

# x86-64 Programming I

CSE 351 Autumn 2017

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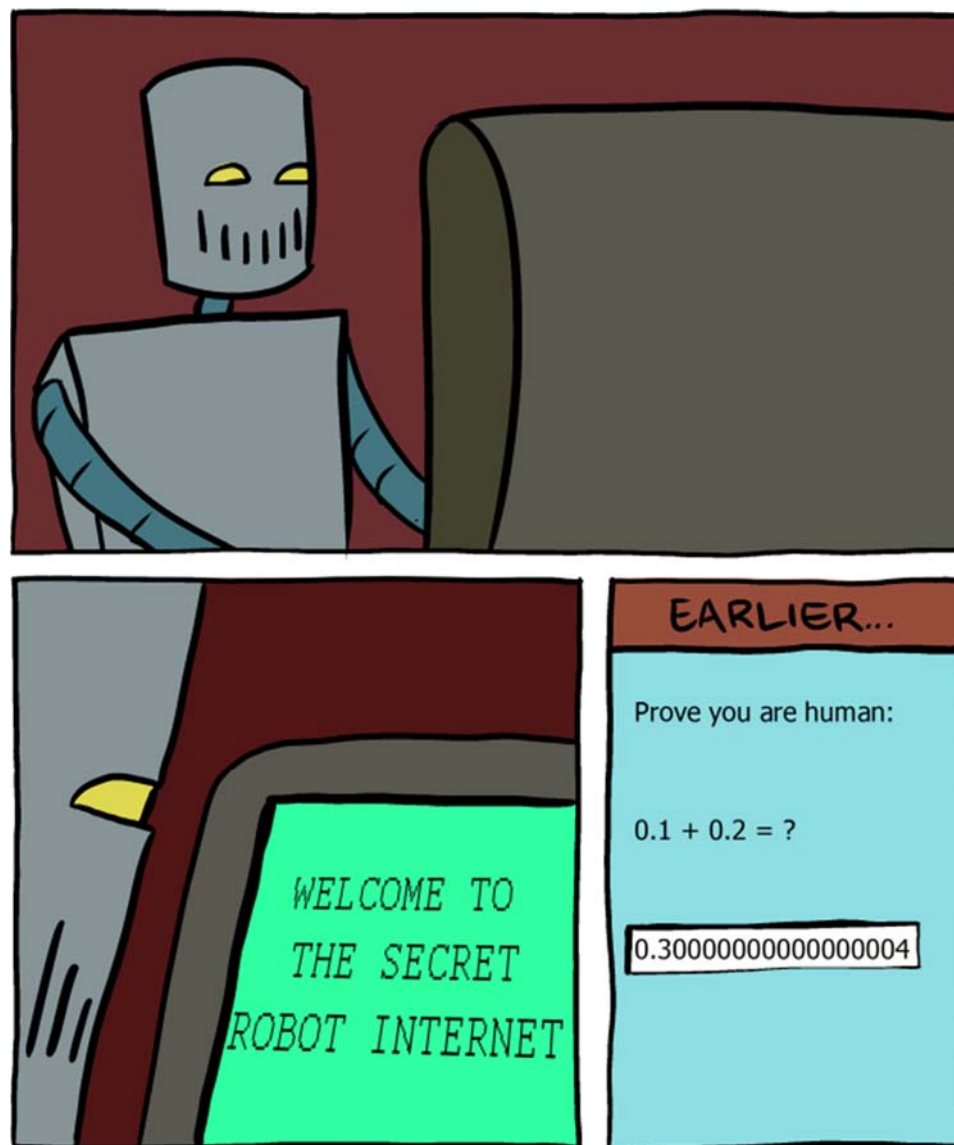
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<http://www.smbc-comics.com/?id=2999>

# Administrivia

- ❖ Lab 1 due tonight at 11:59pm
  - You have *late days* available
- ❖ Homework 2 due next Friday (10/20)
- ❖ Lab 2 (x86-64) released on Tuesday (10/17)
  - Due on 10/27

