Caches IV

CSE 351 Summer 2020

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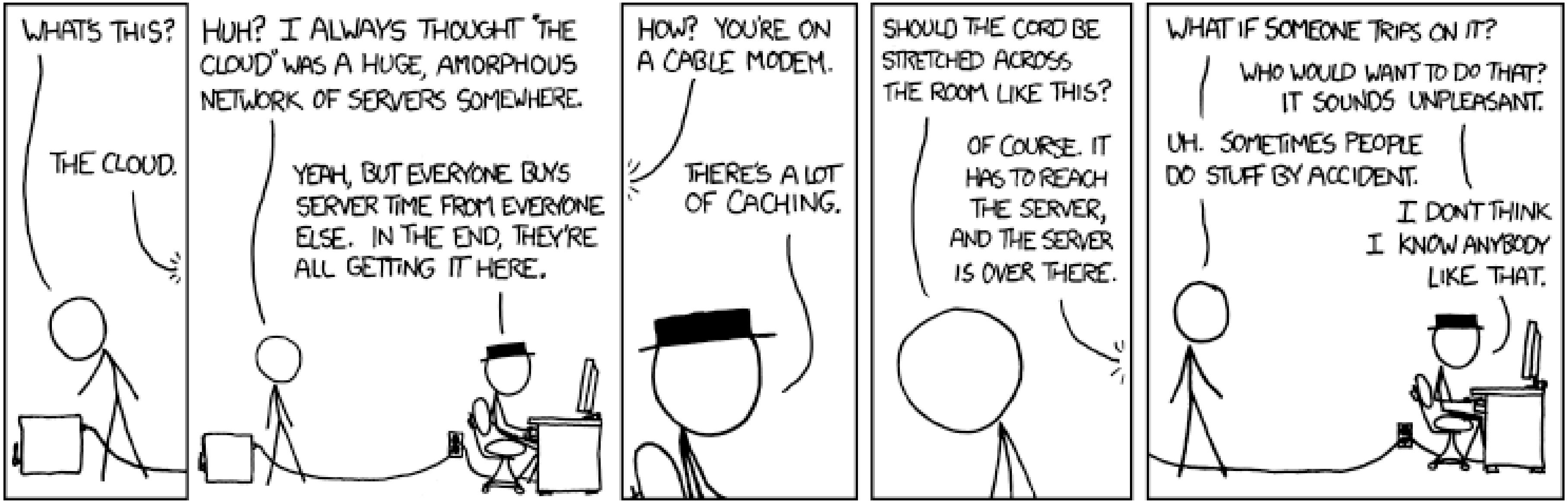


http://xkcd.com/908/

Administrivia

- Questions doc: https://tinyurl.com/CSE351-8-3

- hw16 due Wednesday (8/5) – 10:30am
- hw17 due Friday (8/7) – 10:30am

- Lab 4 due Wednesday (8/12) – 11:59pm
  - All about caches!

- Unit Summary 2 Due Wednesday (8/5) – 11:59pm

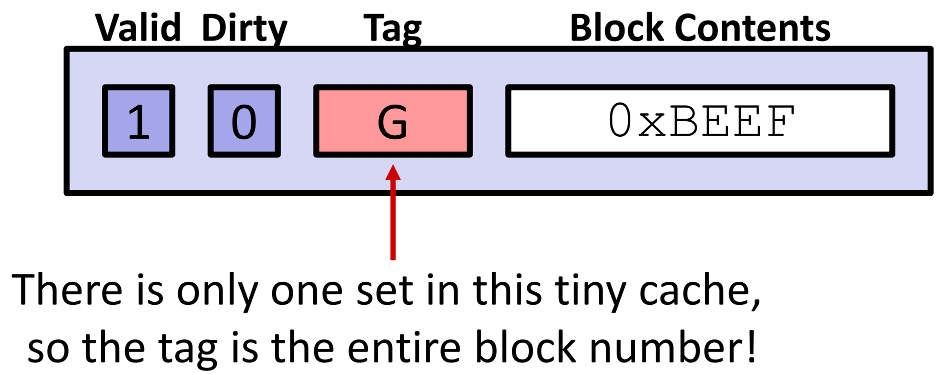
What about writes?

- Multiple copies of data may exist:
  - multiple levels of cache and main memory
- What to do on a write-hit?
  - Write-through: write immediately to next level
  - Write-back: defer write to next level until line is evicted (replaced)
    - Must track which cache lines have been modified ("*dirty bit*")
- What to do on a write-miss?
  - Write allocate: ("fetch on write") load into cache, then execute the write-hit policy
    - Good if more writes or reads to the location follow
  - No-write allocate: ("write around") just write immediately to next level
- Typical caches:
  - Write-back + Write allocate, usually
  - Write-through + No-write allocate, occasionally

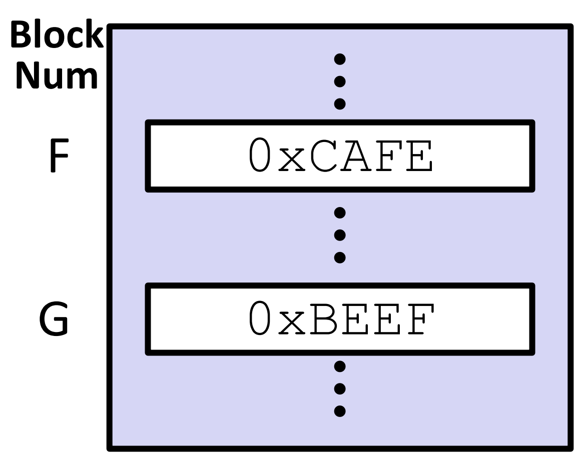
Write-back, Write Allocate Example

Note: While unrealistic, this example assumes that all requests have offset 0 and are for a block's worth of data.

Cache:



Memory:

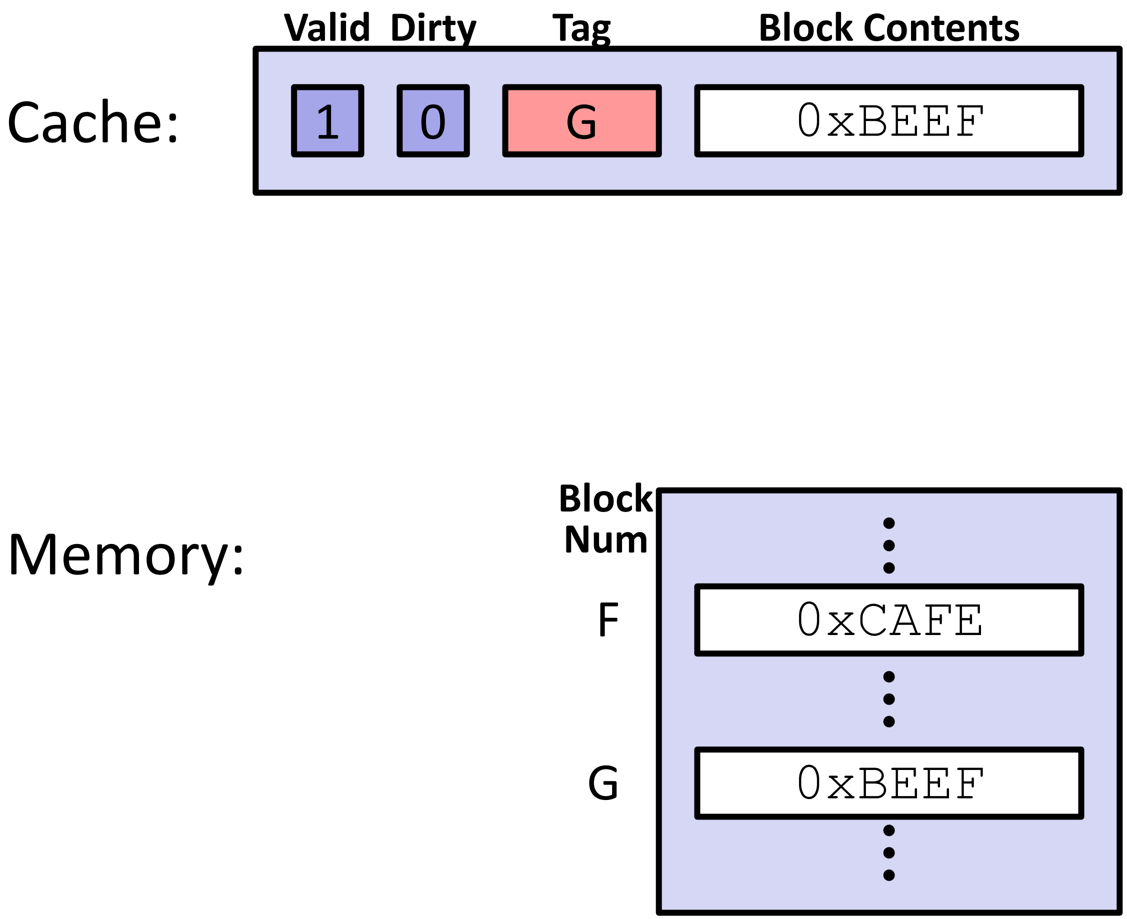


Write-back, Write Allocate Example

Not valid x86, just using block num instead of full byte address to keep the example simple

1) `mov $0xFACE, (F)`

Write Miss!



Step 1: Bring **F** into cache

Write-back, Write Allocate Example

1) `mov $0xFACE, (F)`

Write Miss

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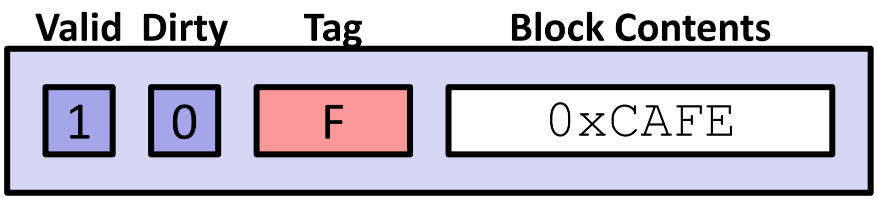
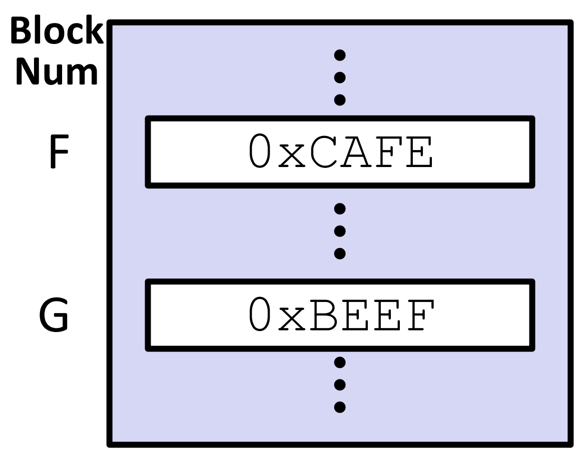


Memory:

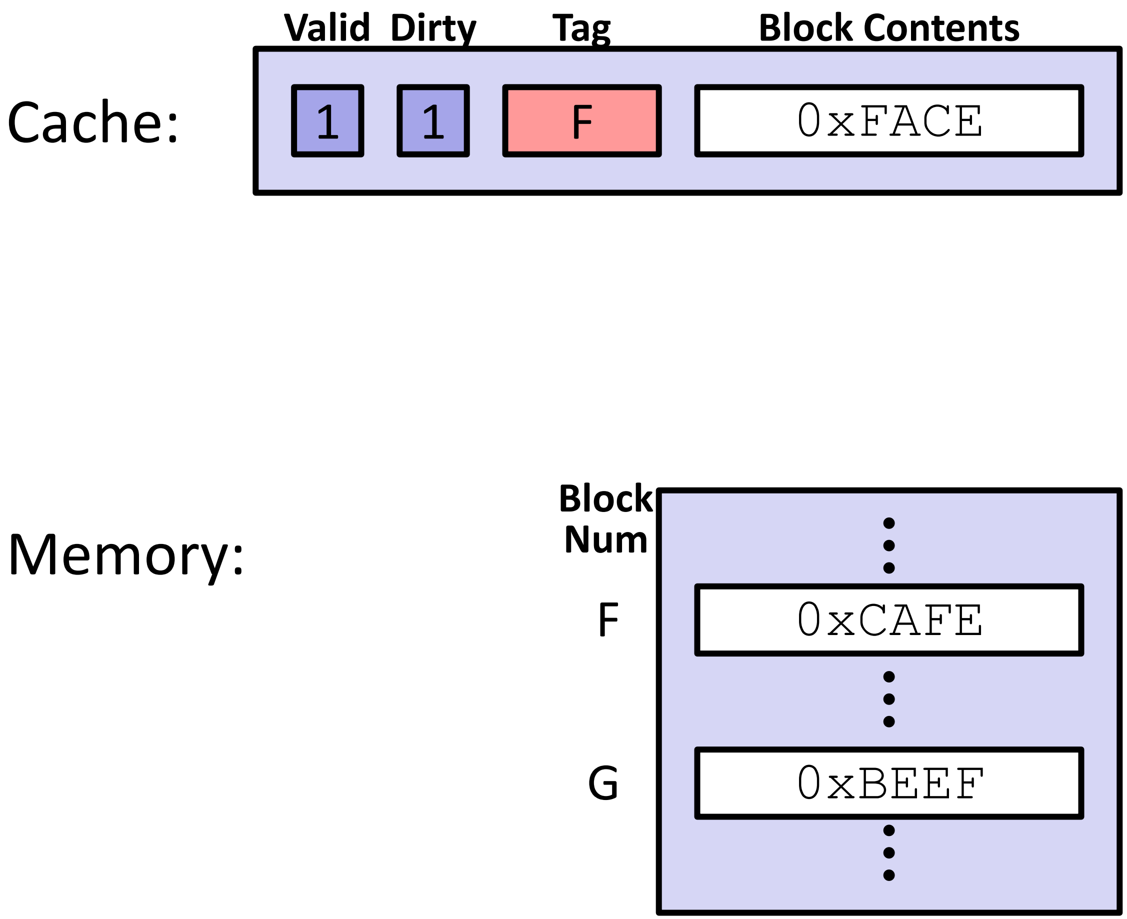


Step 1: Bring **F** into cache

Step 2: Write `0xFACE` to cache only and set the dirty bit

Write-back, Write Allocate Example

1) `mov $0xFACE, (F)`
   Write Miss



Step 1: Bring **F** into cache

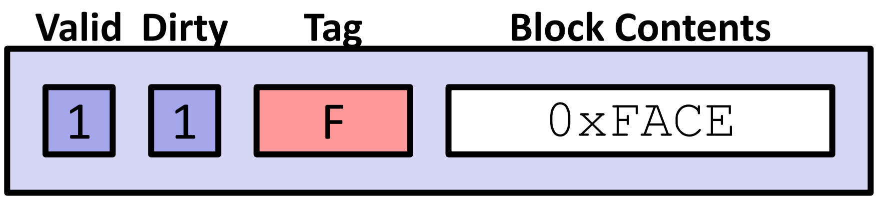
Step 2: Write `0xFACE` to cache only and set the dirty bit

