Floating Point I

CSE 351 Autumn 2018

Instructor: Teaching Assistants:

Justin Hsia Akshat Aggarwal

Brian Dai

Kevin Bi

Sophie Tian

An Wang

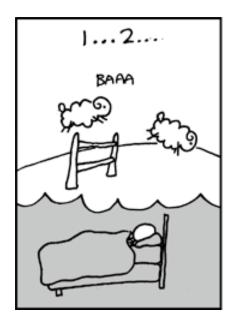
Britt Henderson

Kory Watson

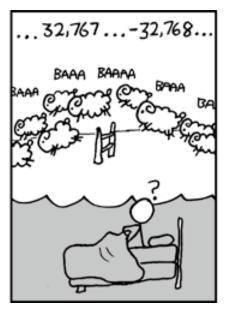
Teagan Horkan

Andrew Hu James Shin

Riley Germundson









http://xkcd.com/571/

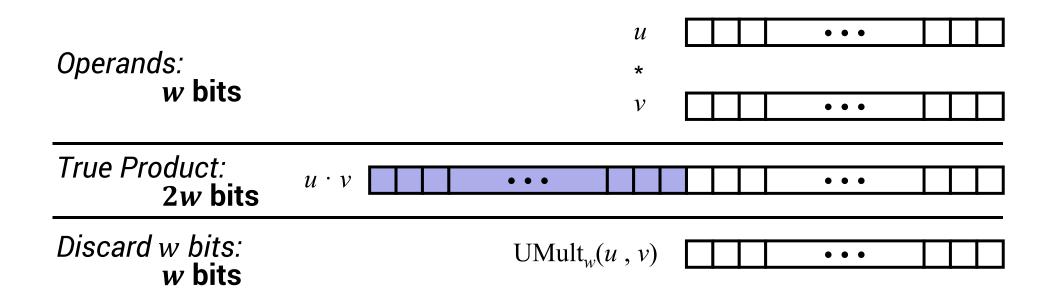


Administrivia

- Lab 1a due tonight at 11:59 pm
 - Submit pointer.c and lab1Areflect.txt
- Lab 1b due Friday (10/8)
 - Submit bits.c and lab1Breflect.txt
- Homework 2 released today, due 10/19
 - On Integers, Floating Point, and x86-64

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Unsigned Multiplication in C

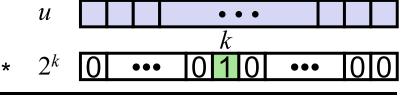


- Standard Multiplication Function
 - Ignores high order w bits
- Implements Modular Arithmetic
 - UMult_w $(u, v) = u \cdot v \mod 2^w$

Multiplication with shift and add

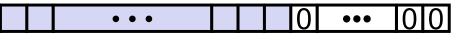
- ❖ Operation u<<k gives u*2^k
 - Both signed and unsigned

Operands: w bits



True Product: w + k bits

 $u\cdot 2^k$



Discard k bits: w bits

$$\mathrm{UMult}_{w}(u, 2^{k})$$



Examples:

u<<3

- u << 5 u << 3 == u * 24
- Most machines shift and add faster than multiply
 - Compiler generates this code automatically

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Number Representation Revisited

- What can we represent in one word?
 - Signed and Unsigned Integers
 - Characters (ASCII)
 - Addresses
- How do we encode the following:
 - Real numbers (e.g. 3.14159)
 - Very large numbers (e.g. 6.02×10²³)
 - Very small numbers (e.g. 6.626×10⁻³⁴)
 - Special numbers (e.g. ∞, NaN)



Floating Point Topics

- Fractional binary numbers
- IEEE floating-point standard
- Floating-point operations and rounding
- Floating-point in C

- There are many more details that we won't cover
 - It's a 58-page standard...

