

Memory, Data, & Addressing II

CSE 351 Winter 2020

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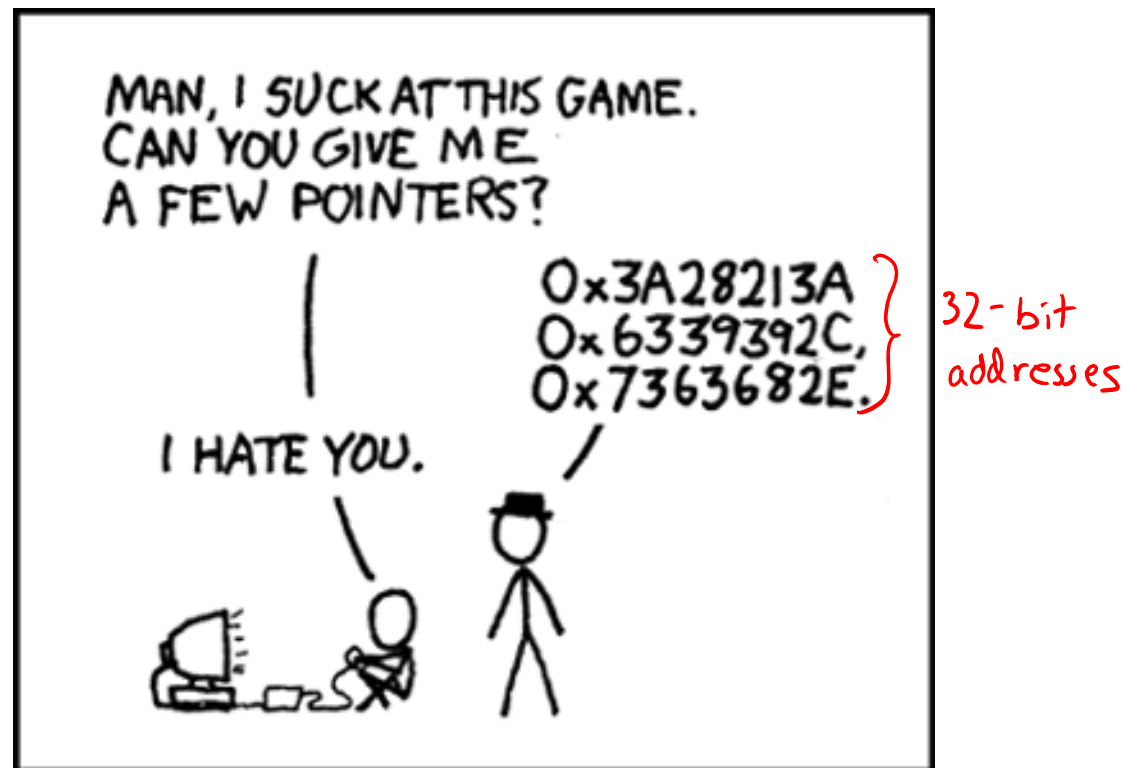
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Josie Lee

Jeffery Tian

Callum Walker

Eddy (Tianyi) Zhou



<http://xkcd.com/138/>

Administrivia

- ❖ Lab 0 due today @ 11:59 pm
 - *You will be revisiting this program throughout this class!*
- ❖ hw2 due Monday, hw3 due Wednesday @ 11:00 am
 - Autograded, unlimited tries, no late submissions
- ❖ Lab 1a released today, due next Friday (1/17)
 - Pointers in C
 - Reminder: last submission graded, *individual* work

Late Days

- ❖ You are given **5 late lab days** for the whole quarter
 - Tokens can only apply to Labs
 - No benefit to having leftover tokens
- ❖ Count lateness in *days* (even if just by a second)
 - Special: weekends count as *one day*
 - No submissions accepted more than two days late
- ❖ Late penalty is 20% deduction of your score per day
 - Only late labs are eligible for penalties
 - Penalties applied at end of quarter to *maximize* your grade
- ❖ Use at own risk – don't want to fall too far behind
 - Intended to allow for unexpected circumstances

Review Questions

- 1) If the word size of a machine is 64-bits, which of the following is usually true? (pick all that apply)
 - a) 64 bits is the size of a pointer T
 - b) 64 bits is the size of an integer F
 - c) 64 bits is the width of a register T
- 2) (True/~~False~~) By looking at the bits stored in memory, I can tell if a particular 4-bytes is being used to represent an integer, floating point number, or instruction.
- 3) If the size of a pointer on a machine is 6 bits, the address space is how many bytes? 0b -----

represent 2^6 things \rightarrow 2^6 addresses \rightarrow 2^6 bytes of data
64 B

Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
 - Binary, hexadecimal, fixed-widths
- ❖ Organizing and addressing data in memory
 - Memory is a byte-addressable array
 - Machine “word” size = address size = register size
 - Endianness – ordering bytes in memory
- ❖ **Manipulating data in memory using C**
 - **Assignment**
 - **Pointers, pointer arithmetic, and arrays**
- ❖ Boolean algebra and bit-level manipulations

Addresses and Pointers in C

* is also used with variable declarations

- ❖ & = “address of” operator
- ❖ * = “value at address” or “dereference” operator

equivalent
 datatype
 `int* ptr;`
 `int *ptr;`

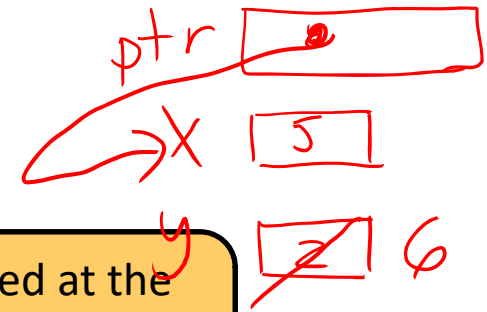
Declares a variable, `ptr`, that is a pointer to (i.e. holds the address of) an `int` in memory

`int x = 5;`
 `int y = 2;`

Declares two variables, `x` and `y`, that hold `ints`, and *initializes* them to 5 and 2, respectively

`ptr = &x;`

Sets `ptr` to the address of `x` (“`ptr` points to `x`”)



`y = 1 + *ptr;`

“Dereference `ptr`”

Sets `y` to “1 plus the value stored at the address held by `ptr`.” Because `ptr` points to `x`, this is equivalent to `y=1+x;`

What is `*(&y)` ?

returns value stored in `y` (equivalent to just using `y`)

Assignment in C

- ❖ A variable is represented by a location
- ❖ Declaration ≠ initialization (initially holds “garbage”)
- ❖ `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

	0x00	0x01	0x02	0x03	
0x00	A7	00	32	00	
<u>0x04</u>	00	01	29	F3	⊗
0x08	EE	EE	EE	EE	
0x0C	FA	CE	CA	FE	
0x10	26	00	00	00	
0x14	00	00	10	00	
<u>0x18</u>	01	00	00	00	y
0x1C	FF	00	F4	96	
0x20	DE	AD	BE	EF	
0x24	00	00	00	00	

Assignment in C

32-bit example
(pointers are 32-bits wide)

little-endian

- ❖ A variable is represented by a location
- ❖ Declaration ≠ initialization (initially holds “garbage”)
- ❖ `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

	0x00	0x01	0x02	0x03	
0x00					
0x04	00	01	29	F3	x
0x08					
0x0C					
0x10					
0x14					
0x18	01	00	00	00	y
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ `int x, y;`

❖ `x = 0;` *pad* → `0x00 00 00 00`
 ↑ *int (4 bytes)*

	0x00	0x01	0x02	0x03	
0x00					
0x04	00	00	00	00	x
0x08					
0x0C					
0x10					
0x14					
0x18	01	00	00	00	y
0x1C					
0x20					
0x24					

Assignment in C

32-bit example
(pointers are 32-bits wide)

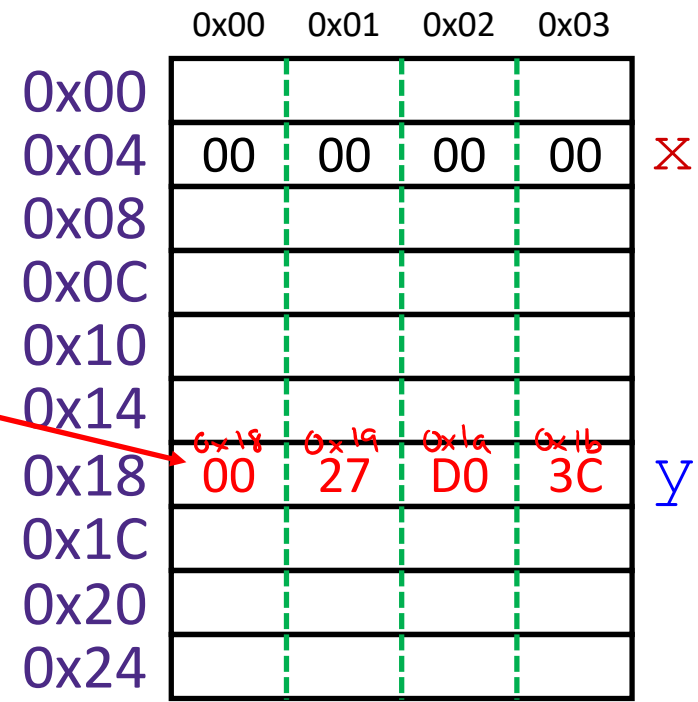
& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ `int x, y;`

❖ `x = 0;`

❖ `y = 0x3CD02700;`
 least significant byte
 little endian!



Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

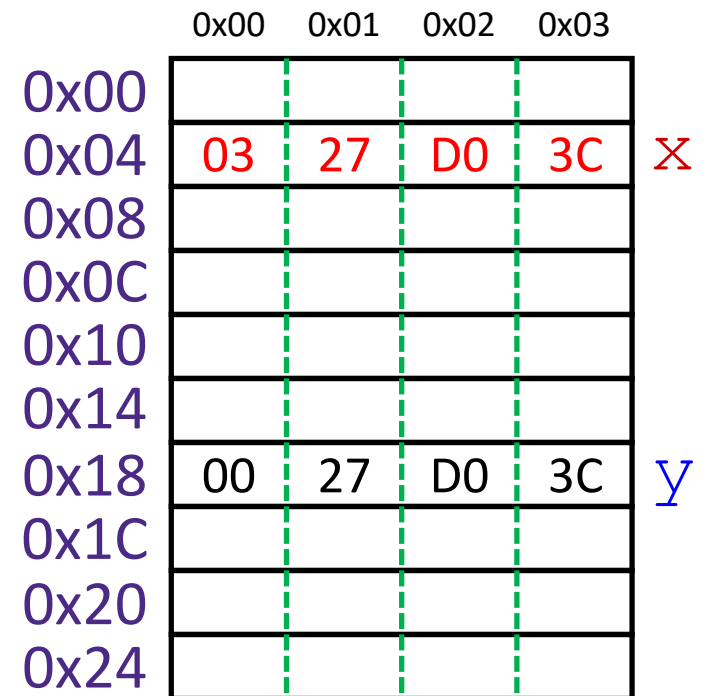
❖ `int x, y;`

❖ `x = 0;`

❖ `y = 0x3CD02700;`

❖ `x = y + 3;`

- Get value at `y`, add 3, store in `x`



Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

❖ `int x, y;`

❖ `x = 0;`

❖ `y = 0x3CD02700;`

❖ `x = y + 3;`

- Get value at `y`, add 3, store in `x`

❖ `int*` `z;` ← pointer to an int

- `z` is at address `0x20`

	0x00	0x01	0x02	0x03	
0x00					
<u>0x04</u>	03	27	D0	3C	X
0x08					
0x0C					
0x10					
0x14					
<u>0x18</u>	00	27	D0	3C	Y
0x1C					
<u>0x20</u>	DE	AD	BE	EF	Z
0x24					

initial value is whatever bits were already there! ("garbage")

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

```
❖ int x, y;
```

```
❖ x = 0;
```

```
❖ y = 0x3CD02700;
```

```
❖ x = y + 3;
```

- Get value at *y*, add 3, store in *x*

```
❖ int* z = 0x18&y + 3; // expect 0x1b
```

- Get address of *y*, "add 3", store in *z*

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	X
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	Y
0x1C					
0x20	24	00	00	00	Z
0x24					

↑ get this instead

Pointer arithmetic

Pointer Arithmetic

- ❖ Pointer arithmetic is scaled by the size of target type
 - In this example, `sizeof(int) = 4`
- ❖ `int* z = &y + 3;` *4 bytes*
 - Get address of `y`, add $3 * \text{sizeof}(\text{int})$, store in `z`
 - $\&y = 0x18 = 1 * 16^1 + 8 * 16^0 = 24$
 - $24 + 3 * (4) = 36 = 2 * 16^1 + 4 * 16^0 = 0x24$
- ❖ **Pointer arithmetic can be dangerous!**
 - Can easily lead to bad memory accesses
 - Be careful with data types and *casting*

Assignment in C

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ `int x, y;`
- ❖ `x = 0;`
- ❖ `y = 0x3CD02700;`
- ❖ `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- ❖ `int* z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`
- ❖ `*z = y;`
 - What does this do?