Write-back, Write Allocate Example

1) `mov $0xFACE, (F)`
Write Miss

2) `mov $0xFEED, (F)`
Write Hit!

Cache:



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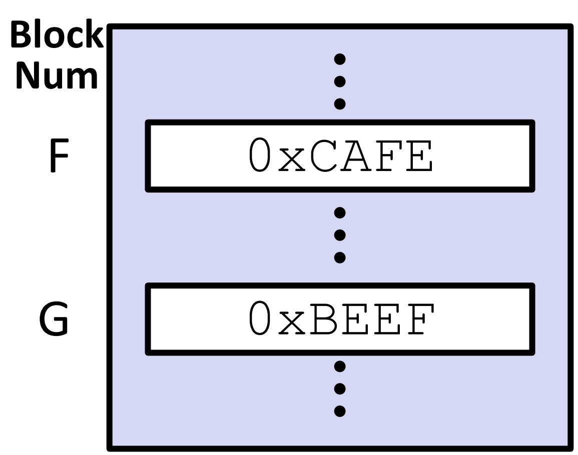
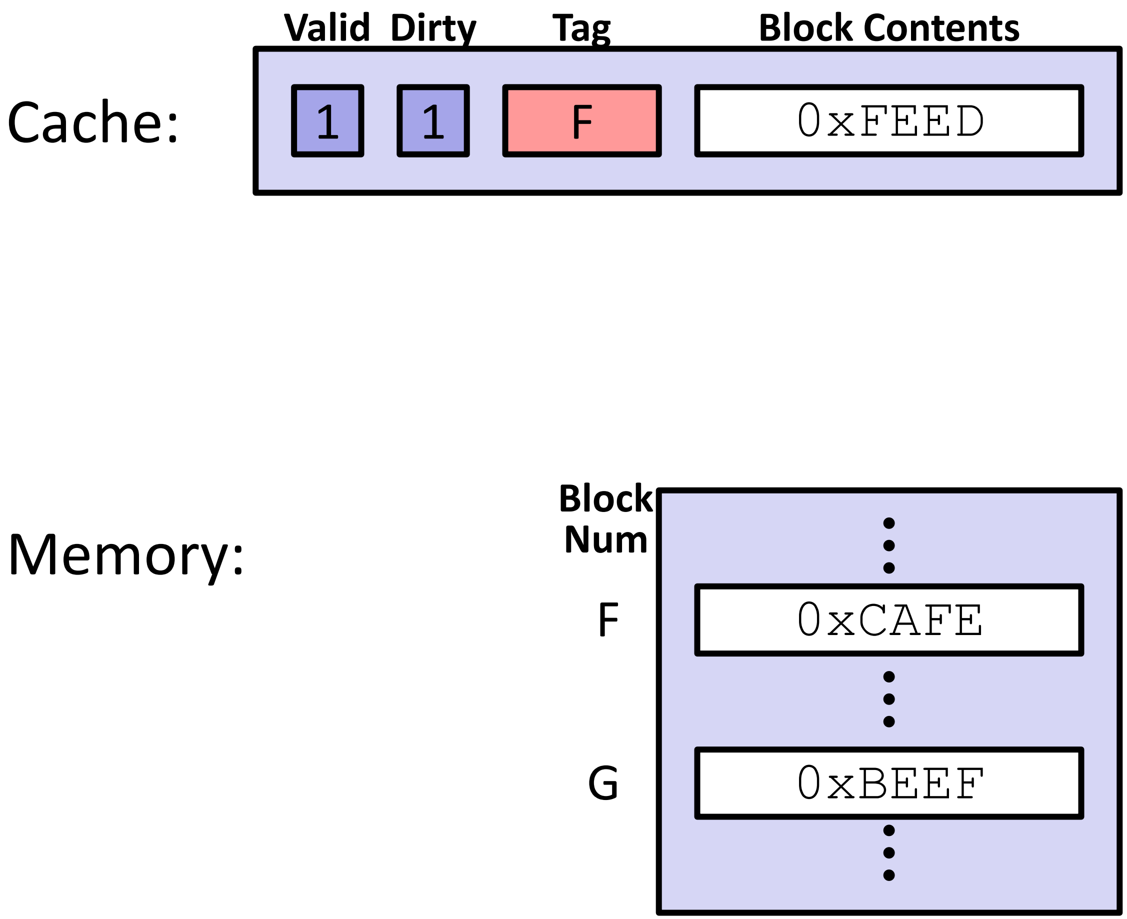


Step: Write `0xFEED` to cache only (and set the dirty bit)

