

Data III & Integers I

CSE 351 Autumn 2017

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<http://xkcd.com/257/>

Administrivia

- ❖ Homework 1 due tonight
- ❖ Lab 1 released
 - *This is considered the hardest assignment of the class by many students*
 - Some progress due Monday 10/9, Lab 1 due Friday 1/13
- ❖ Poll Everywhere: you can change your vote

Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
- ❖ Organizing and addressing data in memory
- ❖ Manipulating data in memory using C
- ❖ **Boolean algebra and bit-level manipulations**

Boolean Algebra

- ❖ Developed by George Boole in 19th Century
 - Algebraic representation of logic (True → 1, False → 0)
 - AND: $A \& B = 1$ when both A is 1 and B is 1
 - OR: $A | B = 1$ when either A is 1 or B is 1
 - XOR: $A ^ B = 1$ when either A is 1 or B is 1, but not both
 - NOT: $\sim A = 1$ when A is 0 and vice-versa
 - DeMorgan's Law:
$$\sim (A | B) = \sim A \& \sim B$$
$$\sim (A \& B) = \sim A | \sim B$$

AND		OR		XOR		NOT	
&	0 1		0 1	^	0 1	~	
0	0 0	0	0 1	0	0 1	0	1
1	0 1	1	1 1	1	1 0	1	0

General Boolean Algebras

- ❖ Operate on bit vectors
 - Operations applied bitwise
 - All of the properties of Boolean algebra apply

$$\begin{array}{r} 01101001 \\ \& 01010101 \\ \hline \end{array} \quad \begin{array}{r} 01101001 \\ + 01010101 \\ \hline \end{array} \quad \begin{array}{r} 01101001 \\ \wedge 01010101 \\ \hline \end{array} \quad \begin{array}{r} 01010101 \\ \sim 01010101 \\ \hline \end{array}$$

- ❖ Examples of useful operations:

$$x \wedge x = 0$$

$$\begin{array}{r} 01010101 \\ \wedge 01010101 \\ \hline 00000000 \end{array}$$

$$x | 1 = 1, \quad x | 0 = x$$

$$\begin{array}{r} 01010101 \\ | 11110000 \\ \hline 11110101 \end{array}$$

Bit-Level Operations in C

- ❖ & (AND) , | (OR) , ^ (XOR) , ~ (NOT)
 - View arguments as bit vectors, apply operations bitwise
 - Apply to any “integral” data type
 - long, int, short, char, unsigned
- ❖ Examples with `char a, b, c;`
 - `a = (char) 0x41;` // 0x41->0b 0100 0001
 - `b = ~a;` // 0b ->0x
 - `a = (char) 0x69;` // 0x69->0b 0110 1001
 - `b = (char) 0x55;` // 0x55->0b 0101 0101
 - `c = a & b;` // 0b ->0x
 - `a = (char) 0x41;` // 0x41->0b 0100 0001
 - `b = a;` // 0b 0100 0001
 - `c = a ^ b;` // 0b ->0x

Contrast: Logic Operations

- ❖ Logical operators in C: `&&` (AND), `||` (OR), `!` (NOT)
 - 0 is False, anything nonzero is True
 - Always return 0 or 1
 - **Early termination** (a.k.a. short-circuit evaluation) of `&&`, `||`
- ❖ Examples (char data type)
 - `! 0x41` → `0x00`
 - `! 0x00` → `0x01`
 - `!! 0x41` → `0x01`
 - `p && *p++`
 - Avoids **null pointer** (`0x0`) access via *early termination*
 - Short for: `if (p) { *p++; }`
 - `0xCC && 0x33` → `0x01`
 - `0x00 || 0x33` → `0x01`

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

Assembly language:

```
get_mpg:
    pushq  %rbp
    movq   %rsp, %rbp
    ...
    popq   %rbp
    ret
```

Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

Computer system:



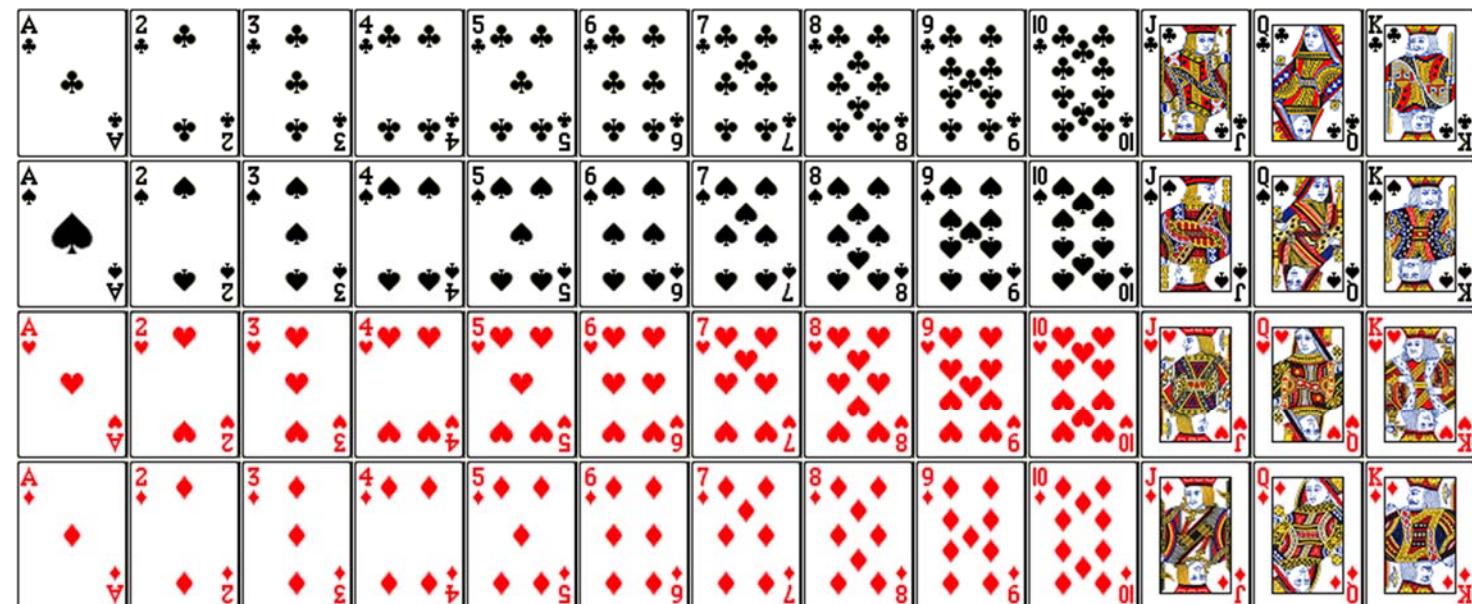
Memory & data
Integers & floats
Machine code & C
x86 assembly
Procedures & stacks
Arrays & structs
Memory & caches
Processes
Virtual memory
Operating Systems

OS:



But before we get to integers....

- ❖ Encode a standard deck of playing cards
- ❖ 52 cards in 4 suits
 - How do we encode suits, face cards?
- ❖ What operations do we want to make easy to implement?
 - Which is the higher value card?
 - Are they the same suit?



Two possible representations

1) 1 bit per card (52): bit corresponding to card set to 1



low-order 52 bits of 64-bit word

- “One-hot” encoding (similar to set notation)
- Drawbacks:
 - Hard to compare values and suits
 - Large number of bits required

2) 1 bit per suit (4), 1 bit per number (13): 2 bits set

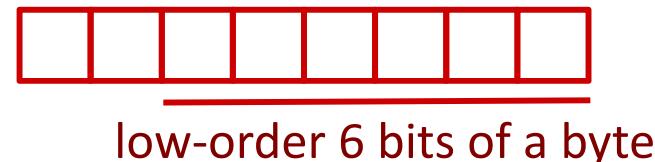


- Pair of one-hot encoded values
- Easier to compare suits and values, but still lots of bits used
- ❖ Can we do better?

Two better representations

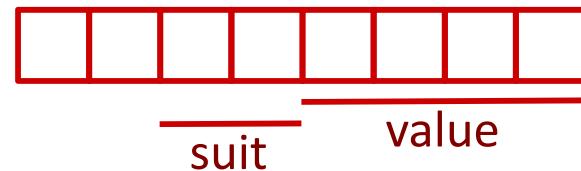
3) Binary encoding of all 52 cards – only 6 bits needed

- $2^6 = 64 \geq 52$



- Fits in one byte (smaller than one-hot encodings)
- How can we make value and suit comparisons easier?

4) Separate binary encodings of suit (2 bits) and value (4 bits)