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	x
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	y
0x1C					
0x20	24	00	00	00	z
0x24					

Assignment in C

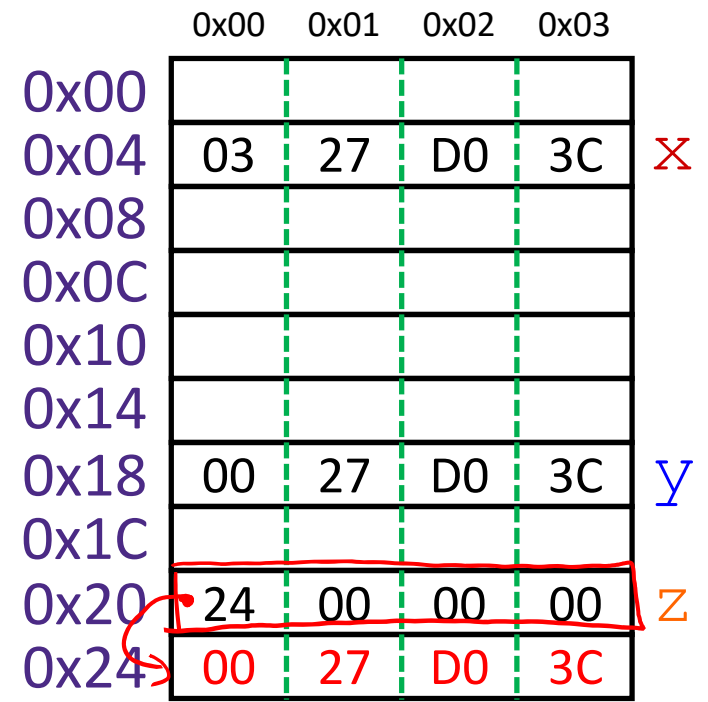
- ❖ `int x, y;`
- ❖ `x = 0;`
- ❖ `y = 0x3CD02700;`
- ❖ `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- ❖ `int* z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`
- ❖ `*z = y;`

The target of a pointer is also a location

 - Get value of `y`, put in address stored in `z`

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"



Arrays in C

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

Declaration: `int a[6]; // &a is 0x10`

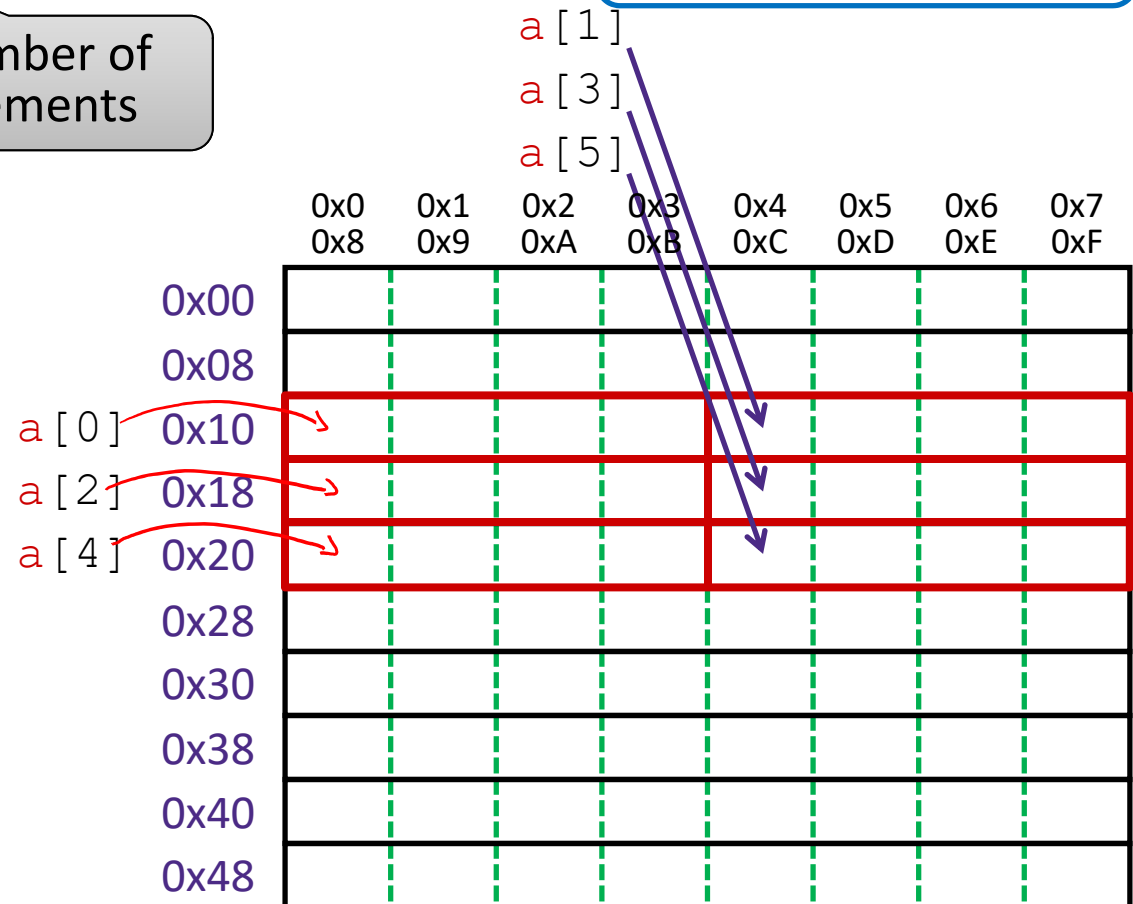
4 bytes each

element type

name

number of elements

64-bit example
(pointers are 64-bits wide)