Write-back, Write Allocate Example

1) `mov $0xFACE, (F)` Write Miss

2) `mov $0xFEED, (F)` Write Hit

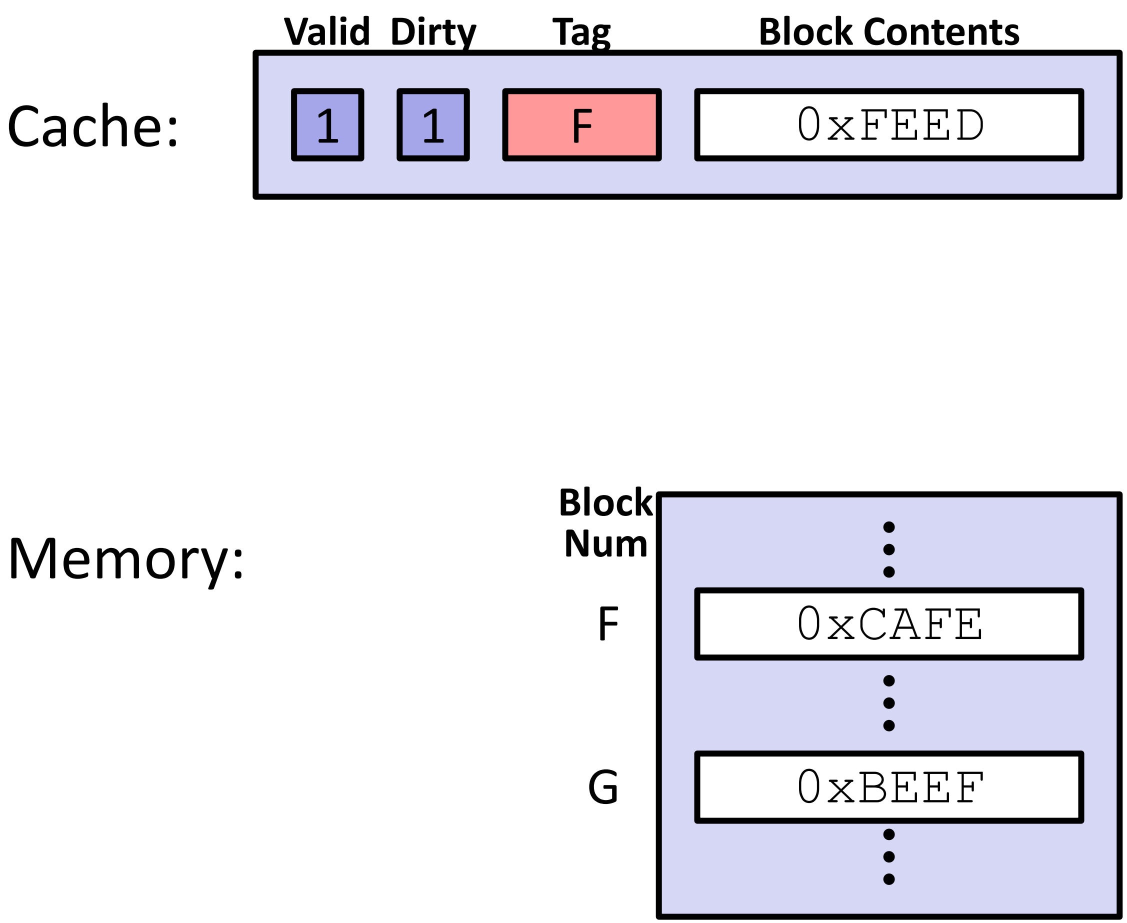


Write-back, Write Allocate Example

1) `mov $0xFACE, (F)`
Write Miss

2) `mov $0xFEED, (F)`
Write Hit

3) `mov (G), %ax`
Read Miss!



Step 1: Write **F** back to memory since it is dirty

Write-back, Write Allocate Example

1) `mov $0xFACE, (F)` Write Miss

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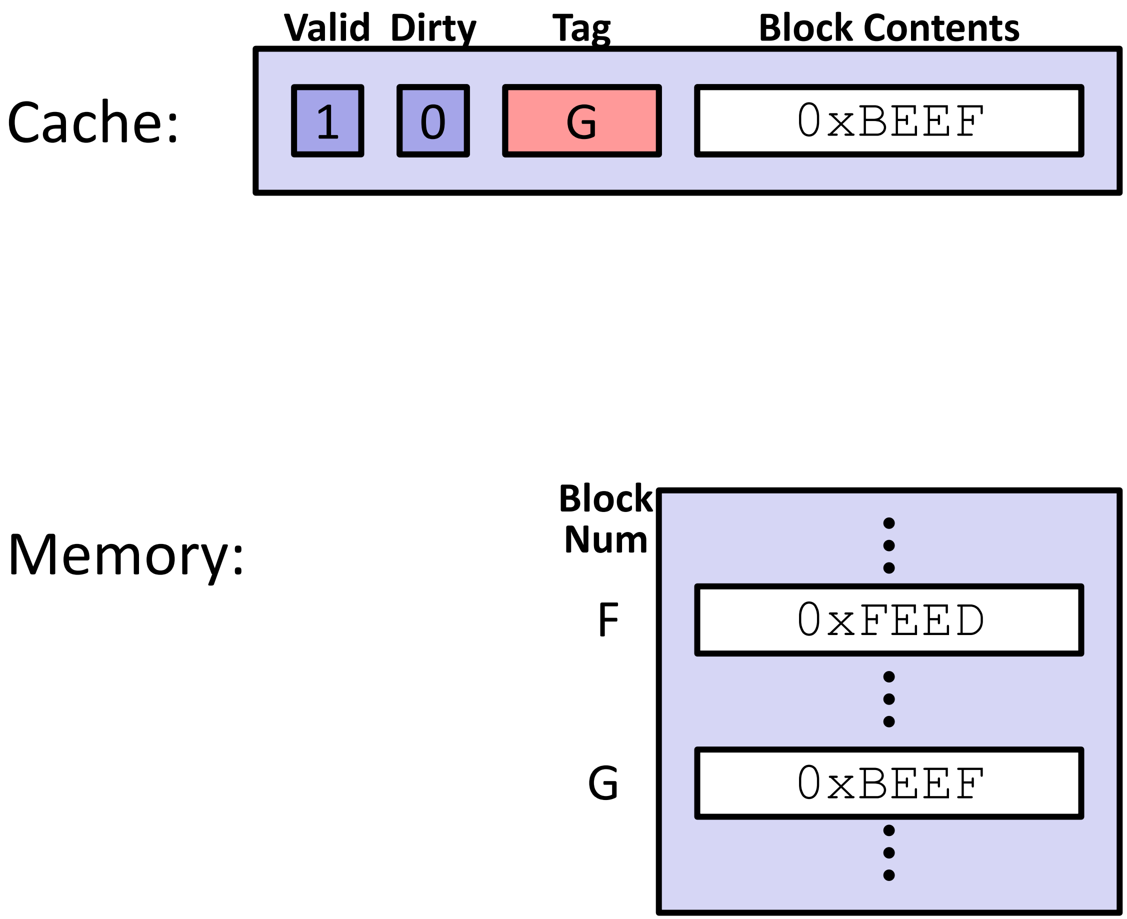
3) `mov (G), %ax` Read Miss



Step 1: Write **F** back to memory since it is dirty

Step 2: Bring **G** into the cache so that we can copy it into `%ax`

Cache Simulator

- Want to play around with cache parameters and policies? Check out our cache simulator!
  - https://courses.cs.washington.edu/courses/cse351/cachesim/
- Way to use:
  - Take advantage of "explain mode" and navigable history to test your own hypotheses and answer your own questions
  - Self-guided Cache Sim Demo posted along with Section 6
  - Will be used in hw17 – Lab 4 Preparation

Polling Question [Cache IV]

- Which of the following cache statements is FALSE?
  - Vote at http://pollev.com/pbjones

