

Memory, Data, & Addressing II

CSE 351 Autumn 2018

Instructor:

Justin Hsia

Teaching Assistants:

Akshat Aggarwal

An Wang

Andrew Hu

Brian Dai

Britt Henderson

James Shin

Kevin Bi

Kory Watson

Riley Germundson

Sophie Tian

Teagan Horkan



<http://xkcd.com/138/>

Administrivia

- ❖ Lab 0 due today @ 11:59 pm
 - *You will be revisiting this program throughout this class!*
- ❖ Homework 1 due Wednesday
 - Reminder: autograded, 20 tries, no late submissions
- ❖ Lab 1a released today, due next Monday (10/8)
 - Pointers in C
 - Reminder: last submission graded, *individual* work

Late Days

- ❖ All submissions due at 11:59 pm
 - Count lateness in *days* (even if just by a second)
 - Special: weekends count as *one day*
 - No submissions accepted more than two days late
- ❖ You are given **5 late day tokens** for the whole quarter
 - Tokens can only apply to Labs (not HW)
 - No benefit to having leftover tokens
- ❖ Late penalty is 20% deduction of your score per day
 - Only late labs are eligible for penalties
 - Penalties applied at end of quarter to *maximize* your grade
- ❖ Use at own risk – don't want to fall too far behind
 - Intended to allow for unexpected circumstances

Review Questions

- 1) If the word size of a machine is 64-bits, which of the following is usually true? (pick all that apply)
 - a) 64 bits is the size of a pointer T
 - b) 64 bits is the size of an integer F (32 bits = 4 bytes)
 - c) 64 bits is the width of a register T
- 2) (True/False) By looking at the bits stored in memory, I can tell if a particular 4-bytes is being used to represent an integer, floating point number, or instruction.
- 3) If the size of a pointer on a machine is 6 bits, the address space is how many bytes? Ob _____
represent 2^6 things $\rightarrow 2^6$ addresses \rightarrow

2^6 bytes of data
64 B

4

Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
- ❖ Organizing and addressing data in memory
- ❖ **Manipulating data in memory using C**
- ❖ Boolean algebra and bit-level manipulations

Addresses and Pointers in C

- ❖ `&` = “address of” operator
- ❖ `*` = “value at address” or “dereference” operator

`*` is also used with variable declarations

datatype
+
`int *ptr;`
`int *ptr;`

Declares a variable, `ptr`, that is a pointer to (i.e. holds the address of) an `int` in memory

`int x = 5;`
`int y = 2;`

Declares two variables, `x` and `y`, that hold `ints`, and initializes them to 5 and 2, respectively

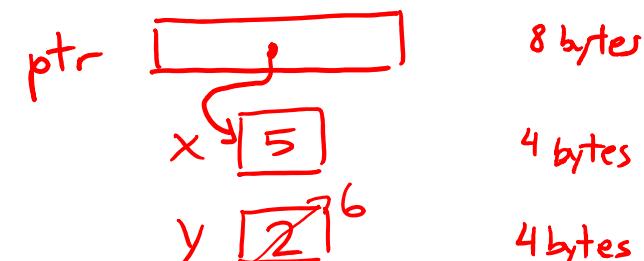
`ptr = &x;`

Sets `ptr` to the address of `x` (“`ptr` points to `x`”)

`y = 1 + *ptr;`

“Dereference `ptr`”

What is `*(&y)` ?



Sets `y` to “1 plus the value stored at the address held by `ptr`. Because `ptr` points to `x`, this is equivalent to `y=1+x`;

`returns value stored in y (equivalent to just using y)`

Assignment in C

- ❖ A variable is represented by a location
- ❖ Declaration ≠ initialization (initially holds “garbage”)
- ❖ **int x, y;**
 - x is at address 0x04, y is at 0x18

0x00	0x01	0x02	0x03
A7	00	32	00
00	01	29	F3
EE	EE	EE	EE
FA	CE	CA	FE
26	00	00	00
00	00	10	00
01	00	00	00
FF	00	F4	96
DE	AD	BE	EF
00	00	00	00

Current state
of memory

Assignment in C

32-bit example
(pointers are 32-bits wide)

little-endian

- ❖ A variable is represented by a location
- ❖ Declaration ≠ initialization (initially holds “garbage”)
- ❖ **int x, y;**
 - x is at address 0x04, y is at 0x18

