

# x86-64 Programming I

CSE 351 Winter 2018

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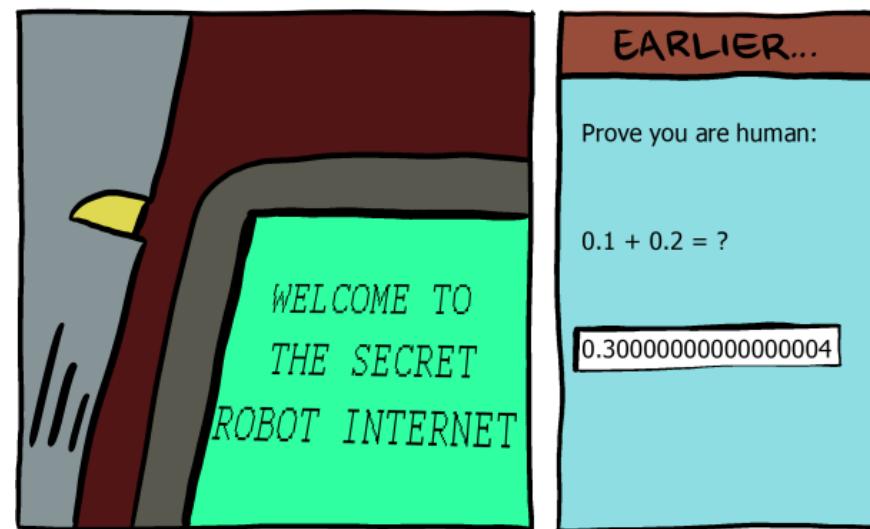
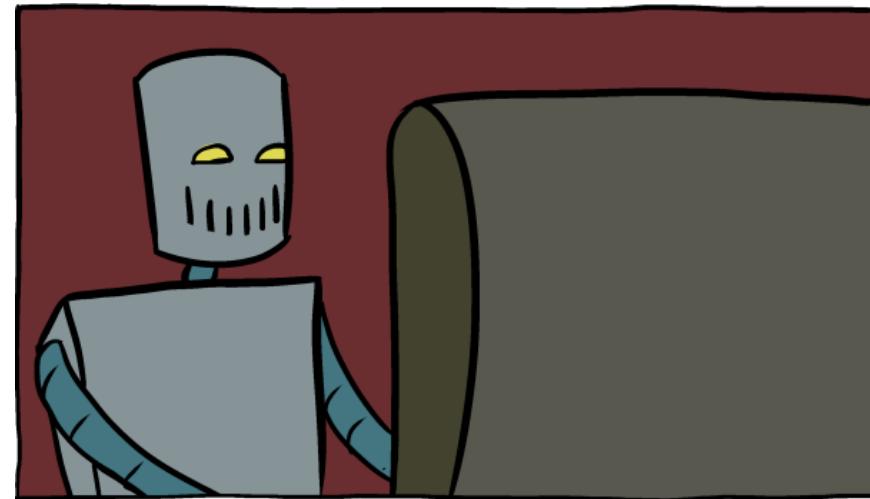
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<http://www.smbc-comics.com/?id=2999>

# Administrivia

- ❖ Homework 2 due Wednesday (1/24)
  - x86-64 part extended until Monday, January 29
- ❖ Lab 2 (x86-64) released today (1/22)
  - Due on Friday, February 2 @ 11:59 pm
- ❖ Don't forget to read the book!
  - Ch. 3 will be very helpful!

# x86-64 Programming I

- ❖ Operand types
- ❖ Moving data
- ❖ Arithmetic operations
- ❖ Memory addressing modes
  - swap example
- ❖ Address computation instruction (`lea`)

# Operand types

- ❖ **Immediate:** Constant integer data
  - Examples: `$0x400`, `$-533`
  - Like C literal, but prefixed with '\$'
  - Encoded with 1, 2, 4, or 8 bytes  
*depending on the instruction*
- ❖ **Register:** 1 of 16 integer registers
  - Examples: `%rax`, `%r13`
  - But `%rsp` reserved for special use
  - Others have special uses for particular instructions
- ❖ **Memory:** Consecutive bytes of memory at a computed address
  - Simplest example: `(%rax)`
  - Various other “address modes”

`%rax`

`%rcx`

`%rdx`

`%rbx`

`%rsi`

`%rdi`

`%rsp`

`%rbp`

`%rN`

# Moving Data

- ❖ General form: `mov_ source, destination`
  - Missing letter (\_) specifies size of operands
  - Note that due to backwards-compatible support for 8086 programs (16-bit machines!), “word” means 16 bits = 2 bytes in x86 instruction names
  - Lots of these in typical code
  