- Also fits in one byte, and easy to do comparisons

K	Q	J	...	3	2	A
1101	1100	1011	...	0011	0010	0001

♣	00
♦	01
♥	10
♠	11

Compare Card Suits

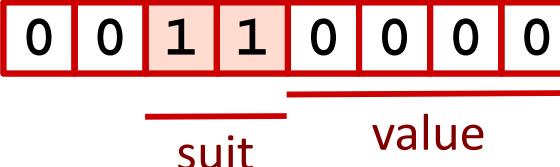
mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .
Here we turns all *but* the bits of interest in v to 0.

```
char hand[ 5 ];           // represents a 5-card hand
char card1, card2;       // two cards to compare
card1 = hand[ 0 ];
card2 = hand[ 1 ];
...
if ( sameSuitP(card1, card2) ) { ... }
```

```
#define SUIT_MASK 0x30
```

```
int sameSuitP(char card1, char card2) {
    return !( (card1 & SUIT_MASK) ^ (card2 & SUIT_MASK) );
    //return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

returns int

SUIT_MASK = 0x30 = 

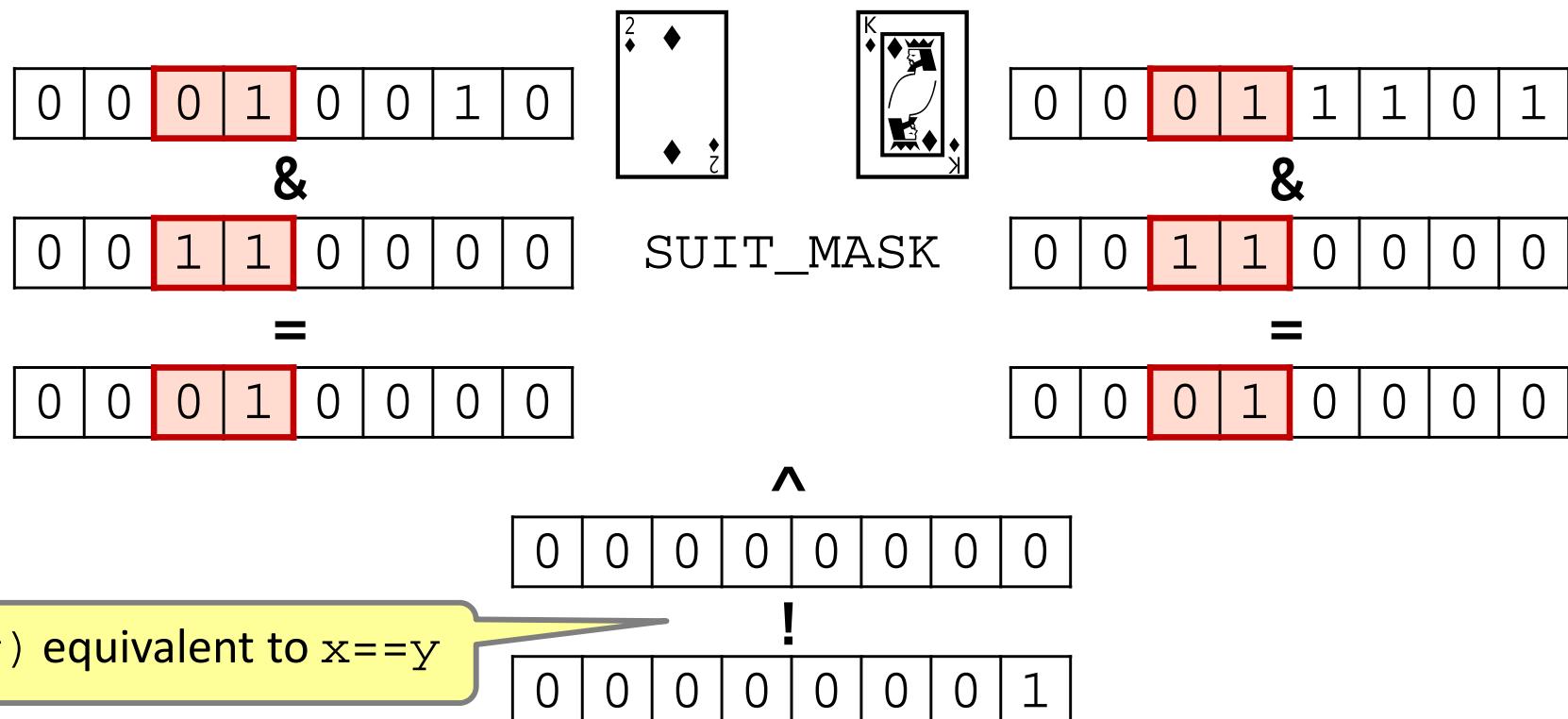
equivalent

Compare Card Suits

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}
```



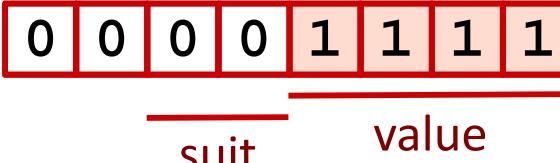
Compare Card Values

```
char hand[ 5 ];           // represents a 5-card hand
char card1, card2;       // two cards to compare
card1 = hand[ 0 ];
card2 = hand[ 1 ];
...
if ( greaterValue(card1, card2) ) { ... }
```

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .

VALUE_MASK = 0x0F = 
The diagram shows the binary value 0x0F (0000 1111) in a row of eight boxes. A red horizontal line under the first four boxes is labeled "suit", and a red horizontal line under the last four boxes is labeled "value".