# Review: Operand types

❖ **Immediate:** Constant integer data

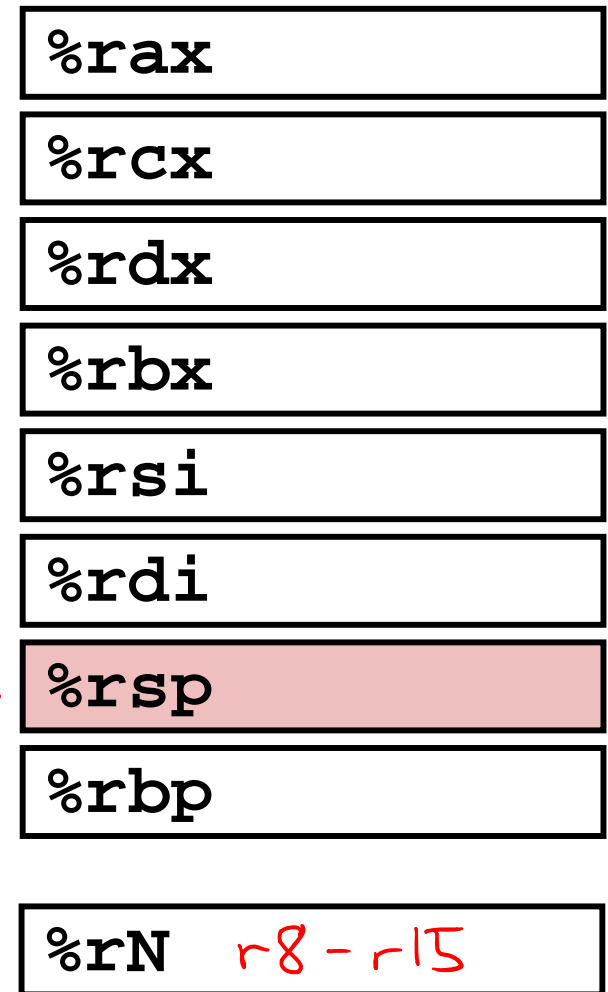
- Examples: \$0x400, \$-533  
hex decimal
- Like C literal, but prefixed with '\$'
- Encoded with 1, 2, 4, or 8 bytes  
*depending on the instruction*

❖ **Register:** 1 of 16 integer registers

- Examples: %rax, %r13
- But %rsp reserved for special use
- Others have special uses for particular instructions

❖ **Memory:** Consecutive bytes of memory at a computed address

- Simplest example: (%rax) ← take data in %rax, treat as address, pull data at that address
- Various other "address modes"



# Moving Data

- ❖ General form: `mov_ source, destination`
  - Missing letter (`_`) specifies size of operands
  - Note that due to backwards-compatible support for 8086 programs (16-bit machines!), “word” means 16 bits = 2 bytes in x86 instruction names
  - Lots of these in typical code
- ❖ `movb src, dst`
  - Move 1-byte “byte”
- ❖ `movw src, dst`
  - Move 2-byte “word”
- ❖ `movl src, dst`
  - Move 4-byte “long word”
- ❖ `movq src, dst`
  - Move 8-byte “quad word”

# movq Operand Combinations

x86                      C  
 Imm ↔ Constant  
 Reg ↔ Variable  
 Mem ↔ dereferencing  
**C Analog**    a pointer

	Source	Dest	Src, Dest	
movq	Imm	Reg	movq \$0x4, %rax	var_a = 0x4;
		Mem	movq \$-147, (%rax)	*p_a = -147;
	Reg	Reg	movq %rax, %rdx	var_d = var_a;
		Mem	movq %rax, (%rdx)	*p_d = var_a;
	Mem	Reg	movq (%rax), %rdx	var_d = *p_a;

❖ *Cannot do memory-memory transfer with a single instruction*

■ How would you do it?

① Mem → Reg

movq (%rax), %rdx

② Reg → Mem

movq %rdx, (%rbx)

# x86-64 Introduction

- ❖ Arithmetic operations
- ❖ Memory addressing modes
  - `swap` example
- ❖ Address computation instruction (`leaq`)

# Some Arithmetic Operations

## ❖ Binary (two-operand) Instructions:

**Maximum of one memory operand**

- Beware argument order!
- No distinction between signed and unsigned
  - Only arithmetic vs. logical shifts
- How do you implement "r3 = r1 + r2"?