Arrays in C

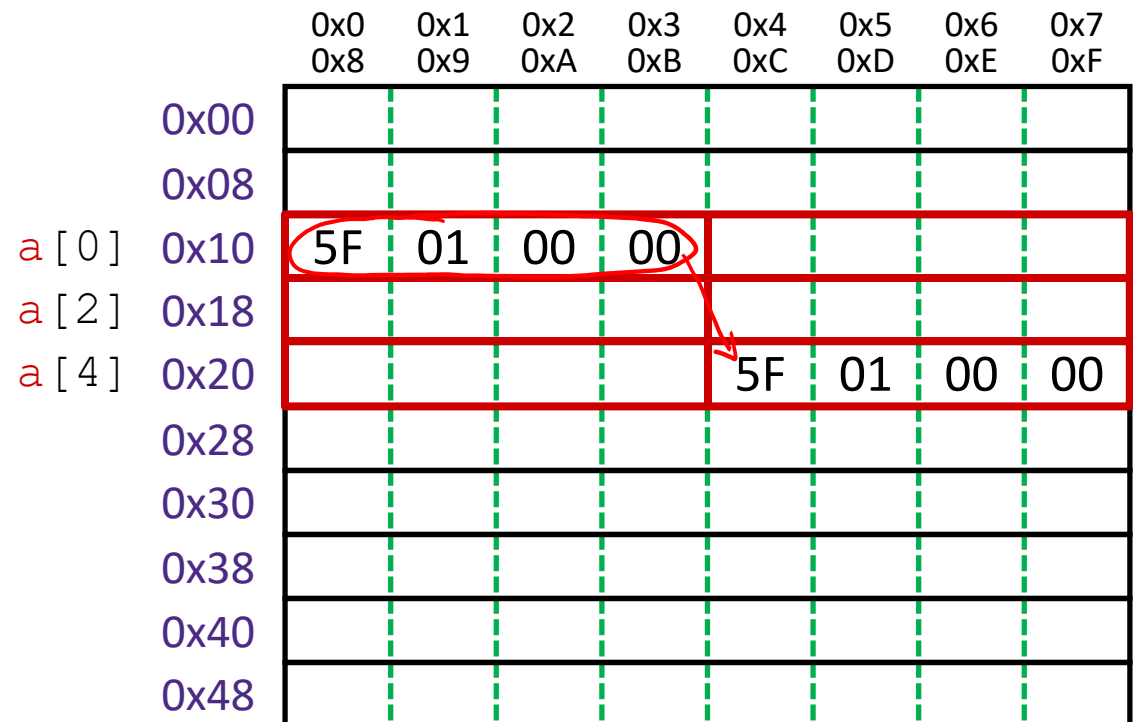
Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes



Arrays in C

Declaration: `int a[6];`

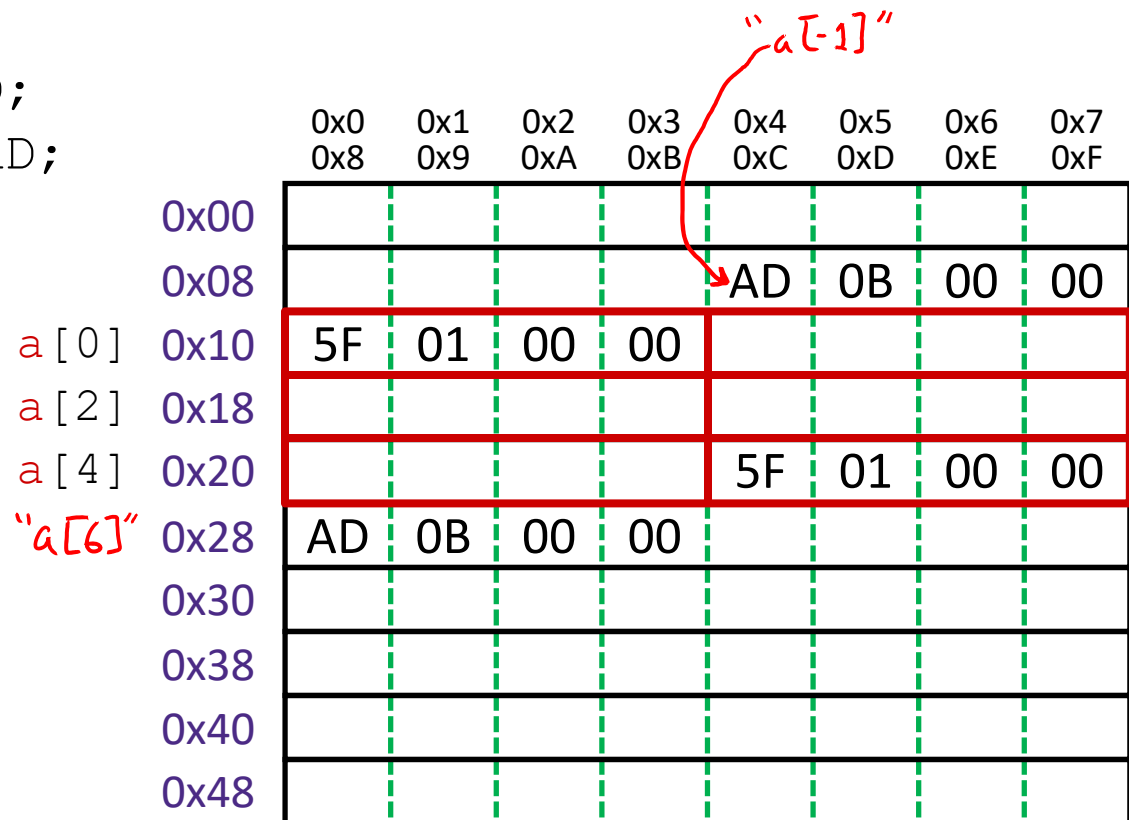
Indexing: \checkmark `a[0] = 0x015f;`
 \checkmark `a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

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Arrays in C

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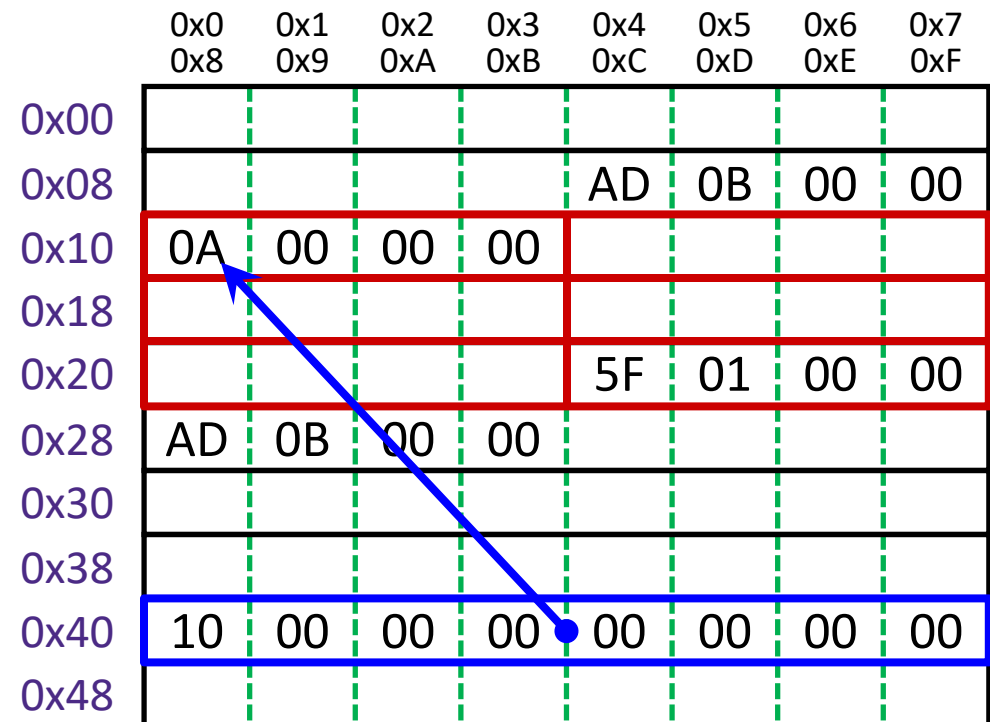
Pointers: `int* p;`

equivalent `[`
`p = a;`
`p = &a[0];`
`*p = 0xA;`
`]` `a[0]`
`a[2]`
`a[4]`

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

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Arrays in C

Declaration: `int a[6];`

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No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \end{array} \right.$

pointer arithmetic: $0x10 + 1 \rightarrow 0x14$
 $p = p + 2;$

$0x10 + 2 \rightarrow 0x18$

`a[0]`
`a[2]`
`a[4]`

`p`

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

$p[i] \iff *(p+i)$

	0x0	0x1	0x2	0x3	0x4	0x5	0x6	0x7
	0x8	0x9	0xA	0xB	0xC	0xD	0xE	0xF
0x00								
0x08					AD	0B	00	00
0x10	0A	00	00	00	0B	00	00	00
0x18								
0x20					5F	01	00	00
0x28	AD	0B	00	00				
0x30								
0x38								
0x40	10	00	00	00	00	00	00	00
0x48								

