



# Administrivia

- ❖ hw3 due Wednesday – 11am
- ❖ hw4 due Friday – 11am
  
- ❖ Lab 1a released
  - Workflow:
    - 1) Edit `pointer.c`
    - 2) Run the Makefile (`make`) and check for compiler errors & warnings
    - 3) Run `ptest` (`./ptest`) and check for correct behavior
    - 4) Run rule/syntax checker (`python dlc.py`) and check output
  - Due Monday 4/13, will overlap a bit with Lab 1b
    - We grade just your *last* submission

# Lab Reflections

- ❖ All subsequent labs (after Lab 0) have a “reflection” portion
  - The Reflection questions can be found on the lab specs and are intended to be done *after* you finish the lab
  - You will type up your responses in a `.txt` file for submission on Gradescope
  - These will be graded “by hand” (read by TAs)
  
- ❖ Intended to check your understand of what you should have learned from the lab

# Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
  - Binary, hexadecimal, fixed-widths
- ❖ Organizing and addressing data in memory
  - Memory is a byte-addressable array
  - Machine “word” size = address size = register size
  - Endianness – ordering bytes in memory
- ❖ Manipulating data in memory using C
  - Assignment
  - Pointers, pointer arithmetic, and arrays
- ❖ **Boolean algebra and bit-level manipulations**

# Boolean Algebra

- ❖ Developed by George Boole in 19th Century
  - Algebraic representation of logic (True  $\rightarrow$  1, False  $\rightarrow$  0)
  - AND:  $A \& B = 1$  when both A is 1 and B is 1
  - OR:  $A | B = 1$  when either A is 1 or B is 1
  - XOR:  $A \wedge B = 1$  when either A is 1 or B is 1, but not both
  - NOT:  $\sim A = 1$  when A is 0 and vice-versa
  - DeMorgan's Law:
    - $\sim (A | B) = \sim A \& \sim B$
    - $\sim (A \& B) = \sim A | \sim B$

AND			OR			XOR			NOT	
&	0	1		0	1	^	0	1	~	
0	0	0	0	0	1	0	0	1	0	1
1	0	1	1	1	1	1	1	0	1	0

# General Boolean Algebras

- ❖ Operate on bit vectors
  - Operations applied bitwise
  - All of the properties of Boolean algebra apply

$$\begin{array}{r}
 01101001 \\
 \& \underline{01010101} \\
 \hline
 \end{array}
 \quad
 \begin{array}{r}
 01101001 \\
 | \underline{01010101} \\
 \hline
 \end{array}
 \quad
 \begin{array}{r}
 01101001 \\
 \wedge \underline{01010101} \\
 \hline
 \end{array}
 \quad
 \begin{array}{r}
 \sim 01010101 \\
 \hline
 \end{array}$$

- ❖ Examples of useful operations:

$$x \wedge x = 0$$

$$\begin{array}{r}
 01010101 \\
 \wedge \underline{01010101} \\
 \hline
 00000000
 \end{array}$$

$$x | 1 = 1, \quad x | 0 = x$$

$$\begin{array}{r}
 01010101 \\
 | \underline{\mathbf{1}1110000} \\
 \hline
 11110101
 \end{array}$$

# Bit-Level Operations in C

- ❖  $\&$  (AND),  $|$  (OR),  $\wedge$  (XOR),  $\sim$  (NOT)
  - View arguments as bit vectors, apply operations bitwise
  - Apply to any “integral” data type
    - long, int, short, char, unsigned
- ❖ Examples with char a, b, c;
  - ```
a = (char) 0x41; // 0x41->0b 0100 0001
b = ~a; // 0b ->0x
```
  - ```
a = (char) 0x69; // 0x69->0b 0110 1001
b = (char) 0x55; // 0x55->0b 0101 0101
c = a & b; // 0b ->0x
```
  - ```
a = (char) 0x41; // 0x41->0b 0100 0001
b = a; // 0b 0100 0001
c = a ^ b; // 0b ->0x
```

# Contrast: Logic Operations

- ❖ Logical operators in C: `&&` (AND), `||` (OR), `!` (NOT)
  - 0 is False, anything nonzero is True
  - Always return 0 or 1
  - **Early termination** (a.k.a. short-circuit evaluation) of `&&`, `||`
- ❖ Examples (`char` data type)
  - `!0x41 -> 0x00`
  - `!0x00 -> 0x01`
  - `!!0x41 -> 0x01`
  - `0xCC && 0x33 -> 0x01`
  - `0x00 || 0x33 -> 0x01`
  - `p && *p`
    - If `p` is the **null pointer** (`0x0`), then `p` is never dereferenced!