- ❖ `movb src, dst`
  - Move 1-byte “byte”
- ❖ `movw src, dst`
  - Move 2-byte “word”
- ❖ `movl src, dst`
  - Move 4-byte “long word”
- ❖ `movq src, dst`
  - Move 8-byte “quad word”

# movq Operand Combinations

Source	Dest	Src, Dest	C Analog
movq	Imm { Reg	movq \$0x4, %rax	var_a = 0x4;
	Mem	movq \$-147, (%rax)	*p_a = -147;
	Reg { Mem	movq %rax, %rdx movq %rax, (%rdx)	var_d = var_a; *p_d = var_a;
Mem	Reg	movq (%rax), %rdx	var_d = *p_a;

- ❖ *Cannot do memory-memory transfer with a single instruction*
  - How would you do it?

# Some Arithmetic Operations

## ❖ Binary (two-operand) Instructions:

- Maximum of one memory operand
- Beware argument order!
- No distinction between signed and unsigned
  - Only arithmetic vs. logical shifts
- How do you implement  
“r3 = r1 + r2”?

Format	Computation	
<b>addq</b> <i>src, dst</i>	$dst = dst + src$	( $dst += src$ )
<b>subq</b> <i>src, dst</i>	$dst = dst - src$	
<b>imulq</b> <i>src, dst</i>	$dst = dst * src$	signed mult
<b>sarq</b> <i>src, dst</i>	$dst = dst >> src$	Arithmetic
<b>shrq</b> <i>src, dst</i>	$dst = dst >> src$	Logical
<b>shlq</b> <i>src, dst</i>	$dst = dst << src$	(same as <b>salq</b> )
<b>xorq</b> <i>src, dst</i>	$dst = dst ^ src$	
<b>andq</b> <i>src, dst</i>	$dst = dst \& src$	
<b>orq</b> <i>src, dst</i>	$dst = dst   src$	



operand size specifier

# Some Arithmetic Operations

- ❖ Unary (one-operand) Instructions:

Format	Computation
<b>incq</b> <i>dst</i>	$dst = dst + 1$
<b>decq</b> <i>dst</i>	$dst = dst - 1$
<b>negq</b> <i>dst</i>	$dst = -dst$
<b>notq</b> <i>dst</i>	$dst = \sim dst$

- ❖ See CSPP Section 3.5.5 for more instructions:  
mulq, cqto, idivq, divq

# Arithmetic Example

```
long simple_arith(long x, long y)
{
    long t1 = x + y;
    long t2 = t1 * 3;
    return t2;
}
```

Register	Use(s)
%rdi	1 <sup>st</sup> argument (x)
%rsi	2 <sup>nd</sup> argument (y)
%rax	return value

```
y += x;
y *= 3;
long r = y;
return r;
```

```
simple_arith:
    addq    %rdi, %rsi
    imulq   $3, %rsi
    movq    %rsi, %rax
    ret
```

# Example of Basic Addressing Modes

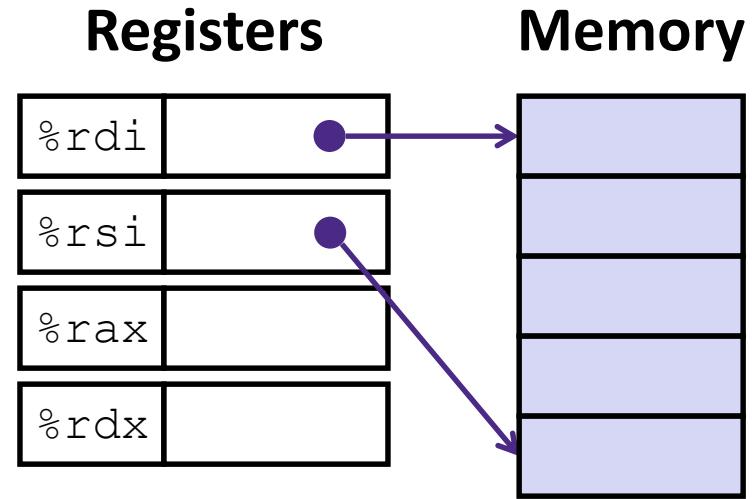
```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movq (%rdi), %rax
    movq (%rsi), %rdx
    movq %rdx, (%rdi)
    movq %rax, (%rsi)
    ret
```