	0x00	0x01	0x02	0x03
0x00				
0x04	00	01	29	F3
0x08				
0x0C				
0x10				
0x14				
0x18	01	00	00	00
0x1C				
0x20				
0x24				

X Y

Assignment in C

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location
- ❖ **int** x, y;
- ❖ x = 0; *int (4 bytes)*

32-bit example
(pointers are 32-bits wide)

& = “address of”

* = “dereference”

	0x00	0x01	0x02	0x03	
0x00					x
0x04	00	00	00	00	
0x08					
0x0C					
0x10					
0x14					
0x18	01	00	00	00	y
0x1C					
0x20					
0x24					

Assignment in C

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location
- ❖ **int** x, y;
- ❖ x = 0;
- ❖ y = 0x3CD02700; least significant byte little endian!

32-bit example
(pointers are 32-bits wide)

`&` = “address of”

`*` = “dereference”

	0x00	0x01	0x02	0x03
0x00				
0x04	00	00	00	00
0x08				
0x0C				
0x10				
0x14				
0x18	00	27	D0	3C
0x1C				
0x20				
0x24				

X Y

Assignment in C

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location
- ❖ **int** x, y;
- ❖ x = 0;
- ❖ y = 0x3CD02700;
- ❖ x = y + 3;
 - Get value at y, add 3, store in x

32-bit example
(pointers are 32-bits wide)

& = “address of”

* = “dereference”

0x00	0x01	0x02	0x03	
03	27	D0	3C	x
0x00	0x08	0x0C	0x10	
0x04				
0x08				
0x0C				
0x10				
0x14				
0x18	00	27	D0	3C
0x1C				
0x20				
0x24				

Assignment in C

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location
- ❖ **int** x, y;
- ❖ x = 0;
- ❖ y = 0x3CD02700;
- ❖ x = y + 3;
 - Get value at y, add 3, store in x
- ❖ **int*** z;
 - z is at address 0x20

32-bit example
(pointers are 32-bits wide)

& = “address of”

* = “dereference”

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	X
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	Y
0x1C					
0x20	DE	AD	BE	EF	Z
0x24					

initial value is whatever bits were already there! (“garbage”)

Assignment in C

- ❖ left-hand side = right-hand side;
 - LHS must evaluate to a *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location
- ❖ **int** x, y;
- ❖ x = 0;
- ❖ y = 0x3CD02700;
- ❖ x = y + 3;
 - Get value at y, add 3, store in x
- ❖ **int*** z = ^{0x18} &y + 3; // expect 0xb
 - Get address of y, “add 3”, store in z

32-bit example
(pointers are 32-bits wide)

& = “address of”

* = “dereference”

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	X
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	Y
0x1C					
0x20	24	00	00	00	Z
0x24					

get this instead

Pointer arithmetic

Pointer Arithmetic

- ❖ Pointer arithmetic is scaled by the size of target type
 - In this example, `sizeof (int) = 4`
- ❖ `int* z = &y + 3;`
 - Get address of `y`, add $3 * \text{sizeof}(\text{int})$, store in `z`
 - $\&y = 0x18 = 1 * 16^1 + 8 * 16^0 = 24$
 - $24 + 3 * (4) = 36 = 2 * 16^1 + 4 * 16^0 = 0x24$
- ❖ Pointer arithmetic can be dangerous!
 - Can easily lead to bad memory accesses
 - Be careful with data types and *casting*

Assignment in C

- ❖ **int** x, y;
- ❖ x = 0;
- ❖ y = 0x3CD02700;
- ❖ x = y + 3;
 - Get value at y, add 3, store in x
- ❖ **int*** z = &y + 3;
 - Get address of y, add 12, store in z
- ❖ *z = y;
 - What does this do?

32-bit example
(pointers are 32-bits wide)

& = “address of”

* = “dereference”

	0x00	0x01	0x02	0x03	
0x00					x
0x04	03	27	D0	3C	
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	RHS y
0x1C					
0x20	24	00	00	00	z
0x24					LHS

Assignment in C

- ❖ **int** `x, y;`
- ❖ `x = 0;`
- ❖ `y = 0x3CD02700;`
- ❖ `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- ❖ **int*** `z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`

The target of a pointer
is also a location
- ❖ ***z = y;**
 - Get value of `y`, put in address stored in `z`

32-bit example
(pointers are 32-bits wide)

`&` = “address of”

`*` = “dereference”

