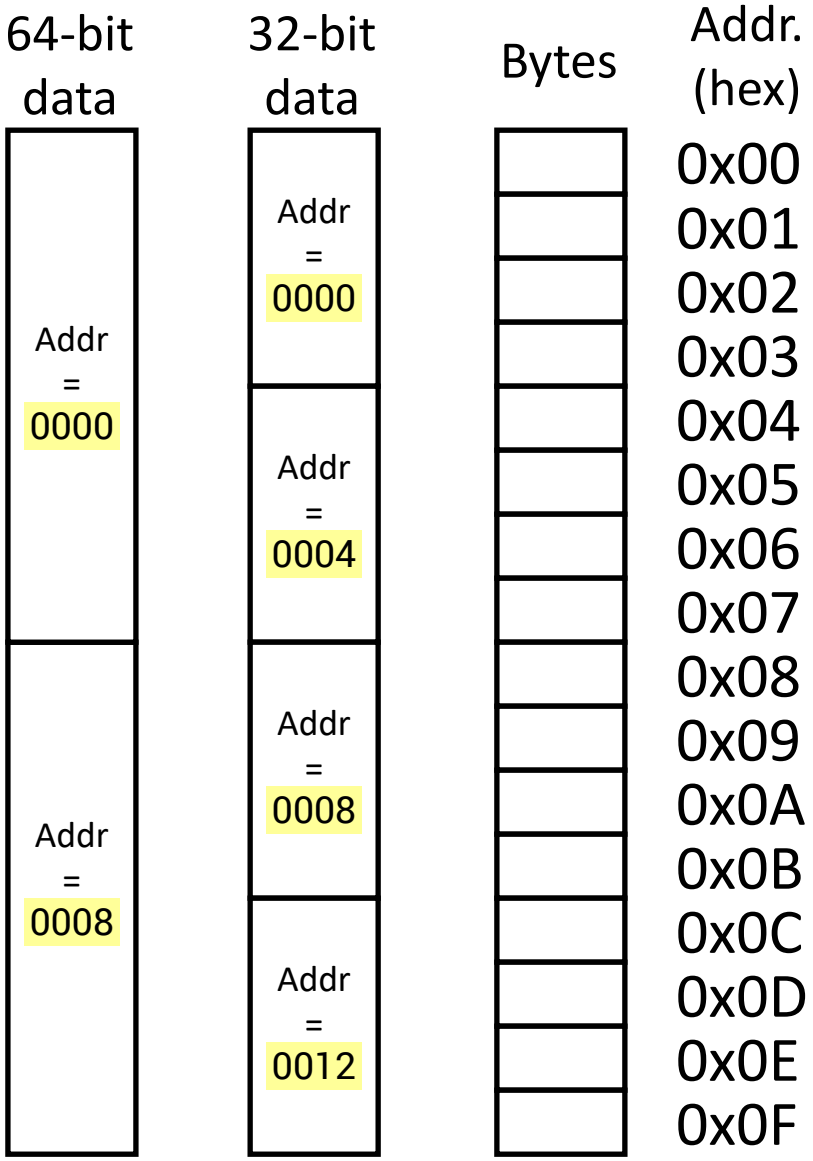
Java Data Type	C Data Type	32-bit (old)	x86-64
`boolean`	`bool`	1	1
`byte`	`char`	1	1
`char`		2	2
`short`	`short int`	2	2
`int`	`int`	4	4
`float`	`float`	4	4
	`long int`	4	8
`double`	`double`	8	8
`long`	`long long`	8	8
	`long double`	8	16
`(reference)`	**`pointer *`**	**4**	**8**