Representation of Fractions

"Binary Point," like decimal point, signifies boundary between integer and fractional parts:

L06: Floating Point I

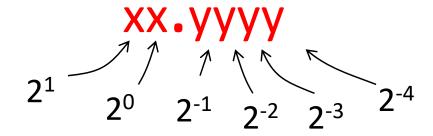
Example 6-bit representation:

* Example: $10.1010_2 = 1 \times 2^1 + 1 \times 2^{-1} + 1 \times 2^{-3} = 2.625_{10}$

Representation of Fractions

"Binary Point," like decimal point, signifies boundary between integer and fractional parts:

Example 6-bit representation:



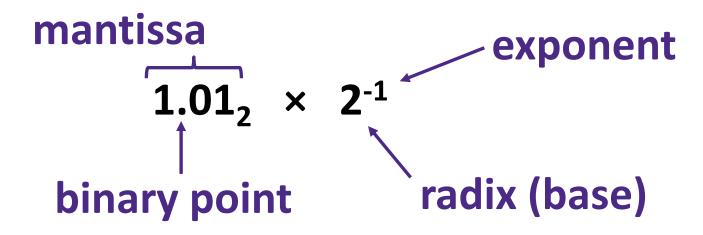
- In this 6-bit representation:
 - What is the encoding and value of the smallest (most negative) number?
 - What is the encoding and value of the largest (most positive) number?
 - What is the smallest number greater than 2 that we can represent?

Scientific Notation (Decimal)

mantissa exponent
$$6.02_{10} \times 10^{23}$$
 decimal point radix (base)

- Normalized form: exactly one digit (non-zero) to left of decimal point
- Alternatives to representing 1/1,000,000,000
 - Normalized: 1.0×10⁻⁹
 - Not normalized: $0.1 \times 10^{-8}, 10.0 \times 10^{-10}$

Scientific Notation (Binary)



- Computer arithmetic that supports this called floating point due to the "floating" of the binary point
 - Declare such variable in C as float (or double)

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Scientific Notation Translation

- Convert from scientific notation to binary point
 - Perform the multiplication by shifting the decimal until the exponent disappears
 - Example: $1.011_2 \times 2^4 = 10110_2 = 22_{10}$
 - Example: $1.011_2 \times 2^{-2} = 0.01011_2 = 0.34375_{10}$
- Convert from binary point to normalized scientific notation
 - Distribute out exponents until binary point is to the right of a single digit
 - Example: $1101.001_2 = 1.101001_2 \times 2^3$
- Practice: Convert 11.375₁₀ to binary scientific notation

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IEEE Floating Point

❖ IEEE 754

- Established in 1985 as uniform standard for floating point arithmetic
- Main idea: make numerically sensitive programs portable
- Specifies two things: representation and result of floating operations
- Now supported by all major CPUs

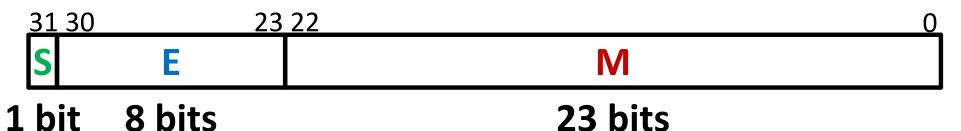
Driven by numerical concerns

- Scientists/numerical analysts want them to be as real as possible
- Engineers want them to be easy to implement and fast
- In the end:
 - Scientists mostly won out
 - Nice standards for rounding, overflow, underflow, but...
 - Hard to make fast in hardware
 - Float operations can be an order of magnitude slower than integer ops



Floating Point Encoding

- Use normalized, base 2 scientific notation:
 - Value: $\pm 1 \times Mantissa \times 2^{Exponent}$
 - Bit Fields: $(-1)^S \times 1.M \times 2^{(E-bias)}$
- Representation Scheme:
 - Sign bit (0 is positive, 1 is negative)
 - Mantissa (a.k.a. significand) is the fractional part of the number in normalized form and encoded in bit vector M
 - Exponent weights the value by a (possibly negative) power of 2 and encoded in the bit vector E



The Exponent Field

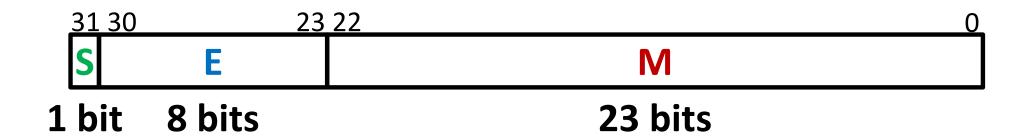
- Use biased notation
 - Read exponent as unsigned, but with *bias* of 2^{w-1}-1 = 127
 - Representable exponents roughly ½ positive and ½ negative
 - Exponent 0 (Exp = 0) is represented as E = 0b 0111 1111
- Why biased?
 - Makes floating point arithmetic easier
 - Makes somewhat compatible with two's complement
- Practice: To encode in biased notation, add the bias then encode in unsigned:

■
$$\mathsf{Exp} = 1 \rightarrow \mathsf{E} = \mathsf{Ob}$$

■
$$Exp = 127 \rightarrow E = 0b$$

Exp =
$$-63 \rightarrow E = 0b$$

The Mantissa (Fraction) Field



$$(-1)^{s} \times (1.M) \times 2^{(E-bias)}$$

- Note the implicit 1 in front of the M bit vector

 - Gives us an extra bit of precision
- Mantissa "limits"
 - Low values near M = 0b0...0 are close to 2^{Exp}
 - High values near M = 0b1...1 are close to 2^{Exp+1}

Peer Instruction Question

- What is the correct value encoded by the following floating point number?
 - 0b 0 10000000 110000000000000000000
 - Vote at http://PollEv.com/justinh

$$A. + 0.75$$

B.
$$+1.5$$

$$C. + 2.75$$

$$D. + 3.5$$

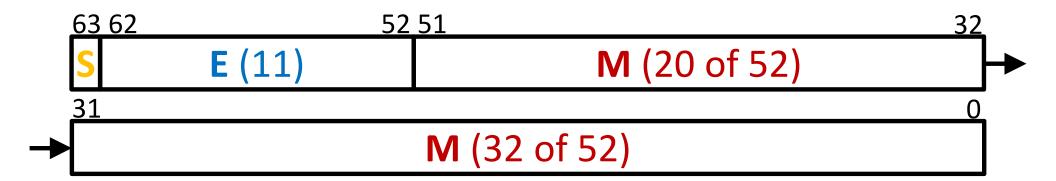
E. We're lost...

Precision and Accuracy

- Precision is a count of the number of bits in a computer word used to represent a value
 - Capacity for accuracy
- Accuracy is a measure of the difference between the actual value of a number and its computer representation
 - High precision permits high accuracy but doesn't guarantee it. It is possible to have high precision but low accuracy.
 - Example: float pi = 3.14;
 - pi will be represented using all 24 bits of the mantissa (highly precise), but is only an approximation (not accurate)

Need Greater Precision?

Double Precision (vs. Single Precision) in 64 bits



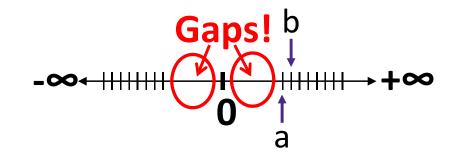
- C variable declared as double
- Exponent bias is now $2^{10}-1 = 1023$
- Advantages: greater precision (larger mantissa), greater range (larger exponent)
- Disadvantages: more bits used, slower to manipulate

Representing Very Small Numbers

- But wait... what happened to zero?
 - Using standard encoding 0x00000000 =
 - Special case: E and M all zeros = 0
 - Two zeros! But at least 0x00000000 = 0 like integers
- New numbers closest to 0:

$$a = 1.0...0_2 \times 2^{-126} = 2^{-126}$$

$$b = 1.0...01_2 \times 2^{-126} = 2^{-126} + 2^{-149}$$



- Normalization and implicit 1 are to blame
- Special case: E = 0, M ≠ 0 are denormalized numbers

Denorm Numbers

This is extra (non-testable) material

- Denormalized numbers
 - No leading 1
 - Uses implicit exponent of -126 even though E = 0x00
- Denormalized numbers close the gap between zero and the smallest normalized number
 - Smallest norm: $\pm 1.0...0_{two} \times 2^{-126} = \pm 2^{-126}$

So much closer to 0

- Smallest denorm: $\pm 0.0...01_{two} \times 2^{-126} = \pm 2^{-149}$
 - There is still a gap between zero and the smallest denormalized number

Other Special Cases

- \star E = 0xFF, M = 0: $\pm \infty$
 - *e.g.* division by 0
 - Still work in comparisons!
- \star E = 0xFF, M \neq 0: Not a Number (NaN)
 - e.g. square root of negative number, 0/0, $\infty-\infty$
 - NaN propagates through computations
 - Value of M can be useful in debugging
- New largest value (besides ∞)?
 - E = 0xFF has now been taken!
 - **E** = 0xFE has largest: $1.1...1_2 \times 2^{127} = 2^{128} 2^{104}$