	0x00	0x01	0x02	0x03	
0x00					
0x04	03	27	D0	3C	X
0x08					
0x0C					
0x10					
0x14					
0x18	00	27	D0	3C	Y
0x1C					
0x20	24	00	00	00	Z
0x24	00	27	D0	3C	

Arrays in C

Declaration: `int a [6] ;` // &a is 0x10

element type

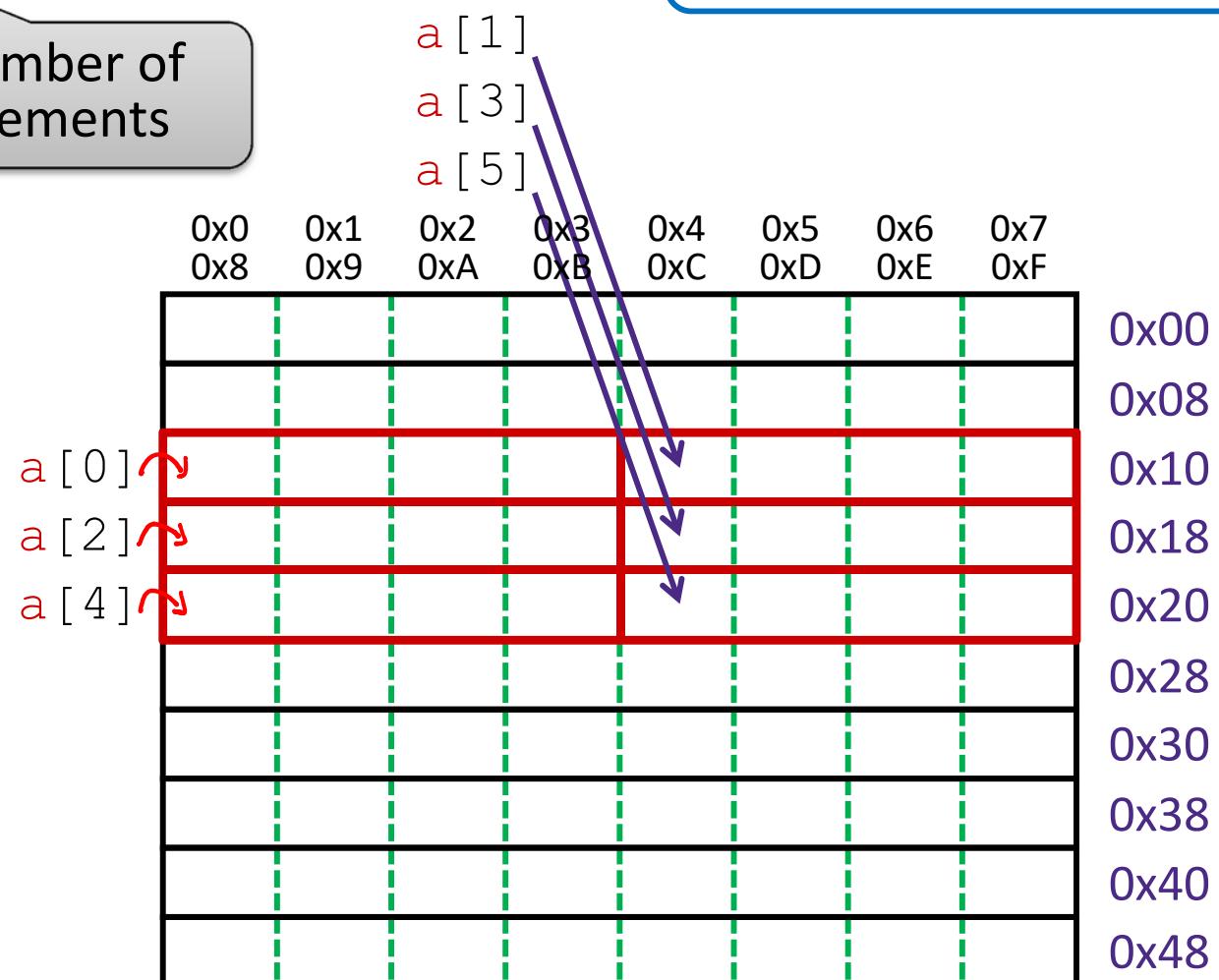
name

4 bytes each

Arrays are adjacent locations in memory
storing the same type of data object

a (array name) returns the array's address

64-bit example
(pointers are 64-bits wide)



