## **Course Wrap-Up**

CSE 351 Spring 2020

**Instructor:** Teaching Assistants:

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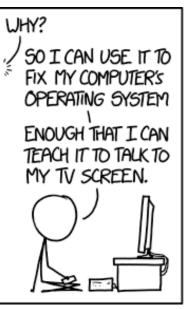
Connie Wang Diya Joy Edan Sneh

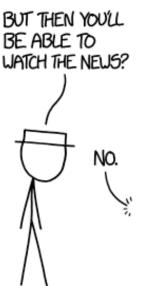
Eddy (Tianyi) Zhou Eric Fan Jeffery Tian

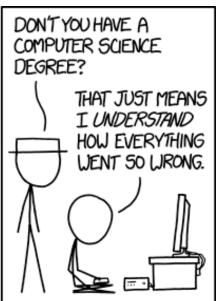
Jonathan Chen Joseph Schafer Melissa Birchfield

Millicent Li Porter Jones Rehaan Bhimani









https://xkcd.com/1760/

#### **Administrivia**

- Lab 5 (on Mem Alloc) NOW due Mon (6/08)
- Unit Summary #4 due Wed (6/10)
- ♦ hw23 on Java and C NOW due Thurs (6/11)
- Course evaluations now open
  - Please fill these out! Close Sunday (6/07)
  - Separate ones for Lecture and Section
- You must log on with your @uw google account to access!!
  - Google doc for 11:30 Lecture: <a href="https://tinyurl.com/351-06-05A">https://tinyurl.com/351-06-05A</a>
  - Google doc for 2:30 Lecture: <a href="https://tinyurl.com/351-06-05B">https://tinyurl.com/351-06-05B</a>

# **Today**

- End-to-end Review
  - What happens after you write your source code?
    - How code becomes a program
    - How your computer executes your code
- Victory lap and high-level concepts (key points)
  - More useful for "5 years from now" than "next week's final"

# C: The Low-Level High-Level Language

- C is a "hands-off" language that "exposes" more of hardware (especially memory)
  - Weakly-typed language that stresses data as bits
    - Anything can be represented with a number!
  - Unconstrained pointers can hold address of anything
    - And no bounds checking buffer overflow possible!
  - Efficient by leaving everything up to the programmer

## C Data Types

#### C Primitive types

- Fixed sizes and alignments
- Characters (char), Integers (short, int, long), Floating Point (float, double)

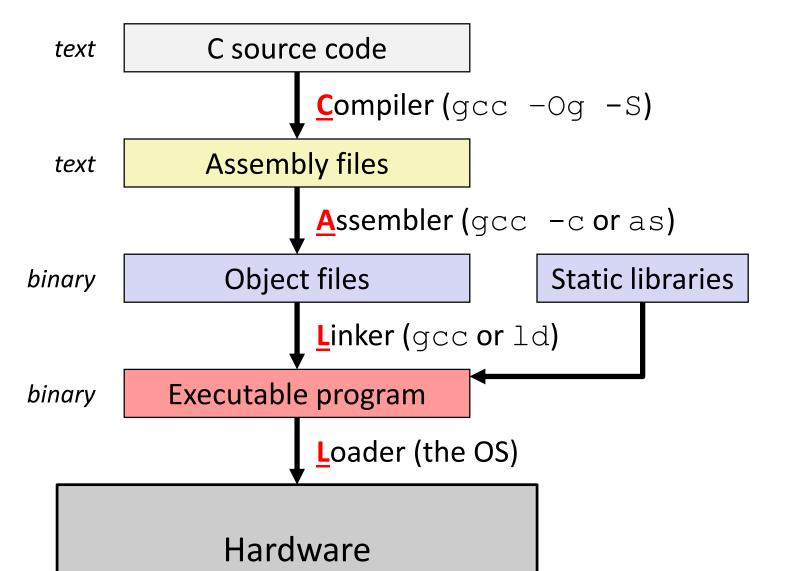
#### C Data Structures

- Arrays contiguous chunks of memory
  - Multidimensional arrays = still one continuous chunk, but row-major
  - Multi-level arrays = array of pointers to other arrays
- Structs structured group of variables
  - Struct fields are ordered according to declaration order
  - Internal fragmentation: space between members to satisfy member alignment requirements (aligned for each primitive element)
  - External fragmentation: space after last member to satisfy overall struct alignment requirement (largest primitive member)

