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## Data III & Integers I

CSE 351 Autumn 2017

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<http://xkcd.com/257/>

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## Administrivia

- ❖ Homework 1 due tonight
- ❖ Lab 1 released
  - This is considered the hardest assignment of the class by many students
  - Some progress due Monday 10/9, Lab 1 due Friday 1/13
- ❖ Poll Everywhere: you can change your vote

2

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## Memory, Data, and Addressing

- ❖ Representing information as bits and bytes
- ❖ Organizing and addressing data in memory
- ❖ Manipulating data in memory using C
- ❖ **Boolean algebra and bit-level manipulations**

3

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## Boolean Algebra

- ❖ Developed by George Boole in 19th Century
  - Algebraic representation of logic (True  $\rightarrow$  1, False  $\rightarrow$  0)
  - AND:  $A \& B = 1$  when both A is 1 and B is 1
  - OR:  $A | B = 1$  when either A is 1 or B is 1
  - XOR:  $A \wedge B = 1$  when either A is 1 or B is 1, but not both
  - NOT:  $\sim A = 1$  when A is 0 and vice-versa
  - DeMorgan's Law:  $\sim(A | B) = \sim A \& \sim B$   
 $\sim(A \& B) = \sim A | \sim B$

	AND	OR	XOR	NOT
$\&$	$\begin{array}{c c} 0 & 1 \\ \hline 0 & 0 \\ 1 & 0 \end{array}$	$\begin{array}{c c} 0 & 1 \\ \hline 0 & 1 \\ 1 & 1 \end{array}$	$\begin{array}{c c} 0 & 1 \\ \hline 0 & 0 \\ 1 & 1 \end{array}$	$\begin{array}{c c} 0 & 1 \\ \hline 0 & 1 \\ 1 & 0 \end{array}$

4

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## General Boolean Algebras

- ❖ Operate on bit vectors
  - Operations applied bitwise
  - All of the properties of Boolean algebra apply

```

01101001  01101001  01101001
& 01010101 | 01010101 ^ 01010101 ~ 01010101
  
```

- ❖ Examples of useful operations:

$$x \wedge x = 0$$

$$x | 1 = 1, \quad x | 0 = x$$

5

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## Bit-Level Operations in C

- ❖  $\&$  (AND),  $|$  (OR),  $\wedge$  (XOR),  $\sim$  (NOT)
  - View arguments as bit vectors, apply operations bitwise
  - Apply to any "integral" data type
    - long, int, short, char, unsigned
- ❖ Examples with char a, b, c;
  - `a = (char) 0x41; // 0x41->0b 0100 0001`  
`b = ~a; // 0b ->0x`
  - `a = (char) 0x69; // 0x69->0b 0110 1001`  
`b = (char) 0x55; // 0x55->0b 0101 0101`  
`c = a & b; // 0b ->0x`
  - `a = (char) 0x41; // 0x41->0b 0100 0001`  
`b = a; // 0b 0100 0001`  
`c = a ^ b; // 0b ->0x`

6

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## Contrast: Logic Operations

- Logical operators in C: && (AND), || (OR), ! (NOT)
  - 0 is False, **anything nonzero** is True
  - Always** return 0 or 1
  - Early termination** (a.k.a. short-circuit evaluation) of &&, ||
- Examples (char data type)
  - !0x41 -> 0x00      0xCC && 0x33 -> 0x01
  - !0x00 -> 0x01      0x00 || 0x33 -> 0x01
  - !!0x41 -> 0x01
  - p && \*p++
    - Avoids **null pointer** (0x0) access via **early termination**
    - Short for: `if (p) { *p++; }`

7

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## Roadmap

**C:**

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

**Java:**

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

Memory & data  
**Integers & floats**  
Machine code & C  
x86 assembly  
Procedures & stacks  
Arrays & structs  
Memory & caches  
Processes  
Virtual memory  
Operating Systems

**Assembly language:**

```
get_mpg:
    pushq   %rbp
    movq   %rsp, %rbp
    ...
    popq   %rbp
    ret
```

**Machine code:**

```
0111010000011000
100011010000010000000010
1000100111000010
11000001111101000011111
```

**OS:** Windows 8, Mac, Linux

**Computer system:**

8

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## But before we get to integers....

