# CSE 351 Section 6 – Buffer Overflow and Caches

Hi there! Welcome back to section, we're happy that you're here  $\ensuremath{\textcircled{\sc o}}$ 

# **Buffer Overflow!**

Consider the following C program:

```
void main() {
   read_input();
}
int read_input() {
   char buf[8];
   gets(buf);
   return 0;
}
```

Here is a diagram of the stack in read\_input() right before the call to gets():

a) What is the value of the return address stored on the stack?

Assume that the user inputs the string "jklmnopqrs"

- b) Write the values in the stack before the "return 0;" statement is executed. Cross out the values that were overwritten and write in their new values. (Hint: use the ASCII table at the bottom to convert from letters to bytes)
- c) What is the new return address after the call to gets ()?
- d) Where will execution jump to after the "return 0;"?
- e) How many characters would we have to enter into the command line to overwrite the return address to 0x6A6B6C6D6E6F?
- f) Create a string that will overwrite the return address, setting it to 0x6A6B6C6D6E6F

In Lab 3, we are given a tool called sendstring, which converts hex digits into the actual bytes

```
>echo ``61 62 63" | ./sendstring
abc
```

g) If we want to overwrite the return address to a stack address like 0x7FFFFAB1234, we need to use a tool like sendstring to send the correct bytes. Why can't we just manually type the characters like we did earlier with "jklmnopgrs"?

> Check out the Lab 3 video on Phase 0 before you start the lab! It's linked on the Lab 3 page

Address	Value
	(hex)
%rsp+15	00
%rsp+14	00
%rsp+13	00
%rsp+12	00
%rsp+11	00
%rsp+10	40
%rsp+9	AF
%rsp+8	3B
%rsp+7	
%rsp+6	
%rsp+5	
%rsp+4	
%rsp+3	
%rsp+2	
%rsp+1	
%rsp+0	

	Char	Hex
	а	61
	b	62
	С	63
	d	64
	е	65
	f	66
	g h i j	67
	h	68
	i	69
	j	6A
	K	6B
	1	6C
	m	6D
	n	6E
s	0	6F
	р	70
	q	71
	r	72
	S	73
ē	t	74
-	u	75
?	V	76
:	W	77
	х	78

# Caches: Locality!

Recall that we have two types of locality that we can have in code:

**Temporal locality**: when recently referenced items are likely to be referenced again in the near future. **Spatial locality**: when nearby addresses tend to be referenced close together in time.

For each type of locality, can you give an example of when we might see it in code?

Temporal Locality:

Spatial Locality:

B3

\_

34

ВC

54

\_

AA \_

00

## Accessing a Cache (Hit or Miss?)

Assume the following caches all have block size K = 4 and are in the current state shown (you can ignore "-"). All values are shown in hex. Tag fields are padded, and bytes of the cache blocks are shown in full. The word size for the machine with these caches is 12 bits (i.e. addresses are 12 bits long)

#### Direct-Mapped:

Set	Valid	Tag	B0	B1	B2	B3	Set	Valid	Tag	B0	B1	B2
0	1	15	63	В4	C1	A4	8	0			-	
1	0	1	-	_		_	9	1	00	01	12	23
2	0	1				-	А	1	01	98	89	СВ
3	1	0 D	DE	AF	BA	DE	В	0	1E	4B	33	10
4	0	-	_	—	_	_	С	0	_	_	_	_
5	0	1	-	_	-	_	D	1	11	C 0	04	39
6	1	13	31	14	15	93	Е	0			-	
7	0			_		_	F	1	ΟF	FF	6F	30

Index bits:

Tag bits:

	Hit or Miss?	Data returned
a) Read 1 byte at 0x7AC		
b) Read 1 byte at 0x024		
c) Read 1 byte at 0x99F		

#### 2-way Set Associative:

Set	Valid	Tag	<b>B0</b>	B1	B2	B3	
0	0	—	—	-	—	_	
1	0	1		1		1	
2	1	03	4 F	D4	A1	3в	
3	0	_		1		_	
4	0	06	CA	FΕ	FO	0 D	
5	1 21		DE	AD	BE	ΕF	
6	0	_	_	_	_	-	
7	1	11	00	12	51	55	

Set	Valid	Tag	<b>B0</b>	B1	B2	B3	
0	0						
1	1	2F	01	20	40	03	
2	1	ΟE	99	09	87	56 -	
3	0	_	—	—	—		
4	0	_	_	—	—		
5	0	_	—	—	—		
6	1	37	22	Bб	DB	AA	
7	0	_	_	_	_	_	

Offset bits:

Index bits:

Tag bits: \_\_\_\_

	Hit or Miss?	Data returned
a) Read 1 byte at 0x435		
b) Read 1 byte at 0x388		
c) Read 1 byte at 0x0D3		

#### Fully Associative:

Set	Valid	Tag	<b>B0</b>	B1	B2	B3	Set	Valid	Tag	B0	B1	B2	B3	
0	1	1F4	00	01	02	03	0	0			—		—	Offset bits:
0	0	-		-	-	-	0	1	0AB	02	30	44	67	
0	1	100	F4	4 D	ΕE	11	0	1	034	FD	EC	BA	23	
0	1	077	12	23	34	45	0	0	_		_		-	Index bits:
0	0		-	—	_	_	0	1	1C6	00	11	22	33	
0	1	101	DA	14	ΕE	22	0	1	045	67	78	89	9A	
0	0	_		-			0	1	001	70	00	44	A6	Tag bits:
0	1	016	90	32	AC	24	0	0	_	_	—	_	—	
														-

	Hit or Miss?	Data returned
a) Read 1 byte at 0x1DD		
b) Read 1 byte at 0x719		
c) Read 1 byte at 0x2AA		

# Cache Sim

If you need help on using the cache sim, take a look at additional supplemental material that will guide you through using the cache sim (posted with today's section handouts)! The cache sim is very useful for lab 4 and corresponding homework assignments.