Format	Computation	
<code>addq src, dst</code>	$dst = dst + src$	( $dst \neq src$ )
<code>subq src, dst</code>	$dst = dst - src$	
<code>imulq src, dst</code>	$dst = dst * src$	signed mult
<code>sarq src, dst</code>	$dst = dst \gg src$	Arithmetic
<code>shrq src, dst</code>	$dst = dst \gg src$	Logical
<code>shlq src, dst</code>	$dst = dst \ll src$	(same as <code>salq</code> )
<code>xorq src, dst</code>	$dst = dst \wedge src$	
<code>andq src, dst</code>	$dst = dst \& src$	
<code>orq src, dst</code>	$dst = dst \vee src$	

Imm, Reg, or Mem

operation      operand size specifier (b, w, l, q)

① `movq r2, r3`      # r3 = r2

② `addq r1, r3`      # r3 = r2 + r1

# Some Arithmetic Operations

## ❖ Unary (one-operand) Instructions:

Format	Computation	
<code>incq dst</code>	$dst = dst + 1$	increment
<code>decq dst</code>	$dst = dst - 1$	decrement
<code>negq dst</code>	$dst = -dst$	negate
<code>notq dst</code>	$dst = \sim dst$	bitwise complement

## ❖ See CSPP Section 3.5.5 for more instructions: `mulq`, `cqto`, `idivq`, `divq`



# Arithmetic Example

Register	Use(s)
<u>%rdi</u>	1 <sup>st</sup> argument (x)
<u>%rsi</u>	2 <sup>nd</sup> argument (y)
<u>%rax</u>	return value

Convention!

```

long simple_arith(long x, long y)
{
    long t1 = x + y;
    long t2 = t1 * 3;
    return t2;
}
    
```

*don't actually need new variables!*

```

y += x;
y *= 3;
long r = y;
return r;
    
```

*must return in %rax*


```

simple_arith:
    addq    %rdi, %rsi
    imulq   $3, %rsi
    movq    %rsi, %rax
    ret
    
```

*# return*

# Example of Basic Addressing Modes

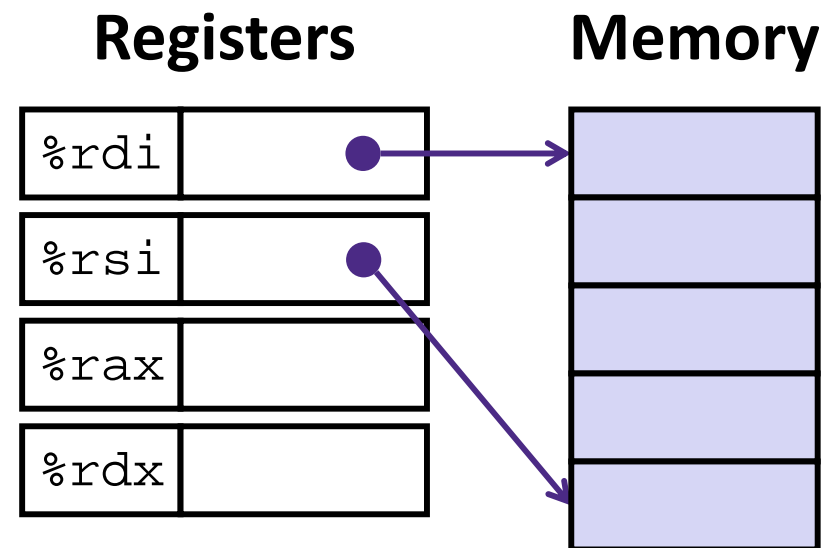
```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```