Arrays in C

Declaration: **int** **a** [6] ;

Indexing: **a** [0] = 0x015f ;
 a [5] = **a** [0] ;

Arrays are adjacent locations in memory
storing the same type of data object

a (array name) returns the array's address

&a [i] is the address of **a** [0] plus **i** times
the element size in bytes

0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	0x00
								0x00
								0x08
								0x10
a [0]	5F	01	00	00				0x18
a [2]								0x20
a [4]				5F	01	00	00	0x28
								0x30
								0x38
								0x40
								0x48

Arrays in C

Declaration: `int a [6] ;`

Indexing: `a [0] = 0x015f ;`
`a [5] = a [0] ;`

No bounds
checking: `a [6] = 0xBAD ;`
`a [-1] = 0xBAD ;`

Arrays are adjacent locations in memory
storing the same type of data object

`a` (array name) returns the array's address

`&a [i]` is the address of `a [0]` plus `i` times
the element size in bytes

The diagram illustrates an array `a` in memory. The array has 8 elements, indexed from 0 to 7. The memory addresses for each element are shown above the array, and the corresponding byte values are shown in the cells. A red arrow points from the label "a[6]" to the value "AD" at address 0x28. Another red arrow points from the label "a[-1]" to the value "AD" at address 0x40. The array elements are:

Address	Value								
0x00	00	0x08	00	0x0F	00	0x10	00	0x18	00
0x01	00	0x09	00	0x0A	00	0x0B	00	0x0C	00
0x02	00	0x0D	00	0x0E	00	0x0F	00	0x10	00
0x03	00	0x11	00	0x12	00	0x13	00	0x14	00
0x04	00	0x15	00	0x16	00	0x17	00	0x18	00
0x05	00	0x19	00	0x1A	00	0x1B	00	0x1C	00
0x06	00	0x1D	00	0x1E	00	0x1F	00	0x20	00
0x07	00	0x21	00	0x22	00	0x23	00	0x24	00
0x08	00	0x25	00	0x26	00	0x27	00	0x28	AD
0x09	00	0x29	00	0x2A	00	0x2B	00	0x2C	00
0x0A	00	0x2D	00	0x2E	00	0x2F	00	0x30	00
0x0B	00	0x31	00	0x32	00	0x33	00	0x34	00
0x0C	00	0x35	00	0x36	00	0x37	00	0x38	00
0x0D	00	0x39	00	0x3A	00	0x3B	00	0x3C	00
0x0E	00	0x3D	00	0x3E	00	0x3F	00	0x40	5F
0x0F	00	0x3F	00	0x3F	00	0x3F	00	0x41	01
0x10	00	0x3F	00	0x3F	00	0x3F	00	0x42	00
0x11	00	0x3F	00	0x3F	00	0x3F	00	0x43	00
0x12	00	0x3F	00	0x3F	00	0x3F	00	0x44	00
0x13	00	0x3F	00	0x3F	00	0x3F	00	0x45	00
0x14	00	0x3F	00	0x3F	00	0x3F	00	0x46	00
0x15	00	0x3F	00	0x3F	00	0x3F	00	0x47	00
0x16	00	0x3F	00	0x3F	00	0x3F	00	0x48	00
0x17	00	0x3F	00	0x3F	00	0x3F	00		
0x18	00	0x3F	00	0x3F	00	0x3F	00		
0x19	00	0x3F	00	0x3F	00	0x3F	00		
0x1A	00	0x3F	00	0x3F	00	0x3F	00		
0x1B	00	0x3F	00	0x3F	00	0x3F	00		
0x1C	00	0x3F	00	0x3F	00	0x3F	00		
0x1D	00	0x3F	00	0x3F	00	0x3F	00		
0x1E	00	0x3F	00	0x3F	00	0x3F	00		
0x1F	00	0x3F	00	0x3F	00	0x3F	00		

Arrays in C

Declaration: **int** a[6];

Indexing: a [0] = 0x015f;
 a [5] = a [0];

