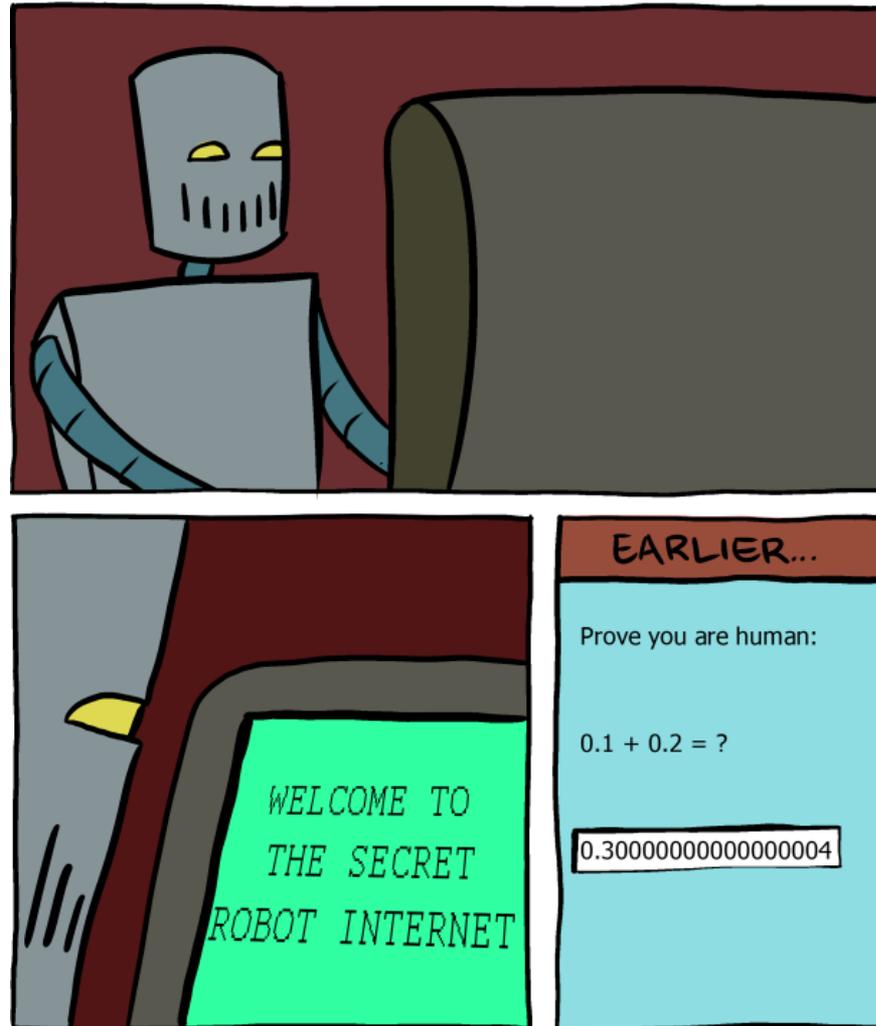


# x86-64 Programming I

CSE 351 Spring 2018



<http://www.smbc-comics.com/?id=2999>

# Administrivia

- ❖ Continue to check the course calendar please 😊

# Review: Operand types

- ❖ **Immediate:** Constant integer data
  - Examples: `$0x400`, `$-533`
  - Like C literal, but prefixed with ``$'`
  - Encoded with 1, 2, 4, or 8 bytes *depending on the instruction*
- ❖ **Register:** 1 of 16 integer registers
  - Examples: `%rax`, `%r13`
  - But `%rsp` reserved for special use
  - Others have special uses for particular instructions
- ❖ **Memory:** Consecutive bytes of memory at a computed address
  - Simplest example: `(%rax)`
  - Various other “address modes”

`%rax``%rcx``%rdx``%rbx``%rsi``%rdi``%rsp``%rbp``%rN`

# Moving Data

- ❖ General form: `mov_ source, destination`
  - Missing letter (`_`) specifies size of operands
  - Note that due to backwards-compatible support for 8086 programs (16-bit machines!), “word” means 16 bits = 2 bytes in x86 instruction names
  - Lots of these in typical code
- ❖ `movb src, dst`
  - Move 1-byte “**byte**”
- ❖ `movw src, dst`
  - Move 2-byte “**word**”
- ❖ `movl src, dst`
  - Move 4-byte “**long word**”
- ❖ `movq src, dst`
  - Move 8-byte “**quad word**”

# movq Operand Combinations

	Source	Dest	Src, Dest	C Analog
movq	Imm	Reg	movq \$0x4, %rax	var_a = 0x4;
		Mem	movq \$-147, (%rax)	*p_a = -147;
	Reg	Reg	movq %rax, %rdx	var_d = var_a;
		Mem	movq %rax, (%rdx)	*p_d = var_a;
	Mem	Reg	movq (%rax), %rdx	var_d = *p_a;

❖ *Cannot do memory-memory transfer with a single instruction*

- How would you do it?