A. We can reduce compulsory misses by decreasing our block size

B. We can reduce conflict misses by increasing associativity

C. A write-back cache will save time for code with good temporal locality on writes

D. A write-through cache will always match data with the memory hierarchy level below it

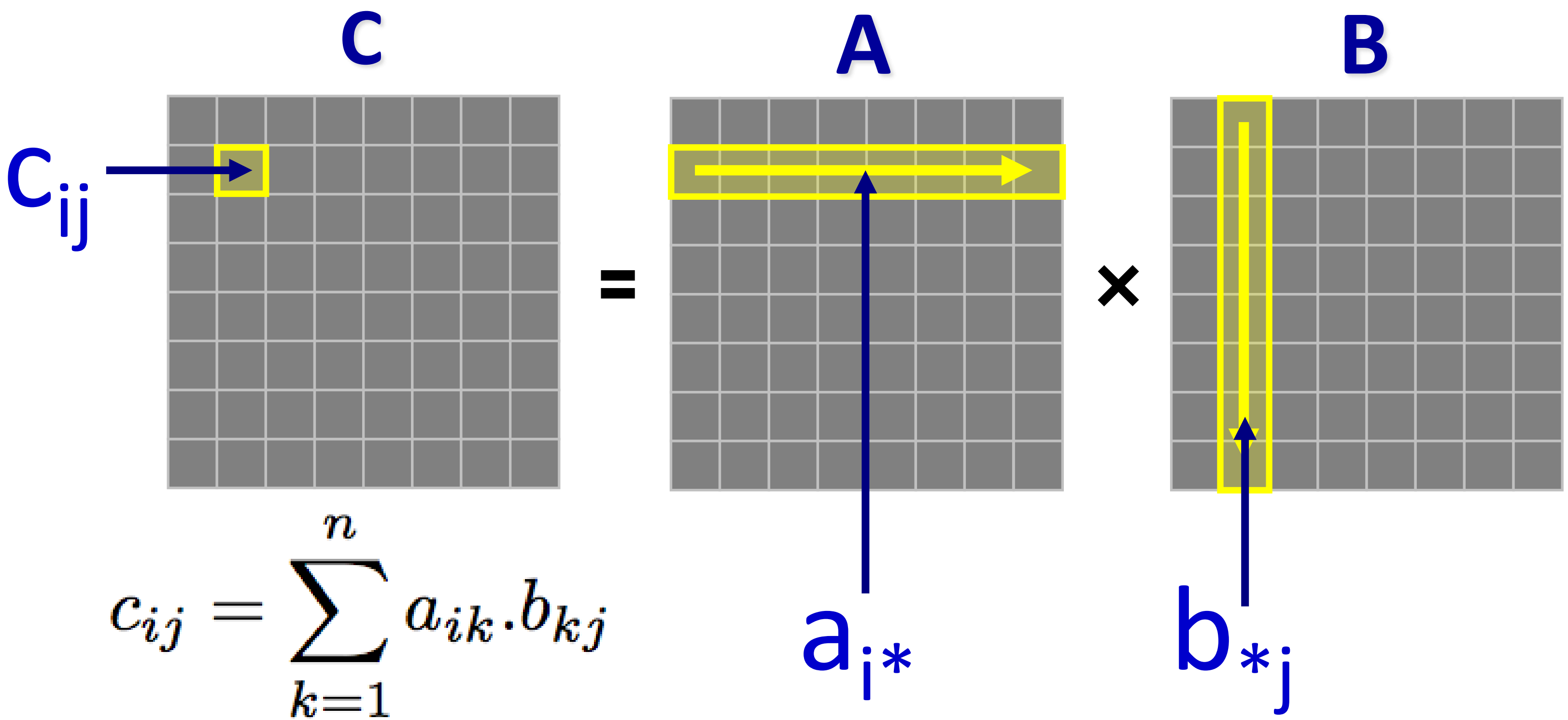
E. We're lost…

Optimizations for the Memory Hierarchy

- Write code that has locality!
  - Spatial: access data contiguously
  - Temporal: make sure access to the same data is not too far apart in time

- How can you achieve locality?
  - Adjust memory accesses in *code* (software) to improve miss rate (MR)
    - Requires knowledge of *both* how caches work as well as your system's parameters
  - Proper choice of algorithm
  - Loop transformations

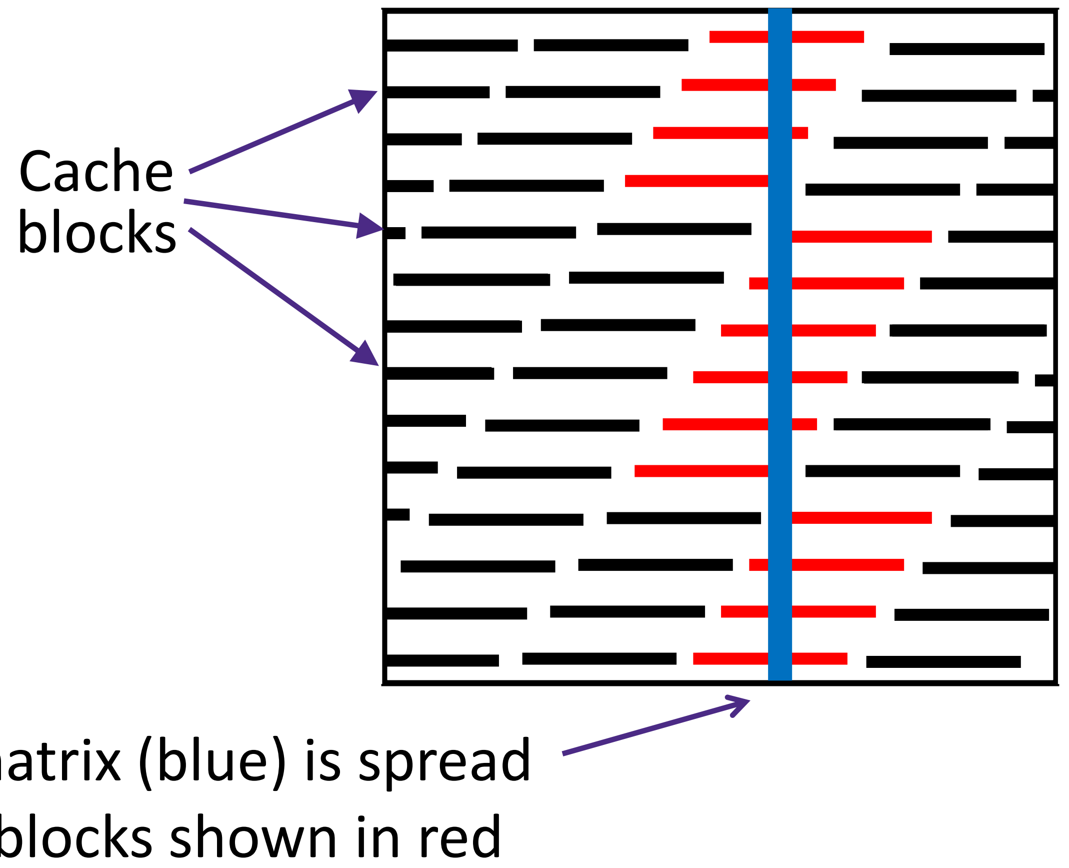
Example: Matrix Multiplication



$$c_{ij} = \sum_{k=1}^{n} a_{ik}.b_{kj}$$

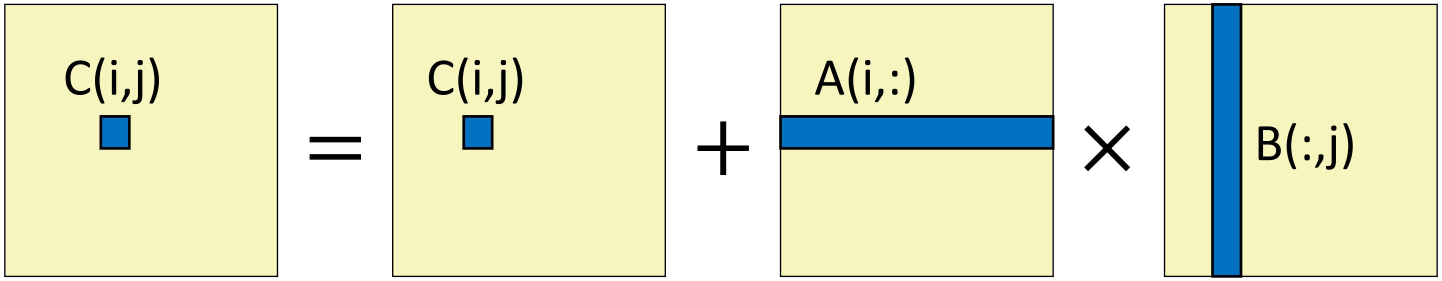
Matrices in Memory

- How do cache blocks fit into this scheme?
  - Row major matrix in memory:



Naïve Matrix Multiply

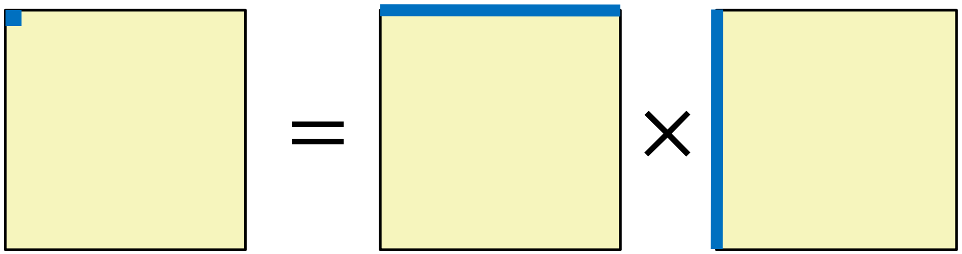
```
# move along rows of A
for (i = 0; i < n; i++)
  # move along columns of B
  for (j = 0; j < n; j++)
    # EACH k loop reads row of A, col of B
    # Also read & write c(i,j) n times
    for (k = 0; k < n; k++)
      c[i*n+j] += a[i*n+k] * b[k*n+j];
```



Cache Miss Analysis (Naïve)

Ignoring matrix `c`

- Scenario Parameters:
  - Square matrix ($n \times n$), elements are `double`s
  - Cache block size K = 64 B = 8 `double`s
  - Cache size $C \ll n$ (much smaller than n)
- Each iteration:
  - $\frac{n}{8} + n = \frac{9n}{8}$ misses



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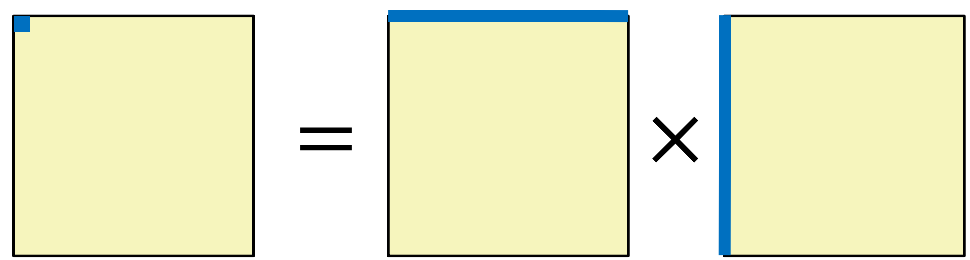
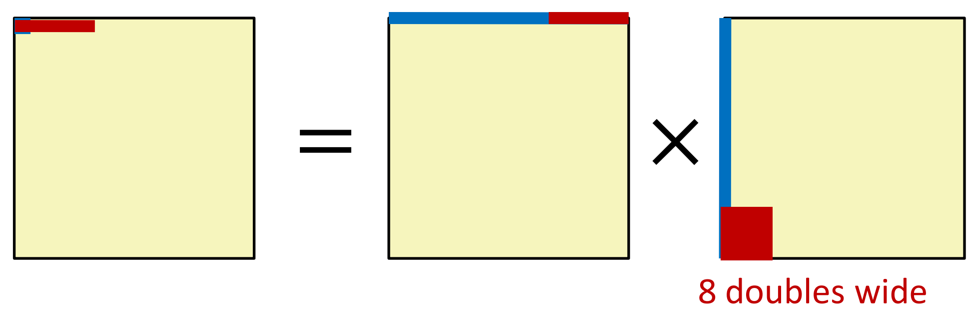


  - Afterwards in cache: (schematic)



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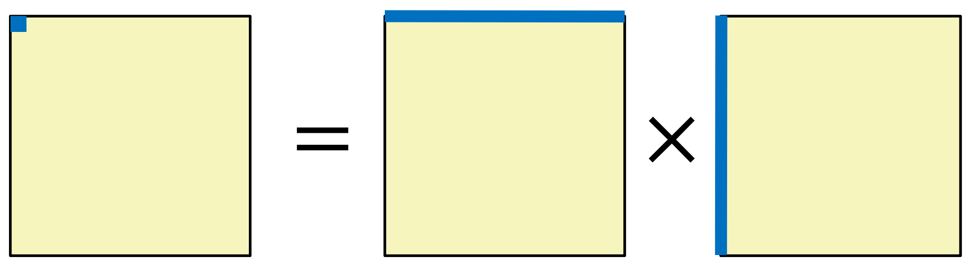


- Total misses: $\frac{9n}{8} \times n^2 = \frac{9}{8}n^3$

once per product matrix element

Linear Algebra to the Rescue (1)

This is extra (non-testable) material

- Can get the same result of a matrix multiplication by splitting the matrices into smaller submatrices (matrix "blocks")

- For example, multiply two 4×4 matrices:

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}, \text{ with } B \text{ defined similarly.}$$

$$AB = \begin{bmatrix} (A_{11}B_{11} + A_{12}B_{21}) & (A_{11}B_{12} + A_{12}B_{22}) \\ (A_{21}B_{11} + A_{22}B_{21}) & (A_{21}B_{12} + A_{22}B_{22}) \end{bmatrix}$$

This is extra (non-testable) material

Linear Algebra to the Rescue (2)

C_{11}	C_{12}	C_{13}	C_{14}
C_{21}	C_{22}	C_{23}	C_{24}
C_{31}	C_{32}	C_{43}	C_{34}
C_{41}	C_{42}	C_{43}	C_{44}

A_{11}	A_{12}	A_{13}	A_{14}
A_{21}	A_{22}	A_{23}	A_{24}
A_{31}	A_{32}	A_{33}	A_{34}
A_{41}	A_{42}	A_{43}	A_{144}

B_{11}	B_{12}	B_{13}	B_{14}
B_{21}	B_{22}	B_{23}	B_{24}
B_{32}	B_{32}	B_{33}	B_{34}
B_{41}	B_{42}	B_{43}	B_{44}

Matrices of size $n \times n$, split into 4 blocks of size r (n=4r)

$$C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} = \sum_k A_{2k}{*}B_{k2}$$

- Multiplication operates on small "block" matrices
  - Choose size so that they fit in the cache!
  - This technique called *"cache blocking"*

