

x86-64 Programming II

CSE 351 Winter 2018

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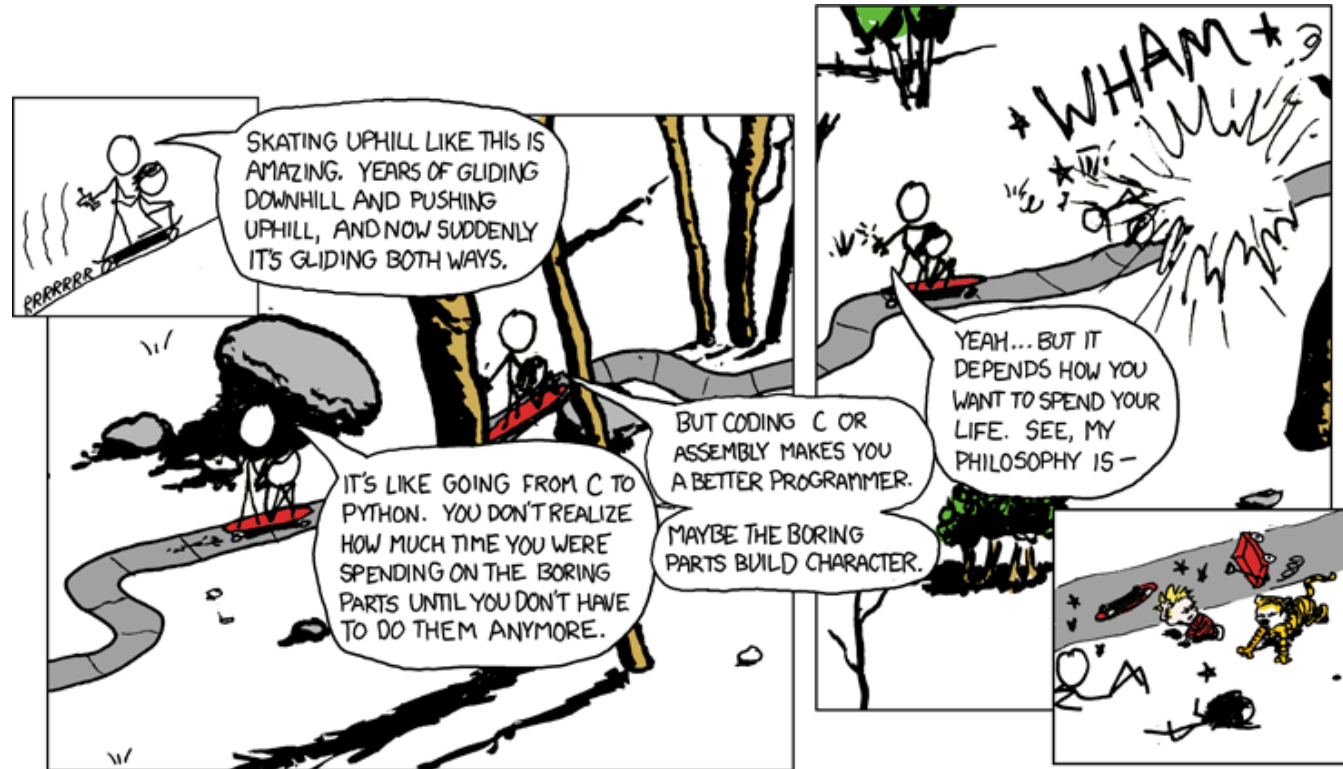
Parker DeWilde

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Waylon Huang

Vinny Palaniappan



<http://xkcd.com/409/>

Administrative

- ❖ Homework 2 (Ints and Floats) due Today!
 - By 11:59 pm, no late submissions!
 - x86-64 part due 1/29
- ❖ GDB Tutorial Session
 - Tomorrow, Thursday 1/25, 5 – 6 pm
 - Room : SIG 134

x86 Control Flow

- ❖ **Condition codes**
- ❖ **Conditional and unconditional branches**
- ❖ **Loops**
- ❖ **Switches**