# x86-64 Introduction

- ❖ Arithmetic operations
- ❖ Memory addressing modes
  - `swap` example
- ❖ Address computation instruction (`leaq`)

# Some Arithmetic Operations

## ❖ Binary (two-operand) Instructions:

- **Maximum of one memory operand**
- Beware argument order!
- No distinction between signed and unsigned
  - Only arithmetic vs. logical shifts
- How do you implement “ $r3 = r1 + r2$ ”?

Format	Computation	
<code>addq src, dst</code>	$dst = dst + src$	( $dst += src$ )
<code>subq src, dst</code>	$dst = dst - src$	
<code>imulq src, dst</code>	$dst = dst * src$	signed mult
<code>sarq src, dst</code>	$dst = dst \gg src$	Arithmetic
<code>shrq src, dst</code>	$dst = dst \gg src$	Logical
<code>shlq src, dst</code>	$dst = dst \ll src$	(same as <code>salq</code> )
<code>xorq src, dst</code>	$dst = dst \wedge src$	
<code>andq src, dst</code>	$dst = dst \& src$	
<code>orq src, dst</code>	$dst = dst   src$	

↑ operand size specifier

# Some Arithmetic Operations

## ❖ Unary (one-operand) Instructions:

Format	Computation	
<code>incq dst</code>	$dst = dst + 1$	increment
<code>decq dst</code>	$dst = dst - 1$	decrement
<code>negq dst</code>	$dst = -dst$	negate
<code>notq dst</code>	$dst = \sim dst$	bitwise complement

- ❖ See CSPP Section 3.5.5 for more instructions:  
`mulq`, `cqto`, `idivq`, `divq`

# Arithmetic Example

```
long simple_arith(long x, long y)
{
    long t1 = x + y;
    long t2 = t1 * 3;
    return t2;
}
```

Register	Use(s)
%rdi	1 <sup>st</sup> argument (x)
%rsi	2 <sup>nd</sup> argument (y)
%rax	return value

```
y += x;
y *= 3;
long r = y;
return r;
```

```
simple_arith:
    addq    %rdi, %rsi
    imulq   $3, %rsi
    movq    %rsi, %rax
    ret
```

# Example of Basic Addressing Modes

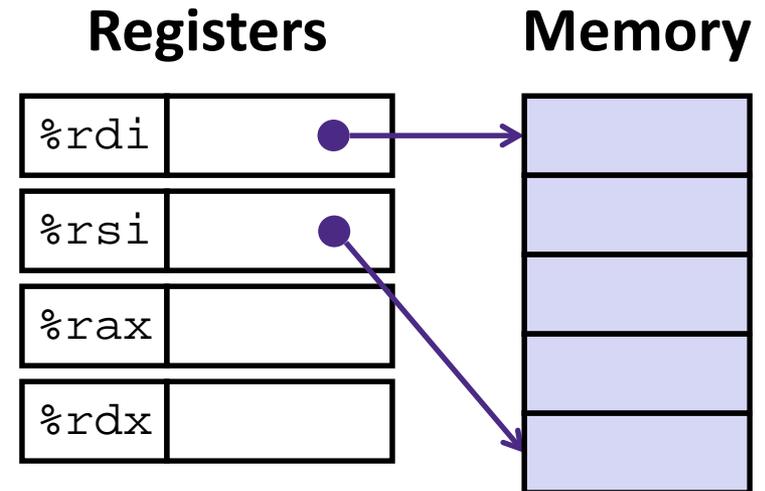
```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movq    (%rdi), %rax
    movq    (%rsi), %rdx
    movq    %rdx, (%rdi)
    movq    %rax, (%rsi)
    ret
```

# Understanding swap( )

```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movq    (%rdi), %rax
    movq    (%rsi), %rdx
    movq    %rdx, (%rdi)
    movq    %rax, (%rsi)
    ret
```



<u>Register</u>		<u>Variable</u>
%rdi	↔	xp
%rsi	↔	yp
%rax	↔	t0
%rdx	↔	t1

# Understanding `swap()`

## Registers

<code>%rdi</code>	<code>0x120</code>
<code>%rsi</code>	<code>0x100</code>
<code>%rax</code>	
<code>%rdx</code>	

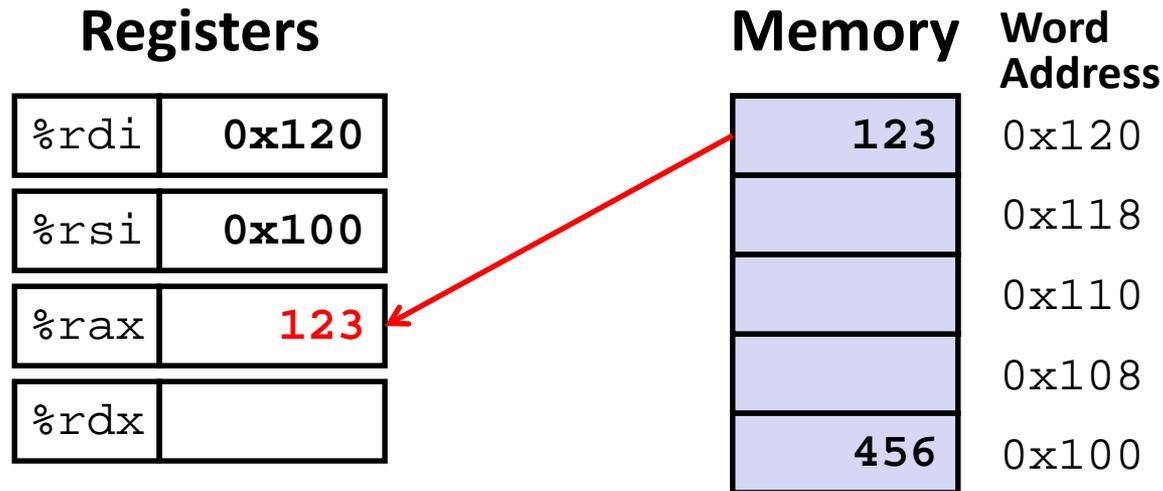
## Memory Word Address

123	0x120
	0x118
	0x110
	0x108
456	0x100

```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)   # *xp = t1
    movq    %rax, (%rsi)   # *yp = t0
    ret
```

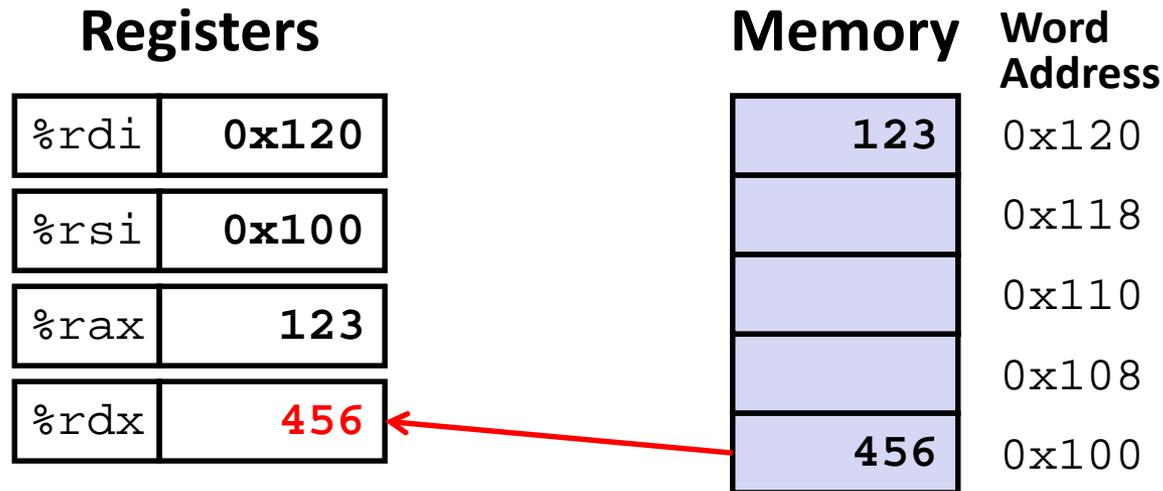
# Understanding swap( )



swap:

```
movq (%rdi), %rax # t0 = *xp
movq (%rsi), %rdx # t1 = *yp
movq %rdx, (%rdi) # *xp = t1
movq %rax, (%rsi) # *yp = t0
ret
```