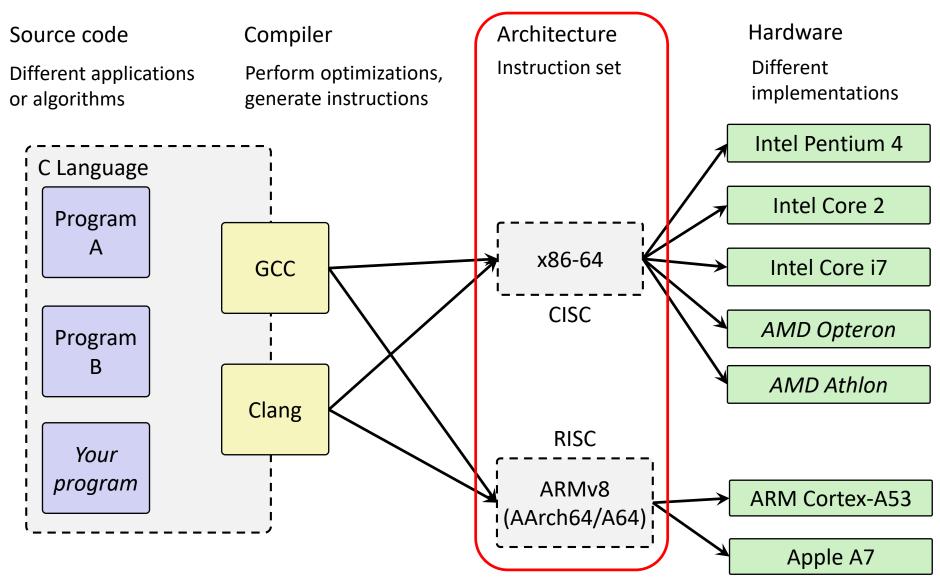
# C and Memory

- Using C allowed us to examine how we store and access data in memory
  - Endianness (only applies to memory)
    - Is the first byte (lowest address) the least significant (little endian) or most significant (big endian) of your data?
  - Array indices and struct fields result in calculating proper addresses to access
- Consequences of your code:
  - Affects performance (locality)
  - Affects security
- But to understand these effects better, we had to dive deeper...