Control Flow

Register	Use(s)
%rdi	1 st argument (x)
%rsi	2 nd argument (y)
%rax	return value

```
long max(long x, long y)
{
    long max;
    if (x > y) {
        max = x;
    } else {
        max = y;
    }
    return max;
}
```

```
max:
    ???
    movq    %rdi, %rax
    ???
    ???
    movq    %rsi, %rax
    ???
    ret
```

Control Flow

Register	Use(s)
%rdi	1 st argument (x)
%rsi	2 nd argument (y)
%rax	return value

```
long max(long x, long y)
{
    long max;
    if (x > y) {
        max = x;
    } else {
        max = y;
    }
    return max;
}
```

Conditional jump

Unconditional jump

```
max:
    if x <= y then jump to else
    movq    %rdi, %rax
    jump to done
else:
    movq    %rsi, %rax
done:
    ret
```

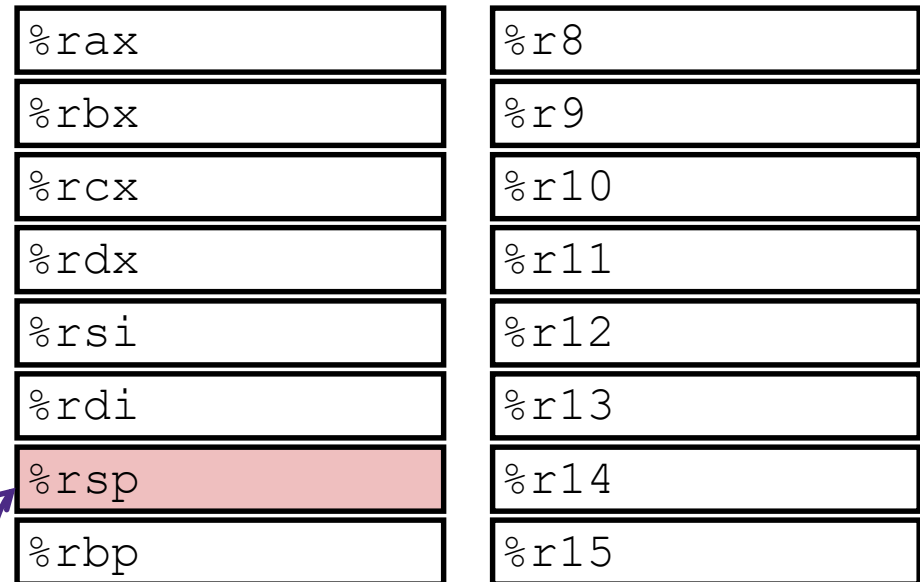
Conditionals and Control Flow

- ❖ Conditional branch/*jump*
 - Jump to somewhere else if some *condition* is true, otherwise execute next instruction
- ❖ Unconditional branch/*jump*
 - *Always* jump when you get to this instruction
- ❖ Together, they can implement most control flow constructs in high-level languages:
 - **if** (*condition*) **then** {...} **else** {...}
 - **while** (*condition*) {...}
 - **do** {...} **while** (*condition*)
 - **for** (*initialization*; *condition*; *iterative*) {...}
 - **switch** {...}

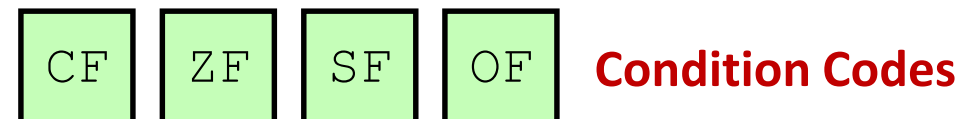
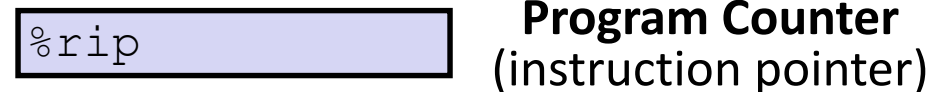
Processor State (x86-64, partial)

- ❖ Information about currently executing program
 - Temporary data (`%rax`, ...)
 - Location of runtime stack (`%rsp`)
 - Location of current code control point (`%rip`, ...)
 - Status of recent tests (**CF**, **ZF**, **SF**, **OF**)
 - Single bit registers:

Registers

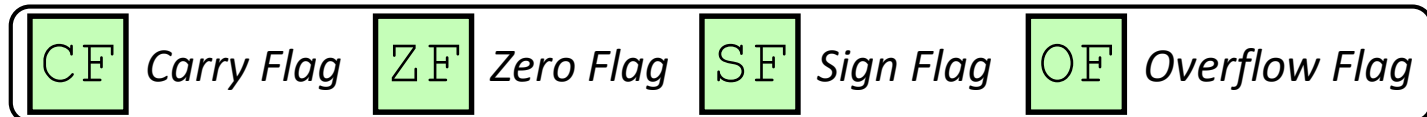


current top of the Stack



Condition Codes (Implicit Setting)

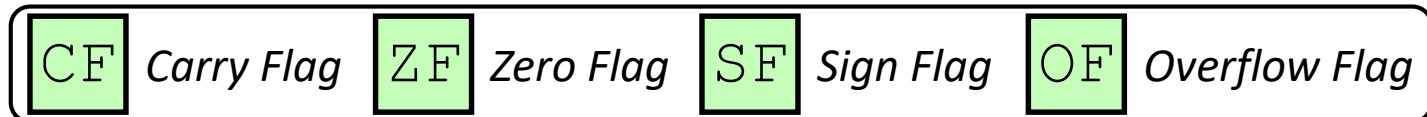
- ❖ *Implicitly* set by **arithmetic** operations
 - (think of it as side effects)
 - Example: **addq** src, dst \leftrightarrow $r = d+s$
 - **CF=1** if carry out from MSB (unsigned overflow)
 - **ZF=1** if $r==0$
 - **SF=1** if $r<0$ (assuming signed, actually just if MSB is 1)
 - **OF=1** if two's complement (signed) overflow
($s>0 \ \&\& \ d>0 \ \&\& \ r<0$) || ($s<0 \ \&\& \ d<0 \ \&\& \ r>=0$)
 - **Not set by lea instruction (beware!)**



Condition Codes (Explicit Setting: Compare)

❖ Explicitly set by **Compare** instruction

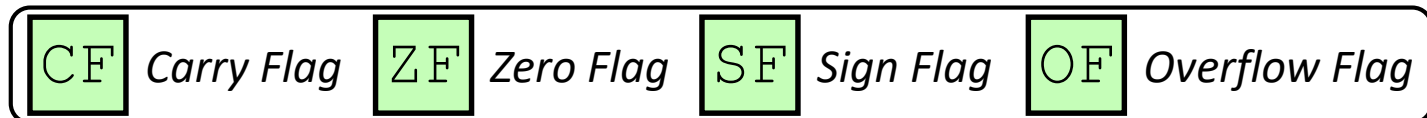
- `cmpq src1, src2`
- `cmpq a, b` sets flags based on $b-a$, but doesn't store
- **CF=1** if carry out from MSB (used for unsigned comparison)
- **ZF=1** if $a==b$
- **SF=1** if $(b-a) < 0$ (signed)
- **OF=1** if two's complement (signed) overflow
 $(a > 0 \ \&\& \ b < 0 \ \&\& \ (b-a) > 0) \ ||$
 $(a < 0 \ \&\& \ b > 0 \ \&\& \ (b-a) < 0)$



Condition Codes (Explicit Setting: Test)

❖ Explicitly set by **Test** instruction

- **testq** `src2, src1`
- **testq** `a, b` sets flags based on `a&b`, but doesn't store
 - Useful to have one of the operands be a *mask*
- **ZF=1** if `a&b==0`
- **SF=1** if `a&b<0` (signed)
- **CF** and **OF** set to 0
- Example: `testq %rax, %rax`
 - Tells you if (+), 0, or (-) based on ZF and SF