address size = word size

To use "`bool`" in C, you must `#include <stdbool.h>`

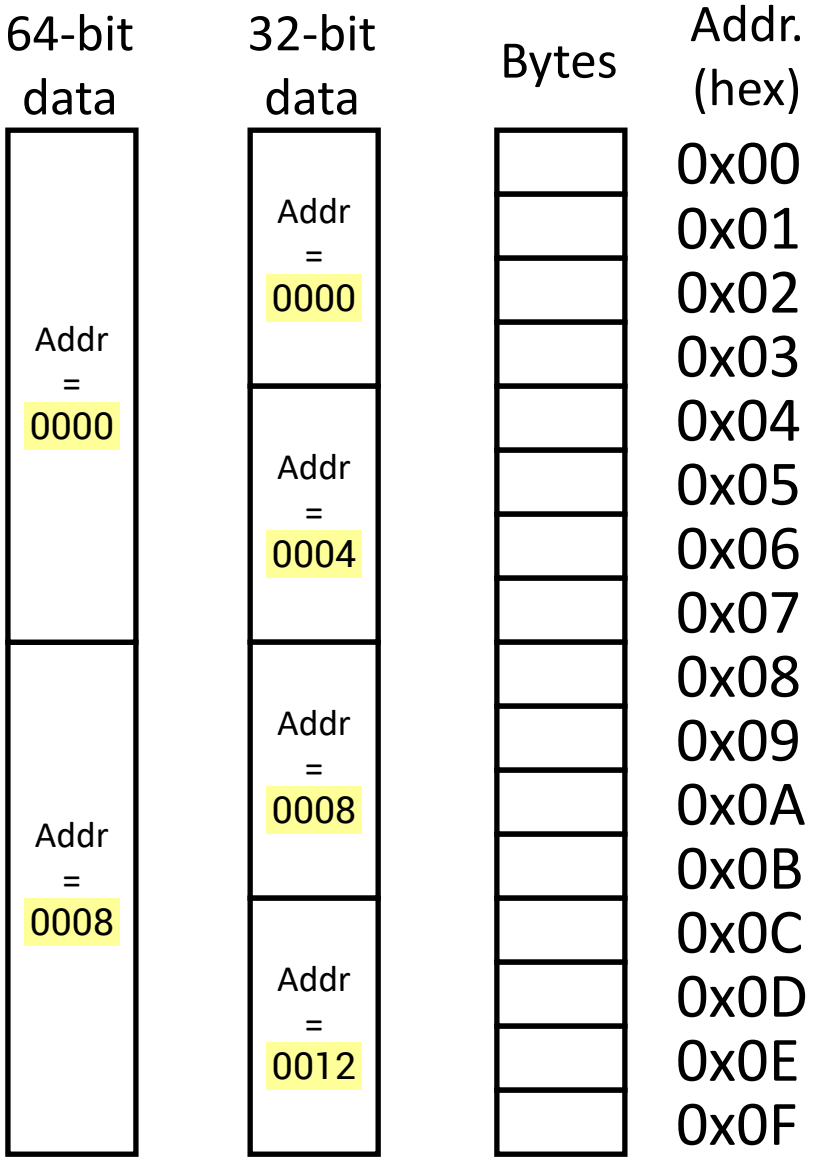
Address of Multibyte Data

- Addresses still specify locations of bytes in memory, but we can choose to *view* memory as a series of chunks of fixed-sized data instead
  - Addresses of successive chunks differ by data size
  - Which byte's address should we use for each word?
- The address of *any* chunk of memory is given by the address of the first byte
  - To specify a chunk of memory, need *both* its **address** and its **size**