# Polling Question

- ❖ Given the bitwise vectors  $\mathbf{x} = 0\mathbf{xBA}$ ;  $\mathbf{y} = 0\mathbf{xE3}$ ;
- ❖ Compute the result of  $(\mathbf{x} \mid \mid \mathbf{y}) \wedge \sim(\mathbf{x} \& \mathbf{y})$ 
  - Vote at <http://pollev.com/pbjones>
- A. **0xA6**
- B. **0xA3**
- C. **0x5C**
- D. **0xFF**
- E. **We're lost...**

# Puppy break



# Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

- Memory & data
- Integers & floats**
- x86 assembly
- Procedures & stacks
- Executables
- Arrays & structs
- Memory & caches
- Processes
- Virtual memory
- Memory allocation
- Java vs. C

Assembly language:

```
get_mpg:
    pushq    %rbp
    movq    %rsp, %rbp
    ...
    popq    %rbp
    ret
```

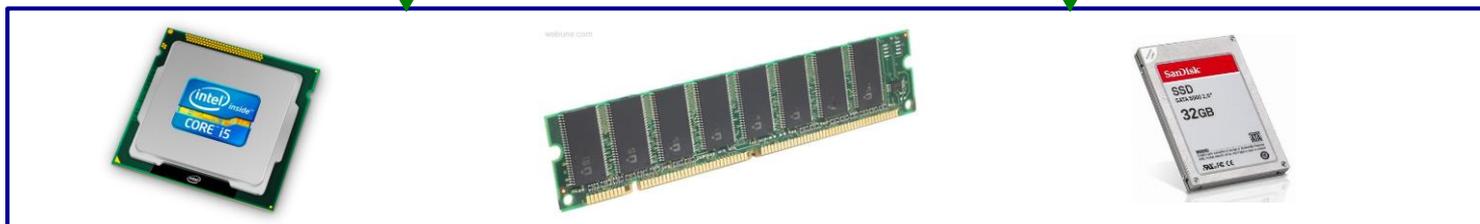
Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

OS:

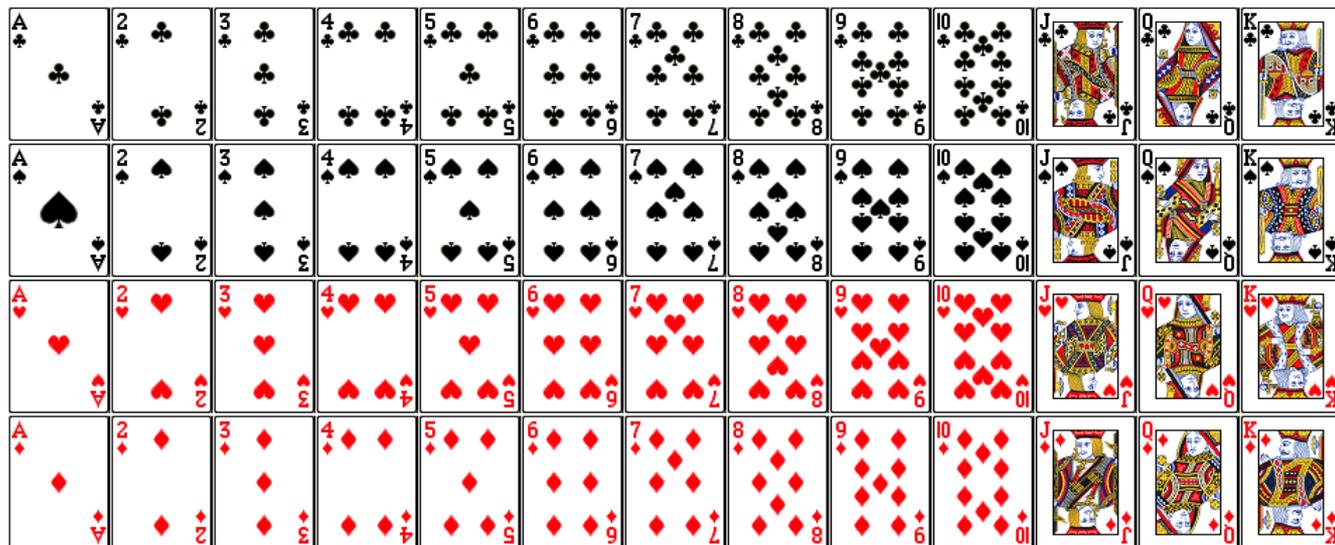


Computer system:



# But before we get to integers....

- ❖ Encode a standard deck of playing cards
- ❖ 52 cards in 4 suits
  - How do we encode suits, face cards?
- ❖ What operations do we want to make easy to implement?
  - Which is the higher value card?
  - Are they the same suit?