Compare Card Values

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v .

```
#define VALUE_MASK 0x0F

int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```

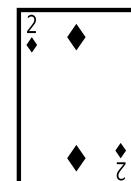
0	0	1	0	0	0	1	0
---	---	---	---	---	---	---	---

&

0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---

=

0	0	0	0	0	0	1	0
---	---	---	---	---	---	---	---



VALUE_MASK

0	0	1	0	1	1	0	1
---	---	---	---	---	---	---	---

&

0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---

=

0	0	0	0	1	1	0	1
---	---	---	---	---	---	---	---

$$2_{10} > 13_{10}$$

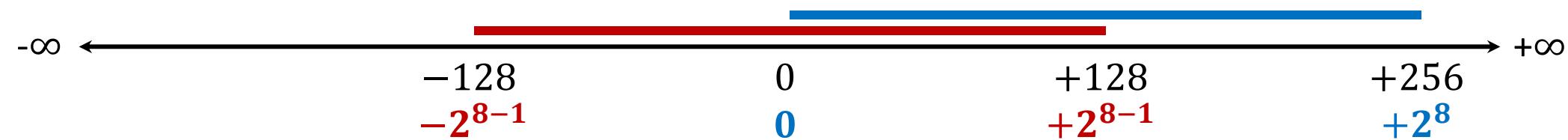
0 (false)

Integers

- ❖ **Binary representation of integers**
 - Unsigned and signed
 - Casting in C
- ❖ Consequences of finite width representation
 - Overflow, sign extension
- ❖ Shifting and arithmetic operations

Encoding Integers

- ❖ The hardware (and C) supports two flavors of integers
 - *unsigned* – only the non-negatives
 - *signed* – both negatives and non-negatives