Blocked Matrix Multiply

- Blocked version of the naïve algorithm:

```
# move by rxr BLOCKS now
for (i = 0; i < n; i += r)
  for (j = 0; j < n; j += r)
    for (k = 0; k < n; k += r)
      # block matrix multiplication
      for (ib = i; ib < i+r; ib++)
        for (jb = j; jb < j+r; jb++)
          for (kb = k; kb < k+r; kb++)
            c[ib*n+jb] += a[ib*n+kb]*b[kb*n+jb];
```

loop over block matrices

loop within block matrices

  - r = block matrix size (assume r divides n evenly)

Cache Miss Analysis (Blocked)

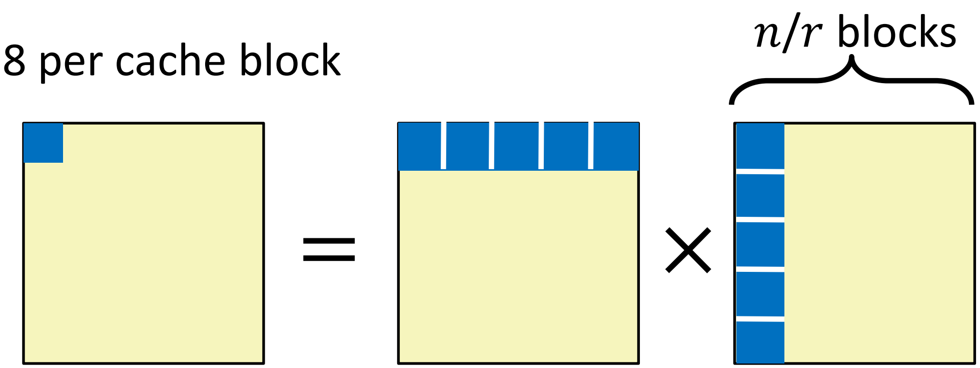
Ignoring matrix c

- Scenario Parameters:
  - Cache block size K = 64 B = 8 doubles
  - Cache size $C \ll n$ (much smaller than n)
  - Three blocks ■ $(r \times r)$ fit into cache: $3r^2 < C$

r^2 elements per block, 8 per cache block

- Each block iteration:
  - $r^2/8$ misses per block
  - $2n/r \times r^2/8 = nr/4$

n/r blocks in row and column



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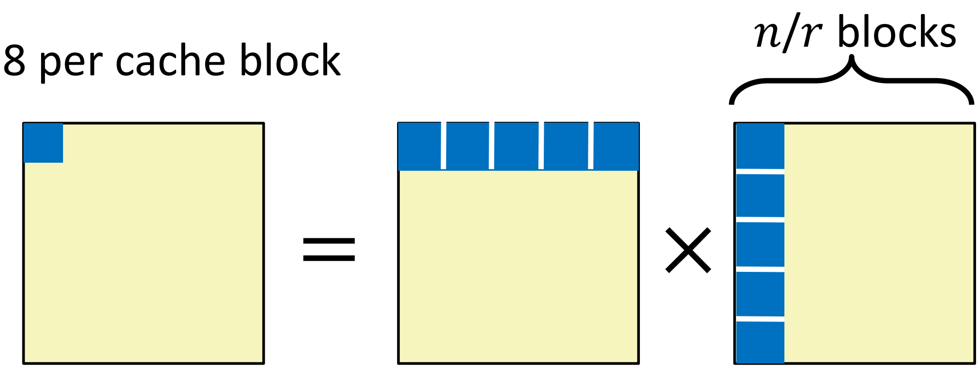
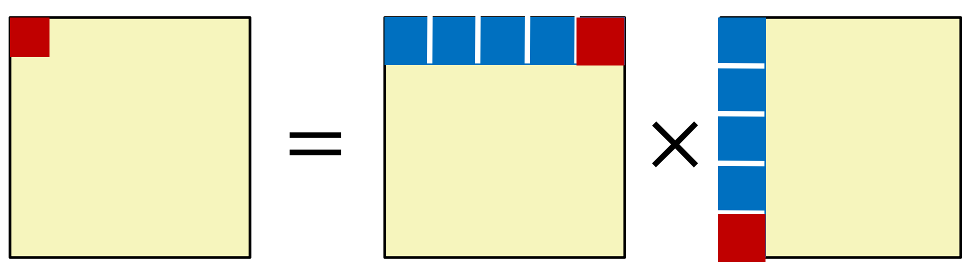


- Afterwards in cache (schematic)



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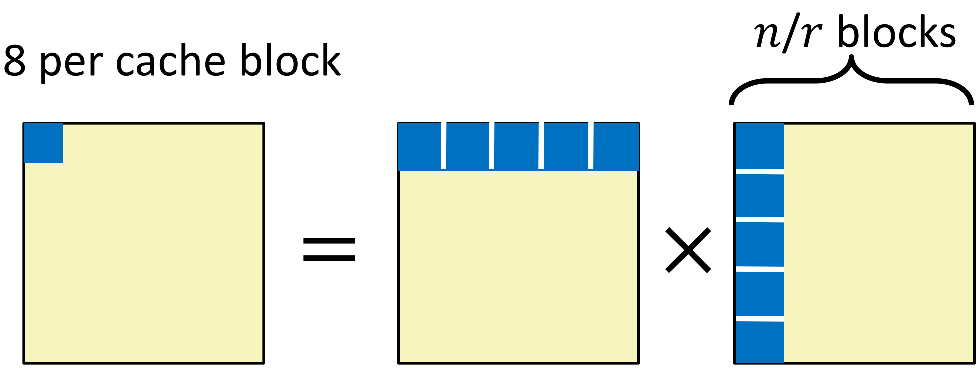