# Two possible representations

1) 1 bit per card (52): bit corresponding to card set to 1



low-order 52 bits of 64-bit word

- “One-hot” encoding (similar to set notation)
- Drawbacks:
  - Hard to compare values and suits
  - Large number of bits required

2) 1 bit per suit (4), 1 bit per number (13): 2 bits set

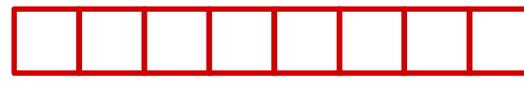


- Pair of one-hot encoded values
- Easier to compare suits and values, but still lots of bits used

# Two better representations

## 3) Binary encoding of all 52 cards – only 6 bits needed

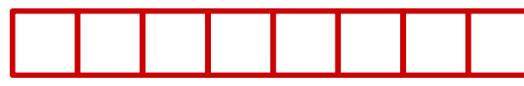
- $2^6 = 64 \geq 52$



low-order 6 bits of a byte

- Fits in one byte (smaller than one-hot encodings)
- How can we make value and suit comparisons easier?

## 4) Separate binary encodings of suit (2 bits) and value (4 bits)



suit value

- Also fits in one byte, and easy to do comparisons

|          |          |          |            |          |          |          |
|----------|----------|----------|------------|----------|----------|----------|
| <b>K</b> | <b>Q</b> | <b>J</b> | <b>...</b> | <b>3</b> | <b>2</b> | <b>A</b> |
| 1101     | 1100     | 1011     | ...        | 0011     | 0010     | 0001     |

|                                                                                       |    |
|---------------------------------------------------------------------------------------|----|
|  | 00 |
|  | 01 |
|  | 10 |
|  | 11 |

# Compare Card Suits

**mask:** a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector  $v$ .

Here we turn all *but* the bits of interest in  $v$  to 0.

```
char hand[5];           // represents a 5-card hand
char card1, card2;     // two cards to compare
card1 = hand[0];
card2 = hand[1];
...
if ( sameSuitP(card1, card2) ) { ... }
```

```
#define SUIT_MASK 0x30
```

```
int sameSuitP(char card1, char card2) {
    return (!(card1 & SUIT_MASK) ^ (card2 & SUIT_MASK));
    //return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

returns **int**

SUIT\_MASK = 0x30 =

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|---|---|

suit

value

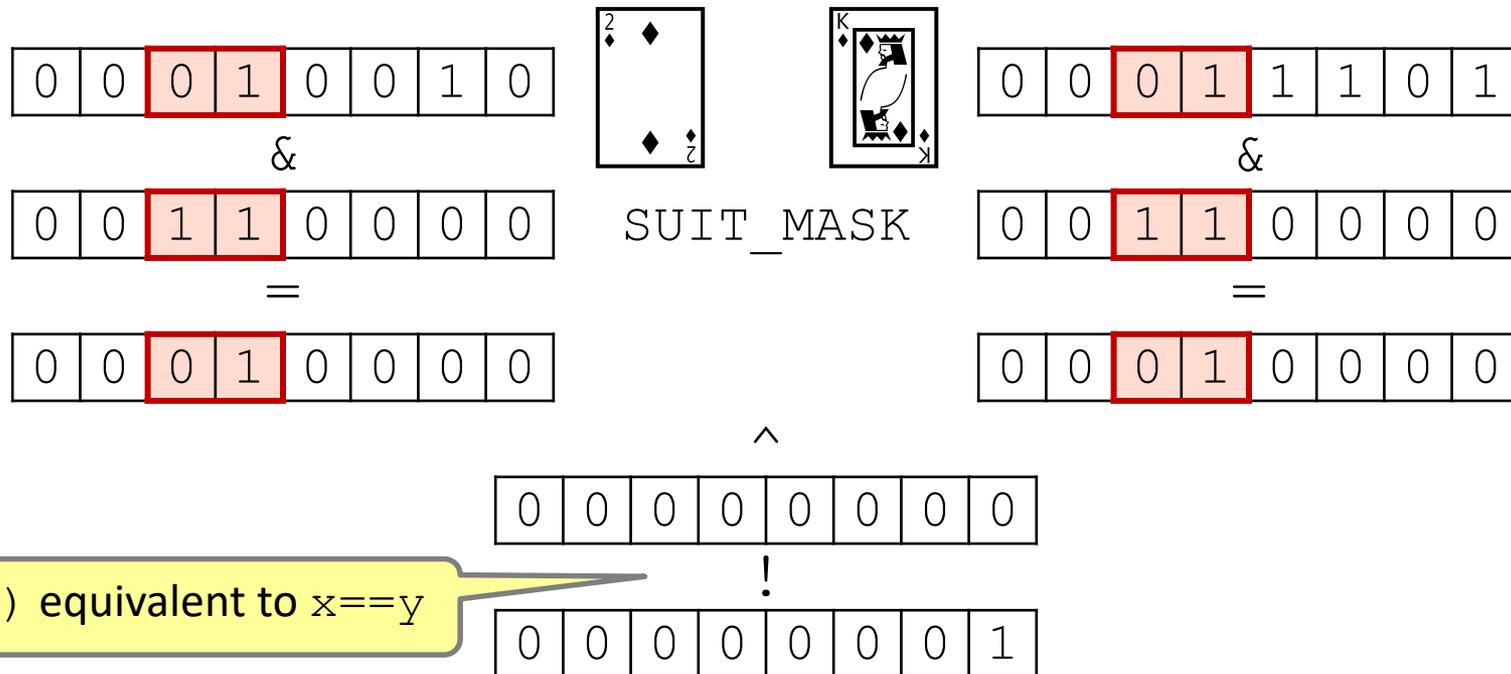
equivalent

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# Compare Card Suits