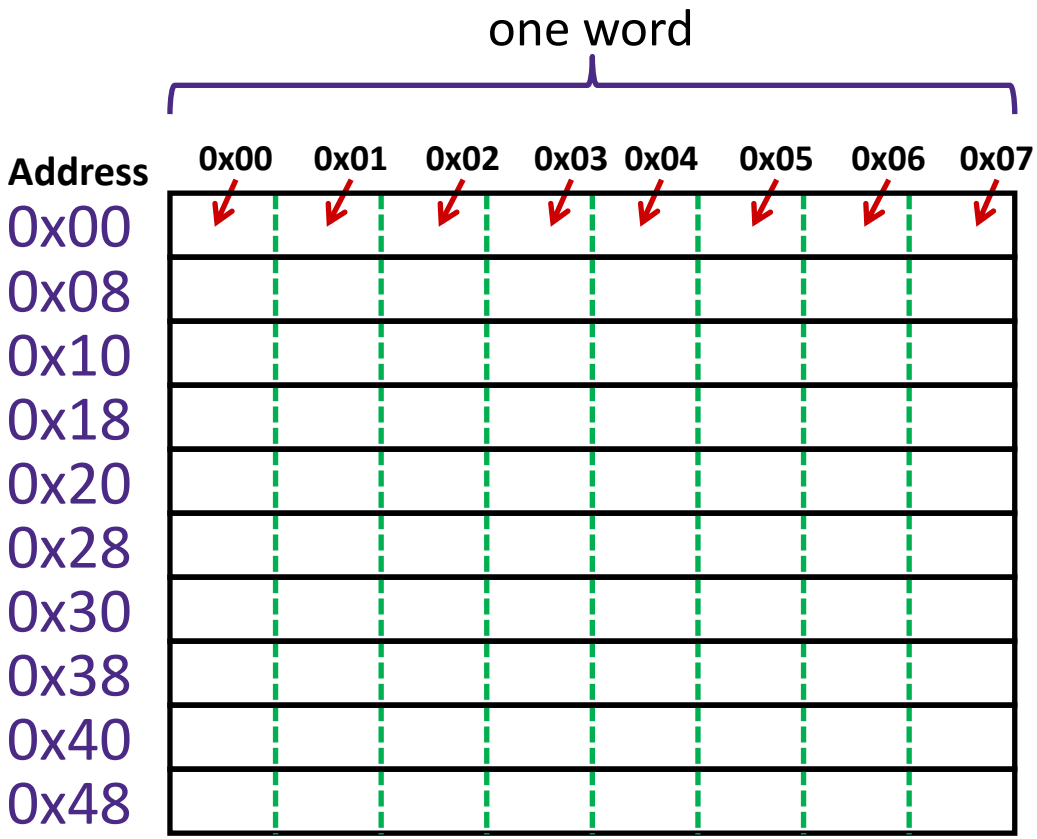
Alignment

- The address of a chunk of memory is considered aligned if its address is a multiple of its size
  - View memory as a series of consecutive chunks of this particular size and see if your chunk doesn't cross a boundary



A Picture of Memory (64-bit view)

- A "64-bit (8-byte) word-aligned" view of memory:
  - In this type of picture, each row is composed of 8 bytes
  - Each cell is a byte
  - An aligned, 64-bit chunk of data will fit on one row



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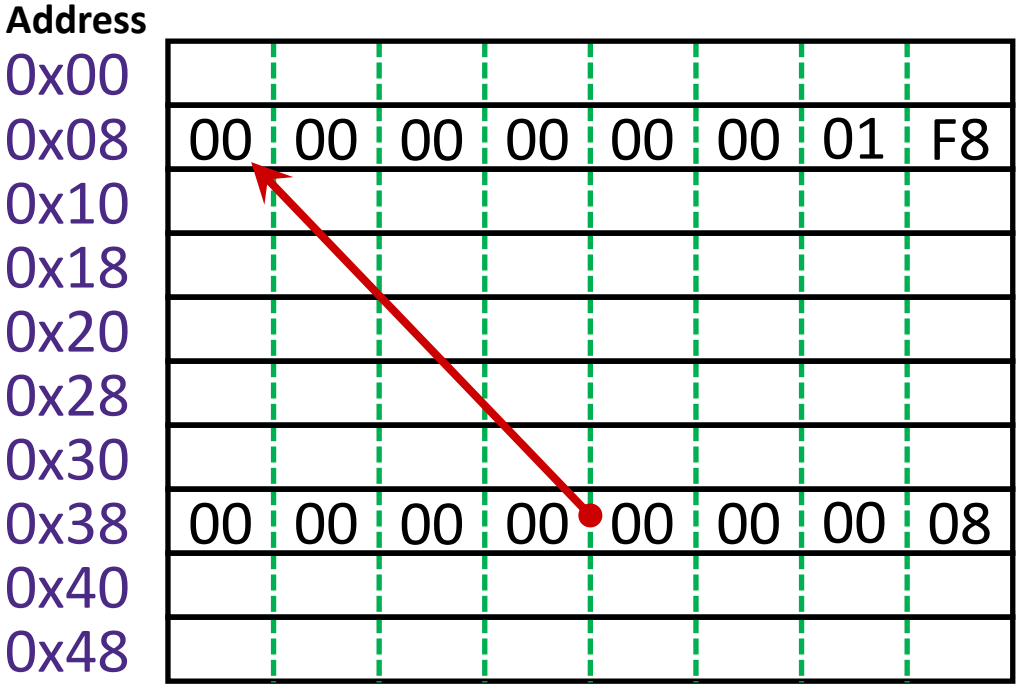
Addresses and Pointers