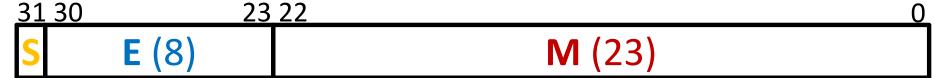


Floating Point Encoding Summary

| E | M | Meaning |
|-------------|----------|--------------|
| 0x00 | 0 | ± 0 |
| 0x00 | non-zero | ± denorm num |
| 0x01 – 0xFE | anything | ± norm num |
| OxFF | 0 | ± 8 |
| OxFF | non-zero | NaN |

Summary

Floating point approximates real numbers:



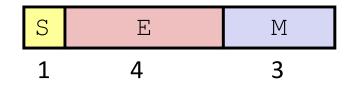
- Handles large numbers, small numbers, special numbers
- Exponent in biased notation (bias = 2^{w-1}-1)
 - Outside of representable exponents is overflow and underflow
- Mantissa approximates fractional portion of binary point
 - Implicit leading 1 (normalized) except in special cases
 - Exceeding length causes rounding

| E | M | Meaning | |
|-------------|----------|--------------|--|
| 0x00 | 0 | ± 0 | |
| 0x00 | non-zero | ± denorm num | |
| 0x01 – 0xFE | anything | ± norm num | |
| 0xFF | 0 | ± ∞ | |
| 0xFF | non-zero | NaN | |

BONUS SLIDES

An example that applies the IEEE Floating Point concepts to a smaller (8-bit) representation scheme. These slides expand on material covered today, so while you don't need to read these, the information is "fair game."

Tiny Floating Point Example



- 8-bit Floating Point Representation
 - The sign bit is in the most significant bit (MSB)
 - The next four bits are the exponent, with a bias of $2^{4-1}-1=7$
 - The last three bits are the mantissa

- Same general form as IEEE Format
 - Normalized binary scientific point notation
 - Similar special cases for 0, denormalized numbers, NaN, ∞

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Dynamic Range (Positive Only)

| | SE | M | Exp | Value | |
|--------------|--------|-------|----------------|---|--------------------|
| | 0 0000 | | -6 -6 | 0 1/8*1/64 = 1/512 | closost to zoro |
| Denormalized | 0 0000 | | - 6 | $\frac{1}{6} \frac{1}{64} = \frac{1}{512}$ $\frac{2}{8} \frac{1}{64} = \frac{2}{512}$ | Closest to Zero |
| numbers | 0 0000 | | -6 | 6/8*1/64 = 6/512 | |
| | 0 0000 |) 111 | - 6 | 7/8*1/64 = 7/512 | O |
| | 0 0001 | 000 | - 6 | 8/8*1/64 = 8/512 | smallest norm |
| | 0 0001 | 001 | - 6 | 9/8*1/64 = 9/512 | |
| | ••• | | | | |
| NI II I | 0 0110 | 110 | -1 | 14/8*1/2 = 14/16 | |
| | 0 0110 | 111 | -1 | 15/8*1/2 = 15/16 | closest to 1 below |
| Normalized | 0 0111 | 000 | 0 | 8/8*1 = 1 | |
| numbers | 0 0111 | 001 | 0 | 9/8*1 = 9/8 | closest to 1 above |
| | 0 0111 | 010 | 0 | 10/8*1 = 10/8 | |
| | ••• | | | | |
| | 0 1110 | 110 | 7 | 14/8*128 = 224 | |
| | 0 1110 |) 111 | 7 | 15/8*128 = 240 | largest norm |
| | 0 1111 | 000 | n/a | inf | |

Special Properties of Encoding

- Floating point zero (0+) exactly the same bits as integer zero
 - All bits = 0
- Can (Almost) Use Unsigned Integer Comparison
 - Must first compare sign bits
 - Must consider $0^- = 0^+ = 0$
 - NaNs problematic
 - Will be greater than any other values
 - What should comparison yield?
 - Otherwise OK
 - Denorm vs. normalized
 - Normalized vs. infinity