

Java and C

CSE 351 Autumn 2017

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<https://xkcd.com/801/>

Administrivia

- ❖ Lab 5 due Friday (12/8)
 - Hard deadline on Sunday (12/10)
- ❖ Course evaluations now open
 - See Piazza post @366 for links (separate for Lec A/B)
- ❖ **Final Exam:** Wed, 12/13, 12:30-2:20pm in KNE 120
 - Review Session: Mon, 12/11, 5-8pm in EEB 105
 - You get TWO double-sided handwritten 8.5×11" cheat sheets
 - Additional practice problems on website

Roadmap

C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
c.setGals(17);  
float mpg =  
    c.getMPG();
```

Memory & data
Integers & floats
x86 assembly
Procedures & stacks
Executables
Arrays & structs
Memory & caches
Processes
Virtual memory
Memory allocation

Java vs. C

Assembly
language:

```
get_mpg:  
    pushq    %rbp  
    movq     %rsp, %rbp  
    ...  
    popq     %rbp  
    ret
```

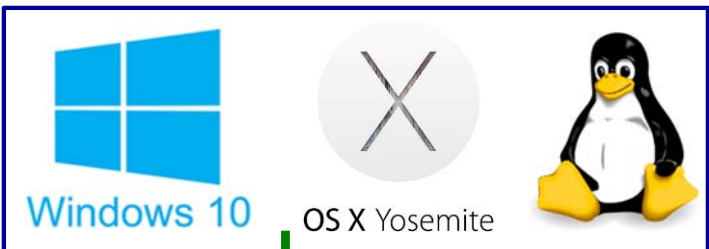
Machine
code:

```
0111010000011000  
100011010000010000000010  
1000100111000010  
110000011111101000011111
```

Computer
system:



OS:



Java vs. C

- ❖ Reconnecting to Java (hello CSE143!)
 - But now you know a lot more about what really happens when we execute programs
- ❖ We've learned about the following items in C; now we'll see what they look like for Java:
 - Representation of data
 - Pointers / references
 - Casting
 - Function / method calls including dynamic dispatch

Worlds Colliding

- ❖ CSE351 has given you a “really different feeling” about what computers do and how programs execute
- ❖ We have occasionally contrasted to Java, but CSE143 may still feel like “a different world”
 - It’s not – it’s just a higher-level of abstraction
 - Connect these levels via how-one-could-implement-Java in 351 terms

Meta-point to this lecture

- ❖ None of the data representations we are going to talk about are guaranteed by Java
- ❖ In fact, the language simply provides an abstraction (Java language specification)
 - Tells us how code should behave for different language constructs, but we can't easily tell how things are really represented
 - But it is important to understand an implementation of the lower levels – useful in thinking about your program

Data in Java

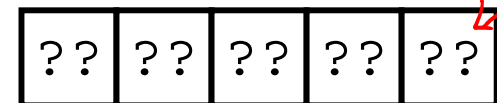
- ❖ Integers, floats, doubles, pointers – same as C
 - “Pointers” are called “references” in Java, but are much more constrained than C’s general pointers
 - Java’s portability-guarantee fixes the sizes of all types
 - Example: `int` is 4 bytes in Java regardless of machine
 - No unsigned types to avoid conversion pitfalls
 - Added some useful methods in Java 8 (also use bigger signed types)
- ❖ `null` is typically represented as 0 but “you can’t tell”
- ❖ Much more interesting:
 - **Arrays**
 - **Characters and strings**
 - **Objects**

Data in Java: Arrays

- ❖ Every element initialized to 0 or `null`
- ❖ Length specified in immutable field at start of array (`int` – 4 bytes)
 - `array.length` returns value of this field
- ❖ *Since it has this info, what can it do?*

C:

```
int array[5];
```

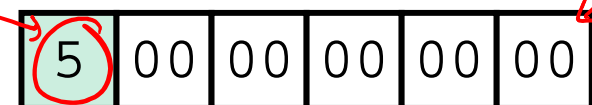


0 4 20

garbage!

Java:

```
int[] array = new int[5];
```



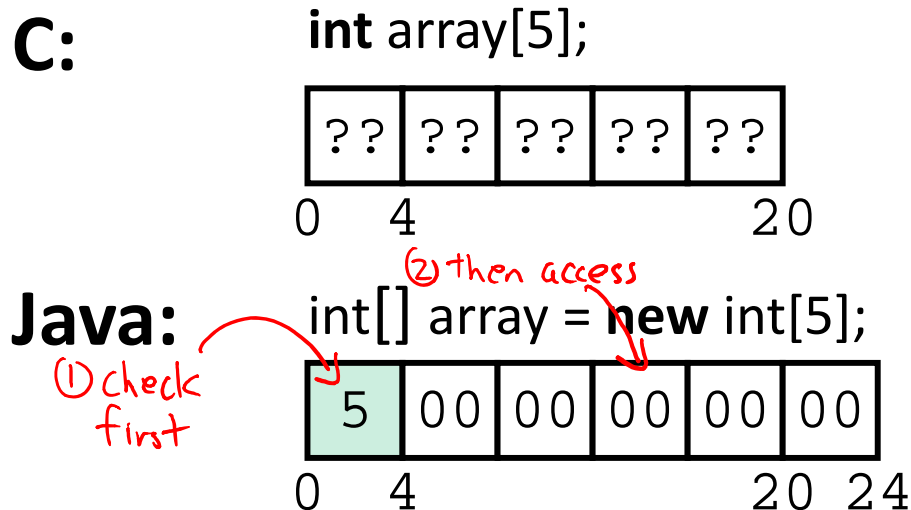
0 4 20 24

size/length
info only

guaranteed zeros

Data in Java: Arrays

- ❖ Every element initialized to 0 or `null`
- ❖ Length specified in immutable field at start of array (`int` – 4 bytes)
 - `array.length` returns value of this field
- ❖ Every access triggers a bounds-check
 - Code is added to ensure the index is within bounds
 - Exception if out-of-bounds



To speed up bounds-checking:

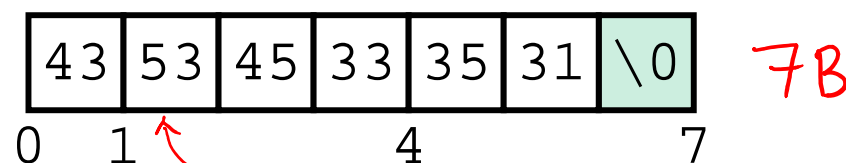
- Length field is likely in cache
- Compiler may store length field in register for loops
- Compiler may prove that some checks are redundant

Data in Java: Characters & Strings

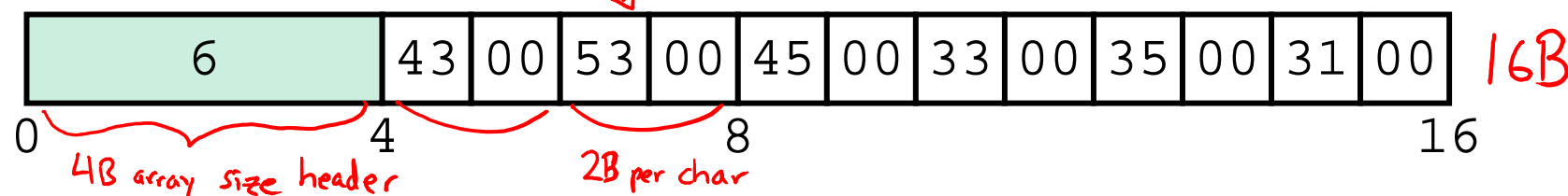
- ❖ Two-byte Unicode instead of ASCII
 - Represents most of the world's alphabets
- ❖ String not bounded by a '\0' (null character)
 - Bounded by hidden length field at beginning of string
- ❖ All String objects read-only (vs. StringBuffer)

Example: the string "CSE351"

C:
(ASCII)



Java:
(Unicode)



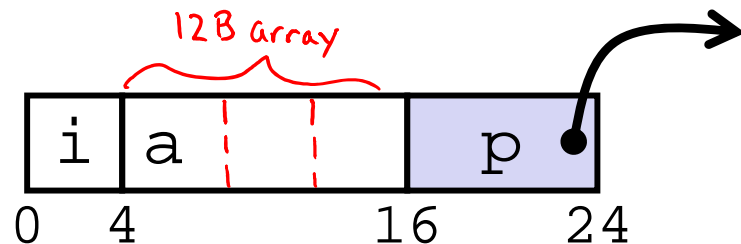
Data in Java: Objects

- ❖ Data structures (objects) are always stored by reference, never stored “inline”
 - Include complex data types (arrays, other objects, etc.) using references

C:

```
struct rec {
    int i;
    int a[3];
    struct rec *p;
};
```

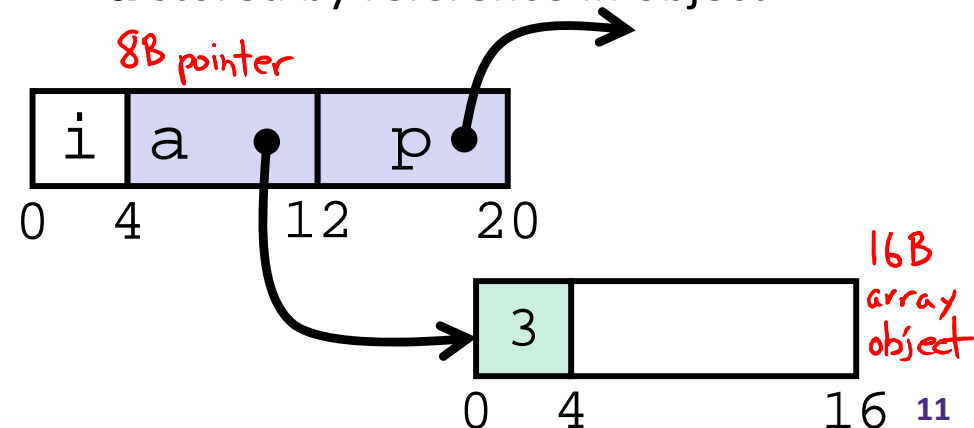
- a[] stored “inline” as part of struct



Java:

```
class Rec {
    int i;
    int[] a = new int[3];
    Rec p;
    ...
}
```

- a stored by reference in object



Pointer/reference fields and variables

- ❖ In C, we have “->” and “.” for field selection depending on whether we have a pointer to a struct or a struct
 - $(*r).a$ is so common it becomes $r->a$
- ❖ In Java, *all non-primitive variables are references to objects*
 - We always use $r.a$ notation
 - But really follow reference to r with offset to a , just like $r->a$ in C
 - So no Java field needs more than 8 bytes !

C:

```
struct rec *r = malloc(...);
struct rec r2;
r->i = val;
r->a[2] = val;
r->p = &r2;
```

Java:

```
r = new Rec();
r2 = new Rec();
r.i = val;
r.a[2] = val;
r.p = r2;
```

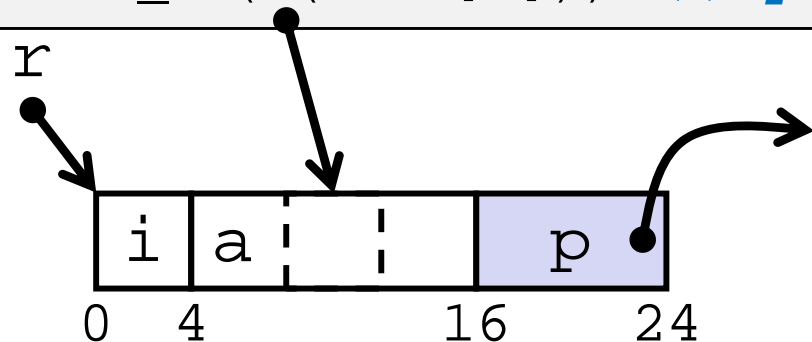
references

Pointers/References

- ❖ *Pointers* in C can point to any memory address
- ❖ *References* in Java can only point to [the starts of] objects
 - Can only be dereferenced to access a field or element of that object

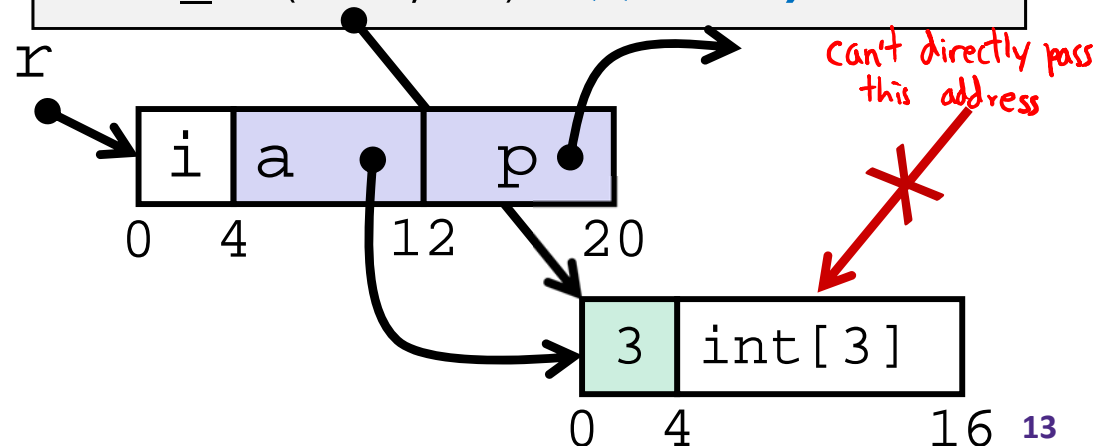
C:

```
struct rec {  
    int i;  
    int a[3];  
    struct rec *p;  
};  
struct rec* r = malloc(...);  
some_fn(&(r->a[1])); // ptr
```



Java:

```
class Rec {  
    int i;  
    int[] a = new int[3];  
    Rec p;  
}  
Rec r = new Rec();  
some_fn(r.a, 1); // ref, index
```



Casting in C (example from Lab 5)

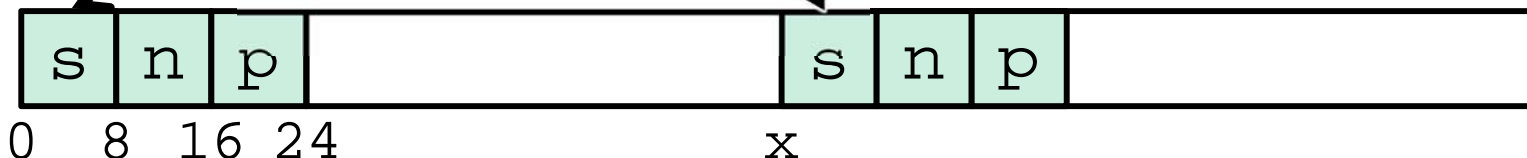
- ❖ Can cast any pointer into any other pointer
 - Changes dereference and arithmetic behavior