64-bit example
(pointers are 64-bits wide)

big-endian

- An *address* refers to a location in memory
- A *pointer* is a data object that holds an address
  - Address can point to *any* data
- Value 504 stored at address 0x08
  - $504_{10} = 1F8_{16}$
    = 0x 00 … 00 01 F8
- Pointer stored at 0x38 points to address 0x08

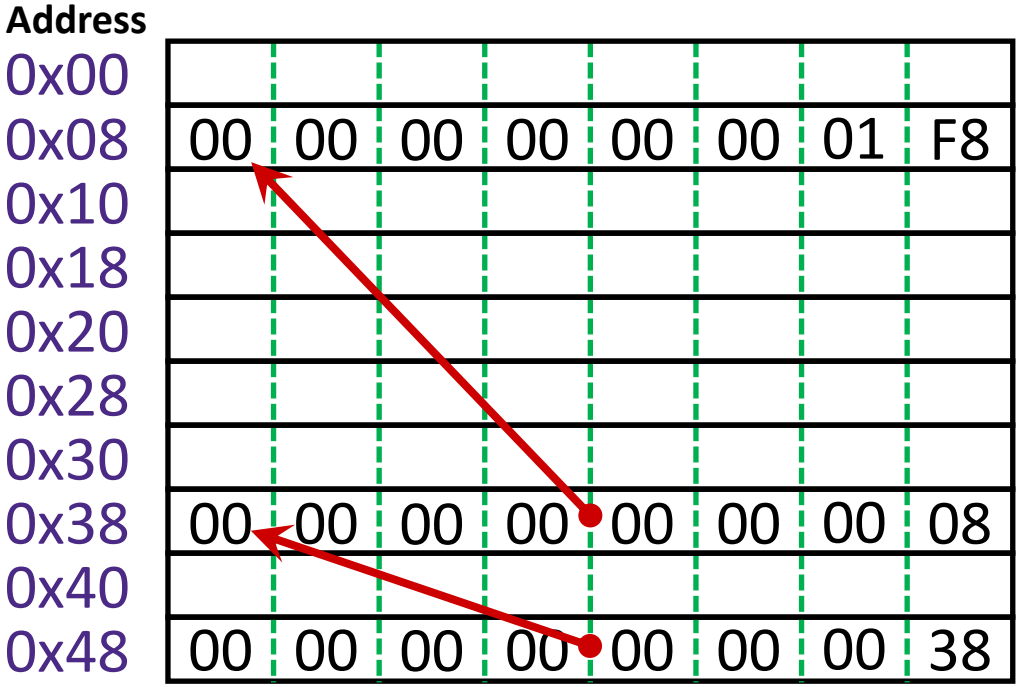


Addresses and Pointers

64-bit example
(pointers are 64-bits wide)

big-endian

- An *address* refers to a location in memory
- A *pointer* is a data object that holds an address
  - Address can point to *any* data
- Pointer stored at 0x48 points to address 0x38
  - Pointer to a pointer!
- Is the data stored at 0x08 a pointer?
  - Could be, depending on how you use it