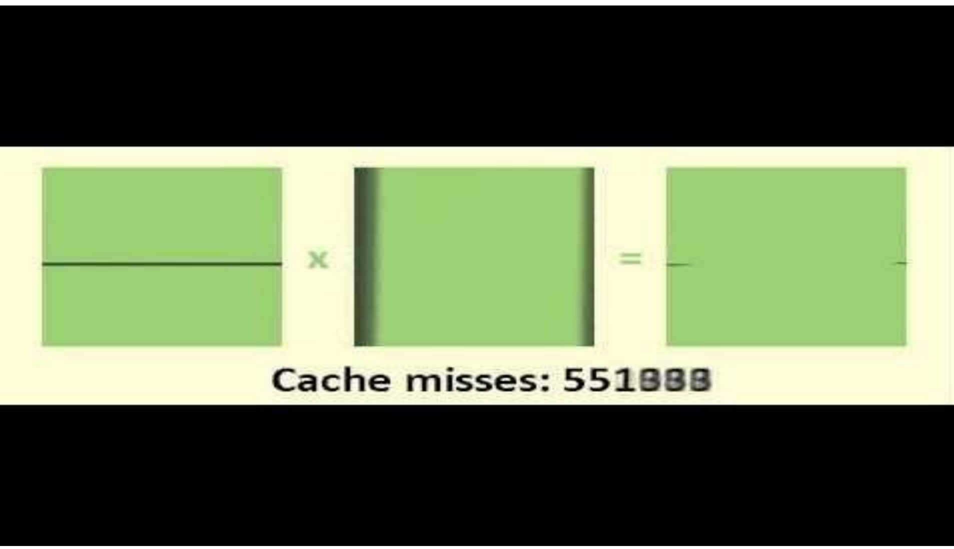
n/r blocks in row and column

- Total misses:
  - $nr/4 \times (n/r)2 = n^3/(4r)$

Matrix Multiply Visualization

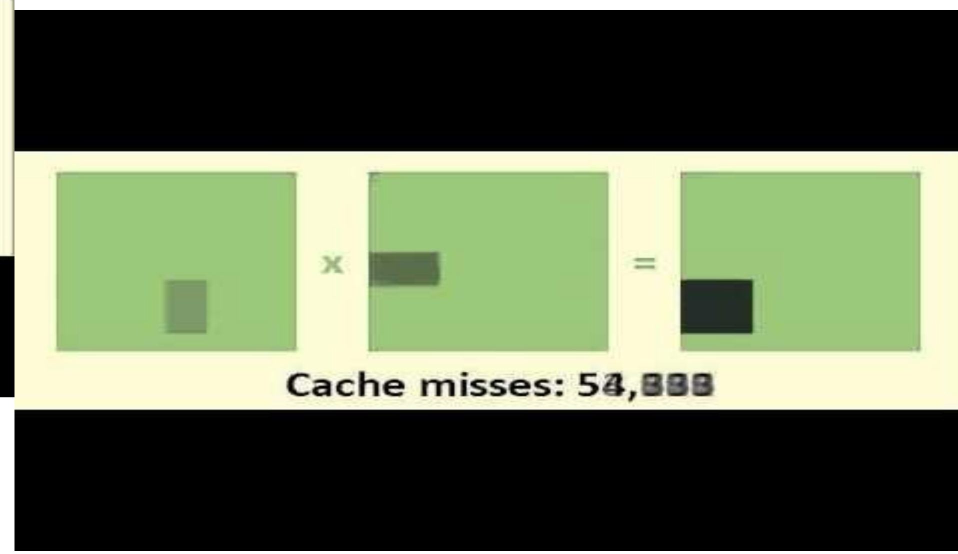
- Here n = 100, C = 32 KiB, r = 30

Naïve:



≈ 1,020,000
cache misses

Blocked:



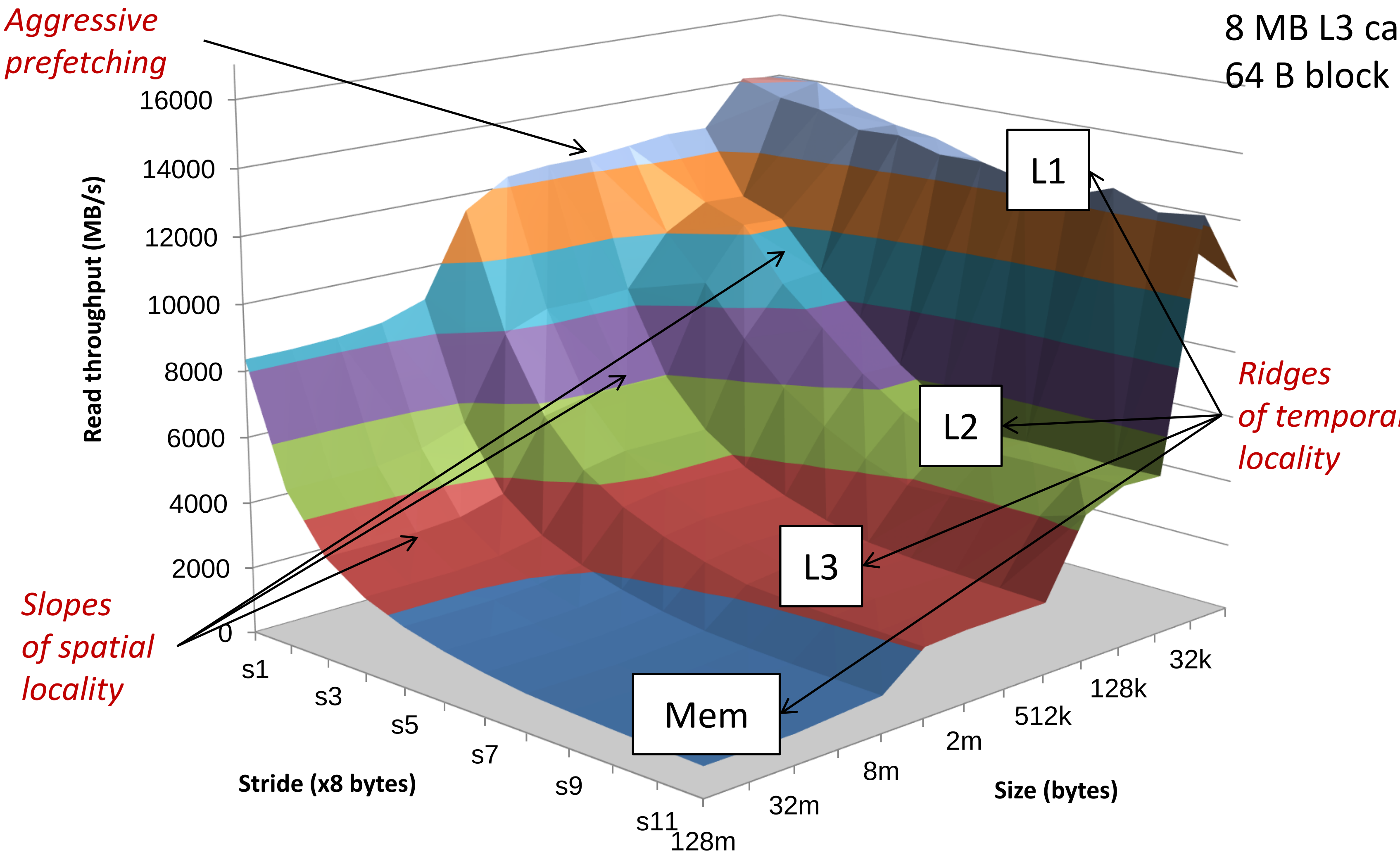
≈ 90,000
cache misses

Cache-Friendly Code

- Programmer can optimize for cache performance
  - How data structures are organized
  - How data are accessed
    - Nested loop structure
    - Blocking is a general technique
- All systems favor "cache-friendly code"
  - Getting absolute optimum performance is very platform specific
    - Cache size, cache block size, associativity, etc.
  - Can get most of the advantage with generic code
    - Keep working set reasonably small (temporal locality)
    - Use small strides (spatial locality)
    - Focus on inner loop code

The Memory Mountain

Core i7 Haswell
2.1 GHz
32 KB L1 d-cache
256 KB L2 cache
8 MB L3 cache
64 B block size



Learning About Your Machine

- **Linux:**
  - `lscpu`
  - ls /sys/devices/system/cpu/cpu0/cache/index0/
    - Example: cat /sys/devices/system/cpu/cpu0/cache/index*/size
- **Windows:**
  - `wmic memcache get <query>` (all values in KB)
  - Example: `wmic memcache get MaxCacheSize`

- Modern processor specs: http://www.7-cpu.com/