```
struct BlockInfo {  
    size_t sizeAndTags;  
    struct BlockInfo* next;  
    struct BlockInfo* prev;  
};  
typedef struct BlockInfo BlockInfo;  
...  
int x;  
BlockInfo *b;  
BlockInfo *newBlock;  
...  
newBlock = (BlockInfo *) ( (char *) b + x );  
...
```

Cast b into char * to
do unscaled addition

Cast back into
BlockInfo * to use
as BlockInfo struct

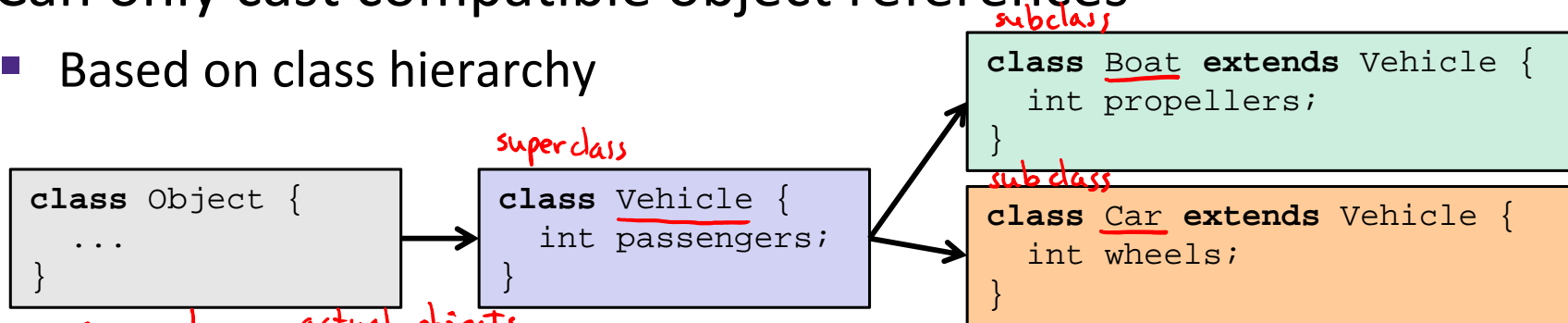
move by x bytes



Type-safe casting in Java

❖ Can only cast compatible object references

- Based on class hierarchy



references! *actual objects*

```
Vehicle v = new Vehicle(); // super class of Boat and Car
Boat b1 = new Boat(); // |--> sibling
Car c1 = new Car(); // |--> sibling
```

```
Vehicle v1 = new Car();
```

```
Vehicle v2 = v1;
```

```
Car c2 = new Boat();
```

```
Car c3 = new Vehicle();
```

```
Boat b2 = (Boat) v;
```

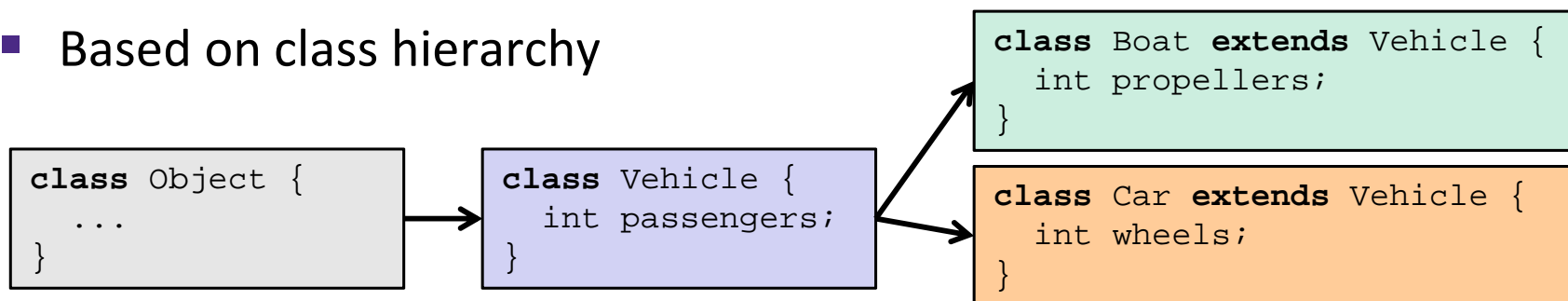
```
Car c4 = (Car) v2;
```

```
Car c5 = (Car) b1;
```

Type-safe casting in Java

❖ Can only cast compatible object references

■ Based on class hierarchy



```
Vehicle v = new Vehicle(); // super class of Boat and Car
```

```
Boat b1 = new Boat(); // |--> sibling
```

```
Car c1 = new Car(); // |--> sibling
```

```
Vehicle v1 = new Car();
```

← ✓ Everything needed for `Vehicle` also in `Car`