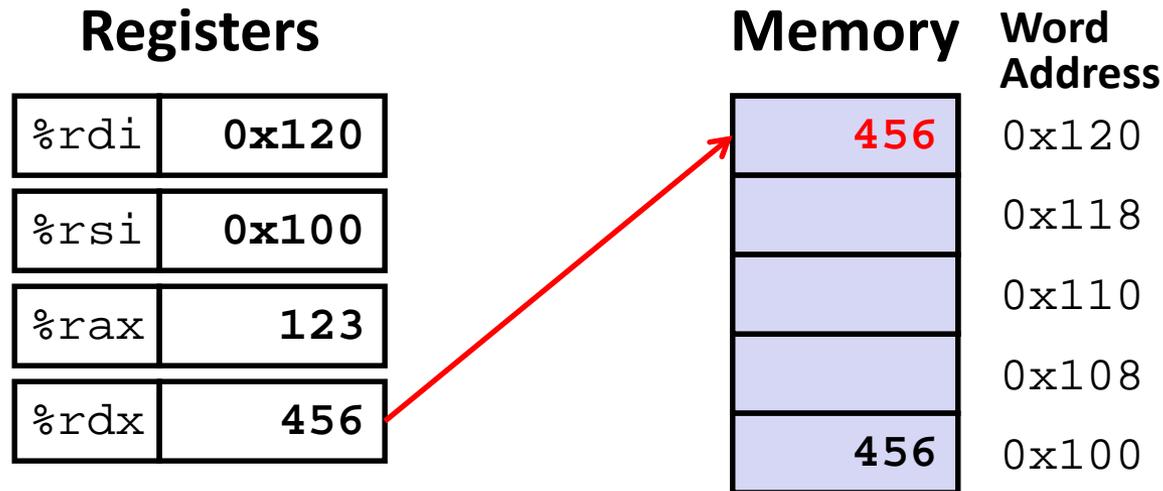
# Understanding swap ( )



```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp  
    movq    (%rsi), %rdx    # t1 = *yp  
    movq    %rdx, (%rdi)   # *xp = t1  
    movq    %rax, (%rsi)   # *yp = t0  
    ret
```

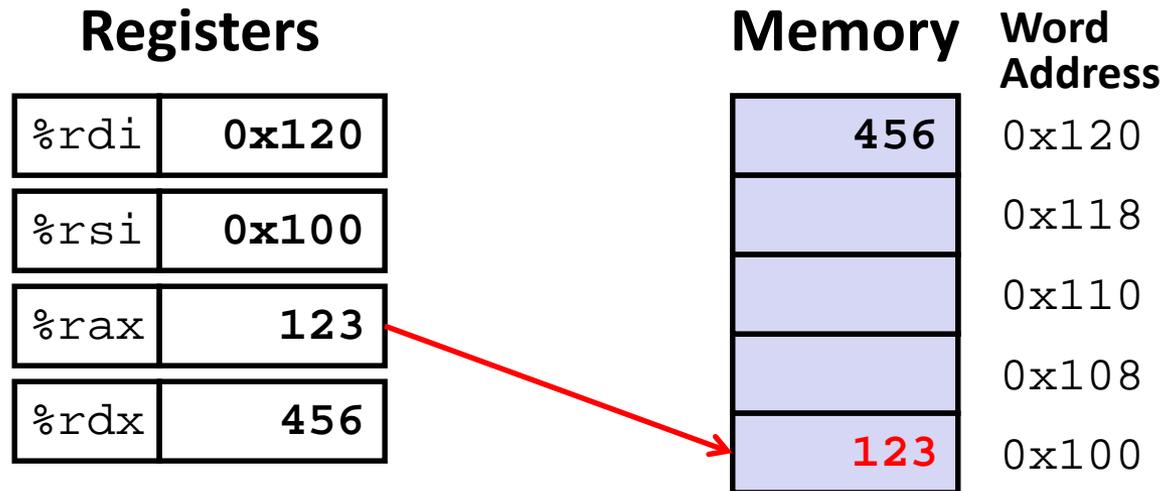
# Understanding `swap()`



`swap:`

```
movq    (%rdi), %rax    # t0 = *xp
movq    (%rsi), %rdx    # t1 = *yp
movq    %rdx, (%rdi)  # *xp = t1
movq    %rax, (%rsi)    # *yp = t0
ret
```

# Understanding swap ( )



```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)    # *xp = t1
    movq    %rax, (%rsi)    # *yp = t0
    ret
```

# Memory Addressing Modes: Basic

❖ **Indirect:**  $(R)$   $\text{Mem}[\text{Reg}[R]]$

- Data in register  $R$  specifies the memory address
- Like pointer dereference in C
- Example: `movq (%rcx), %rax`