Using Condition Codes: Jump

❖ j^* Instructions

- Jumps to **target** (an address) based on condition codes

Instruction	Condition	Description
<code>jmp target</code>	1	Unconditional
<code>je target</code>	ZF	Equal / Zero
<code>jne target</code>	\sim ZF	Not Equal / Not Zero
<code>js target</code>	SF	Negative
<code>jns target</code>	\sim SF	Nonnegative
<code>jg target</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
<code>jge target</code>	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
<code>jl target</code>	$(SF \wedge OF)$	Less (Signed)
<code>jle target</code>	$(SF \wedge OF) \ \ ZF$	Less or Equal (Signed)
<code>ja target</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned ">")
<code>jb target</code>	CF	Below (unsigned "<")

Using Condition Codes: Set

❖ `set*` Instructions

- Set low-order byte of `dst` to 0 or 1 based on condition codes
- Does **not** alter remaining 7 bytes

Instruction	Condition	Description
<code>sete dst</code>	ZF	Equal / Zero
<code>setne dst</code>	\sim ZF	Not Equal / Not Zero
<code>sets dst</code>	SF	Negative
<code>setns dst</code>	\sim SF	Nonnegative
<code>setg dst</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
<code>setge dst</code>	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
<code>setl dst</code>	$(SF \wedge OF)$	Less (Signed)
<code>setle dst</code>	$(SF \wedge OF) \ \ ZF$	Less or Equal (Signed)
<code>seta dst</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned ">")
<code>setb dst</code>	CF	Below (unsigned "<")

Reminder: x86-64 Integer Registers

❖ Accessing the low-order byte:

<code>%rax</code>	<code>%al</code>
<code>%rbx</code>	<code>%bl</code>
<code>%rcx</code>	<code>%cl</code>
<code>%rdx</code>	<code>%dl</code>
<code>%rsi</code>	<code>%sil</code>
<code>%rdi</code>	<code>%dil</code>
<code>%rsp</code>	<code>%spl</code>
<code>%rbp</code>	<code>%bpl</code>

<code>%r8</code>	<code>%r8b</code>
<code>%r9</code>	<code>%r9b</code>
<code>%r10</code>	<code>%r10b</code>
<code>%r11</code>	<code>%r11b</code>
<code>%r12</code>	<code>%r12b</code>
<code>%r13</code>	<code>%r13b</code>
<code>%r14</code>	<code>%r14b</code>
<code>%r15</code>	<code>%r15b</code>

Reading Condition Codes

Register	Use(s)
%rdi	1 st argument (x)
%rsi	2 nd argument (y)
%rax	return value

❖ set* Instructions

- Set a low-order byte to 0 or 1 based on condition codes
- Operand is byte register (e.g. al, dl) or a byte in memory
- Do not alter remaining bytes in register
 - Typically use `movzbl` (zero-extended `mov`) to finish job

```
int gt(long x, long y)
{
    return x > y;
}
```

```
cmpq    %rsi, %rdi    #
setg    %al           #
movzbl  %al, %eax     #
ret
```

Reading Condition Codes

Register	Use(s)
%rdi	1 st argument (x)
%rsi	2 nd argument (y)
%rax	return value

❖ set* Instructions

- Set a low-order byte to 0 or 1 based on condition codes
- Operand is byte register (e.g. al, dl) or a byte in memory
- Do not alter remaining bytes in register
 - Typically use `movzbl` (zero-extended `mov`) to finish job

```
int gt(long x, long y)
{
    return x > y;
}
```

```
cmpq    %rsi, %rdi    # Compare x to y (x - y)
setg    %al           # Set when >
movzbl  %al, %eax     # Zero rest of %rax
ret
```

Aside: movz and movs

`movz __ src, regDest`

Move with zero extension

`movs __ src, regDest`

Move with sign extension

- Copy from a *smaller* source value to a *larger* destination
- Source can be memory or register; Destination *must* be a register
- Fill remaining bits of dest with **zero** (`movz`) or **sign bit** (`movs`)

`movzSD` / `movsSD`:

S – size of source (**b** = 1 byte, **w** = 2)

D – size of dest (**w** = 2 bytes, **l** = 4, **q** = 8)

Example:

`movzbq %al, %rbx`

0x??	0x??	0x??	0x??	0x??	0x??	0x??	0xFF	←%rax
------	------	------	------	------	------	------	------	-------

0x00	0x00	0x00	0x00	0x00	0x00	0x00	0xFF	←%rbx
------	------	------	------	------	------	------	------	-------