# Understanding swap()

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movq (%rdi), %rax
    movq (%rsi), %rdx
    movq %rdx, (%rdi)
    movq %rax, (%rsi)
    ret
```



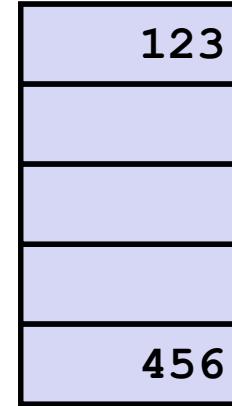
<u>Register</u>	<u>Variable</u>
%rdi	↔ xp
%rsi	↔ yp
%rax	↔ t0
%rdx	↔ t1

# Understanding swap()

Registers

%rdi	0x120
%rsi	0x100
%rax	
%rdx	

Memory



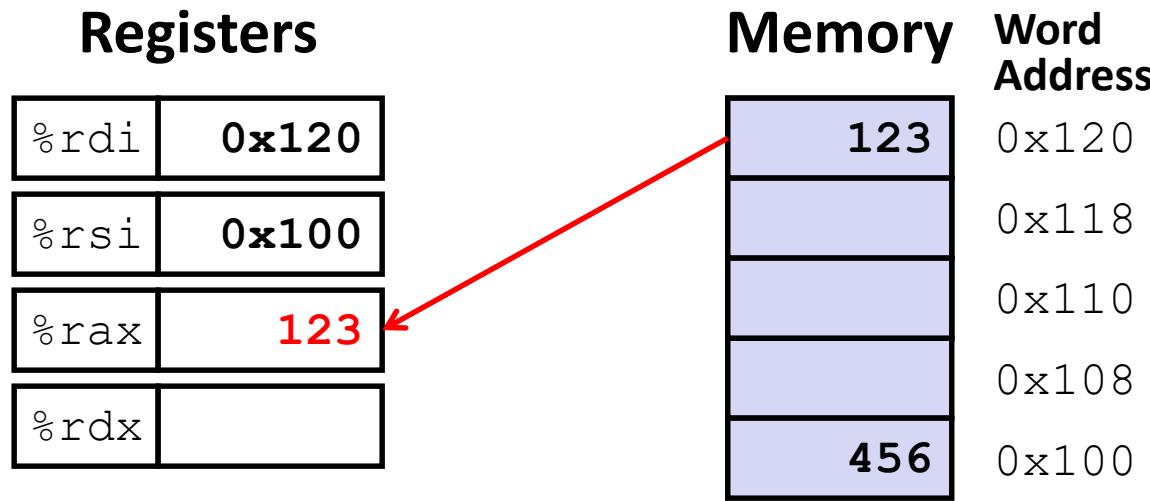
Word Address

0x120  
0x118  
0x110  
0x108  
0x100

Swap:

```
movq (%rdi), %rax    # t0 = *xp
movq (%rsi), %rdx    # t1 = *yp
movq %rdx, (%rdi)    # *xp = t1
movq %rax, (%rsi)    # *yp = t0
ret
```

# Understanding swap()



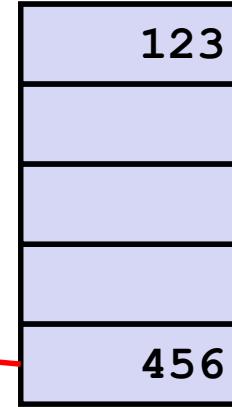
Swap:

```
movq (%rdi), %rax    # t0 = *xp
movq (%rsi), %rdx    # t1 = *yp
movq %rdx, (%rdi)    # *xp = t1
movq %rax, (%rsi)    # *yp = t0
ret
```

# Understanding swap()

**Registers**

%rdi	0x120
%rsi	0x100
%rax	123
%rdx	456

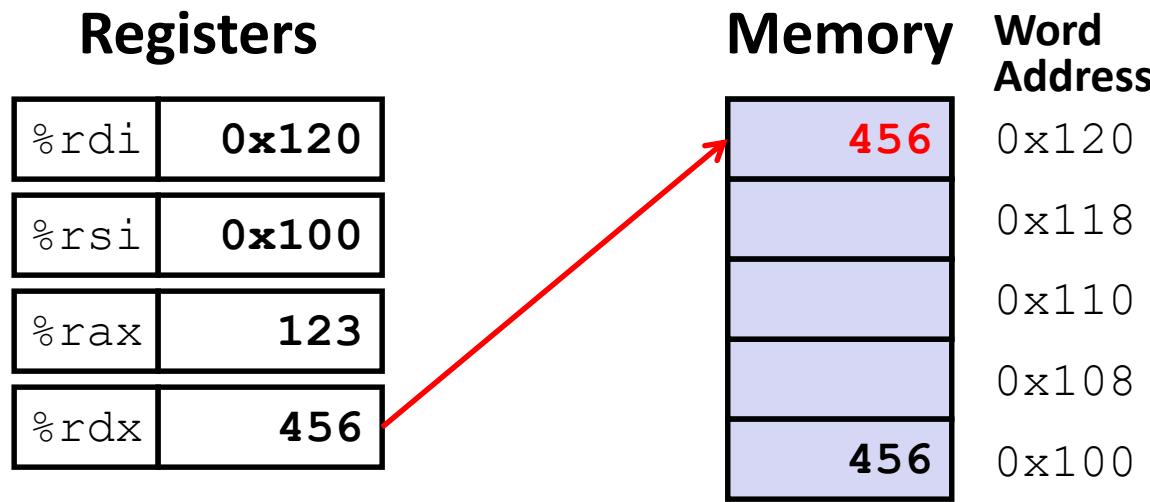
**Memory****Word Address**

0x120  
0x118  
0x110  
0x108  
0x100

**Swap:**

```
movq (%rdi), %rax    # t0 = *xp
movq (%rsi), %rdx    # t1 = *yp
movq %rdx, (%rdi)    # *xp = t1
movq %rax, (%rsi)    # *yp = t0
ret
```

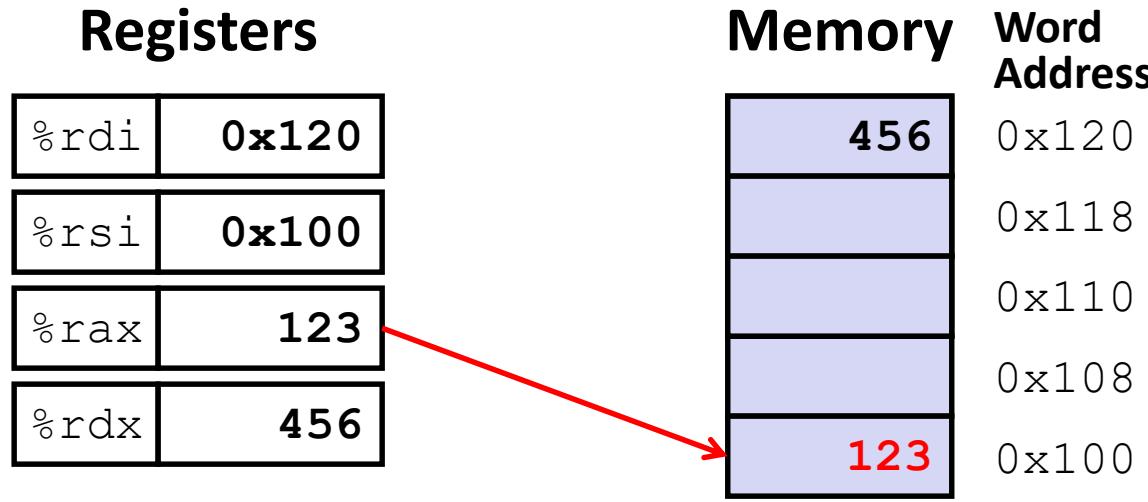
# Understanding swap()



Swap:

```
movq (%rdi), %rax    # t0 = *xp
movq (%rsi), %rdx    # t1 = *yp
movq %rdx, (%rdi)    # *xp = t1
movq %rax, (%rsi)    # *yp = t0
ret
```

# Understanding swap()



Swap:

```
movq (%rdi), %rax    # t0 = *xp
movq (%rsi), %rdx    # t1 = *yp
movq %rdx, (%rdi)    # *xp = t1
movq %rax, (%rsi)    # *yp = t0
ret
```

# Memory Addressing Modes: Basic

❖ **Indirect:** (R) Mem[Reg[R]]

- Data in register R specifies the memory address
- Like pointer dereference in C
- Example: `movq (%rcx), %rax`

❖ **Displacement:** D (R) Mem[Reg[R]+D]

- Data in register R specifies the *start* of some memory region
- Constant displacement D specifies the offset from that address
- Example: `movq 8(%rbp), %rdx`

# Complete Memory Addressing Modes

## ❖ General:

- $D(Rb, Ri, S) \quad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S + D]$ 
  - $Rb$ : Base register (any register)
  - $Ri$ : Index register (any register except  $\%rsp$ )
  - $S$ : Scale factor (1, 2, 4, 8) – *why these numbers?*
  - $D$ : Constant displacement value (a.k.a. immediate)

## ❖ Special cases (see CSPP Figure 3.3 on p.181)

- $D(Rb, Ri) \quad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D] \quad (S=1)$
- $(Rb, Ri, S) \quad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S] \quad (D=0)$
- $(Rb, Ri) \quad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]] \quad (S=1, D=0)$
- $(, Ri, S) \quad \text{Mem}[\text{Reg}[Ri] * S] \quad (Rb=0, D=0)$

# Address Computation Examples

%rdx	0xf000
%rcx	0x0100

$D(Rb, Ri, S) \rightarrow$   
 $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S + D]$

Expression	Address Computation	Address
$0x8(%rdx)$		
$(%rdx, %rcx)$		
$(%rdx, %rcx, 4)$		
$0x80(,%rdx,2)$		

# Address Computation Instruction

- ❖ `leaq src, dst`
  - “`lea`” stands for *load effective address*
  - `src` is address expression (any of the formats we’ve seen)
  - `dst` is a register
  - Sets `dst` to the *address* computed by the `src` expression  
**(does not go to memory! – it just does math)**
  - Example: `leaq (%rdx,%rcx,4), %rax`
- ❖ Uses:
  - Computing addresses without a memory reference
    - e.g. translation of `p = &x[i];`
  - Computing arithmetic expressions of the form `x+k*i+d`
    - Though `k` can only be 1, 2, 4, or 8

# Example: lea vs. mov

**Registers**

%rax	
%rbx	
%rcx	<b>0x4</b>
%rdx	<b>0x100</b>
%rdi	
%rsi	

**Memory**

<b>0x400</b>
<b>0xF</b>
<b>0x8</b>
<b>0x10</b>
<b>0x1</b>

**Word Address**

0x120  
0x118  
0x110  
0x108  
0x100

```
leaq (%rdx,%rcx,4), %rax
movq (%rdx,%rcx,4), %rbx
leaq (%rdx), %rdi
movq (%rdx), %rsi
```

# Arithmetic Example

```
long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
leaq    (%rdi,%rsi), %rax
addq    %rdx, %rax
leaq    (%rsi,%rsi,2), %rdx
salq    $4, %rdx
leaq    4(%rdi,%rdx), %rcx
imulq   %rcx, %rax
ret
```

Register	Use(s)
%rdi	1 <sup>st</sup> argument (x)
%rsi	2 <sup>nd</sup> argument (y)
%rdx	3 <sup>rd</sup> argument (z)

- ❖ Interesting Instructions
  - leaq: “address” computation
  - salq: shift
  - imulq: multiplication
    - Only used once!

# Arithmetic Example

```
long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

Register	Use(s)
%rdi	x
%rsi	y
%rdx	z, t4
%rax	t1, t2, rval
%rcx	t5

arith:

<b>leaq</b>	(%rdi,%rsi), %rax	# rax/t1	= x + y
<b>addq</b>	%rdx, %rax	# rax/t2	= t1 + z
<b>leaq</b>	(%rsi,%rsi,2), %rdx	# rdx	= 3 * y
<b>salq</b>	\$4, %rdx	# rdx/t4	= (3*y) * 16
<b>leaq</b>	4(%rdi,%rdx), %rcx	# rcx/t5	= x + t4 + 4
<b>imulq</b>	%rcx, %rax	# rax/rval	= t5 * t2
<b>ret</b>			

# Peer Instruction Question

- ❖ Which of the following x86-64 instructions correctly calculates  $\%rax = 9 * \%rdi$ ?
  - A. **leaq (,%rdi,9), %rax**
  - B. **movq (,%rdi,9), %rax**
  - C. **leaq (%rdi,%rdi,8), %rax**
  - D. **movq (%rdi,%rdi,8), %rax**
  - E. **We're lost...**

# Summary

- ❖ **Operands:** Assembly operands include *immediates*, *registers*, and *data at a specified memory location*.
- ❖ **Memory Addressing Modes:** The addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways
  - *Base register, index register, scale factor, and displacement* map well to pointer arithmetic operations
- ❖ `lea` is address calculation instruction
  - Does NOT actually go to memory
  - Used to compute addresses or some arithmetic expressions