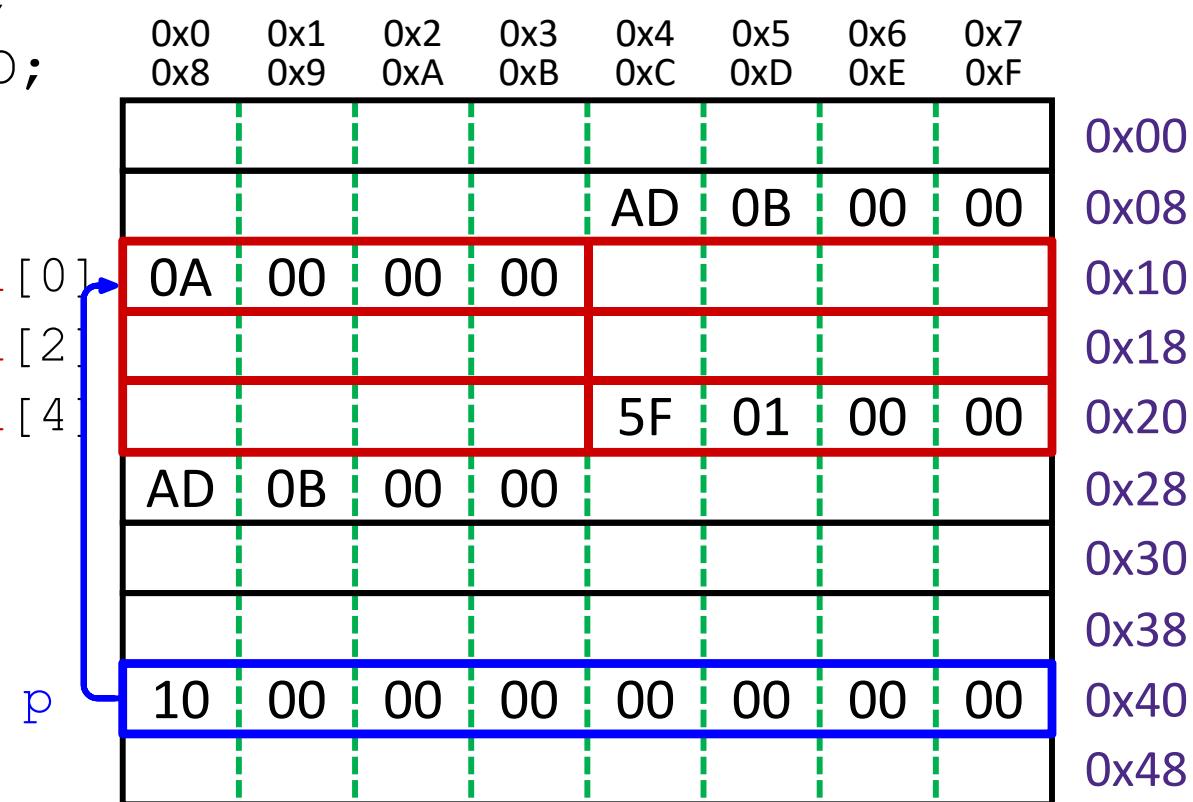
No bounds checking: `a[-1] = 0xBAD;`

Pointers: **int*** p;
equivalent {
 p = a;
 p = &a[0];
 *p = 0xA;

Arrays are adjacent locations in memory
storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes



Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`
 equivalent $\begin{cases} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{cases}$

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\begin{cases} p[1] = 0xB; \\ * (p+1) = 0xB; \end{cases}$
 pointer arithmetic: $0x10 + 1 \rightarrow 6 \times 4$
 $p = p + 2;$