```
#define SUIT_MASK 0x30

int sameSuitP(char card1, char card2) {
    return (!(card1 & SUIT_MASK) ^ (card2 & SUIT_MASK));
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}
```



! (x ^ y) equivalent to x == y

# Compare Card Values

**mask:** a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector  $v$ .

```
char hand[5];           // represents a 5-card hand
char card1, card2;     // two cards to compare
card1 = hand[0];
card2 = hand[1];

...

if ( greaterValue(card1, card2) ) { ... }
```

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```

VALUE\_MASK = 0x0F = 

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
|---|---|---|---|---|---|---|---|

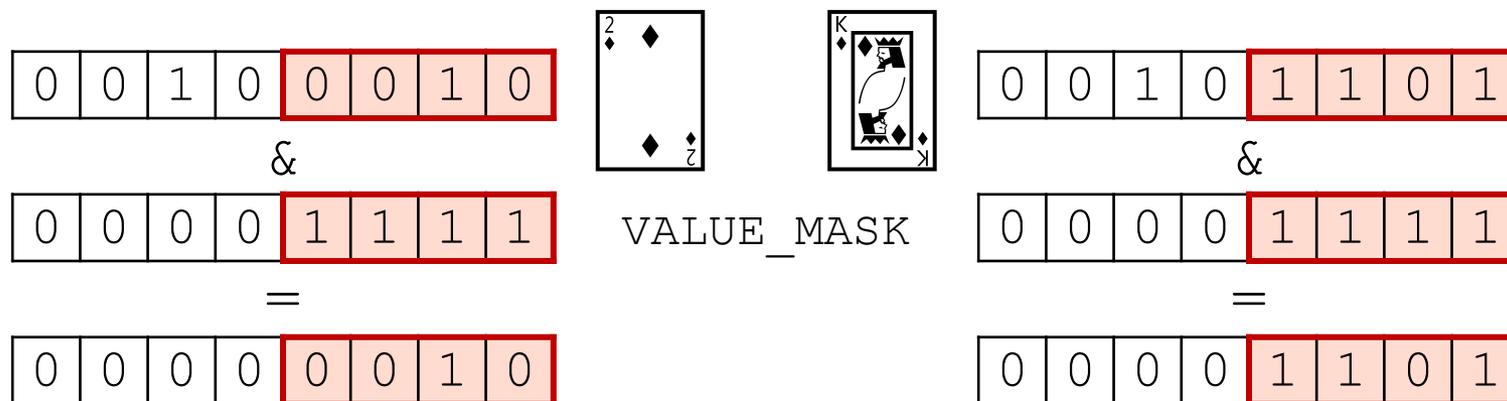
└── suit ─┘      └── value ─┘

# Compare Card Values

**mask:** a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector  $v$ .

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
           (unsigned int)(card2 & VALUE_MASK));
}
```



$$2_{10} > 13_{10}$$

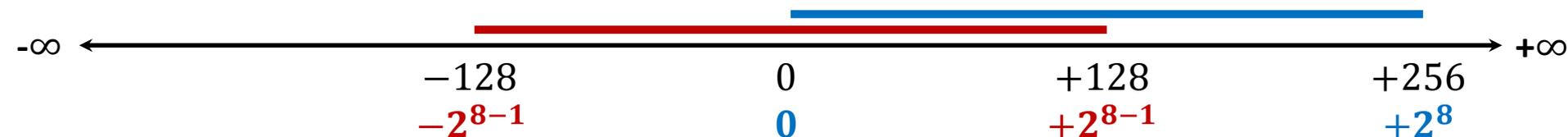
0 (false)

# Integers

- ❖ **Binary representation of integers**
  - **Unsigned and signed**
  - Casting in C
- ❖ Consequences of finite width representation
  - Overflow, sign extension
- ❖ Shifting and arithmetic operations

# Encoding Integers

- ❖ The hardware (and C) supports two flavors of integers
  - *unsigned* – only the non-negatives
  - *signed* – both negatives and non-negatives
- ❖ Cannot represent all integers with  $w$  bits
  - Only  $2^w$  distinct bit patterns
  - Unsigned values:  $0 \dots 2^w - 1$
  - Signed values:  $-2^{w-1} \dots 2^{w-1} - 1$
- ❖ **Example:** 8-bit integers (*e.g.* `char`)



# Unsigned Integers

- ❖ Unsigned values follow the standard base 2 system
  - $b_7b_6b_5b_4b_3b_2b_1b_0 = b_72^7 + b_62^6 + \dots + b_12^1 + b_02^0$
- ❖ Add and subtract using the normal “carry” and “borrow” rules, just in binary

|     |
|-----|
| 63  |
| + 8 |
| 71  |

|           |
|-----------|
| 00111111  |
| +00001000 |
| 01000111  |

- ❖ Useful formula:  $2^{N-1} + 2^{N-2} + \dots + 2 + 1 = 2^N - 1$ 
  - *i.e.* N ones in a row =  $2^N - 1$
- ❖ How would you make *signed* integers?