- ❖ Cannot represent all integers with w bits
 - Only 2^w distinct bit patterns
 - Unsigned values: $0 \dots 2^w - 1$
 - Signed values: $-2^{w-1} \dots 2^{w-1} - 1$
- ❖ Example: 8-bit integers (e.g. `char`)



Unsigned Integers

- ❖ Unsigned values follow the standard base 2 system
 - $b_7b_6b_5b_4b_3b_2b_1b_0 = b_72^7 + b_62^6 + \dots + b_12^1 + b_02^0$
- ❖ Add and subtract using the normal “carry” and “borrow” rules, just in binary

63	00111111
+ 8	+ <hr/> 00001000
71	01000111

- ❖ Useful formula: $2^{N-1} + 2^{N-2} + \dots + 2 + 1 = 2^N - 1$
 - i.e. N ones in a row = $2^N - 1$
- ❖ How would you make *signed* integers?

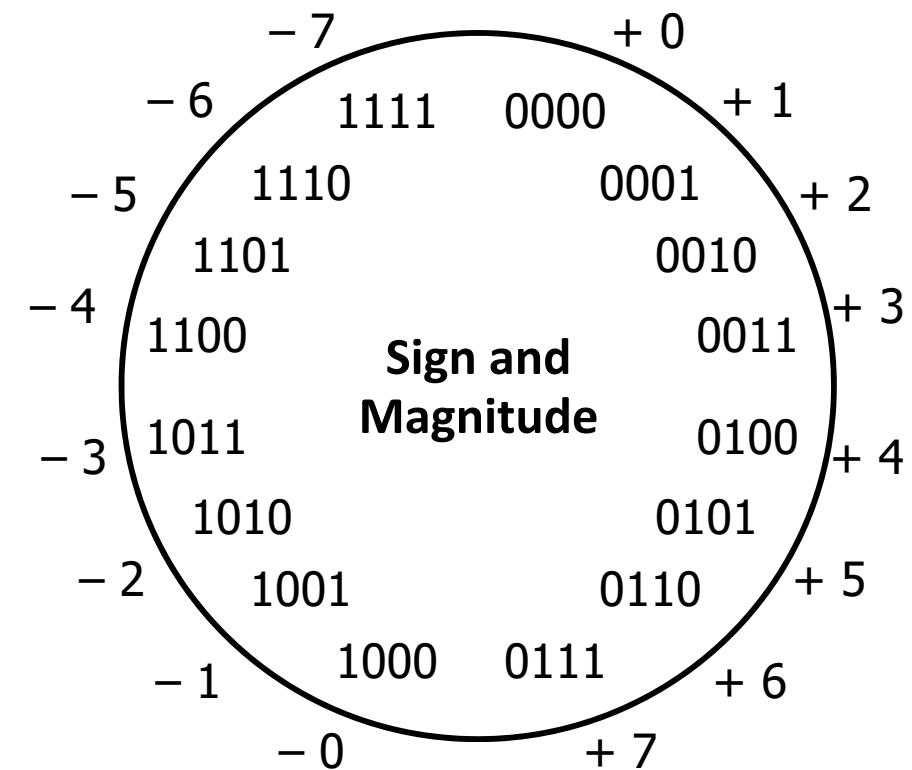
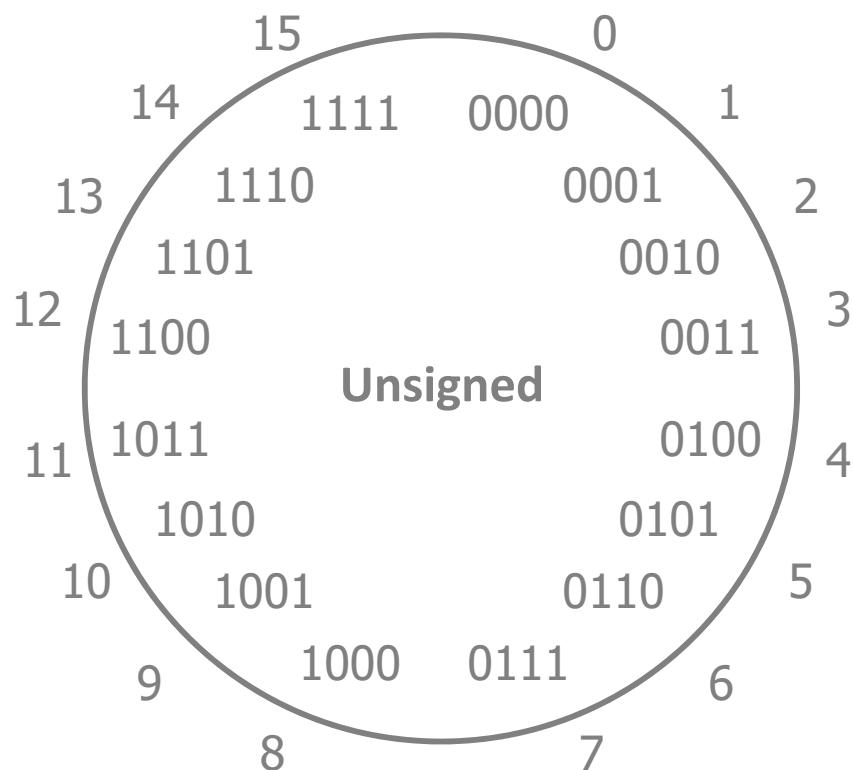
Sign and Magnitude

Most Significant Bit

- ❖ Designate the high-order bit (MSB) as the “sign bit”
 - sign=0: positive numbers; sign=1: negative numbers
- ❖ Benefits:
 - Using MSB as sign bit matches positive numbers with unsigned
 - All zeros encoding is still = 0
- ❖ Examples (8 bits):