$a + 2 * \text{sizeof}(int) = 0x18$

Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$ $a[0]$
 $a[2]$
 $a[4]$

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{array} \right.$

store at 0x18 → $*p = 0xB + 1 = 0xC$

(no pointer arithmetic)

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes

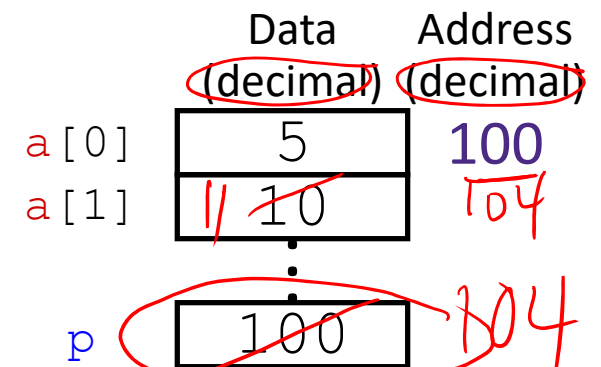
	0x0	0x1	0x2	0x3	0x4	0x5	0x6	0x7
	0x8	0x9	0xA	0xB	0xC	0xD	0xE	0xF
0x00								
0x08					AD	0B	00	00
0x10	0A	00	00	00	0B	00	00	00
0x18	0C	00	00	00				
0x20					5F	01	00	00
0x28	AD	0B	00	00				
0x30								
0x38								
0x40	18	00	00	00	00	00	00	00
0x48								

Question: The variable values after Line 3 executes are shown on the right. What are they after Line 4 & 5?

- Vote at <http://pollev.com/rea>

```

1 void main() {
2   int a[] = {5, 10};
3   int* p = a;
4   p = p + 1;
5   *p = *p + 1;
6 }
```



after line 4

p	*p	a[0]	a[1]
----------	-----------	-------------	-------------

after line 5

p	*p	a[0]	a[1]
----------	-----------	-------------	-------------

(A) 101 10 5 10 then 101 11 5 11

→ (B) 104 10 5 10 then 104 11 5 11

(C) 100 6 6 10 then 101 6 6 10

(D) 100 6 6 10 then 104 6 6 10

Representing strings

- ❖ C-style string stored as an array of bytes (**char***)
 - Elements are one-byte **ASCII codes** for each character
 - No “String” keyword, unlike Java

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	”	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	'	55	7	71	G	87	W	103	g	119	w
40	(56	8	72	H	88	X	104	h	120	x
41)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