Aside: movz and movs

`movz __ src, regDest`

Move with zero extension

`movs __ src, regDest`

Move with sign extension

- Copy from a *smaller* source value to a *larger* destination
- Source can be memory or register; Destination *must* be a register
- Fill remaining bits of dest with **zero** (`movz`) or **sign bit** (`movs`)

`movz`*SD* / `movs`*SD*:

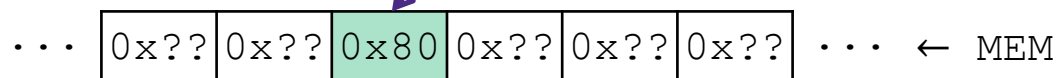
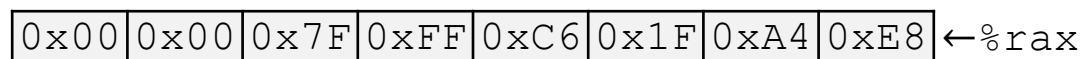
S – size of source (**b** = 1 byte, **w** = 2)

D – size of dest (**w** = 2 bytes, **l** = 4, **q** = 8)

Note: In x86-64, any instruction that generates a 32-bit (long word) value for a register also sets the high-order portion of the register to 0. Good example on p. 184 in the textbook.

Example:

`movsbl (%rax), %ebx`



Copy 1 byte from memory into 8-byte register & sign extend it

Using Condition Codes: Jump

❖ j^* Instructions

- Jumps to **target** (an address) based on condition codes

Instruction	Condition	Description
<code>je target</code>	ZF	Equal / Zero
<code>jne target</code>	$\sim ZF$	Not Equal / Not Zero
<code>js target</code>	SF	Negative
<code>jns target</code>	$\sim SF$	Nonnegative
<code>jg target</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
<code>jge target</code>	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
<code>j1 target</code>	$(SF \wedge OF)$	Less (Signed)
<code>jle target</code>	$(SF \wedge OF) \ \ ZF$	Less or Equal (Signed)
<code>ja target</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned ">")
<code>jb target</code>	CF	Below (unsigned "<")

Choosing instructions for conditionals

- ❖ All arithmetic instructions set condition flags based on result of operation (op)
 - Conditionals are comparisons against 0

```
addq 5, (p)
*p = *p + 5;
```

```
addq 5, (p)
je: *p+5 == 0
jne: *p+5 != 0
jg: *p+5 > 0
jl: *p+5 < 0
```

		addq (op) s, d
je	"Equal"	d (op) s == 0
jne	"Not equal"	d (op) s != 0
js	"Sign" (negative)	d (op) s < 0
jns	(non-negative)	d (op) s >= 0
jg	"Greater"	d (op) s > 0
jge	"Greater or equal"	d (op) s >= 0
jl	"Less"	d (op) s < 0
jle	"Less or equal"	d (op) s <= 0
ja	"Above" (unsigned >)	d (op) s > 0U
jb	"Below" (unsigned <)	d (op) s < 0U

Choosing instructions for conditionals

- ❖ Reminder: `cmp` is like `sub`, `test` is like `and`
 - Result is not stored anywhere

		<code>cmp a,b</code>	<code>test a,b</code>
je	“Equal”	<code>b == a</code>	<code>b&a == 0</code>
jne	“Not equal”	<code>b != a</code>	<code>b&a != 0</code>
js	“Sign” (negative)	<code>b-a < 0</code>	<code>b&a < 0</code>
jns	(non-negative)	<code>b-a >= 0</code>	<code>b&a >= 0</code>
jg	“Greater”	<code>b > a</code>	<code>b&a > 0</code>
jge	“Greater or equal”	<code>b >= a</code>	<code>b&a >= 0</code>
jl	“Less”	<code>b < a</code>	<code>b&a < 0</code>
jle	“Less or equal”	<code>b <= a</code>	<code>b&a <= 0</code>
ja	“Above” (unsigned >)	<code>b > a</code>	<code>b&a > 0U</code>
jb	“Below” (unsigned <)	<code>b < a</code>	<code>b&a < 0U</code>

```

cmpq 5, (p)
je:   *p == 5
jne:  *p != 5
jg:   *p > 5
jl:   *p < 5