- Encode a standard deck of playing cards
- 52 cards in 4 suits
  - How do we encode suits, face cards?
- What operations do we want to make easy to implement?
  - Which is the higher value card?
  - Are they the same suit?

9

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## Two possible representations

- 1 bit per card (52): bit corresponding to card set to 1
  - low-order 52 bits of 64-bit word
  - “One-hot” encoding (similar to set notation)
  - Drawbacks:
    - Hard to compare values and suits
    - Large number of bits required
- 1 bit per suit (4), 1 bit per number (13): 2 bits set
  - Pair of one-hot encoded values
  - 4 suits      13 numbers
  - Easier to compare suits and values, but still lots of bits used

Can we do better?

10

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## Two better representations

- Binary encoding of all 52 cards – only 6 bits needed
  - $2^6 = 64 \geq 52$
  - low-order 6 bits of a byte
  - Fits in one byte (smaller than one-hot encodings)
  - How can we make value and suit comparisons easier?
- Separate binary encodings of suit (2 bits) and value (4 bits)
  - Also fits in one byte, and easy to do comparisons

K	Q	J	...	3	2	A
1101	1100	1011	...	0011	0010	0001

♠	00
♠	01
♥	10
♥	11

11

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## Compare Card Suits

**mask:** a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector v. Here we turns all *but* the bits of interest in v to 0.

```
char hand[5]; // represents a 5-card hand
char card1, card2; // two cards to compare
card1 = hand[0];
card2 = hand[1];
...
if ( sameSuitP(card1, card2) ) { ... }
```

```
#define SUIT_MASK 0x30

int sameSuitP(char card1, char card2) {
    return (!(card1 & SUIT_MASK) ^ (card2 & SUIT_MASK));
    return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

returns int      SUIT\_MASK = 0x30 = **0001100000** equivalent

suit      value

12

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## Compare Card Suits

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector *v*. Here we turn all *but* the bits of interest in *v* to 0.

```
#define SUIT_MASK 0x30
int sameSuitP(char card1, char card2) {
    return (!((card1 & SUIT_MASK) ^ (card2 & SUIT_MASK)));
    //return (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

! (x^y) equivalent to x=y

13

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## Compare Card Values

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector *v*.

```
char hand[5]; // represents a 5-card hand
char card1, card2; // two cards to compare
card1 = hand[0];
card2 = hand[1];
...
if ( greaterValue(card1, card2) ) { ... }
```

```
#define VALUE_MASK 0x0F
int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```

VALUE\_MASK = 0x0F = 0000011111

suit value

14

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## Compare Card Values

mask: a bit vector designed to achieve a desired behavior when used with a bitwise operator on another bit vector *v*.

```
#define VALUE_MASK 0x0F
int greaterValue(char card1, char card2) {
    return ((unsigned int)(card1 & VALUE_MASK) >
            (unsigned int)(card2 & VALUE_MASK));
}
```

$2_{10} > 13_{10}$   
0 (false)

15

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## Integers

- ❖ Binary representation of integers
  - Unsigned and signed
  - Casting in C
- ❖ Consequences of finite width representation
  - Overflow, sign extension
- ❖ Shifting and arithmetic operations

16

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## Encoding Integers

- ❖ The hardware (and C) supports two flavors of integers
  - *unsigned* – only the non-negatives
  - *signed* – both negatives and non-negatives
- ❖ Cannot represent all integers with *w* bits
  - Only  $2^w$  distinct bit patterns
  - Unsigned values:  $0 \dots 2^w - 1$
  - Signed values:  $-2^{w-1} \dots 2^{w-1} - 1$
- ❖ Example: 8-bit integers (e.g. char)

17

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## Unsigned Integers

- ❖ Unsigned values follow the standard base 2 system
  - $b_7b_6b_5b_4b_3b_2b_1b_0 = b_72^7 + b_62^6 + \dots + b_12^1 + b_02^0$
- ❖ Add and subtract using the normal “carry” and “borrow” rules, just in binary

63	00111111
+ 8	+00001000
71	01000111

- ❖ Useful formula:  $2^{N-1} + 2^{N-2} + \dots + 2 + 1 = 2^N - 1$ 
  - i.e. *N* ones in a row =  $2^N - 1$
- ❖ How would you make *signed* integers?

18

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## Sign and Magnitude

Most Significant Bit

- Designate the high-order bit (MSB) as the "sign bit"
  - $sign=0$ : positive numbers;  $sign=1$ : negative numbers
- Benefits:
  - Using MSB as sign bit matches positive numbers with unsigned
  - All zeros encoding is still = 0
- Examples (8 bits):
  - $0x00 = 00000000_2$  is non-negative, because the sign bit is 0
  - $0x7F = 01111111_2$  is non-negative ( $+127_{10}$ )
  - $0x85 = 10000101_2$  is negative ( $-5_{10}$ )
  - $0x80 = 10000000_2$  is negative... zero???

19

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## Sign and Magnitude

- MSB is the sign bit, rest of the bits are magnitude
- Drawbacks?

20

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## Sign and Magnitude

- MSB is the sign bit, rest of the bits are magnitude
- Drawbacks:
  - Two representations of 0 (bad for checking equality)

21

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## Sign and Magnitude

- MSB is the sign bit, rest of the bits are magnitude
- Drawbacks:
  - Two representations of 0 (bad for checking equality)
  - Arithmetic is cumbersome
    - Example:  $4-3 \neq 4+(-3)$

4	0100
-3	-0011
1	0001

✓

4	0100
+ -3	+ 1011
-7	1111

✗

- Negatives "increment" in wrong direction!

22

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## Two's Complement

- Let's fix these problems:
  - "Flip" negative encodings so incrementing works

23

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## Two's Complement

- Let's fix these problems:
  - "Flip" negative encodings so incrementing works
  - "Shift" negative numbers to eliminate -0
- MSB still indicates sign!
  - This is why we represent one more negative than positive number ( $-2^{N-1}$  to  $2^{N-1}-1$ )

24

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## Two's Complement Negatives

- Accomplished with one neat mathematical trick!

$b_{w-1}$  has weight  $-2^{w-1}$ , other bits have usual weights  $+2^i$

- 4-bit Examples:
  - $1010_2$ , unsigned:  $1*2^3+0*2^2+1*2^1+0*2^0 = 10$
  - $1010_2$ , two's complement:  $-1*2^3+0*2^2+1*2^1+0*2^0 = -6$
- 1 represented as:  $1111_2 = -2^3+(2^3-1)$ 
  - MSB makes it super negative, add up all the other bits to get back up to -1

Two's Complement

25

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## Why Two's Complement is So Great

- Roughly same number of (+) and (-) numbers
- Positive number encodings match unsigned
- Single zero
- All zeros encoding = 0
- Simple negation procedure:
  - Get negative representation of any integer by taking bitwise complement and then adding one!
  - $(\sim x + 1 == -x)$

Two's Complement

26

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## Peer Instruction Question

- Take the 4-bit number encoding  $x = 0b1011$
- Which of the following numbers is NOT a valid interpretation of  $x$  using any of the number representation schemes discussed today?
  - Unsigned, Sign and Magnitude, Two's Complement
  - Vote at <http://PollEv.com/justinh>

A. -4  
 B. -5  
 C. 11  
 D. -3  
 E. We're lost...

27

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## Summary

- Bit-level operators allow for fine-grained manipulations of data
  - Bitwise AND (&), OR (|), and NOT (~) different than logical AND (&&), OR (||), and NOT (!)
  - Especially useful with bit masks
- Choice of *encoding scheme* is important
  - Tradeoffs based on size requirements and desired operations
- Integers represented using unsigned and two's complement representations
  - Limited by fixed bit width
  - We'll examine arithmetic operations next lecture

29