```
swap:          src , dst (AT &T syntax)
    movq    (%rdi), %rax
    movq    (%rsi), %rdx
    movq    %rdx, (%rdi)
    movq    %rax, (%rsi)
    ret
```

# Understanding swap( )

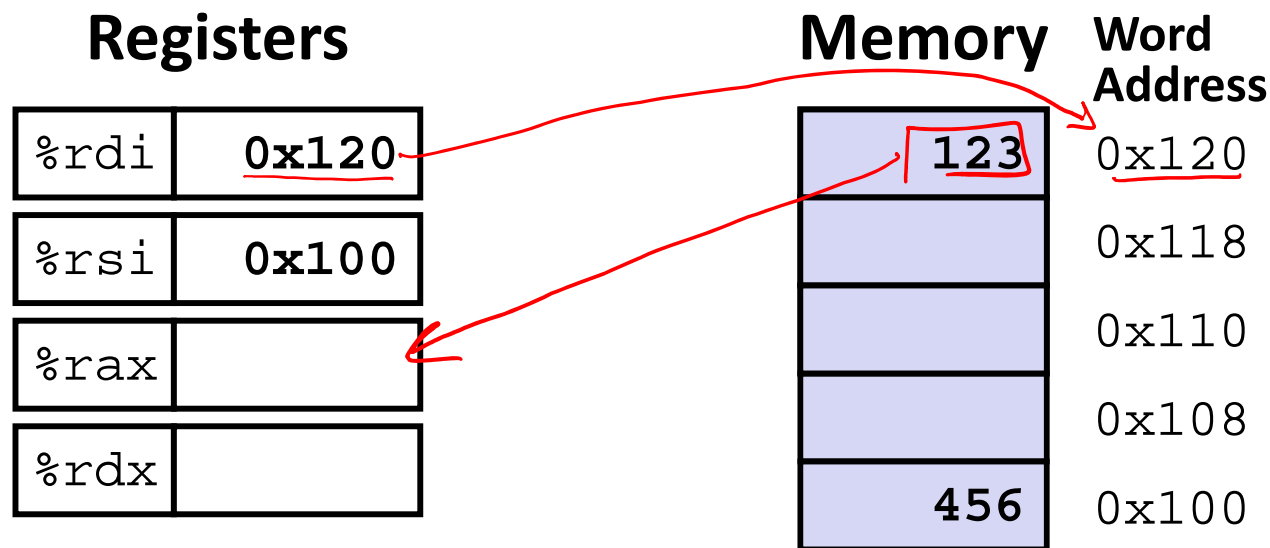
```
void swap(long *xp, long *yp)
{
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```



```
swap:
    movq    (%rdi), %rax
    movq    (%rsi), %rdx
    movq    %rdx, (%rdi)
    movq    %rax, (%rsi)
    ret
```

<u>Register</u>		<u>Variable</u>
%rdi	↔	xp
%rsi	↔	yp
%rax	↔	t0
%rdx	↔	t1

# Understanding swap( )



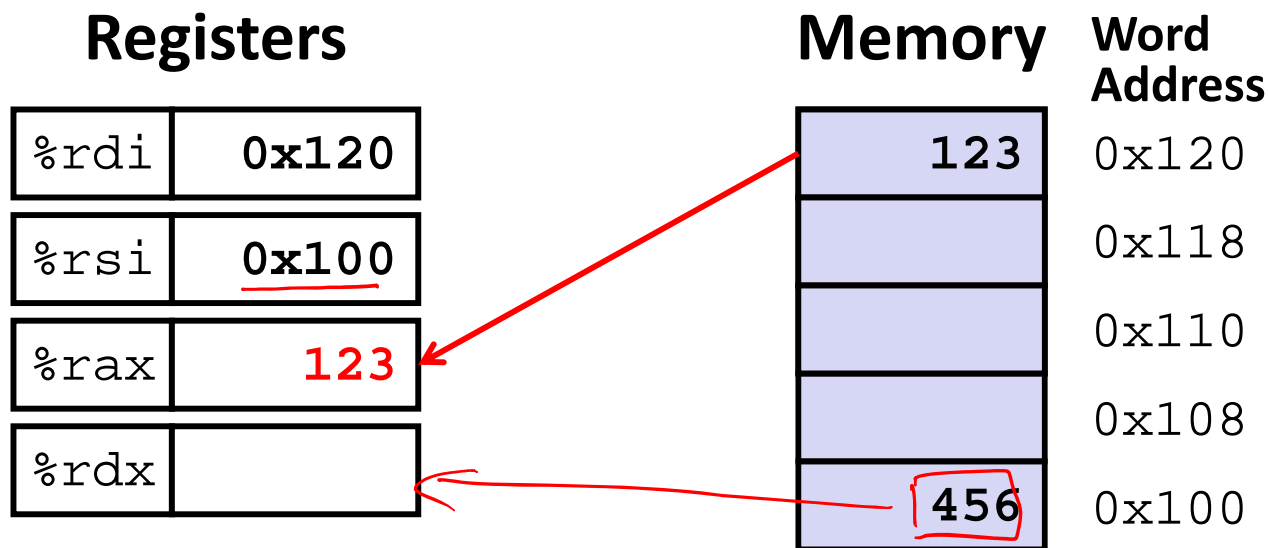
```

swap:
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)    # *xp = t1
    movq    %rax, (%rsi)    # *yp = t0
    ret
    
```