```
Vehicle v2 = v1;
```

← ✓ `v1` is declared as type `Vehicle`

```
Car c2 = new Boat();
```

← X Compiler error: Incompatible type – elements in `Car` that are not in `Boat` (siblings)

```
Car c3 = new Vehicle();
```

← X Compiler error: Wrong direction – elements `Car` not in `Vehicle` (wheels)

what if: `v = new Boat();`

```
Boat b2 = (Boat) v;
```

← X Runtime error: `Vehicle` does not contain all elements in `Boat` (propellers)

```
Car c4 = (Car) v2;
```

← ✓ `v2` refers to a `Car` at runtime

```
Car c5 = (Car) b1;
```

← X Compiler error: Unconvertable types – `b1` is declared as type `Boat`

Java Object Definitions

```
class Point {  
    double x;  
    double y;  
  
    Point ( ) {  
        x = 0;  
        y = 0;  
    }  
  
    boolean samePlace(Point p) {  
        return (x == p.x) && (y == p.y);  
    }  
}  
...  
Point p = new Point ( );  
...
```

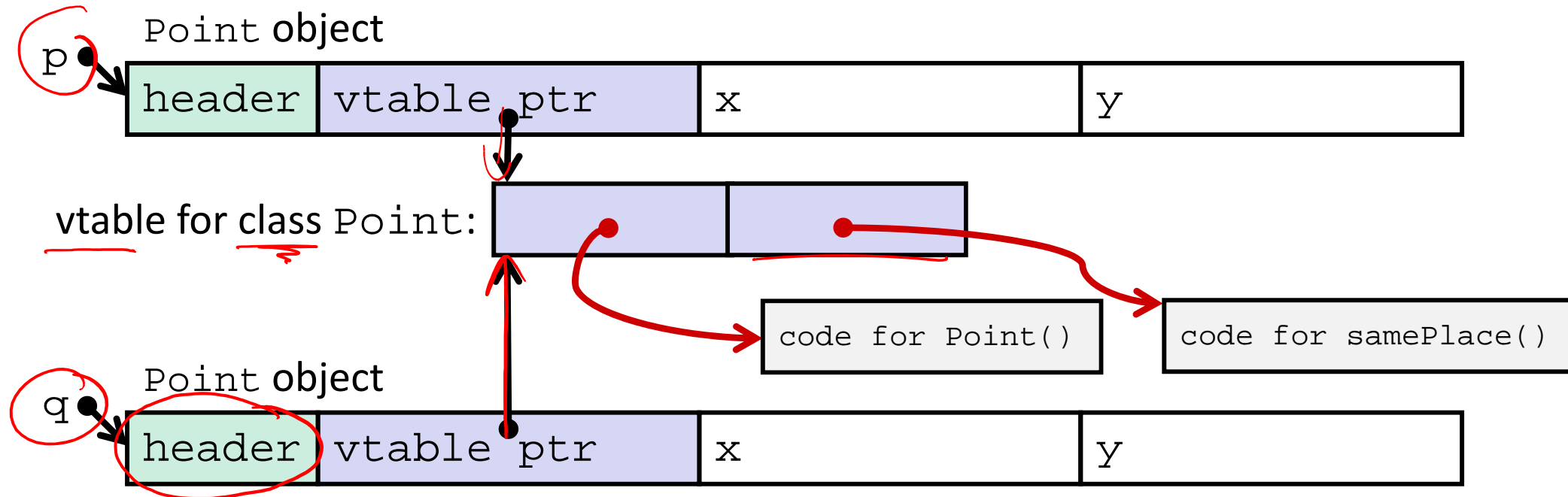
fields

constructor

method(s)

creation

Java Objects and Method Dispatch



❖ *Virtual method table (vtable)*

- Like a jump table for instance (“virtual”) methods plus other class info
- One table per class

❖ *Object header* : GC info, hashing info, lock info, etc.

- Why no size?

Java Constructors

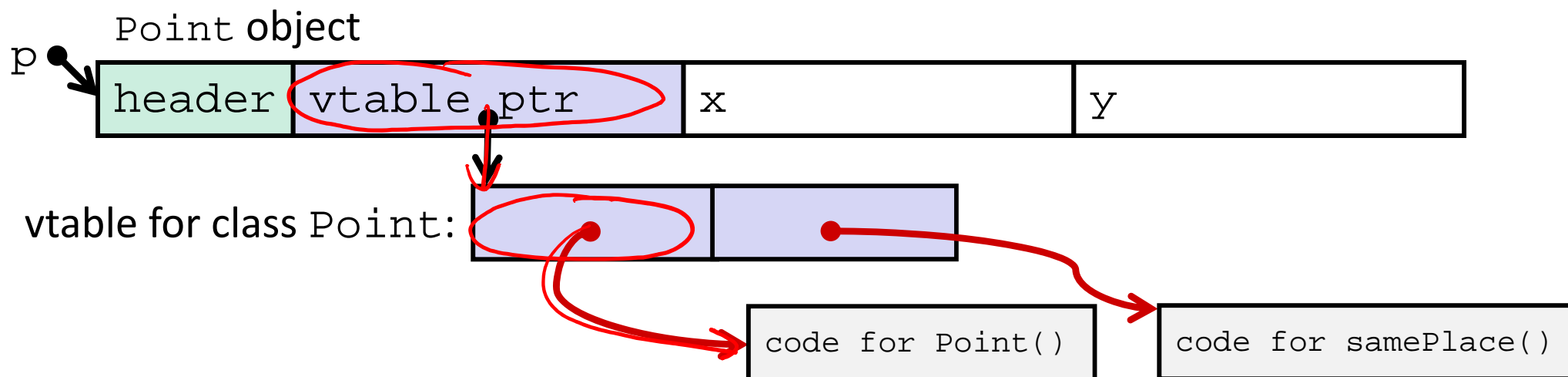
- ❖ When we call **new**: allocate space for object (data fields and references), initialize to zero/null, and run constructor method

Java:

```
Point p = new Point();
```

C pseudo-translation:

```
Point* p = calloc(1, sizeof(Point));  
p->header = ...; // set up header (somehow)  
p->vtable = &Point_vtable; } run the  
p->vtable[0](p);             } constructor
```



Java Methods

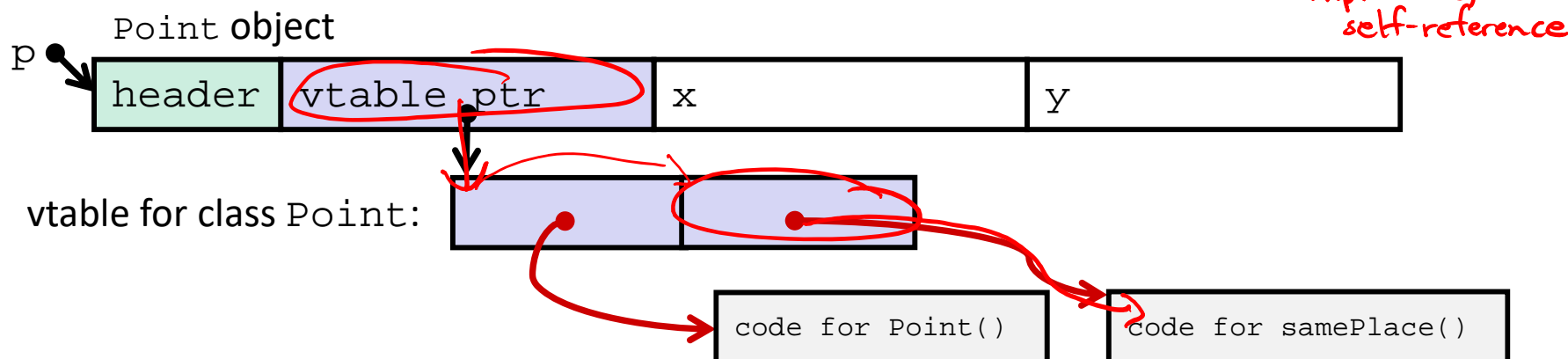
- ❖ Static methods are just like functions
- ❖ Instance methods:
 - Can refer to *this*; reference to particular instance of class
 - Have an implicit first parameter for *this*; and
 - Can be overridden in subclasses
- ❖ The code to run when calling an instance method is chosen *at runtime* by lookup in the vtable (i.e. dispatch)

Java:

```
p.samePlace(q);
```

C pseudo-translation:

```
p->vtable[1](p, q);
```



Subclassing

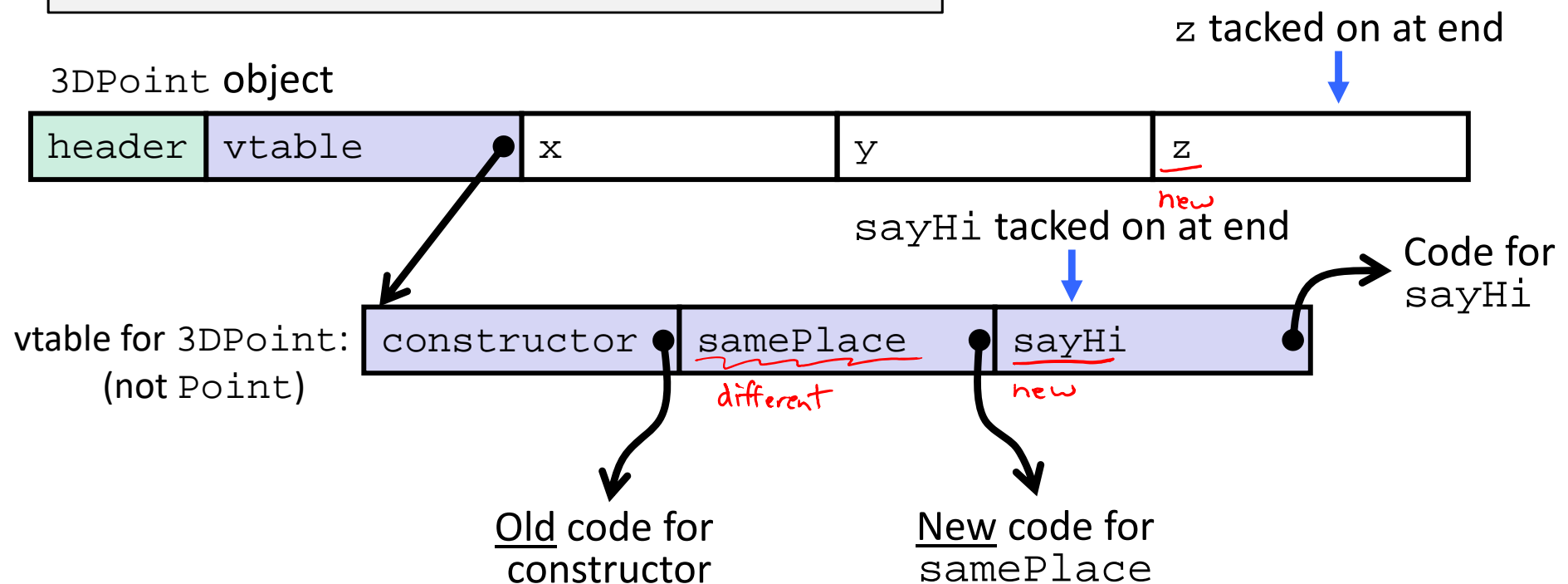
```
class 3DPoint extends Point {  
    double z;  
    boolean samePlace(Point p2) {  
        return false;  
    }  
    void sayHi() {  
        System.out.println("hello");  
    }  
}
```

← new field
} override method
} new method

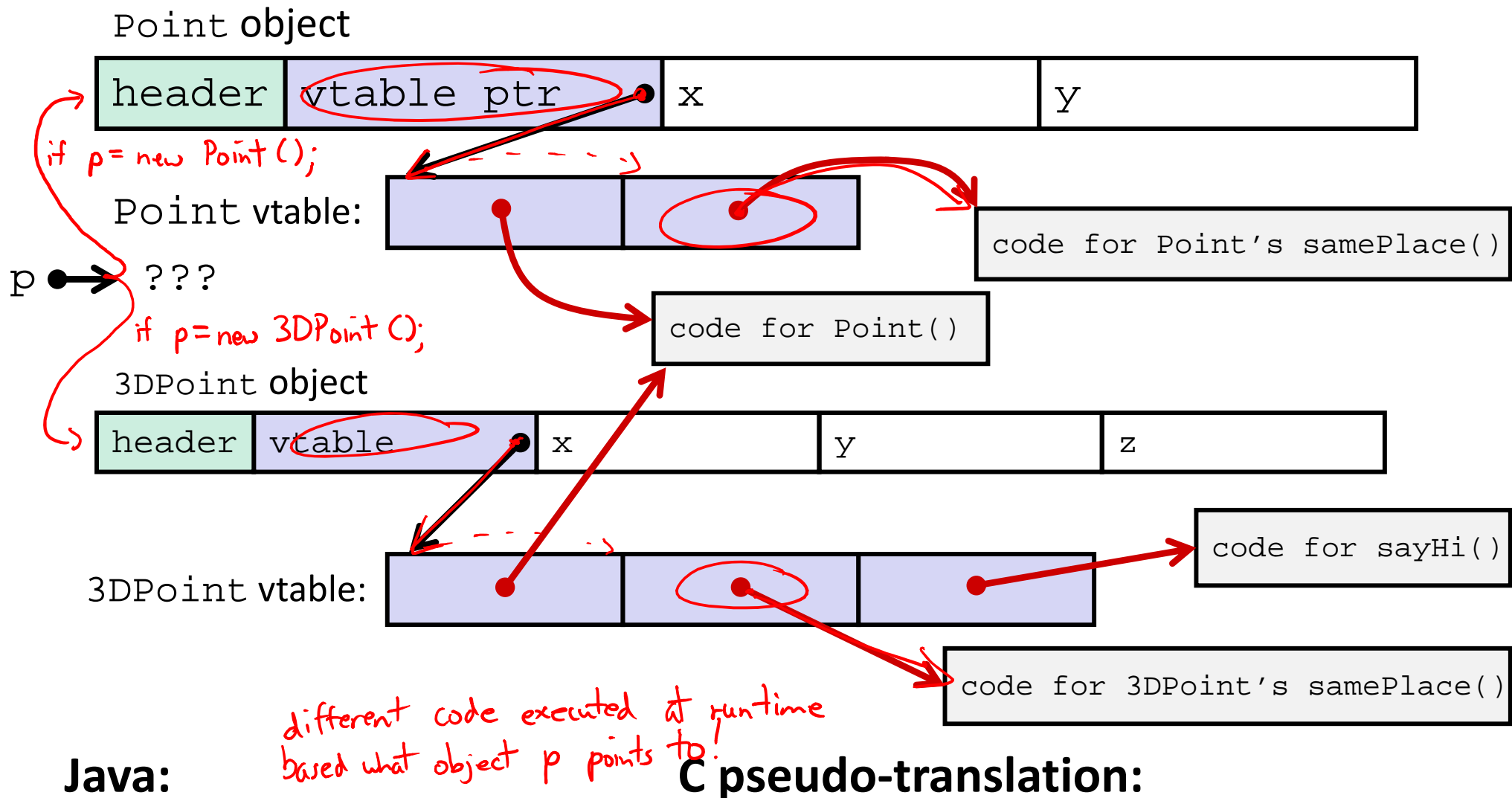
- ❖ Where does “z” go? At end of fields of `Point`
 - `Point` fields are always in the same place, so `Point` code can run on `3DPoint` objects without modification
- ❖ Where does pointer to code for two new methods go?
 - No constructor, so use default `Point` constructor
 - To override “`samePlace`”, use same vtable position
 - Add new pointer at end of vtable for new method “`sayHi`”

Subclassing