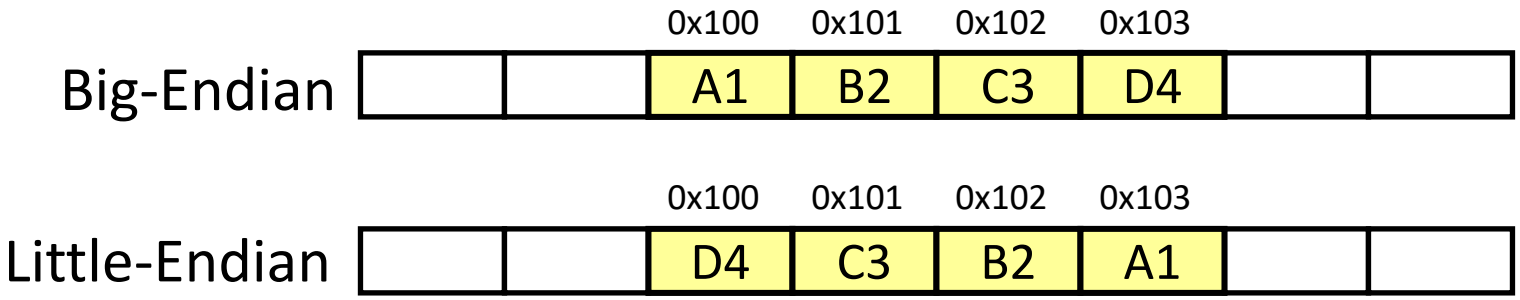
Byte Ordering

- How should bytes within a word be ordered *in memory?*
  - Want to keep consecutive bytes in consecutive addresses
  - **Example:** store the 4-byte (32-bit) `int`:
    0x A1 B2 C3 D4

- By convention, ordering of bytes called *endianness*
  - The two options are big-endian and little-endian
    - In which address does the least significant *byte* go?
    - Based on *Gulliver's Travels*: tribes cut eggs on different sides (big, little)

Byte Ordering

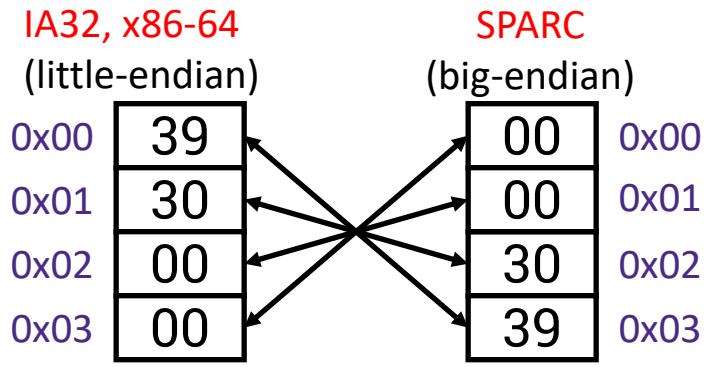
- Big-endian (SPARC, z/Architecture)
  - Least significant byte has highest address
- Little-endian (x86, x86-64)
  - Least significant byte has lowest address
- Bi-endian (ARM, PowerPC)
  - Endianness can be specified as big or little
- **Example:** 4-byte data 0xA1B2C3D4 at address 0x100