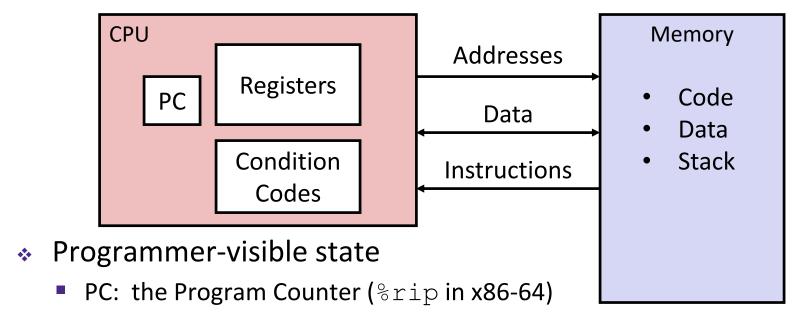
## **How Code Becomes a Program**



### **Instruction Set Architecture**



### **Assembly Programmer's View**

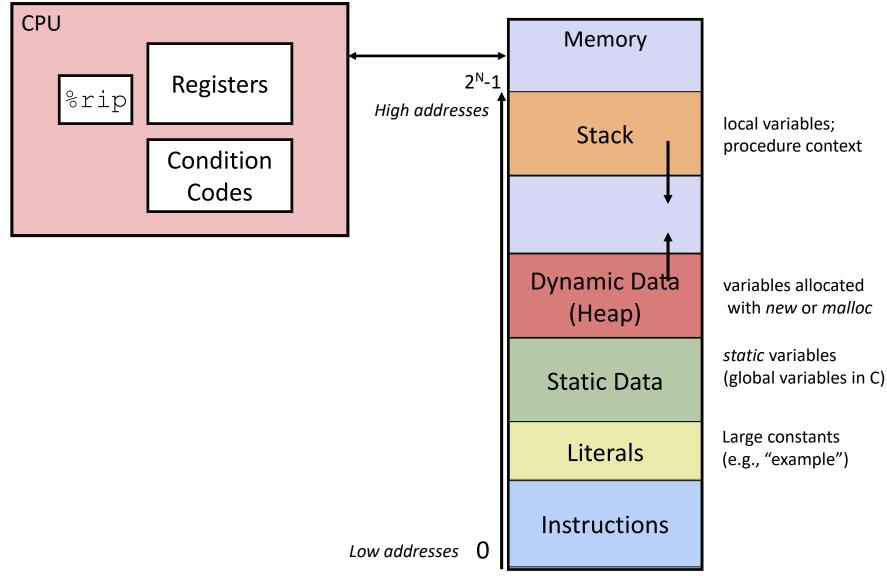


- Address of next instruction
- Named registers
  - Together in "register file"
  - Heavily used program data
- Condition codes
  - Store status information about most recent arithmetic operation
  - Used for conditional branching

#### Memory

- Byte-addressable array
- Huge virtual address space
- Private, all to yourself...

# **Program's View: Parts of Memory**





# **Program's View: Instructions**

#### Instructions

- Data movement
  - mov, movz, movz
  - push, pop
- Arithmetic
  - add, sub, imul
- Control flow
  - · cmp, test
  - jmp, je, jgt, ...
  - call, ret

#### Operand types

- Literal: \$8
- Register: %rdi, %al
- Memory: D(Rb,Ri,S) = D+Rb+Ri\*S
  - lea: not a memory access! Low addresses

Memory  $2^{N}-1$ High addresses Stack

local variables; procedure context

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Dynamic Data (Heap)

Static Data

Literals

**Instructions** 

variables allocated with *new* or *malloc* 

static variables (global variables in C)

Large constants (e.g., "example")

## **Program's View: Procedures & the Stack**

- **Procedures** 
  - **Essential abstraction**
  - Recursion...
- Stack discipline
  - Stack frame per call
  - Local variables
- Calling convention
  - How to pass arguments
    - Diane's Silk Dress Costs \$89
  - How to return data
  - Return address
  - Caller-saved / callee-saved registers

Memory  $2^{N}-1$ High addresses local variables; Stack procedure context Dynamic Data variables allocated with *new* or *malloc* (Heap) static variables (global variables in C) Static Data Large constants Literals (e.g., "example") **Instructions** 

High addresses

### **Program's View: The Heap**



- Variable size
- Variable lifetime

#### Allocator

- Balance throughput and memory utilization
- Data structures to keep track of free blocks

#### Garbage collection

- Must always free memory
- Garbage collectors help by finding anything reachable
- Failing to free results in memory leaks

Low addresses (

 $2^{N}-1$ local variables; Stack procedure context Dynamic Data variables allocated with *new* or *malloc* (Heap) static variables (global variables in C) Static Data Large constants Literals (e.g., "example") **Instructions** 

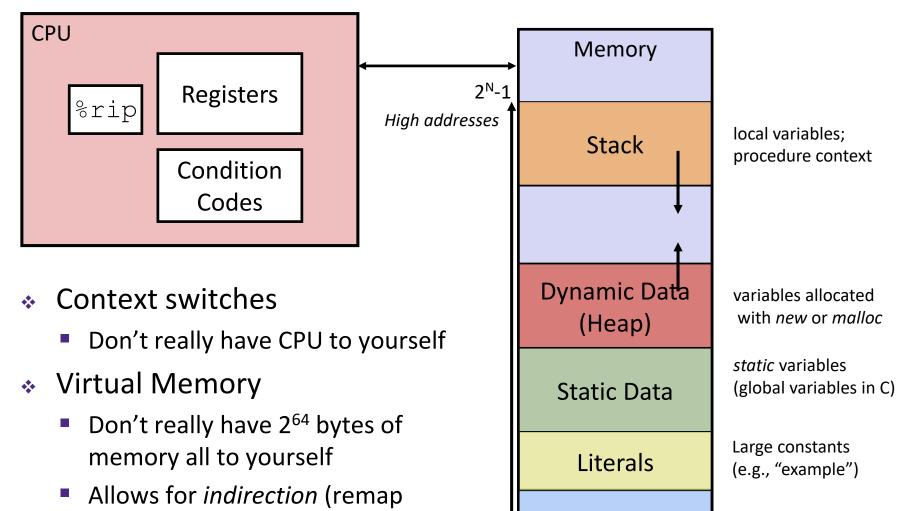
Memory

Instructions

physical pages, sharing...)

