Memory Allocation I

CSE 351 Winter 2018

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WHEN WILL WE FORGET?
BASED ON US CENSUS BUREAU *NATIONAL POPULATION PROJECTIONS*
ASSUMING WE DON'T REMEMBER CULTURAL EVENTS FROM BEFORE AGE 5 OR 6

BY THIS YEAR:	THE MAJORITY OF AMERICANS WILL BE TOO YOUNG TO REMEMBER:
2016	*RETURN OF THE JEDI* RELEASE
2017	THE FIRST APPLE MACINTOSH
2018	NEW COKE
2019	*CHALLENGER*
2020	CHERNOBYL
2021	BLACK MONDAY
2022	THE REAGAN PRESIDENCY
2023	THE BERLIN WALL
2024	HAMMERTIME
2025	THE SOVIET UNION
2026	THE LA RIOTS
2027	LORENA BOBBITT
2028	THE *FORREST GUMP* RELEASE
2029	THE RWANDAN GENOCIDE
2030	OJ SIMPSON'S TRIAL
2038	A TIME BEFORE FACEBOOK
2039	VH1's *I LOVE THE 90s*
2040	HURRICANE KATRINA
2041	THE PLANET PLUTO
2042	THE FIRST iPHONE
2047	ANYTHING EMBARRASSING YOU DO TODAY

Adapted from
https://xkcd.com/1093/

Administrative

- Homework 5 due Wednesday
- Lab 5 (on Mem Alloc), due 3/10

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

Memory & data
Integers & floats
x86 assembly
Procedures & stacks
Executables
Arrays & structs
Memory & caches
Processes
Virtual memory
Memory allocation
Java vs. C

Assembly language:

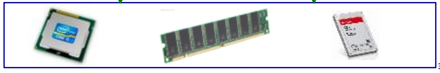
```
get_mpg:
    pushq   %rbp
    movq    %rsp, %rbp
    ...
    popq    %rbp
    ret
```

Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

OS:

Computer system:



Multiple Ways to Store Program Data

- Static global data
  - *Fixed size* at compile-time
  - Entire *lifetime of the program* (loaded from executable)
  - Portion is read-only (e.g. string literals)
- Stack-allocated data
  - Local/temporary variables
    - *Can* be dynamically sized (in some versions of C)
  - *Known lifetime* (deallocated on `return`)
- Dynamic (heap) data
  - Size known only at runtime (i.e. based on user-input)
  - Lifetime known only at runtime (long-lived data structures)

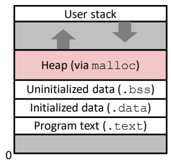
```
int array[1024];

void foo(int n) {
  int tmp;
  int local_array[n];

  int* dyn =
    (int*)malloc(n*sizeof(int));
}
```

Memory Allocation

- **Dynamic memory allocation**
  - **Introduction and goals**
  - **Allocation and deallocation (free)**
  - **Fragmentation**
- Explicit allocation implementation
  - Implicit free lists
  - Explicit free lists (Lab 5)
  - Segregated free lists
- Implicit deallocation: garbage collection
- Common memory-related bugs in C

Dynamic Memory Allocation

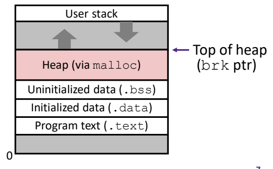
- Programmers use dynamic memory allocators to acquire virtual memory at run time
  - For data structures whose size (or lifetime) is known only at runtime
  - Manage the heap of a process' virtual memory:



- Types of allocators
  - ***Explicit* allocator:** programmer allocates and frees space
    - Example: `malloc` and `free` in C
  - ***Implicit* allocator:** programmer only allocates space (no free)
    - Example: garbage collection in Java, Caml, and Lisp

Dynamic Memory Allocation

- Allocator organizes heap as a collection of variable-sized *blocks*, which are either *allocated* or *free*
  - Allocator requests pages in the heap region; virtual memory hardware and OS kernel allocate these pages to the process
  - Application objects are typically smaller than pages, so the allocator manages blocks *within* pages
    - (Larger objects handled too; ignored here)



Allocating Memory in C

- Need to `#include <stdlib.h>`
- `void* malloc(size_t size)`
  - Allocates a continuous block of `size` bytes of uninitialized memory
  - Returns a pointer to the beginning of the allocated block; NULL indicates failed request
    - Typically aligned to an 8-byte (x86) or 16-byte (x86-64) boundary
    - Returns `NULL` if allocation failed (also sets `errno`) or `size==0`
  - Different blocks not necessarily adjacent
- Good practices:
  - `ptr = (int*) malloc(n*sizeof(int));`
    - `sizeof` makes code more portable
    - `void*` is implicitly cast into any pointer type; explicit typecast will help you catch coding errors when pointer types don't match

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- Related functions:
  - `void* calloc(size_t nitems, size_t size)`
    - "Zeros out" allocated block
  - `void* realloc(void* ptr, size_t size)`
    - Changes the size of a previously allocated block (if possible)
  - `void* sbrk(intptr_t increment)`
    - Used internally by allocators to grow or shrink the heap