*src dst*

*Comment*

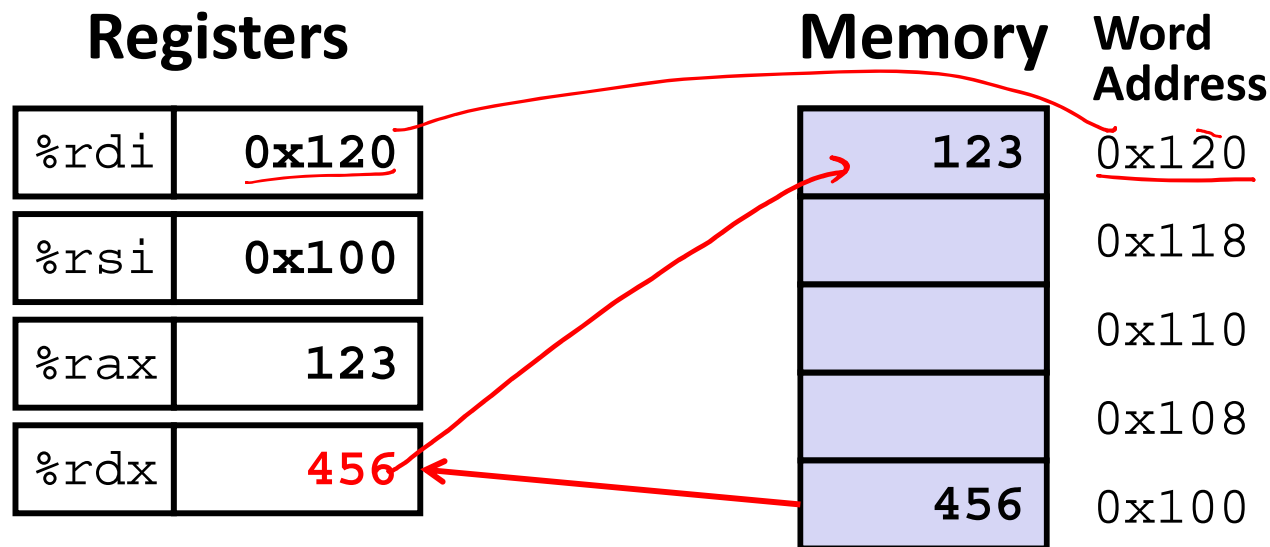
# Understanding swap( )



```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)   # *xp = t1
    movq    %rax, (%rsi)   # *yp = t0
    ret
```

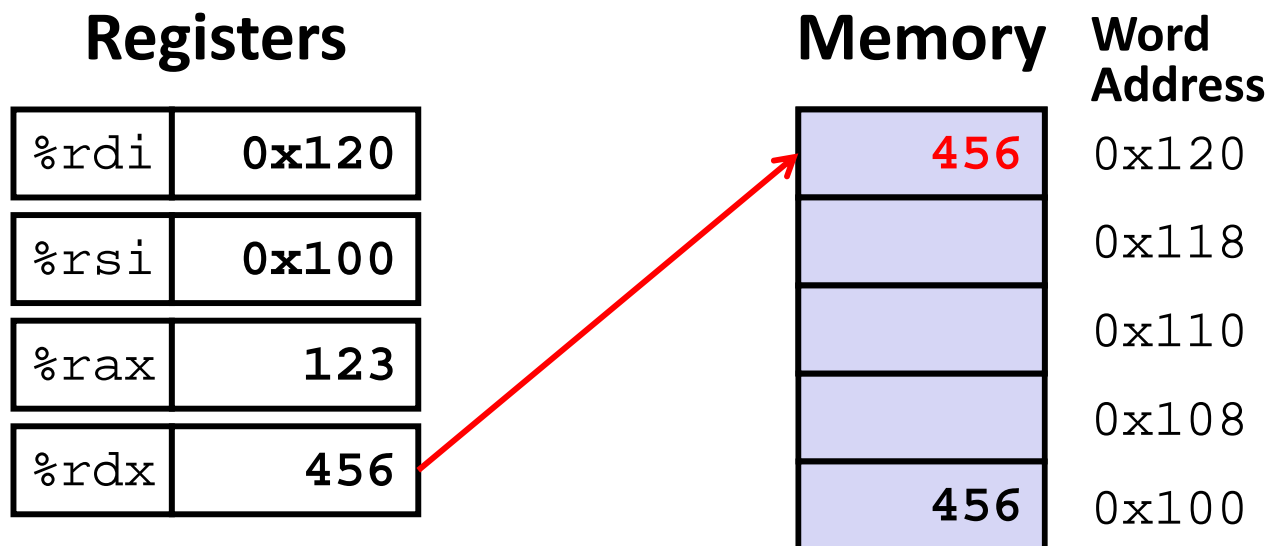
# Understanding swap( )



```

swap:
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi) # *xp = t1
    movq    %rax, (%rsi)   # *yp = t0
    ret
    
```