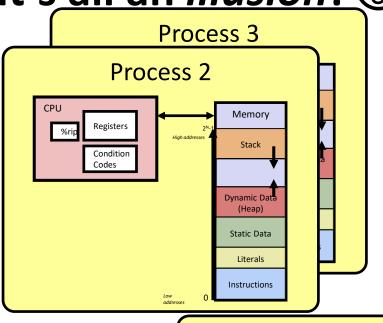


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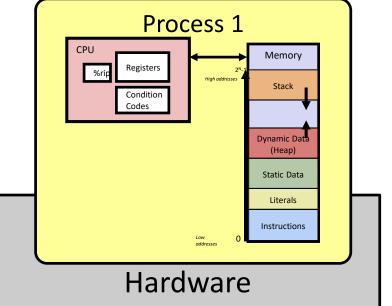


Low addresses

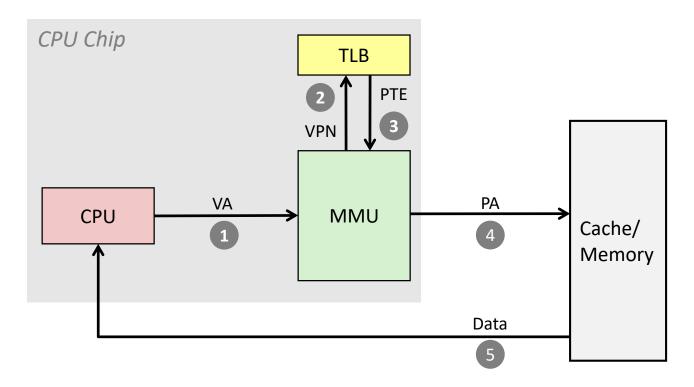
But remember... it's all an illusion! (3)



- \* fork
  - Creates copy of the process
- \* execv
  - Replace with new program
- wait
  - Wait for child to die (to reap it and prevent zombies)



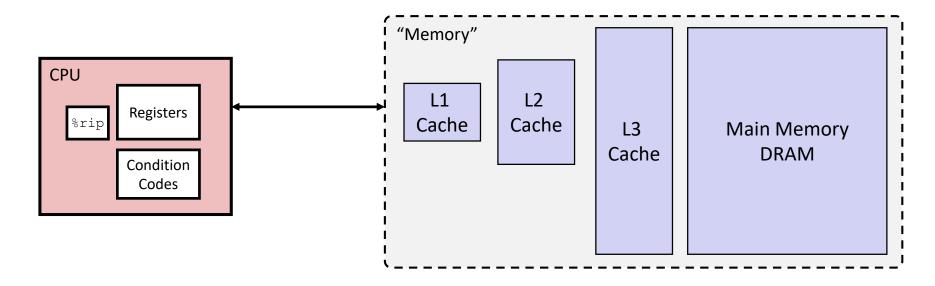
# Virtual Memory



#### Address Translation

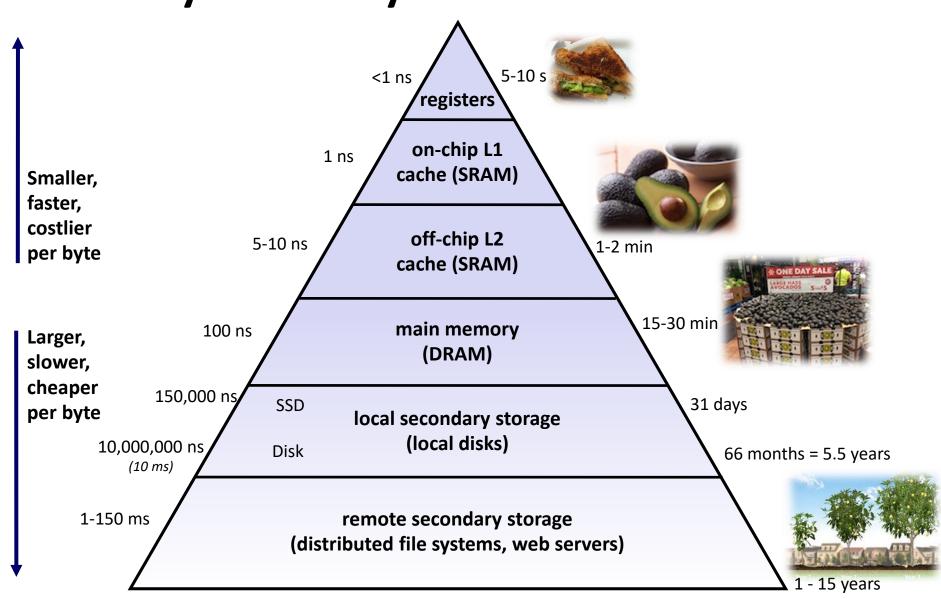
- Every memory access must first be converted from virtual to physical
- Indirection: just change the address mapping when switching processes
- Luckily, TLB (and page size) makes it pretty fast