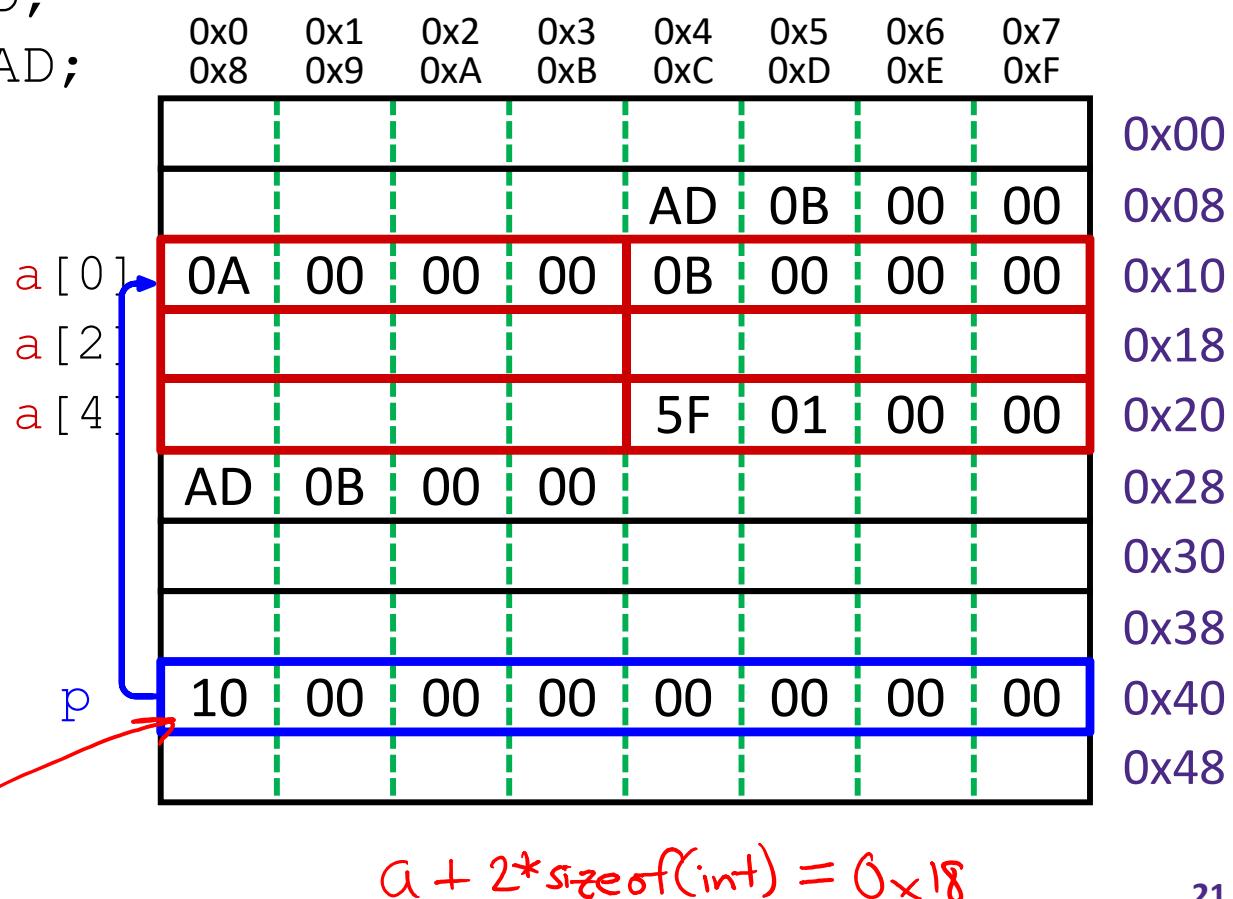
$0x10 + 2 \rightarrow 0x18$

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

`&a[i]` is the address of `a[0]` plus i times the element size in bytes

$$p[i] \iff *(p + i)$$



Arrays in C

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`
 equivalent `{ p = a;`
`p = &a[0];`
`*p = 0xA;`

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent `{ p[1] = 0xB;`
`* (p+1) = 0xB;`
`p = p + 2;`