ASCII: American Standard Code for Information Interchange

Null-Terminated Strings

- ❖ **Example:** "Donald Trump" stored as a 13-byte array

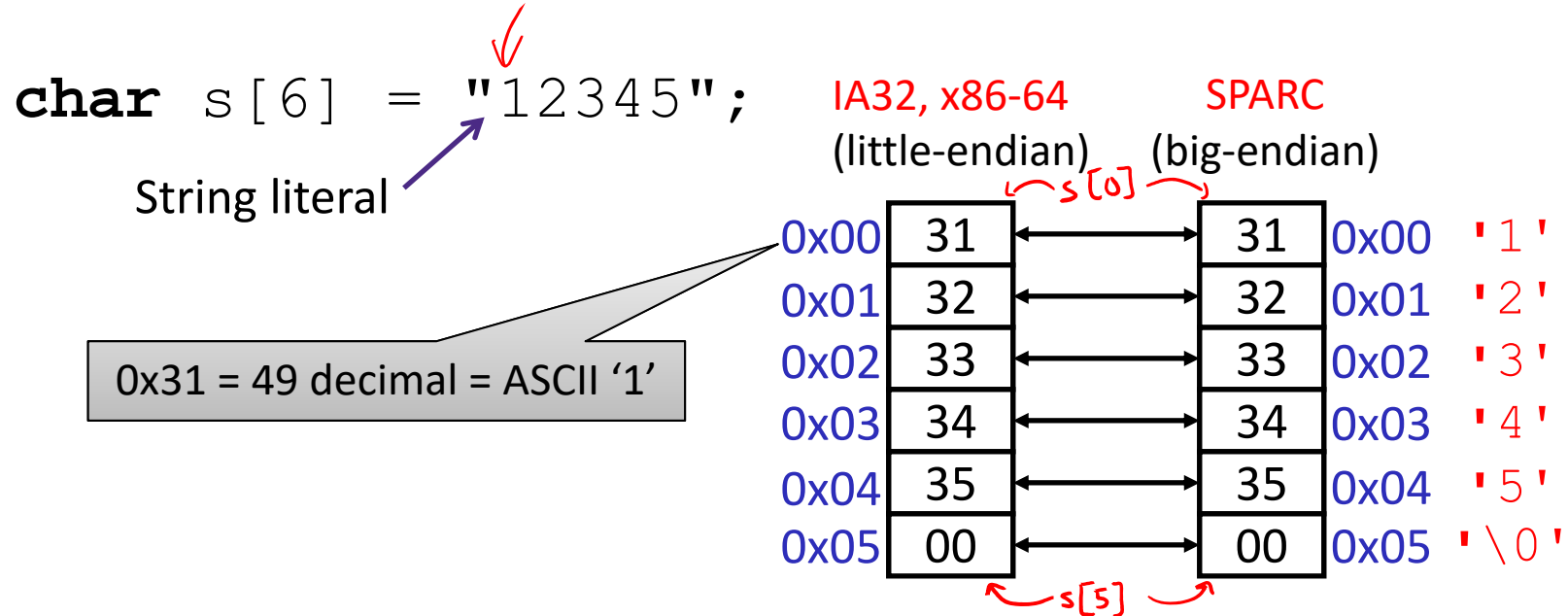
<i>Decimal:</i>	68	111	110	97	108	100	32	84	114	117	109	112	0
<i>Hex:</i>	0x44	0x6F	0x6E	0x61	0x6C	0x64	0x20	0x54	0x72	0x75	0x6D	0x70	0x00
<i>Text:</i>	D	o	n	a	l	d		T	r	u	m	p	\0

13 bytes total!

- ❖ Last character followed by a 0 byte (' \0 ') (a.k.a. "null terminator")
 - Must take into account when allocating space in memory
 - Note that ⁴⁸'0' ≠ ⁰' \0 ' (i.e. character 0 has non-zero value)
- ❖ How do we compute the length of a string?
 - Traverse array until null terminator encountered

C (char = 1 byte)

Endianness and Strings



- ❖ Byte ordering (endianness) is not an issue for 1-byte values
 - The whole array does not constitute a single value
 - Individual elements are values; chars are single bytes

Examining Data Representations

- ❖ Code to print byte representation of data
 - Any data type can be treated as a *byte array* by **casting** it to `char`
 - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {
    int i;
    for (i = 0; i < len; i++)
        printf("%p\t0x%.2x\n", start+i, *(start+i));
    printf("\n");
}
```

`printf` **directives:**

<code>%p</code>	Print pointer
<code>\t</code>	Tab
<code>%x</code>	Print value as hex
<code>\n</code>	New line

Examining Data Representations

- ❖ Code to print byte representation of data
 - Any data type can be treated as a *byte array* by **casting** it to `char`
 - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {
    int i;
    for (i = 0; i < len; i++)
        printf("%p\t0x%.2x\n", start+i, *(start+i));
    printf("\n");
}
```

format string

*pointer arithmetic on char**

```
void show_int(int x) {
    show_bytes( (char *) &x, sizeof(int) );
}
```

*int**

4 bytes

"cast" (treat as)

show_bytes Execution Example

```
int x = 12345; // 0x00003039
printf("int x = %d;\n", x);
show_int(x); // show_bytes((char *) &x, sizeof(int));
```

❖ Result (Linux x86-64):

- **Note:** The addresses will change on each run (try it!), but fall in same general range

```
int x = 12345;
0x7fffb7f71dbc      0x39
0x7fffb7f71dbd      0x30
0x7fffb7f71dbe      0x00
0x7fffb7f71dbf      0x00
```

Summary

- ❖ Assignment in C results in value being put in memory location
- ❖ Pointer is a C representation of a data address
 - $\&$ = “address of” operator
 - $*$ = “value at address” or “dereference” operator
- ❖ Pointer arithmetic scales by size of target type
 - Convenient when accessing array-like structures in memory
 - Be careful when using – particularly when *casting* variables
- ❖ Arrays are adjacent locations in memory storing the same type of data object
 - Strings are null-terminated arrays of characters (ASCII)

Assignment in C - Handout

32-bit example
(pointers are 32-bits wide)

& = "address of"
* = "dereference"

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value*
 - Store RHS value at LHS location

- ❖ `int x, y;`
- ❖ `x = 0;`
- ❖ `y = 0x3CD02700;`
- ❖ `x = y + 3;`
- ❖ `int* z = &y + 3;`
- ❖ `*z = y;`