Byte Ordering Examples

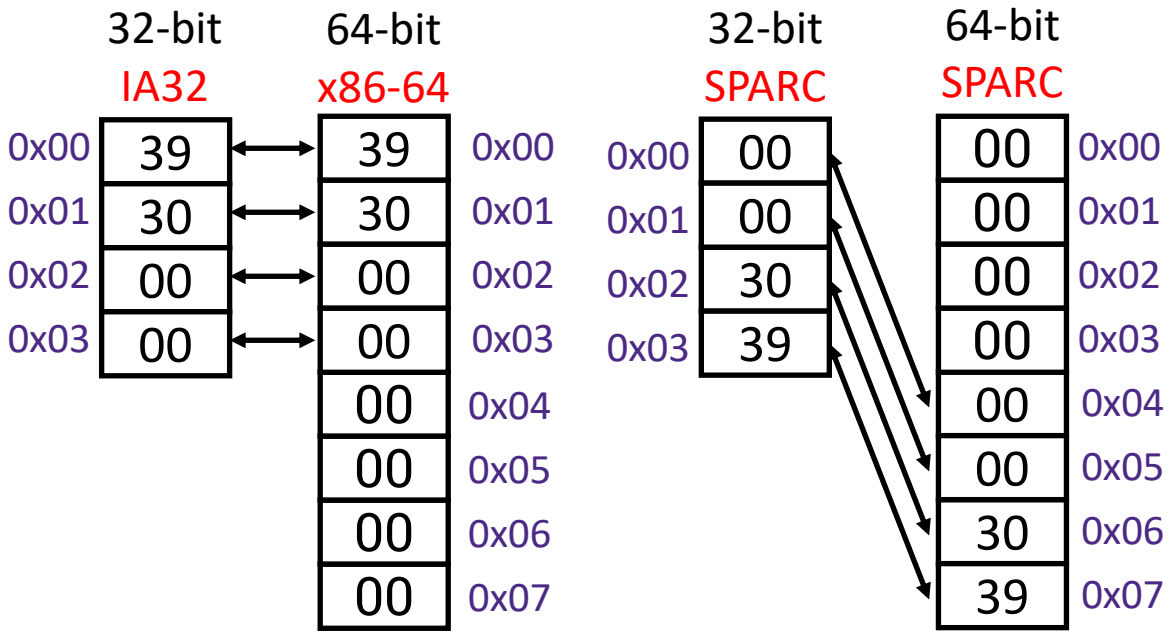
Decimal:	12345
Binary:	0011 0000 0011 1001
Hex:	3 0 3 9

```
int x = 12345;
// or x = 0x3039;
```




```
long int y = 12345;
// or y = 0x3039;
```

(A `long int` is the size of a word)



Polling Question

- We store the value 0x 01 02 03 04 as a ***word*** at address 0x100 in a big-endian, 64-bit machine
- What is the ***byte of data*** stored at address 0x104?
  - Vote in Ed Lessons

A. 0x04

B. 0x40

C. 0x01

D. 0x10

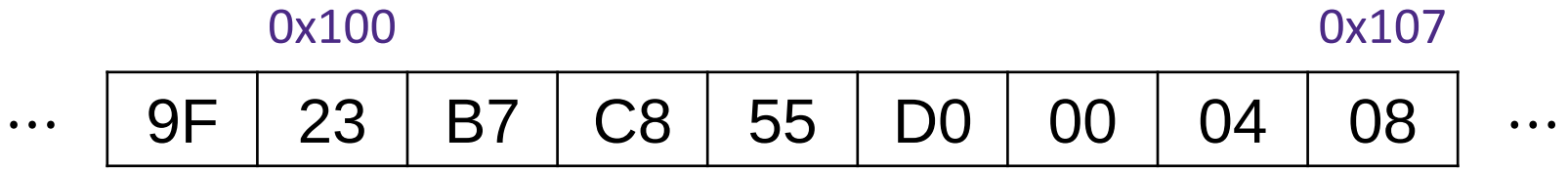
E. We're lost…

Endianness

- *Endianness only applies to memory storage*
- Often programmer can ignore endianness because it is handled for you
  - Bytes wired into correct place when reading or storing from memory (hardware)
  - Compiler and assembler generate correct behavior (software)
- Endianness still shows up:
  - Logical issues: accessing different amount of data than how you stored it (*e.g.*, store `int`, access byte as a `char`)
  - Need to know exact values to debug memory errors
  - Manual translation to and from machine code (in 351)

Challenge Question

- Assume the state of memory is as shown below for a little-endian machine.



- If we (1) *read* the value of an int at address 0x102, (2) add 8 to it, and then (3) store the new value as an int at address 0x104, which of the following addresses retain their original value?

A. 0x102 **B. 0x104** **C. 0x105** **D. 0x107**

Summary

- Memory is a long, *byte-addressed* array
  - Word size bounds the size of the *address space* and memory
  - Different data types use different number of bytes
  - Address of chunk of memory given by address of lowest byte in chunk
  - Object of K bytes is *aligned* if it has an address that is a multiple of K
- Pointers are data objects that hold addresses
- Endianness determines memory storage order for multi-byte data