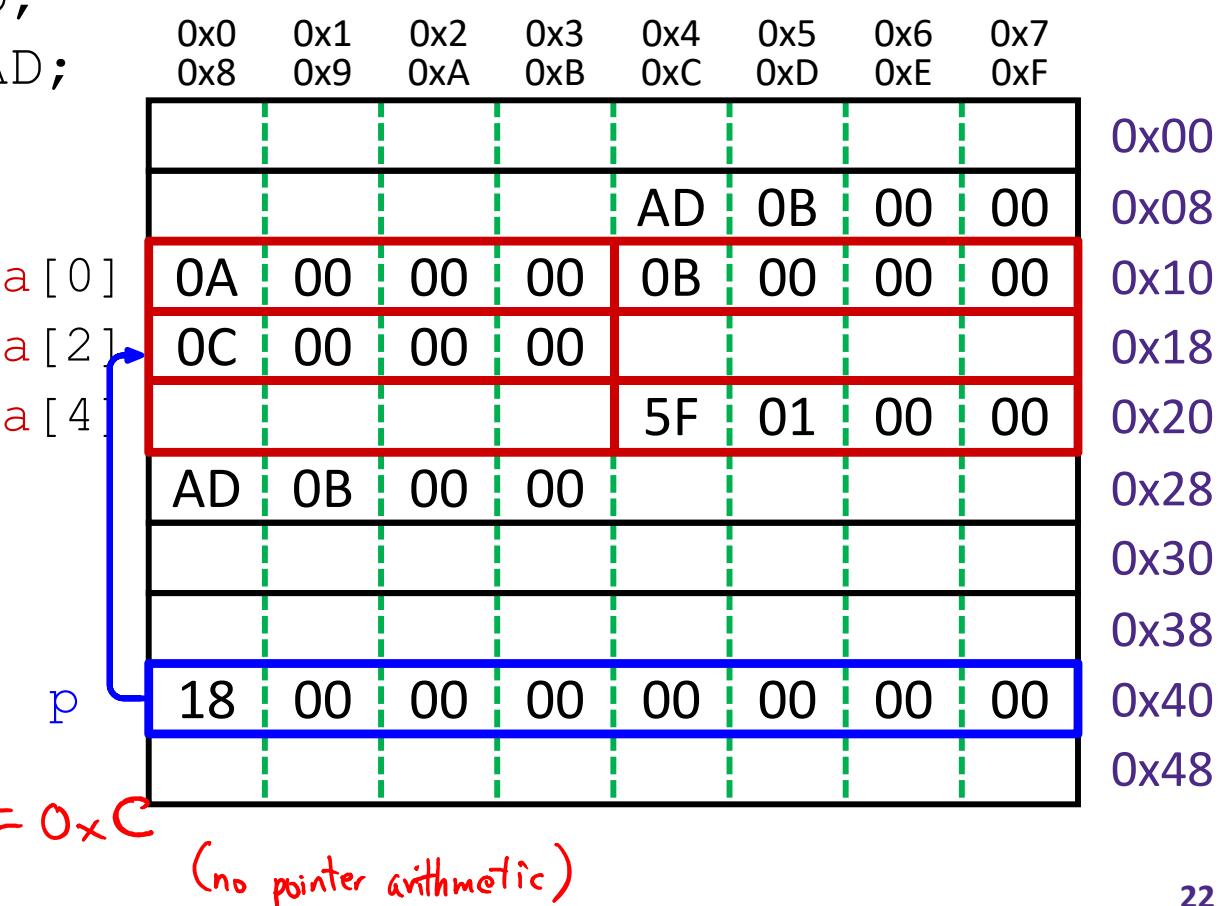
store at `0x18` →

$$*p = a[1] + 1; \quad 0xB + 1 = 0xC$$

Arrays are adjacent locations in memory storing the same type of data object

`a` (array name) returns the array's address

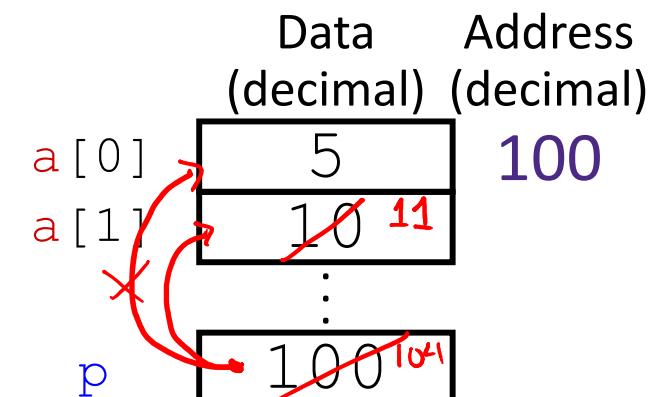
`&a[i]` is the address of `a[0]` plus `i` times the element size in bytes



Question: The variable values after Line 3 executes are shown on the right. What are they after Line 4 & 5?

- Vote at <http://PollEv.com/justinh>

```
1 void main () {  
2     int a [] = { 5, 10 } ;  
3     int* p = a ;  
4     p = p + 1 ; // sizeof(int) = 4  
5     *p = *p + 1 ;  
6 }
```



After Line 4
p *p a[0] a[1]

After Line 5
p *p a[0] a[1]

- (A) 101 10 5 10 then 101 11 5 11
- (B) 104 10 5 10 then 104 11 5 11**
- (C) 100 6 6 10 then 101 6 6 10
- (D) 100 6 6 10 then 104 6 6 10

Representing strings

- ❖ C-style string stored as an array of bytes (**char***)
 - Elements are one-byte **ASCII codes** for each character
 - No “String” keyword, unlike Java

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	"	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	,	55	7	71	G	87	W	103	g	119	w
40	(56	8	72	H	88	X	104	h	120	x
41)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

ASCII: American Standard Code for Information Interchange

Null-Terminated Strings

- ❖ Example: "Donald Trump" stored as a 13-byte array

Decimal:	68	111	110	97	108	100	32	84	114	117	109	112	0
Hex:	0x44	0x6F	0x6E	0x61	0x6C	0x64	0x20	0x54	0x72	0x75	0x6D	0x70	0x00
Text:	D	o	n	a	l	d		T	r	u	m	p	\0

13 bytes total!

- ❖ Last character followed by a 0 byte (' \0 ')
(a.k.a. "null terminator")
 - Must take into account when allocating space in memory
 - Note that '0' ≠ '\0' (i.e. character 0 has non-zero value)
- ❖ How do we compute the length of a string?
 - Traverse array until null terminator encountered

Endianness and Strings

C (char = 1 byte)

```
char s[6] = "12345";
```

String literal

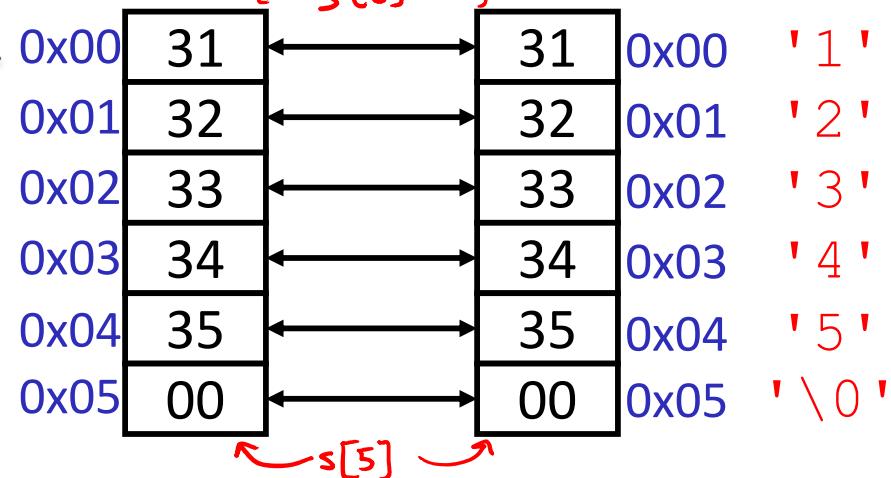
0x31 = 49 decimal = ASCII '1'

IA32, x86-64

(little-endian)

SPARC

(big-endian)



- ❖ Byte ordering (endianness) is not an issue for 1-byte values
 - The whole array does not constitute a single value
 - Individual elements are values; chars are single bytes

Examining Data Representations

- ❖ Code to print byte representation of data
 - Any data type can be treated as a *byte array* by **casting** it to `char`
 - C has **unchecked casts** **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

printf directives:

%p	Print pointer
\t	Tab
%x	Print value as hex
\n	New line

Examining Data Representations

- Code to print byte representation of data

- Any data type can be treated as a *byte array* by **casting** it to `char`
- C has **unchecked casts** **!! DANGER !!**

```
void show_bytes(char* start, int len) {
    int i;
    for (i = 0; i < len; i++)
        printf("%p\t0x%.2x\n", start+i, *(start+i));
    printf("\n");  format string
}
```

pointer arithmetic on
`char*`

```
void show_int(int x) {
    show_bytes( (char *) &x, sizeof(int));
}
```

`int*`
`4 bytes`
"cast"
(treat as)

show_bytes Execution Example

```
int x = 12345; // 0x00003039
printf("int x = %d;\n", x);
show_int(x); // show_bytes((char *) &x, sizeof(int));
```

- ❖ Result (Linux x86-64):

- **Note:** The addresses will change on each run (try it!), but fall in same general range

```
int x = 12345;
0x7fffb7f71dbc      0x39
0x7fffb7f71dbd      0x30
0x7fffb7f71dbe      0x00
0x7fffb7f71dbf      0x00
```

Summary

- ❖ Assignment in C results in value being put in memory location
- ❖ Pointer is a C representation of a data address
 - $\&$ = “address of” operator
 - $*$ = “value at address” or “dereference” operator
- ❖ Pointer arithmetic scales by size of target type
 - Convenient when accessing array-like structures in memory
 - Be careful when using – particularly when *casting* variables
- ❖ Arrays are adjacent locations in memory storing the same type of data object
 - Strings are null-terminated arrays of characters (ASCII)