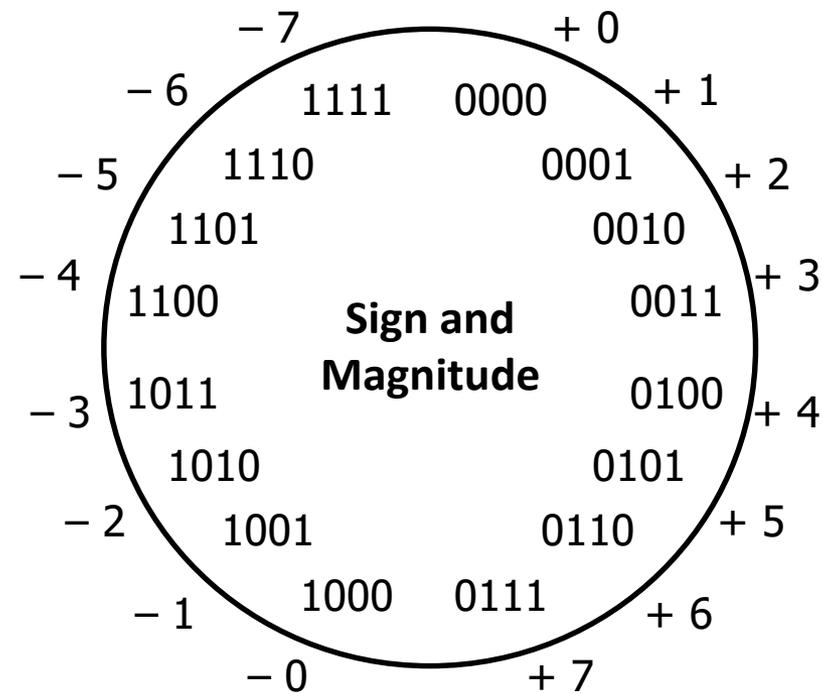
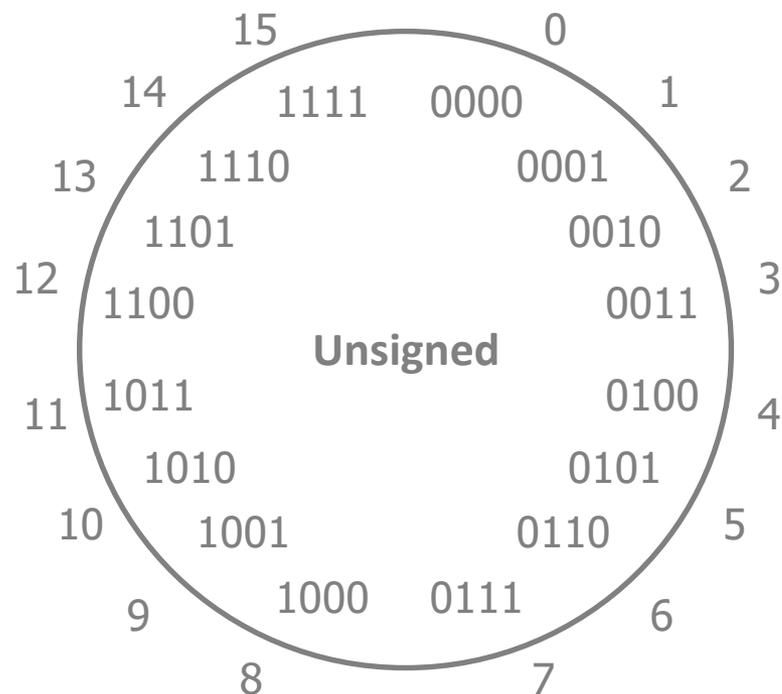
# Sign and Magnitude

Most Significant Bit

- ❖ Designate the high-order bit (MSB) as the “sign bit”
  - $\text{sign}=0$ : positive numbers;  $\text{sign}=1$ : negative numbers
- ❖ Benefits:
  - Using MSB as sign bit matches positive numbers with unsigned
  - All zeros encoding is still = 0
- ❖ Examples (8 bits):
  - $0x00 = 00000000_2$  is non-negative, because the sign bit is 0
  - $0x7F = 01111111_2$  is non-negative ( $+127_{10}$ )
  - $0x85 = 10000101_2$  is negative ( $-5_{10}$ )
  - $0x80 = 10000000_2$  is negative... zero???