```

```

testq a, a
je:   a == 0
jne:  a != 0
jg:   a > 0
jl:   a < 0

```

```

testb a, 0x1
je:   aLSB == 0
jne:  aLSB == 1

```

Choosing instructions for conditionals

		cmp a,b	test a,b
je	"Equal"	b == a	b&a == 0
jne	"Not equal"	b != a	b&a != 0
js	"Sign" (negative)	b-a < 0	b&a < 0
jns	(non-negative)	b-a >= 0	b&a >= 0
jb	"Below" (unsigned <)	b < a	b&a < 0U
jbe	"Below or equal" (unsigned <=)	b <= a	b&a <= 0U
ja	"Above" (unsigned >)	b > a	b&a > 0U
jge	"Greater or equal"	b >= a	b&a >= 0
jae	"Above or equal" (unsigned >=)	b >= a	b&a >= 0U
jg	"Greater"	b > a	b&a > 0
jl	"Less"	b < a	b&a < 0
jle	"Less or equal"	b <= a	b&a <= 0

```

    if (x < 3 && x == y) {
        return 1;
    } else {
        return 2;
    }
    
```

If either %al or %bl are false

```

    cmpq $3, %rdi
    setl %al
    cmpq %rsi, %rdi
    sete %bl
    testb %al, %bl
    je T2
T1: # x < 3 && x == y:
    movq $1, %rax
    ret
T2: # else
    movq $2, %rax
    ret
    
```

Choosing instructions for conditionals

		<code>cmp a,b</code>	<code>test a,b</code>
<code>je</code>	“Equal”	<code>b == a</code>	<code>b&a == 0</code>
<code>jne</code>	“Not equal”	<code>b != a</code>	<code>b&a != 0</code>
<code>js</code>	“Sign” (negative)	<code>b-a < 0</code>	<code>b&a < 0</code>
<code>jns</code>	(non-negative)	<code>b-a >= 0</code>	<code>b&a >= 0</code>
<code>jg</code>	“Greater”	<code>b > a</code>	<code>b&a > 0</code>
<code>jge</code>	“Greater or equal”	<code>b >= a</code>	<code>b&a >= 0</code>
<code>jl</code>	“Less”	<code>b < a</code>	<code>b&a < 0</code>
<code>jle</code>	“Less or equal”	<code>b <= a</code>	<code>b&a <= 0</code>
<code>ja</code>	“Above” (unsigned >)	<code>b > a</code>	<code>b&a > 0U</code>
<code>jb</code>	“Below” (unsigned <)	<code>b < a</code>	<code>b&a < 0U</code>

❖ <https://godbolt.org/g/KntpyG>

```

if (x < 3 && x == y) {
    return 1;
} else {
    return 2;
}

```

```

cmpq $3, %rdi
setl %al

cmpq %rsi, %rdi
sete %bl

testb %al, %bl
je T2

```

T1: # `x < 3 && x == y`:

```

movq $1, %rax
ret

```

T2: # `else`

```

movq $2, %rax
ret

```

Question

Register	Use(s)
%rdi	1 st argument (x)
%rsi	2 nd argument (y)
%rax	return value

- A. `cmpq %rsi, %rdi`
`jle .L4`
- B. `cmpq %rsi, %rdi`
`jg .L4`
- C. `testq %rsi, %rdi`
`jle .L4`
- D. `testq %rsi, %rdi`
`jg .L4`
- E. We're lost...

```
long absdiff(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    _____
    _____
                                     # x > y:
    movq    %rdi, %rax
    subq   %rsi, %rax
    ret
.L4:                                     # x <= y:
    movq   %rsi, %rax
    subq   %rdi, %rax
    ret
```

Summary

- ❖ Control flow in x86 determined by status of Condition Codes
 - Showed **C**arry, **Z**ero, **S**ign, and **O**verflow, though others exist
 - Set flags with arithmetic instructions (implicit) or Compare and Test (explicit)
 - Set instructions read out flag values
 - Jump instructions use flag values to determine next instruction to execute