Freeing Memory in C

- Need to `#include <stdlib.h>`
- `void free(void* p)`
  - Releases whole block pointed to by `p` to the pool of available memory
  - Pointer `p` must be the address *originally* returned by `m/c/realloc` (i.e. beginning of the block), otherwise throws system exception
  - Don't call `free` on a block that has already been released or on `NULL`

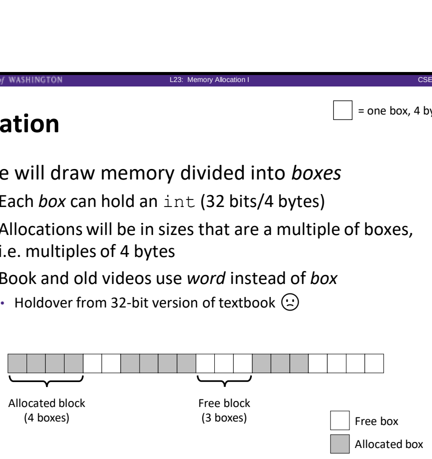
Memory Allocation Example in C

```
void foo(int n, int m) {
  int i, *p;
  p = (int*) malloc(n*sizeof(int));  /* allocate block of n ints  */
  if (p == NULL) {                   /* check for allocation error  */
    perror("malloc");
    exit(0);
  }
  for (i=0; i<n; i++)                /* initialize int array  */
    p[i] = i;
                             /* add space for m ints to end of p block  */
  p = (int*) realloc(p, (n+m)*sizeof(int));
  if (p == NULL) {                   /* check for allocation error  */
    perror("realloc");
    exit(0);
  }
  for (i=n; i < n+m; i++)            /* initialize new spaces  */
    p[i] = i;
  for (i=0; i<n+m; i++)              /* print new array  */
    printf("%d\n", p[i]);
  free(p);                           /* free p */
}
```

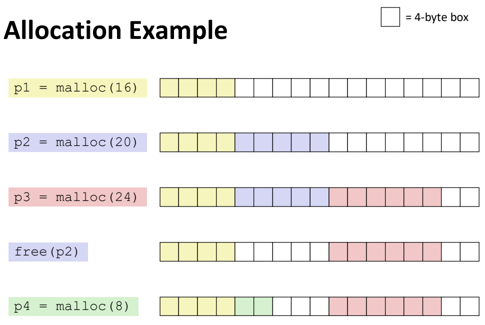
Notation

☐ = one box, 4 bytes

- We will draw memory divided into *boxes*
  - Each *box* can hold an `int` (32 bits/4 bytes)
  - Allocations will be in sizes that are a multiple of boxes, i.e. multiples of 4 bytes
  - Book and old videos use *word* instead of *box*
    - Holdover from 32-bit version of textbook ☹



Allocation Example



Implementation Interface

- **Applications**
  - Can issue arbitrary sequence of `malloc` and `free` requests
  - Must never access memory not currently allocated
  - Must never free memory not currently allocated
    - Also must only use `free` with previously `malloc`'ed blocks
- **Allocators**
  - Can't control number or size of allocated blocks
  - Must respond immediately to `malloc`
  - Must allocate blocks from free memory
  - Must align blocks so they satisfy all alignment requirements
  - Can't move the allocated blocks

Performance Goals

- **Goals:** Given some sequence of `malloc` and `free` requests $R_0, R_1, \ldots, R_k, \ldots, R_{n-1}$, maximize throughput and peak memory utilization
  - These goals are often conflicting

1) Throughput
- Number of completed requests per unit time
- Example:
  - If 5,000 `malloc` calls and 5,000 `free` calls completed in 10 seconds, then throughput is 1,000 operations/second

Performance Goals

- Definition: *Aggregate payload* P_k
  - `malloc(p)` results in a block with a *payload* of `p` bytes
  - After request R_k has completed, the *aggregate payload* P_k is the sum of currently allocated payloads
- Definition: *Current heap size* H_k
  - Assume H_k is monotonically non-decreasing
    - Allocator can increase size of heap using `sbrk`

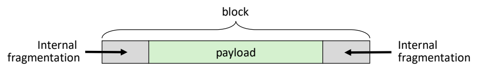
2) Peak Memory Utilization
- Defined as $U_k = (\max_{i \le k} P_i)/H_k$ after k+1 requests
- Goal: maximize utilization for a sequence of requests
- Why is this hard? And what happens to throughput?