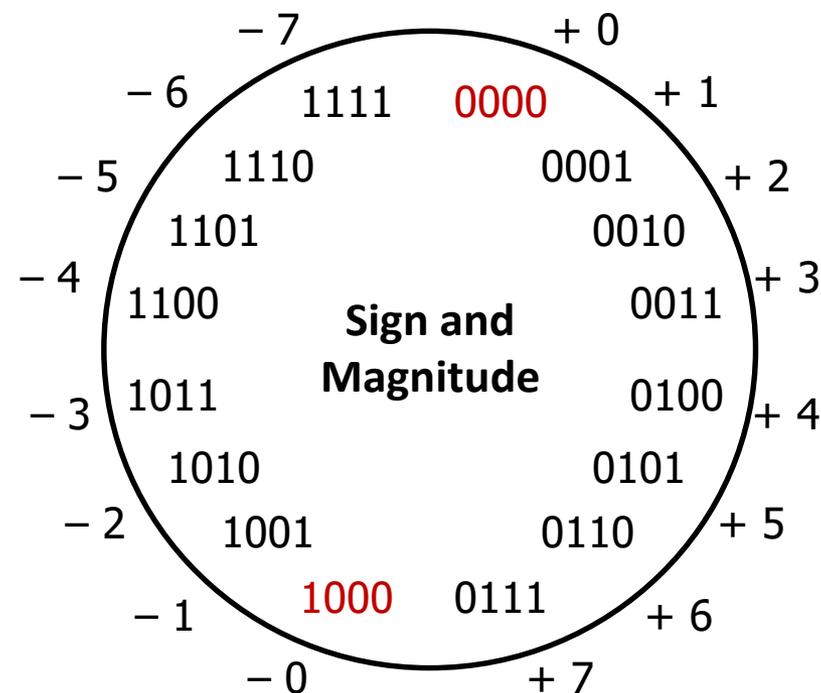
# Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks?



# Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks:
  - **Two representations of 0** (bad for checking equality)



# Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks:
  - Two representations of 0 (bad for checking equality)
  - **Arithmetic is cumbersome**
    - Example:  $4 - 3 \neq 4 + (-3)$

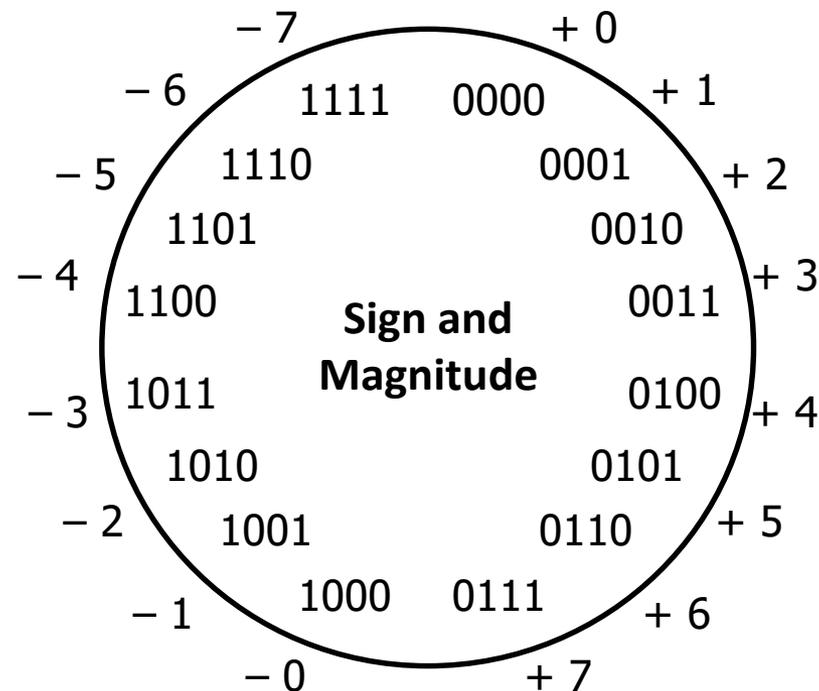
|     |        |
|-----|--------|
| 4   | 0100   |
| - 3 | - 0011 |
| 1   | 0001   |



|      |        |
|------|--------|
| 4    | 0100   |
| + -3 | + 1011 |
| -7   | 1111   |