❖ **Displacement:**  $D(R)$   $\text{Mem}[\text{Reg}[R]+D]$

- Data in register  $R$  specifies the *start* of some memory region
- Constant displacement  $D$  specifies the offset from that address
- Example: `movq 8(%rbp), %rdx`

# Complete Memory Addressing Modes

## ❖ General:

- $D(Rb, Ri, S)$      $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S + D]$ 
  - Rb:        Base register (any register)
  - Ri:        Index register (any register except `%rsp`)
  - S:        Scale factor (1, 2, 4, 8) – *why these numbers?*
  - D:        Constant displacement value (a.k.a. immediate)

## ❖ Special cases (see CSPP Figure 3.3 on p.181)

- $D(Rb, Ri)$          $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D]$     ( $S=1$ )
- $(Rb, Ri, S)$         $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S]$     ( $D=0$ )
- $(Rb, Ri)$           $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]]$         ( $S=1, D=0$ )
- $(, Ri, S)$           $\text{Mem}[\text{Reg}[Ri] * S]$                 ( $Rb=0, D=0$ )

# Address Computation Examples

<code>%rdx</code>	<code>0xf000</code>
<code>%rcx</code>	<code>0x0100</code>

$$D(Rb, Ri, S) \rightarrow$$

$$\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S + D]$$

Expression	Address Computation	Address
<code>0x8(%rdx)</code>		
<code>(%rdx,%rcx)</code>		
<code>(%rdx,%rcx,4)</code>		
<code>0x80(,%rdx,2)</code>		

# Address Computation Instruction

- ❖ `leaq src, dst`
  - “lea” stands for *load effective address*
  - `src` is address expression (any of the formats we’ve seen)
  - `dst` is a register
  - Sets `dst` to the *address* computed by the `src` expression (**does not go to memory! – it just does math**)
  - Example: `leaq (%rdx,%rcx,4), %rax`
- ❖ Uses:
  - Computing addresses without a memory reference
    - e.g. translation of `p = &x[i];`
  - Computing arithmetic expressions of the form  $x+k*i+d$ 
    - Though `k` can only be 1, 2, 4, or 8

# Example: lea vs. mov

Registers		Memory	Word Address
%rax		0x400	0x120
%rbx		0xF	0x118
%rcx	0x4	0x8	0x110
%rdx	0x100	0x10	0x108
%rdi		0x1	0x100
%rsi			

```
leaq (%rdx,%rcx,4), %rax
movq (%rdx,%rcx,4), %rbx
leaq (%rdx), %rdi
movq (%rdx), %rsi
```

# Arithmetic Example

```
long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
    leaq    (%rdi,%rsi), %rax
    addq    %rdx, %rax
    leaq    (%rsi,%rsi,2), %rdx
    salq    $4, %rdx
    leaq    4(%rdi,%rdx), %rcx
    imulq   %rcx, %rax
    ret
```

Register	Use(s)
%rdi	1 <sup>st</sup> argument (x)
%rsi	2 <sup>nd</sup> argument (y)
%rdx	3 <sup>rd</sup> argument (z)

- ❖ Interesting Instructions
  - leaq: “address” computation
  - salq: shift
  - imulq: multiplication
    - Only used once!

# Arithmetic Example

```

long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}

```

Register	Use(s)
%rdi	x
%rsi	y
%rdx	z, t4
%rax	t1, t2, rval
%rcx	t5

```

arith:
    leaq    (%rdi,%rsi), %rax    # rax/t1    = x + y
    addq    %rdx, %rax          # rax/t2    = t1 + z
    leaq    (%rsi,%rsi,2), %rdx  # rdx       = 3 * y
    salq    $4, %rdx           # rdx/t4    = (3*y) * 16
    leaq    4(%rdi,%rdx), %rcx   # rcx/t5    = x + t4 + 4
    imulq   %rcx, %rax          # rax/rval  = t5 * t2
    ret

```

# Peer Instruction Question

- ❖ Which of the following x86-64 instructions correctly calculates `%rax=9 * %rdi`?
- A. `leaq (, %rdi, 9), %rax`
- B. `movq (, %rdi, 9), %rax`
- C. `leaq (%rdi, %rdi, 8), %rax`
- D. `movq (%rdi, %rdi, 8), %rax`
- E. We're lost...

# Summary

- ❖ **Memory Addressing Modes:** The addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways
  - *Base register, index register, scale factor, and displacement* map well to pointer arithmetic operations
- ❖ `lea` is address calculation instruction
  - Does NOT actually go to memory
  - Used to compute addresses or some arithmetic expressions