Fragmentation

- Poor memory utilization is caused by *fragmentation*
  - Sections of memory are not used to store anything useful, but cannot satisfy allocation requests
  - Two types: *internal* and *external*
- **Recall:** Fragmentation in structs
  - Internal fragmentation was wasted space *inside* of the struct (between fields) due to alignment
  - External fragmentation was wasted space *between* struct instances (e.g. in an array) due to alignment
- Now referring to wasted space in the heap *inside* or *between* allocated blocks

Internal Fragmentation

- For a given block, *internal fragmentation* occurs if payload is smaller than the block

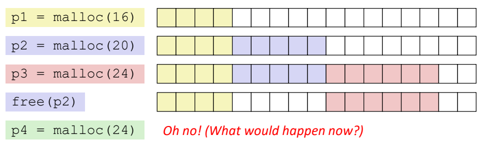


- **Causes:**
  - Padding for alignment purposes
  - Overhead of maintaining heap data structures (inside block, outside payload)
  - Explicit policy decisions (e.g. return a big block to satisfy a small request)
- Easy to measure because only depends on past requests

External Fragmentation

☐ = 4-byte box

- For the heap, *external fragmentation* occurs when allocation/free pattern leaves "holes" between blocks
  - That is, the aggregate payload is non-continuous
  - Can cause situations where there is enough aggregate heap memory to satisfy request, but no single free block is large enough



- Don't know what future requests will be
  - Difficult to impossible to know if past placements will become problematic

Peer Instruction Question

- Which of the following statements is FALSE?

A. Temporary arrays should *not* be allocated on the Heap
B. `malloc` returns an address filled with garbage
C. Peak memory utilization is a measure of both internal and external fragmentation
D. An allocation failure will cause your program to stop
E. We're lost…

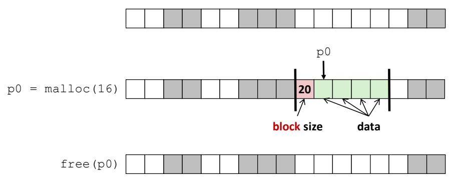
Implementation Issues

- How do we know how much memory to free given just a pointer?
- How do we keep track of the free blocks?
- How do we pick a block to use for allocation (when many might fit)?
- What do we do with the extra space when allocating a structure that is smaller than the free block it is placed in?
- How do we reinsert a freed block into the heap?

Knowing How Much to Free

☐ = 4-byte box (free)
▨ = 4-byte box (allocated)

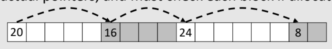
- Standard method
  - Keep the length of a block in the box preceding the block
    - This box is often called the ***header field*** or ***header***
  - Requires an extra box for every allocated block



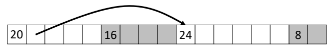
Keeping Track of Free Blocks

☐ = 4-byte box (free)
▨ = 4-byte box (allocated)

1) *Implicit free list* using length – links <u>all</u> blocks using math
   - No actual pointers, and must check each block if allocated or free



2) *Explicit free list* among <u>only the free blocks</u>, using pointers



3) *Segregated free list*
   - Different free lists for different size "classes"

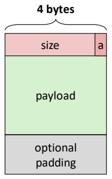
4) *Blocks sorted by size*
   - Can use a balanced binary tree (e.g. red-black tree) with pointers within each free block, and the length used as a key

Implicit Free Lists

e.g. with 8-byte alignment, possible values for size:
00001000 = 8 bytes
00010000 = 16 bytes
00011000 = 24 bytes
. . .

- For each block we need: **`size`, is-allocated?**
  - Could store using two boxes, but wasteful
- Standard trick
  - If blocks are aligned, some low-order bits of `size` are always 0
  - Use lowest bit as an allocated/free flag (fine as long as aligning to K>1)
  - When reading `size`, must remember to mask out this bit!

Format of allocated and free blocks:



a = 1: allocated block
a = 0: free block

size: block size (in bytes)

payload: application data (allocated blocks only)

If `x` is first box (header):

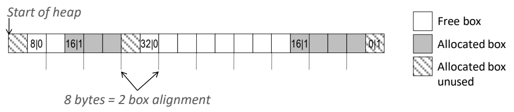
```
x = size | a;

a = x & 1;

size = x & ~1;
```

Implicit Free List Example

- Each block begins with header (size in bytes and allocated bit)
- Sequence of blocks in heap (`size|allocated`):
  8|0, 16|1, 32|0, 16|1



- 8-byte alignment for *payload*
  - May require initial padding (internal fragmentation)
  - Note `size`: padding is considered part of *previous* block
- Special one-box marker (0|1) marks end of list
  - Zero `size` is distinguishable from all other blocks

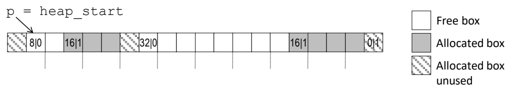
Implicit List: Finding a Free Block

(*p) gets the block *header*
(*p & 1) extracts the `allocated` bit
(*p & -2) extracts the `size`

- *First fit*
  - **Search list from beginning**, choose first free block that fits:

```
p = heap_start;
while ((p < end) &&     // not past end
       ((*p & 1) ||     // already allocated
        (*p <= len))) { // too small
  p = p + (*p & -2);    // go to next block (UNSCALED +)
}                       // p points to selected block or end
```

  - Can take time linear in total number of blocks
  - In practice can cause "splinters" at beginning of list



Implicit List: Finding a Free Block

- *Next fit*
  - Like first-fit, but **search list starting where previous search finished**
  - Should often be faster than first-fit: avoids re-scanning unhelpful blocks
  - Some research suggests that fragmentation is worse
- *Best fit*
  - Search the list, choose the ***best*** free block: large enough AND with fewest bytes left over
  - Keeps fragments small—usually helps fragmentation
  - Usually worse throughput