# But Memory is Also a Lie! 🕲



- Illusion of one flat array of bytes
  - But caches invisibly make accesses to physical addresses faster!
- Caches
  - Associativity tradeoff with miss rate and access time
  - Block size tradeoff with spatial and temporal locality
  - Cache size tradeoff with miss rate and cost

# **Memory Hierarchy**



### **Review of Course Themes**

- Review course goals
  - They should make much more sense now!

## **Big Theme: Abstractions and Interfaces**

- Computing is about abstractions
  - (but we can't forget reality)
- What are the abstractions that we use?
- What do you need to know about them?
  - When do they break down and you have to peek under the hood?
  - What bugs can they cause and how do you find them?
- \* How does the hardware relate to the software?
  - Become a better programmer and begin to understand the important concepts that have evolved in building ever more complex computer systems

### **Little Theme 1: Representation**

- All digital systems represent everything as 0s and 1s
  - The 0 and 1 are really two different voltage ranges in the wires
  - Or magnetic positions on a disc, or hole depths on a DVD, or even DNA...
- "Everything" includes:
  - Numbers integers and floating point
  - Characters the building blocks of strings
  - Instructions the directives to the CPU that make up a program
  - Pointers addresses of data objects stored away in memory
- Encodings are stored throughout a computer system
  - In registers, caches, memories, disks, etc.
- They all need addresses (a way to locate)
  - Find a new place to put a new item
  - Reclaim the place in memory when data no longer needed

#### **Little Theme 2: Translation**

- There is a big gap between how we think about programs and data and the 0s and 1s of computers
  - Need languages to describe what we mean
  - These languages need to be translated one level at a time
- We know Java as a programming language
  - Have to work our way down to the 0s and 1s of computers
  - Try not to lose anything in translation!
  - We encountered C language, assembly language, and machine code (for the x86 family of CPU architectures)

#### **Little Theme 3: Control Flow**

- How do computers orchestrate everything they are doing?
- Within one program:
  - How do we implement if/else, loops, switches?
  - What do we have to keep track of when we call a procedure, and then another, and then another, and so on?
  - How do we know what to do upon "return"?
- Across programs and operating systems:
  - Multiple user programs
  - Operating system has to orchestrate them all
    - Each gets a share of computing cycles
    - They may need to share system resources (memory, I/O, disks)
  - Yielding and taking control of the processor
    - Voluntary or "by force"?

### **Course Perspective**

- CSE351 will make you a better programmer
  - Purpose is to show how software really works
    - Understanding of some of the abstractions that exist between programs and the hardware they run on, why they exist, and how they build upon each other
  - Understanding the underlying system makes you more effective
    - Better debugging
    - Better basis for evaluating performance
    - How multiple activities work in concert (e.g. OS and user programs)
  - "Stuff everybody learns and uses and forgets not knowing"
- CSE351 presents a world-view that will empower you
  - The intellectual and software tools to understand the trillions+ of 1s and 0s that are "flying around" when your program runs

### **Courses: What's Next?**

- Staying near the hardware/software interface:
  - CSE369/EE271: Digital Design basic hardware design using FPGAs
  - CSE474/EE474: Embedded Systems software design for microcontrollers
- Systems software (CSE majors/non-majors courses)
  - CSE341/CSE413: Programming Languages
  - CSE332/CSE373: Data Structures and Parallelism
  - CSE333/CSE374: Systems Programming building well-structured systems in C/C++
- Looking ahead
  - CSE401/CSE413: Compilers (pre-reqs: 332)
  - CSE451: Operating Systems (pre-reqs: 332, 333)
  - CSE461: Networks (pre-reqs: 332, 333)

# Thanks for a great quarter!

Huge thanks to your awesome TAs!































- Don't be a stranger!
  - I hope to see you in a course sometime in the future!