 - 0x00 = 00000000_2 is non-negative, because the sign bit is 0
 - 0x7F = 01111111_2 is non-negative ($+127_{10}$)
 - 0x85 = 10000101_2 is negative (-5_{10})
 - 0x80 = 10000000_2 is negative... zero???

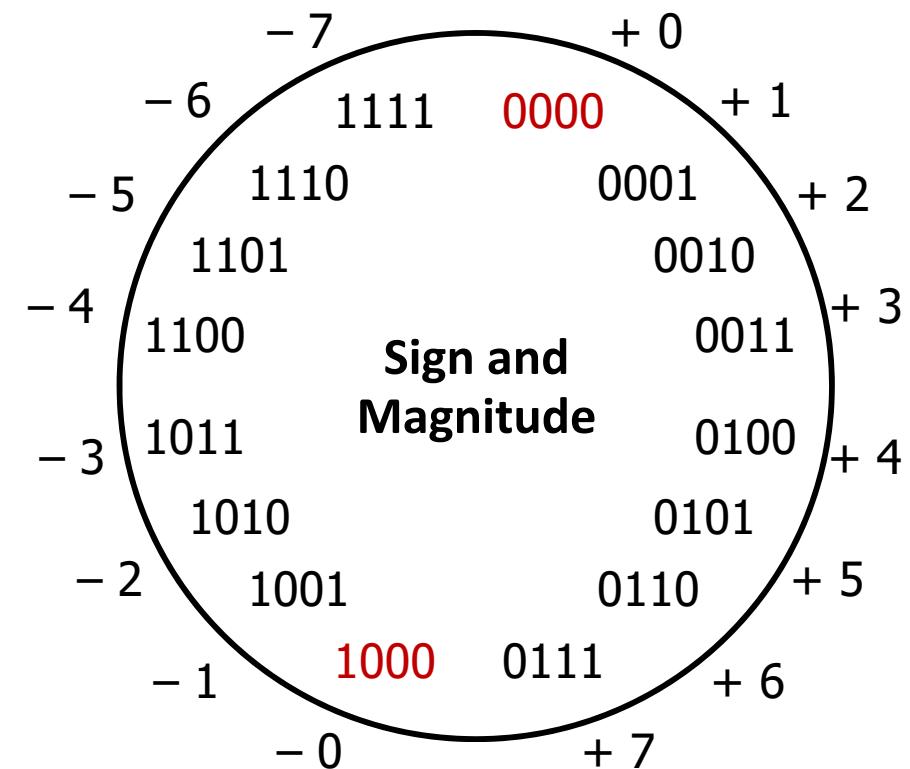
Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks?



Sign and Magnitude

- ❖ MSB is the sign bit, rest of the bits are magnitude
- ❖ Drawbacks:
 - Two representations of 0 (bad for checking equality)



Sign and Magnitude

- MSB is the sign bit, rest of the bits are magnitude
- Drawbacks:
 - Two representations of 0 (bad for checking equality)
 - Arithmetic is cumbersome**
 - Example: $4 - 3 \neq 4 + (-3)$

$$\begin{array}{r} 4 \\ - 3 \\ \hline 1 \end{array}$$

0100
- 0011
0001

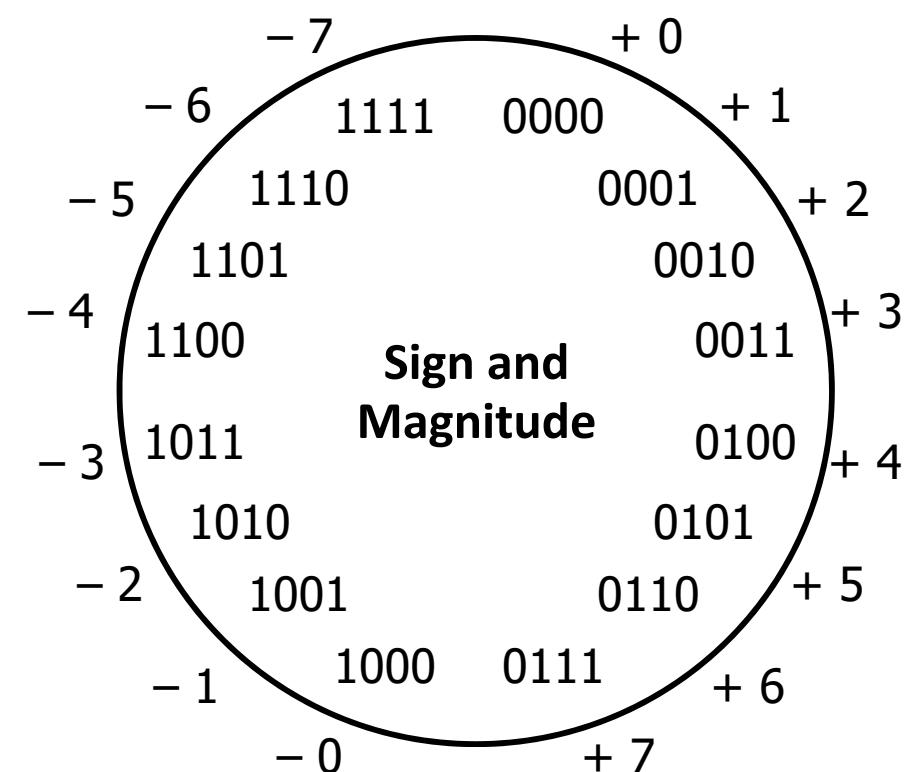
✓

$$\begin{array}{r} 4 \\ + -3 \\ \hline -7 \end{array}$$

0100
+ 1011
1111

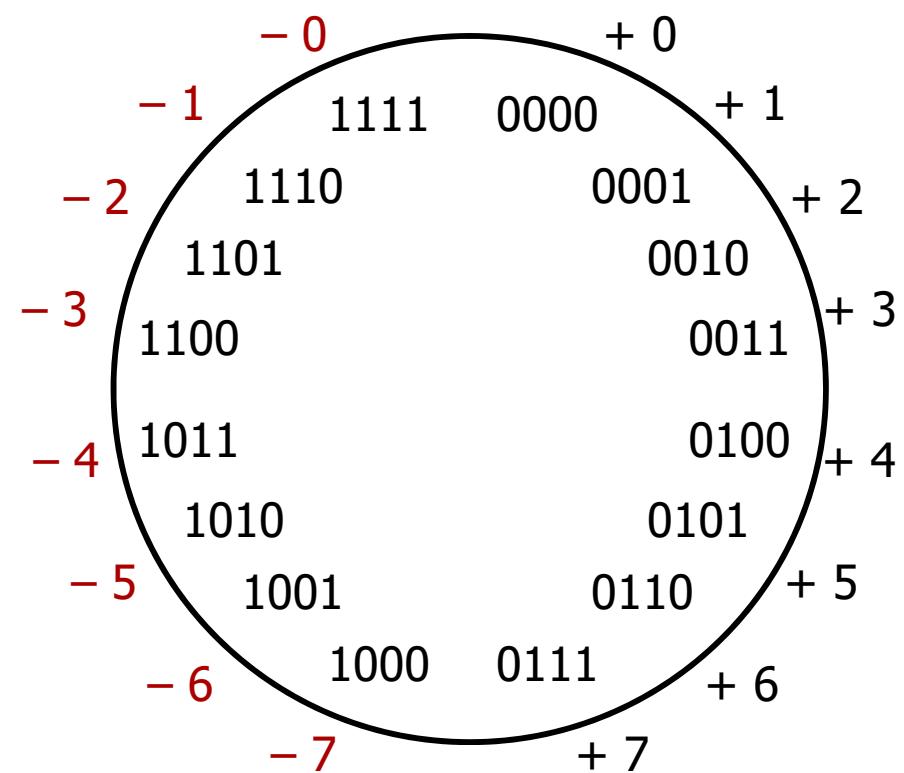
✗

- Negatives “increment” in wrong direction!



Two's Complement

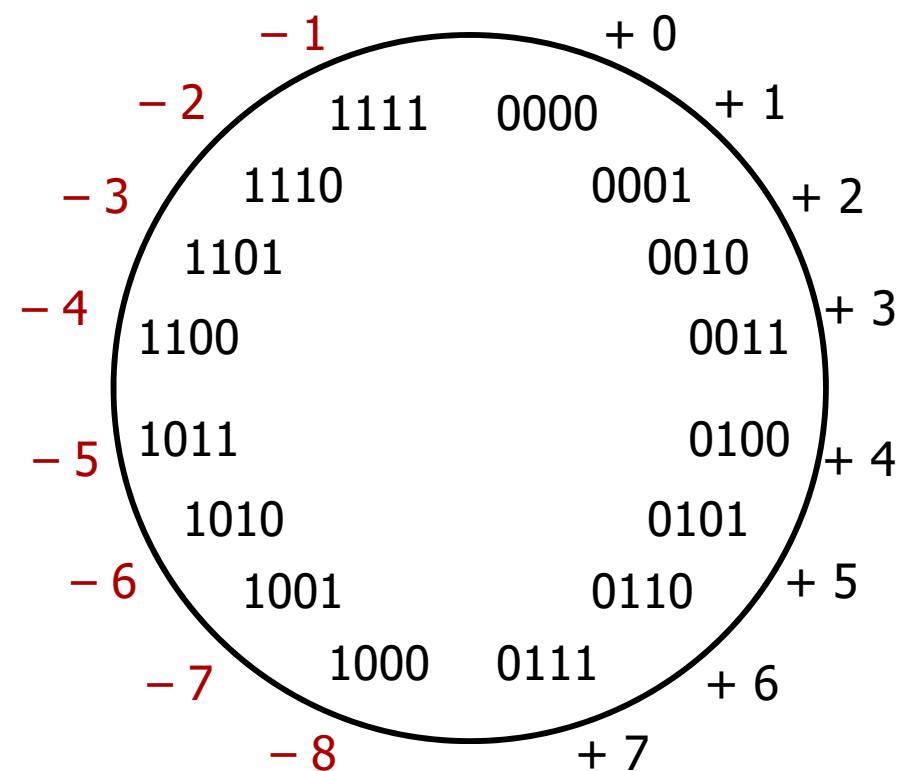
- ❖ Let's fix these problems:
 - 1) “Flip” negative encodings so incrementing works



Two's Complement

- ❖ Let's fix these problems:
 - 1) “Flip” negative encodings so incrementing works
 - 2) “Shift” negative numbers to eliminate -0

- ❖ MSB *still* indicates sign!
 - This is why we represent one more negative than positive number (-2^{N-1} to $2^{N-1} - 1$)



Two's Complement Negatives

- Accomplished with one neat mathematical trick!

b_{w-1} has weight -2^{w-1} , other bits have usual weights $+2^i$



- 4-bit Examples:

- 1010_2 unsigned:

$$1*2^3 + 0*2^2 + 1*2^1 + 0*2^0 = 10$$

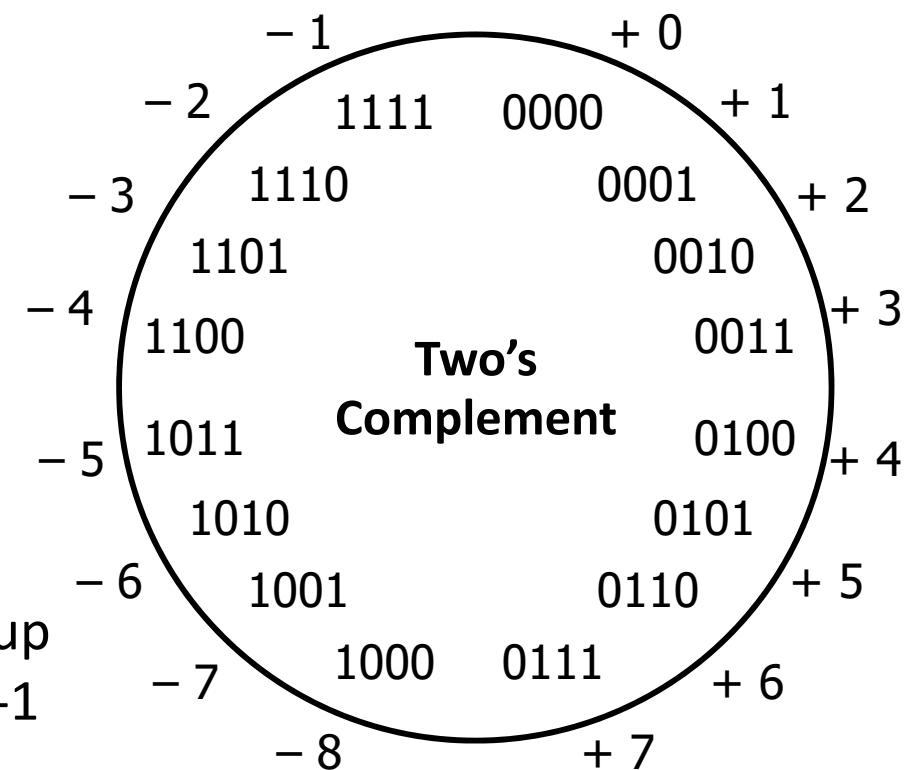
- 1010_2 two's complement:

$$-1*2^3 + 0*2^2 + 1*2^1 + 0*2^0 = -6$$

- 1 represented as:

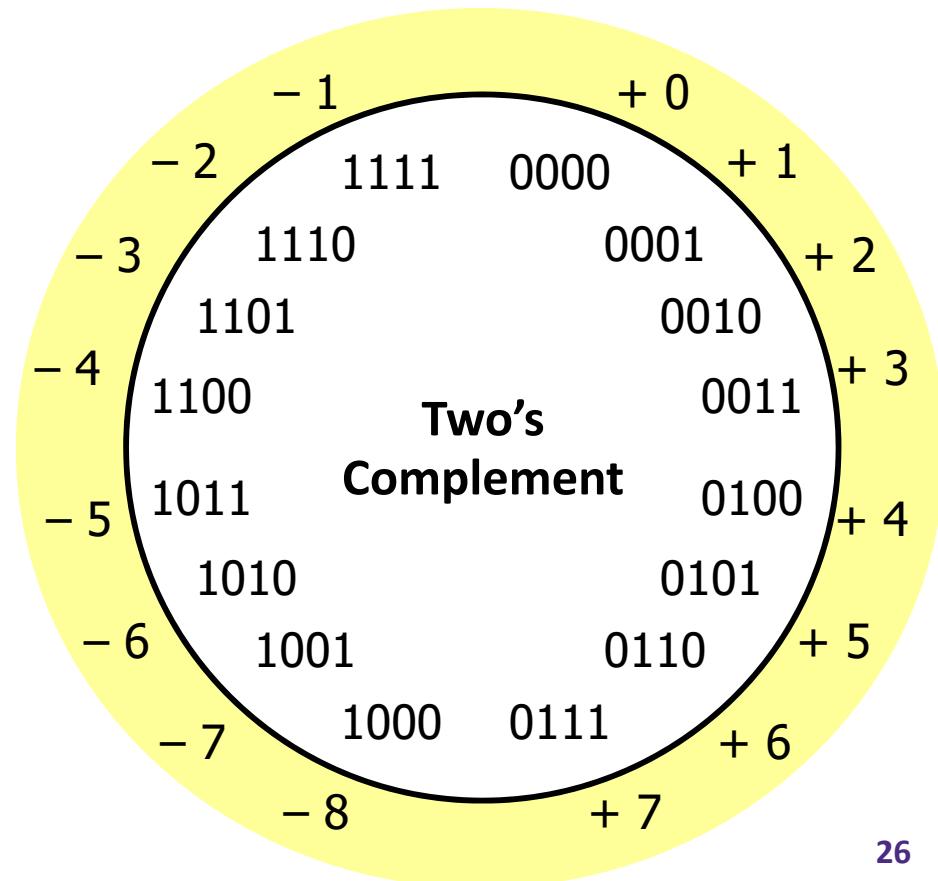
$$1111_2 = -2^3 + (2^3 - 1)$$

- MSB makes it super negative, add up all the other bits to get back up to -1



Why Two's Complement is So Great

- ❖ Roughly same number of (+) and (-) numbers
- ❖ Positive number encodings match unsigned
- ❖ Single zero
- ❖ All zeros encoding = 0
- ❖ Simple negation procedure:
 - Get negative representation of any integer by taking bitwise complement and then adding one!
 $(\sim x + 1 == -x)$



Peer Instruction Question

- ❖ Take the 4-bit number encoding $x = 0b1011$
 - ❖ Which of the following numbers is NOT a valid interpretation of x using any of the number representation schemes discussed today?
 - Unsigned, Sign and Magnitude, Two's Complement
 - Vote at <http://PollEv.com/justinh>
- A. -4
- B. -5
- C. 11
- D. -3
- E. We're lost...

Summary

- ❖ Bit-level operators allow for fine-grained manipulations of data
 - Bitwise AND (`&`), OR (`|`), and NOT (`~`) different than logical AND (`&&`), OR (`||`), and NOT (`!`)
 - Especially useful with bit masks
- ❖ Choice of *encoding scheme* is important
 - Tradeoffs based on size requirements and desired operations
- ❖ Integers represented using unsigned and two's complement representations
 - Limited by fixed bit width
 - We'll examine arithmetic operations next lecture