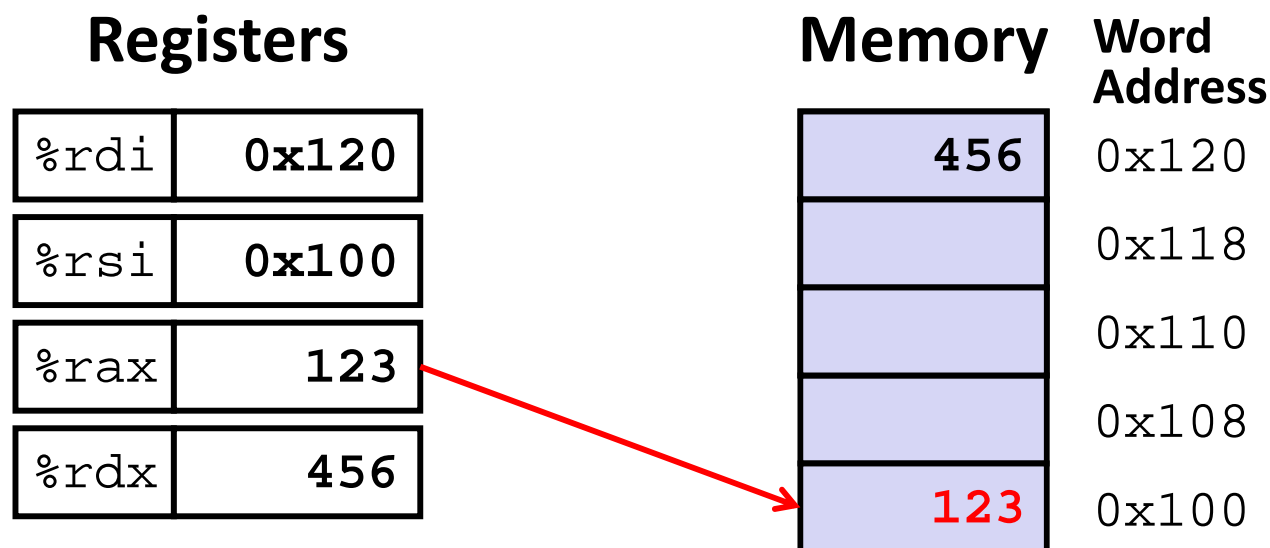
# Understanding swap( )



```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)   # *xp = t1
    movq    %rax, (%rsi)   # *yp = t0
    ret
```

# Understanding swap()



```
swap:
```

```
    movq    (%rdi), %rax    # t0 = *xp
    movq    (%rsi), %rdx    # t1 = *yp
    movq    %rdx, (%rdi)    # *xp = t1
    movq    %rax, (%rsi)    # *yp = t0
    ret
```



# Memory Addressing Modes: Basic

- ❖ **Indirect:**  $(R)$  Mem[Reg[R]]
    - Data in register R specifies the memory address
    - Like pointer dereference in C
    - Example: `movq (%rcx), %rax`
- Handwritten notes:*  
 - "name of register" with an arrow pointing to (R)  
 - "treat Mem as an array" with an arrow pointing to Mem[Reg[R]]  
 - "value stored in register" with an arrow pointing to Reg[R]

- ❖ **Displacement:**  $D(R)$  Mem[Reg[R]+D]
    - Data in register R specifies the *start* of some memory region
    - Constant displacement D specifies the offset from that address
    - Example: `movq 8(%rbp), %rdx`
- Handwritten note:*  
 - "no space" with an arrow pointing to the space between D and (R) in D(R)

# Complete Memory Addressing Modes

$$ar[i] \leftrightarrow *(ar + i) \rightarrow \text{Mem}[ar + i * \text{size of (data type)}]$$

## ❖ General:

- $D(Rb, Ri, S)$      $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S + D]$ 
  - $Rb$ :     Base register (any register)
  - $Ri$ :     Index register (any register except `%rsp`)
  - $S$ :     Scale factor (1, 2, 4, 8) – *why these numbers?*
  - $D$ :     Constant displacement value (a.k.a. immediate)

## ❖ Special cases (see CSPP Figure 3.3 on p.181)

- $D(Rb, Ri)$       $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D]$     ( $S=1$ )
- $(Rb, Ri, S)$      $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] * S]$     ( $D=0$ )
- $(Rb, Ri)$       $\text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]]$      ( $S=1, D=0$ )
- $(, Ri, S)$       $\text{Mem}[\text{Reg}[Ri] * S]$             ( $Rb=0, D=0$ )

↑ so reg name not interpreted as  $Rb$

(if not specified)

default values:

$S = 1$

$D = 0$

$Reg[Rb] = 0$

$Reg[Ri] = 0$

# Address Computation Examples

%rdx	0xf000
%rcx	0x0100

$D(Rb, Ri, S) \rightarrow$

$Mem[Reg[Rb] + Reg[Ri] * S + D]$

Expression	Address Computation	Address
$0x8 ( \overset{D}{}, \overset{Rb}{\%rdx} )$	$Reg[Rb] + D = 0xf000 + 0x8$	$0xf008$
$( \overset{Rb}{\%rdx}, \overset{Ri}{\%rcx} )$	$Reg[Rb] + Reg[Ri] * 1$	$0xf100$
$( \overset{Rb}{\%rdx}, \overset{Ri}{\%rcx}, \overset{S}{4} )$	$*4$	$0xf400$
$0x80 ( \overset{D}{}, \overset{Ri}{\%rdx}, \overset{S}{2} )$	$Reg[Ri] * 2 + 0x80$	$0x1e080$

$0xf000 * 2$

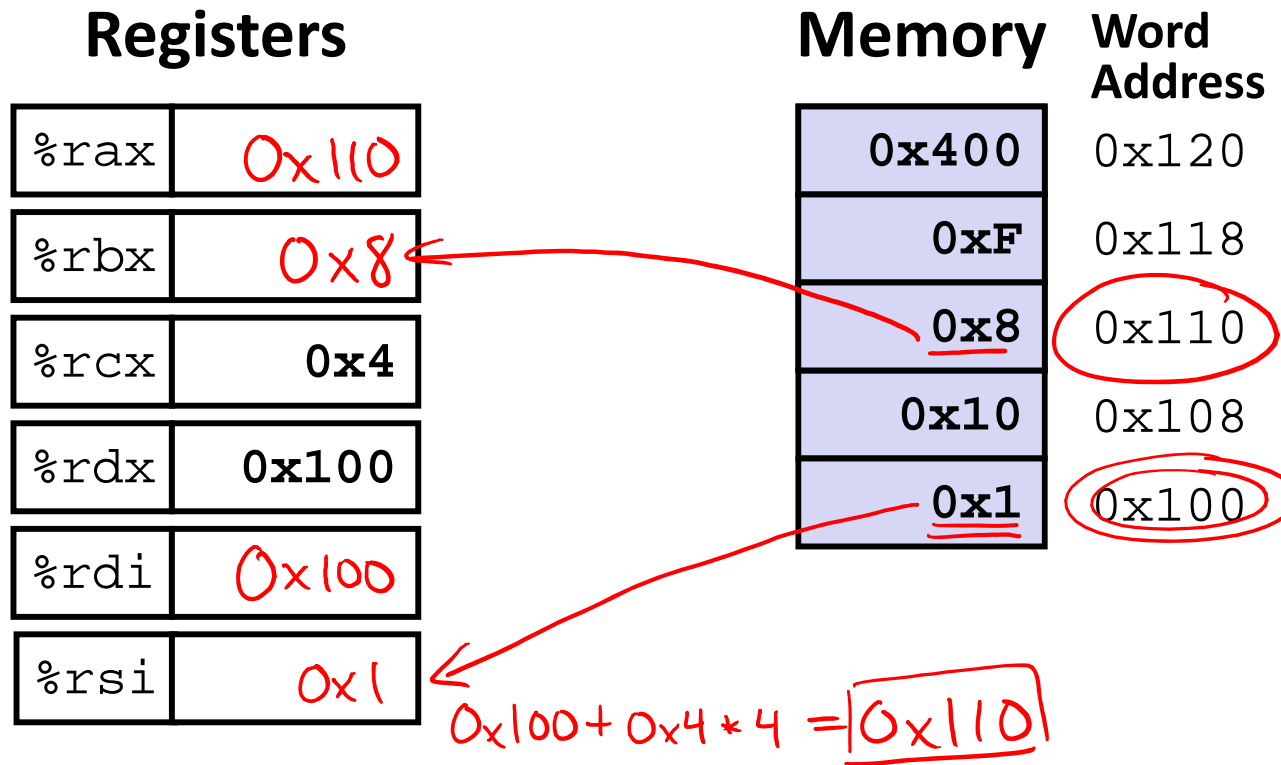
$0xf000 \ll 1 = 0x1e000$

$$\begin{array}{r} 1111\ 0000 \\ 1\ 1110\ 0000 \dots 0 \end{array}$$

# Address Computation Instruction

- ❖  $\overset{\text{"Mem"}}{\text{leaq}} \overset{\text{Reg}}{\text{src}}, \text{dst}$ 
  - "lea" stands for *load effective address*
  - src is address expression (any of the formats we've seen)
  - dst is a register  $\hookrightarrow$  calculates  $\text{Reg}[\text{Rb}] + \text{Reg}[\text{Ri}] * S + D$
  - Sets dst to the *address* computed by the src expression  
(**does not go to memory!** – it just does math) ~~Mem!~~
  - Example: `leaq (%rdx,%rcx,4), %rax`
- ❖ Uses:
  - Computing addresses without a memory reference
    - e.g. translation of `p = &x[i];`  $\leftarrow$  address-of operator
  - Computing arithmetic expressions of the form  $x+k*i+d$ 
    - Though  $k$  can only be 1, 2, 4, or 8

# Example: lea vs. mov



<code>leaq</code> ( <sup>Rb</sup> %rdx, <sup>Ri</sup> %rcx, <sup>S</sup> 4), %rax	→ 0x110	("addr")
<code>movq</code> (%rdx, %rcx, 4), %rbx	→ 0x8	(data)
<code>leaq</code> (%rdx), %rdi	→ 0x100	("addr")
<code>movq</code> (%rdx), %rsi	→ 0x1	(data)

0x100

# Arithmetic Example

```
long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

Register	Use(s)
%rdi	1 <sup>st</sup> argument (x)
%rsi	2 <sup>nd</sup> argument (y)
%rdx	3 <sup>rd</sup> argument (z)

← replaced by `lea` & `shift`

```
arith:
    leaq    (%rdi,%rsi), %rax    # rax = x + y (t1)
    addq   %rdx, %rax           # rax = x + y + z (t2)
    leaq   (%rsi,%rsi,2), %rdx  # rdx = 3y
    salq   $4, %rdx            # rdx = 48y (t4)
    leaq   4(%rdi,%rdx), %rcx
    imulq  %rcx, %rax
    ret
```

← multiplying two variables

## Interesting Instructions

- `leaq`: “address” computation
- `salq`: shift
- `imulq`: multiplication
- Only used once!

# Arithmetic Example

```

long arith(long x, long y, long z)
{
    long t1 = x + y;
    long t2 = z + t1;
    long t3 = x + 4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
    
```

Register	Use(s)
%rdi	x
%rsi	y
%rdx	z, t4
%rax	t1, t2, rval
%rcx	t5

limited registers means they often get reused!

```

arith:
    leaq    (%rdi,%rsi), %rax    # rax/t1    = x + y
    addq   %rdx, %rax           # rax/t2    = t1 + z
    leaq   (%rsi,%rsi,2), %rdx  # rdx      = 3 * y
    salq   $4, %rdx            # rdx/t4    = (3*y) * 16
    leaq   4(%rdi,%rdx), %rcx   # rcx/t5    = x + t4 + 4
    imulq  %rcx, %rax          # rax/rval  = t5 * t2
    ret
    
```

*comment (AT &T syntax)*

*S ∈ {1,2,4,8}*

# Peer Instruction Question

❖ Which of the following x86-64 instructions correctly calculates  $\%rax = 9 * \%rdi$ ?

▪ Vote at <http://PollEv.com/justinh>

~~A.~~ `leaq (, %rdi, 9), %rax` ↖  $s \in \{1, 2, 4, 8\}$

~~B.~~ `movq (, %rdi, 9), %rax`

**C.** `leaq (%rdi, %rdi, 8), %rax`

**D.** `movq (%rdi, %rdi, 8), %rax`

**E.** We're lost...

→  $\%rax = 9 * \%rdi$   
 $\%rax = *(9 * \%rdi)$



# Summary

- ❖ **Memory Addressing Modes:** The addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways
  - *Base register, index register, scale factor, and displacement* map well to pointer arithmetic operations
- ❖ `lea` is address calculation instruction
  - Does NOT actually go to memory
  - Used to compute addresses or some arithmetic expressions