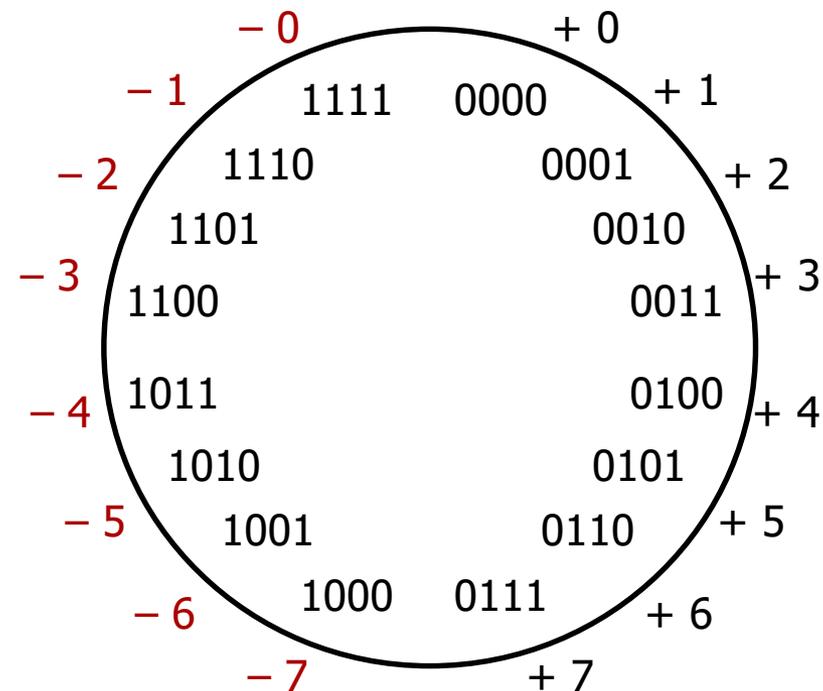
- Negatives “increment” in wrong direction!



# Two's Complement

❖ Let's fix these problems:

1) "Flip" negative encodings so incrementing works



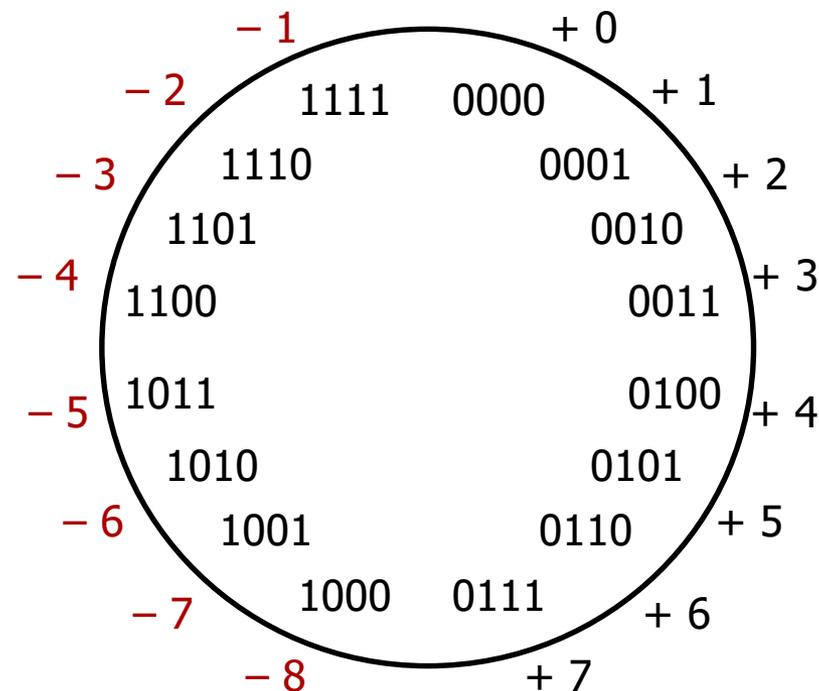
# Two's Complement

❖ Let's fix these problems:

- 1) "Flip" negative encodings so incrementing works
- 2) "Shift" negative numbers to eliminate  $-0$

❖ MSB *still* indicates sign!

- This is why we represent one more negative than positive number ( $-2^{N-1}$  to  $2^{N-1} - 1$ )



# Two's Complement Negatives

- ❖ Accomplished with one neat mathematical trick!

$b_{w-1}$  has weight  $-2^{w-1}$ , other bits have usual weights  $+2^i$



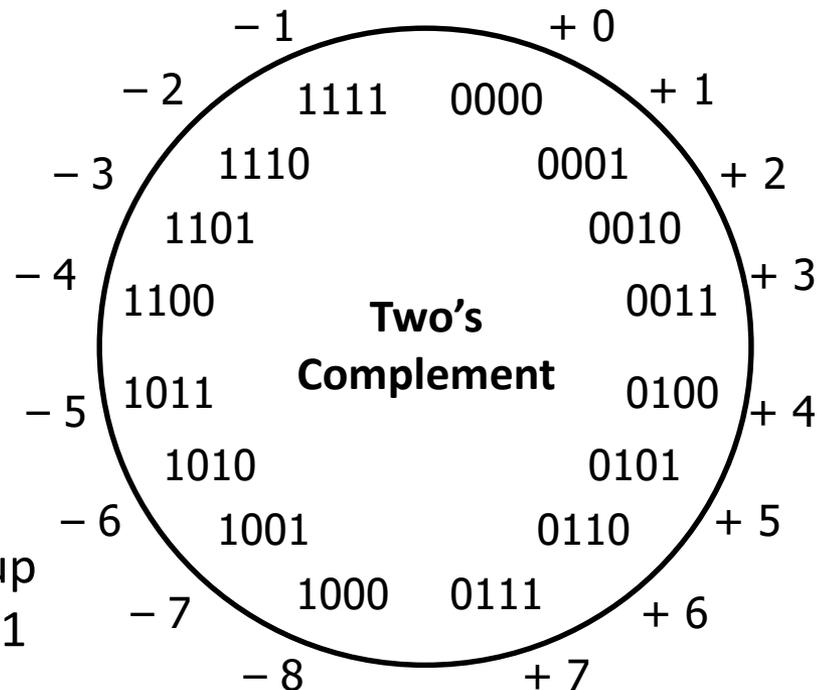
- 4-bit Examples:

- $1010_2$  unsigned:  
 $1*2^3+0*2^2+1*2^1+0*2^0 = 10$
- $1010_2$  two's complement:  
 $-1*2^3+0*2^2+1*2^1+0*2^0 = -6$

- -1 represented as:

$$1111_2 = -2^3 + (2^3 - 1)$$

- MSB makes it super negative, add up all the other bits to get back up to -1



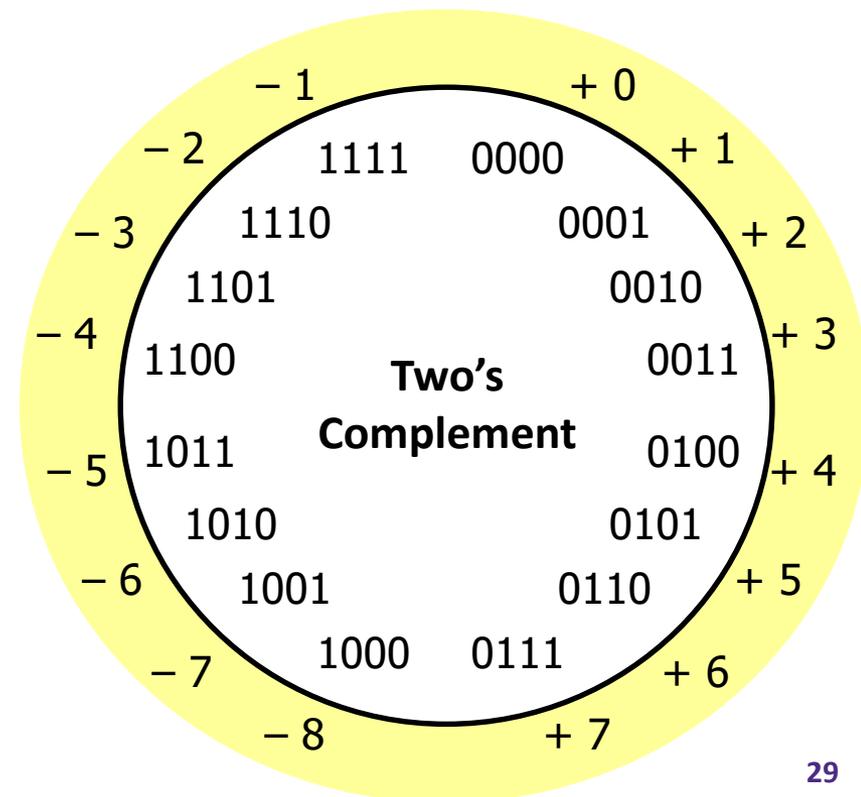
# Why Two's Complement is So Great

- ❖ Roughly same number of (+) and (-) numbers
- ❖ Positive number encodings match unsigned
- ❖ Single zero
- ❖ All zeros encoding = 0

- ❖ Simple negation procedure:

- Get negative representation of any integer by taking bitwise complement and then adding one!

$$(\sim x + 1 == -x)$$



# Polling Question [Int I - b]

- ❖ Take the 4-bit number encoding  $x = 0b1011$
- ❖ Which of the following numbers is NOT a valid interpretation of  $x$  using any of the number representation schemes discussed today?
  - Unsigned, Sign and Magnitude, Two's Complement
  - Vote at <http://pollev.com/rea>
- A. -4
- B. -5
- C. 11
- D. -3
- E. We're lost...

# Summary

- ❖ Bit-level operators allow for fine-grained manipulations of data
  - Bitwise AND ( $\&$ ), OR ( $|$ ), and NOT ( $\sim$ ) different than logical AND ( $\&\&$ ), OR ( $| |$ ), and NOT ( $!$ )
  - Especially useful with bit masks
- ❖ Choice of *encoding scheme* is important
  - Tradeoffs based on size requirements and desired operations
- ❖ Integers represented using unsigned and two's complement representations
  - Limited by fixed bit width
  - We'll examine arithmetic operations next lecture