```
class 3DPoint extends Point {  
    double z;  
    boolean samePlace(Point p2) {  
        return false;  
    }  
    void sayHi() {  
        System.out.println("hello");  
    }  
}
```



Dynamic Dispatch



Java:

```
Point p = ???;
return p.samePlace(q);
```

C pseudo-translation:

```
// works regardless of what p is
return p->vtable[1](p, q);
```

Ta-da!

- ❖ In CSE143, it may have seemed “magic” that an *inherited* method could call an *overridden* method
 - You were tested on this endlessly
- ❖ The “trick” in the implementation is this part:
`p->vtable[i](p,q)`
 - In the body of the pointed-to code, any calls to (other) methods of `this` will use `p->vtable`
 - Dispatch determined by `p`, not the class that defined a method

Practice Question

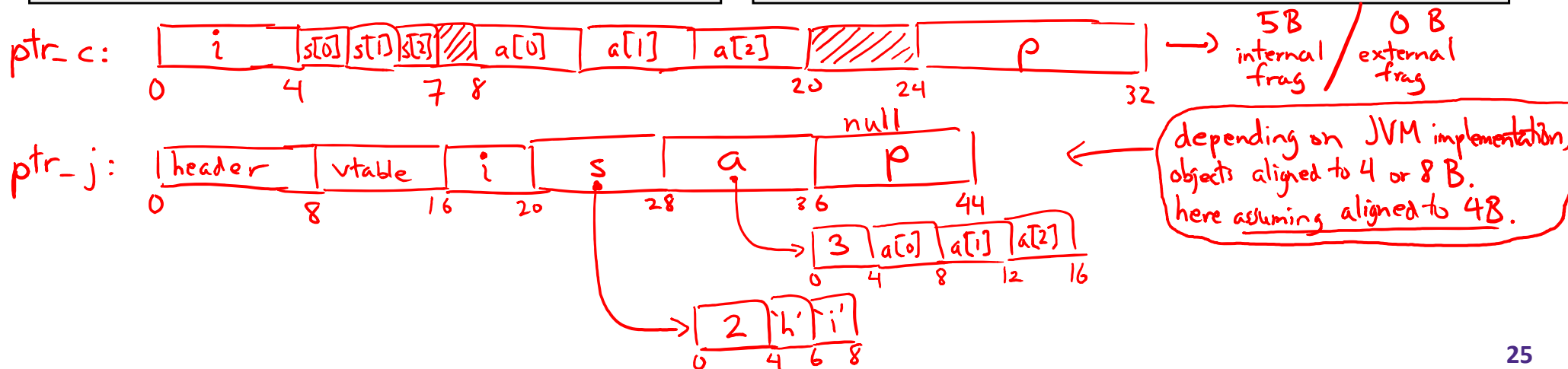
- ❖ Assume: 64-bit pointers and that a Java object header is 8 B
- ❖ What are the sizes of the things being pointed at by `ptr_c` (32 B) and `ptr_j`? (44 B)

```
struct c {
    int i;
    char s[3];
    int a[3];
    struct c *p;
};
struct c* ptr_c;
```

Handwritten notes:
 K = 4 (for int)
 1 (for char)
 4 (for int)
 8 (for struct c*)
 internal frag (for s and a)
 K_{max} = 8 (for struct c*)
 external frag (for p)

```
class jobj {
    int i;
    String s = "hi";
    int[] a = new int[3];
    jobj p;
}
jobj ptr_j = new jobj();
```

Handwritten note:
 no explicit methods, but still inherits constructor & methods from Object class



We made it! 😊 😎 😄

C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
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float mpg =  
    c.getMPG();
```

Assembly
language:

```
get_mpg:  
    pushq    %rbp  
    movq     %rsp, %rbp  
    ...  
    popq     %rbp  
    ret
```

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Java vs. C

OS:

