

The Hardware/Software Interface

CSE351 Winter 2013

Data Structures I: Arrays

Data Structures in Assembly

■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

■ Structs

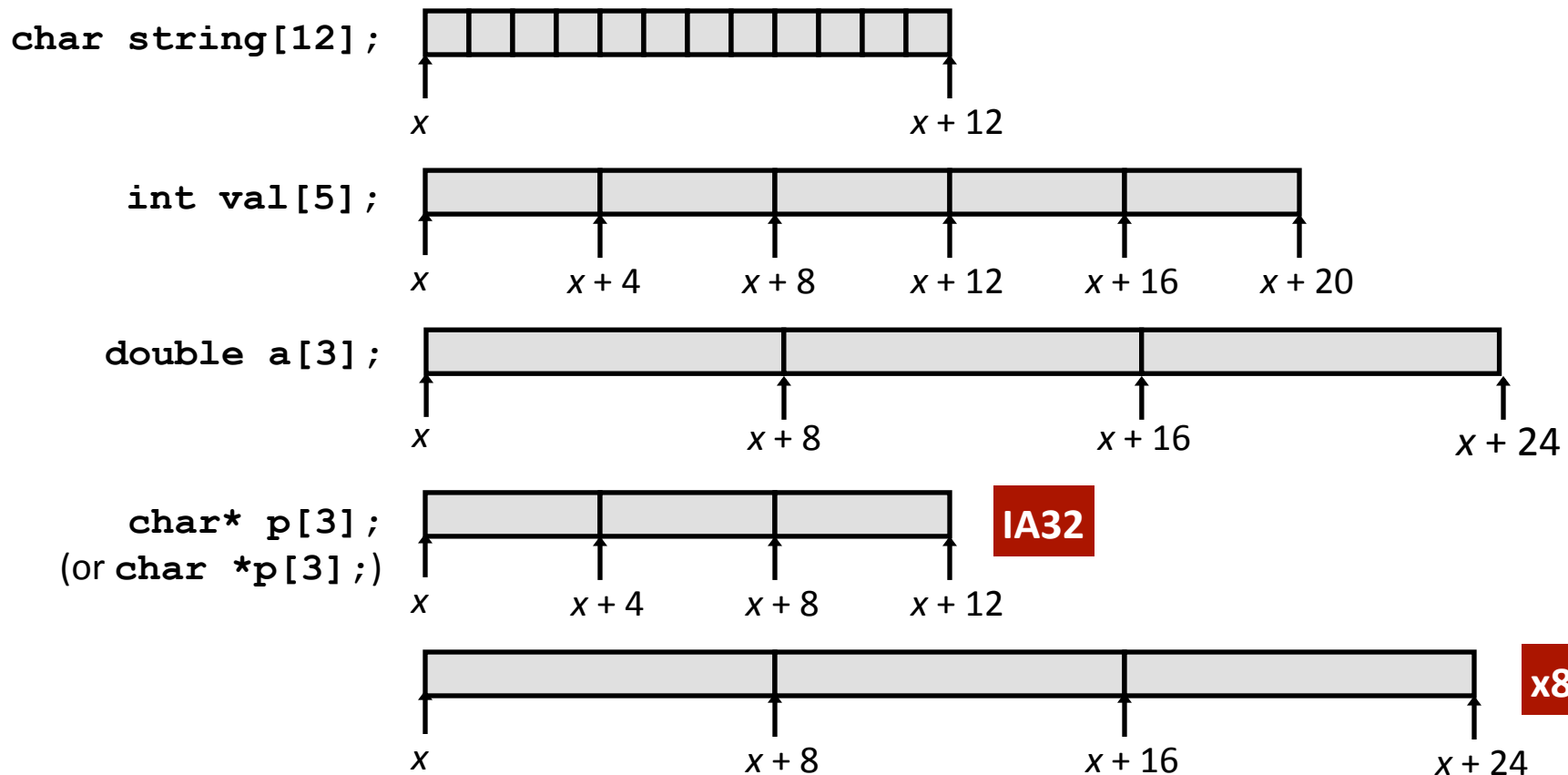
- Alignment

■ Unions

Array Allocation

■ Basic Principle

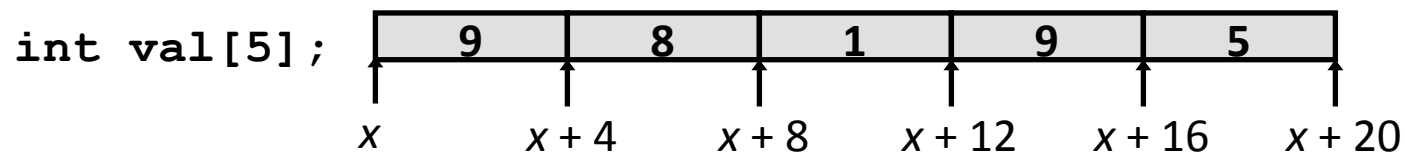
- $T \ A[N];$
- Array of data type T and length N
- *Contiguously* allocated region of $N * \text{sizeof}(T)$ bytes



Array Access

■ Basic Principle

- $T\ A[N];$
- Array of data type T and length N
- Identifier A can be used as a pointer to array element 0: Type T^*



■ Reference Type Value

- `val[4]` `int` 5
- `val` `int *` x
- `val+1` `int *` $x + 4$
- `&val[2]` `int *` $x + 8$
- `val[5]` `int` ??
- `*(val+1)` `int` 8
- `val + i` `int *` $x + 4*i$

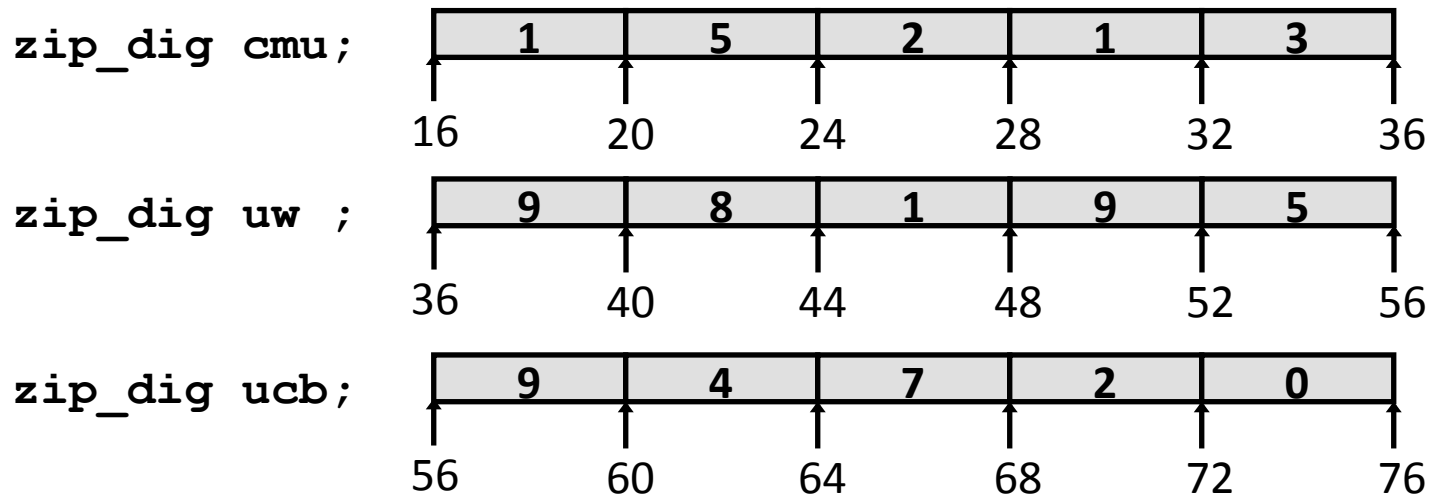
Array Example

```
typedef int zip_dig[5];  
  
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig uw  = { 9, 8, 1, 9, 5 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

Array Example

```
typedef int zip_dig[5];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig uw  = { 9, 8, 1, 9, 5 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “zip_dig uw” equivalent to “int uw[5]”
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



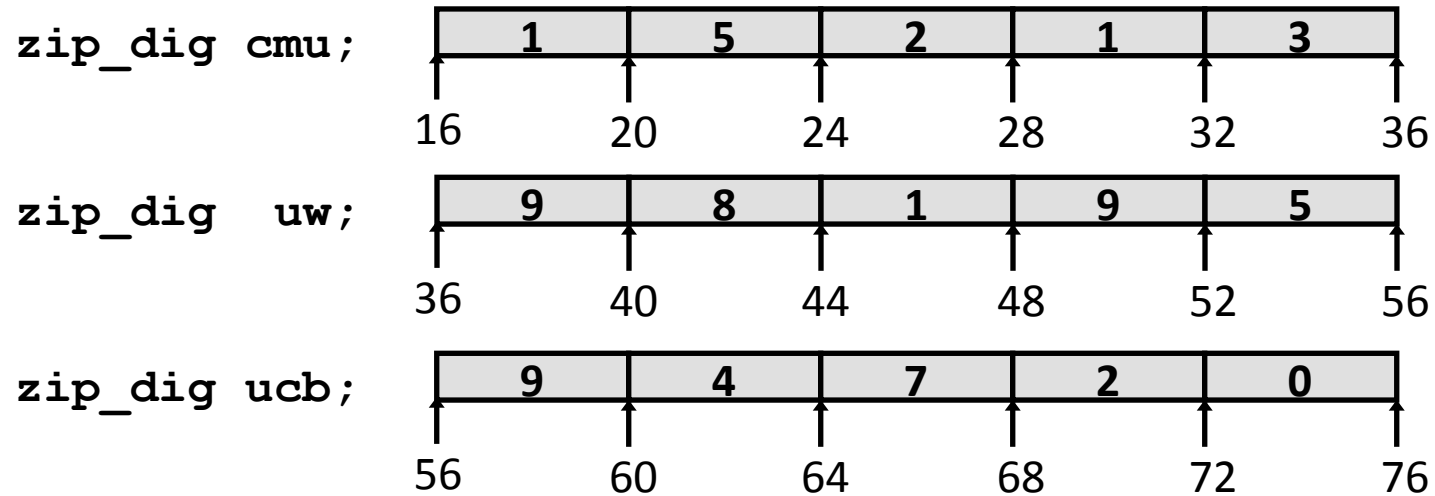
```
int get_digit
  (zip_dig z, int dig)
{
  return z[dig];
}
```

IA32

```
# %edx = z
# %eax = dig
movl (%edx,%eax,4),%eax # z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at $4 * \%eax + \%edx$
- Use memory reference $(\%edx, \%eax, 4)$

Referencing Examples



Reference	Address	Value	Guaranteed?
<code>uw[3]</code>	$36 + 4 * 3 = 48$	9	Yes
<code>uw[6]</code>	$36 + 4 * 6 = 60$	4	No
<code>uw[-1]</code>	$36 + 4 * -1 = 32$	3	No
<code>cmu[15]</code>	$16 + 4 * 15 = 76$??	No

- No bounds checking
- Location of each separate array in memory is not guaranteed

Array Loop Example

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

Array Loop Example

■ Original

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

■ Transformed

- Eliminate loop variable **i**, use pointer **zend** instead
- Convert array code to pointer code
 - Pointer arithmetic on **z**
- Express in do-while form (no test at entrance)

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while (z <= zend);
    return zi;
}
```

Array Loop Implementation (IA32)

■ Registers

```
%ecx  z
%eax  zi
%ebx  zend
```

■ Computations

- $10 * z_i + *z$ implemented as $*z + 2 * (5 * z_i)$
- $z++$ increments by 4

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

```
# %ecx = z
xorl %eax,%eax          # zi = 0
leal 16(%ecx),%ebx      # zend = z+4
.L59:
leal (%eax,%eax,4),%edx # zi + 4*zi = 5*zi
movl (%ecx),%eax       # *z
addl $4,%ecx           # z++
leal (%eax,%edx,2),%eax # zi = *z + 2*(5*zi)
cmpl %ebx,%ecx       # z : zend
jle .L59            # if <= goto loop
```

Nested Array Example

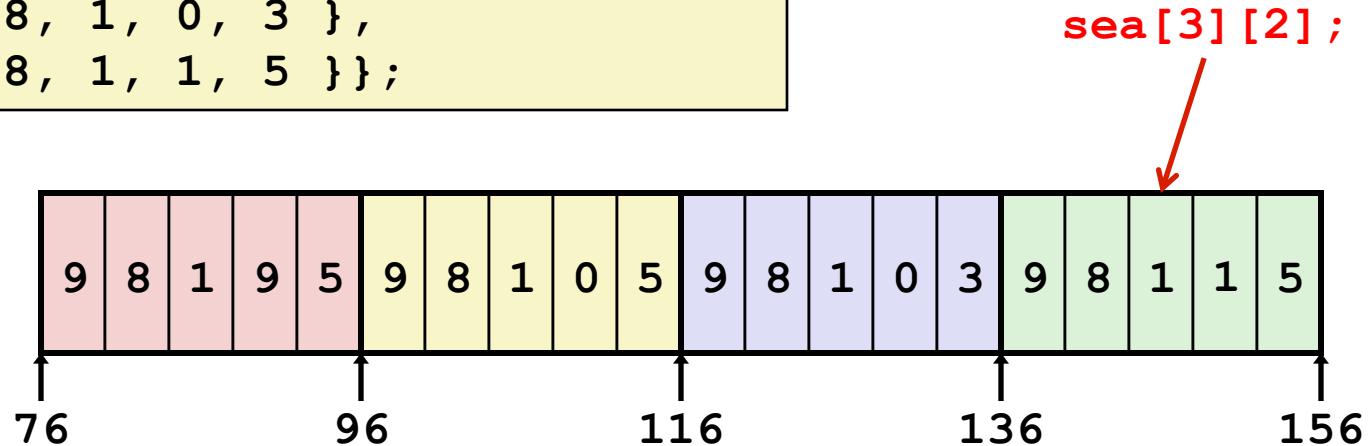
```
#define PCOUNT 4
zip_dig sea[PCOUNT] =
    {{ 9, 8, 1, 9, 5 },
     { 9, 8, 1, 0, 5 },
     { 9, 8, 1, 0, 3 },
     { 9, 8, 1, 1, 5 }};
```

Remember, $\mathbf{T} \mathbf{A}[\mathbf{N}]$ is an array with elements of type \mathbf{T} , with length \mathbf{N}

Nested Array Example

```
#define PCOUNT 4
zip_dig sea[PCOUNT] =
  { { 9, 8, 1, 9, 5 },
    { 9, 8, 1, 0, 5 },
    { 9, 8, 1, 0, 3 },
    { 9, 8, 1, 1, 5 } };
```

Remember, $\mathbf{T} \mathbf{A}[\mathbf{N}]$ is an array with elements of type \mathbf{T} , with length \mathbf{N}



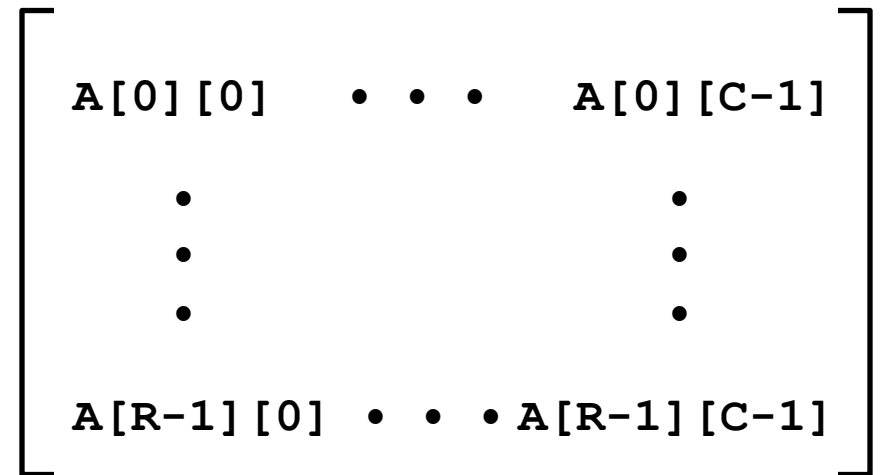
- “Row-major” ordering of all elements
- Guaranteed?

Multidimensional (Nested) Arrays

■ Declaration

- $T \ A[R][C];$
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

■ Array size?



Multidimensional (Nested) Arrays

■ Declaration

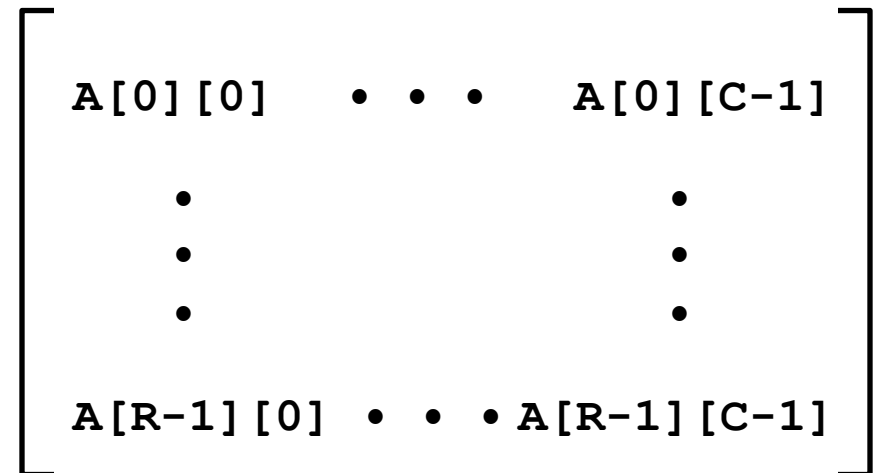
- `T A[R][C];`
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

■ Array size

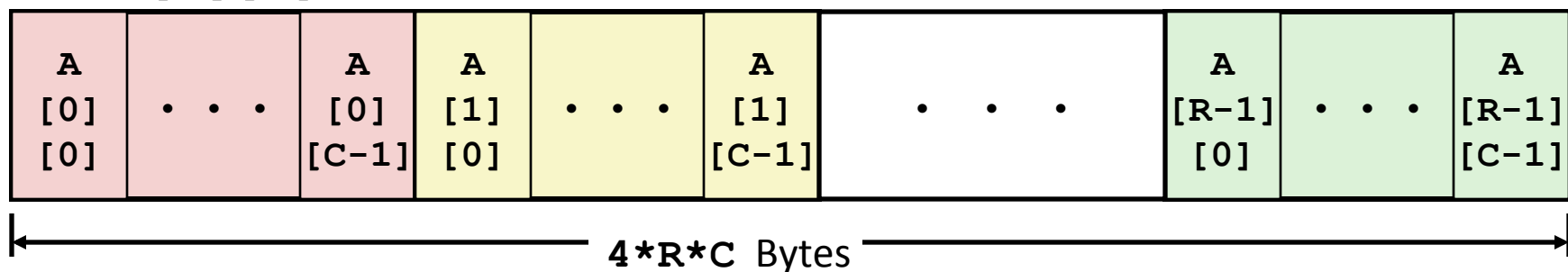
- $R * C * K$ bytes

■ Arrangement

- Row-major ordering



```
int A[R][C];
```

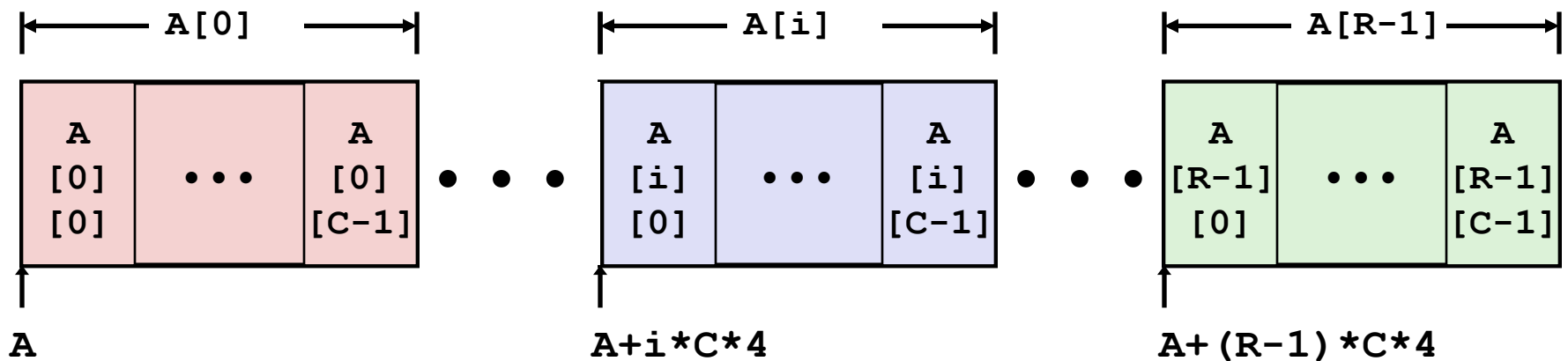


Nested Array Row Access

■ Row vectors

- $T \ A[R][C]$: $A[i]$ is array of C elements
- Each element of type T requires K bytes
- Starting address $A + i * (C * K)$

```
int A[R][C];
```



Nested Array Row Access Code

```
int *get_sea_zip(int index)
{
    return sea[index];
}
```

```
#define PCOUNT 4
zip_dig sea[PCOUNT] =
    {{ 9, 8, 1, 9, 5 },
     { 9, 8, 1, 0, 5 },
     { 9, 8, 1, 0, 3 },
     { 9, 8, 1, 1, 5 }};
```

Nested Array Row Access Code

```
int *get_sea_zip(int index)
{
    return sea[index];
}
```

```
#define PCOUNT 4
zip_dig sea[PCOUNT] =
    {{ 9, 8, 1, 9, 5 },
     { 9, 8, 1, 0, 5 },
     { 9, 8, 1, 0, 3 },
     { 9, 8, 1, 1, 5 }};
```

- What data type is sea[index]?
- What is its starting address?

```
# %eax = index
leal (%eax,%eax,4),%eax
leal sea(,%eax,4),%eax
```

Translation?

Nested Array Row Access Code

```
int *get_sea_zip(int index)
{
    return sea[index];
}
```

```
#define PCOUNT 4
zip_dig sea[PCOUNT] =
    {{ 9, 8, 1, 9, 5 },
     { 9, 8, 1, 0, 5 },
     { 9, 8, 1, 0, 3 },
     { 9, 8, 1, 1, 5 }};
```

```
# %eax = index
leal (%eax,%eax,4),%eax # 5 * index
leal sea(,%eax,4),%eax # sea + (20 * index)
```

■ Row Vector

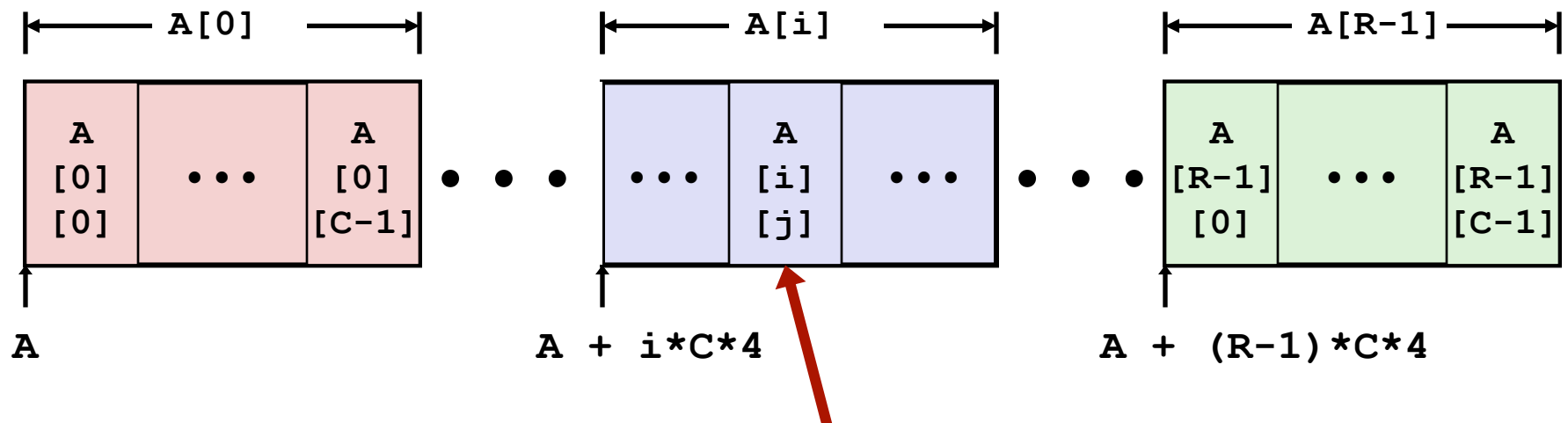
- `sea[index]` is array of 5 ints
- Starting address `sea+20*index`

■ IA32 Code

- Computes and returns address
- Compute as `sea+4*(index+4*index)=sea+20*index`

Nested Array Row Access

```
int A[R][C];
```

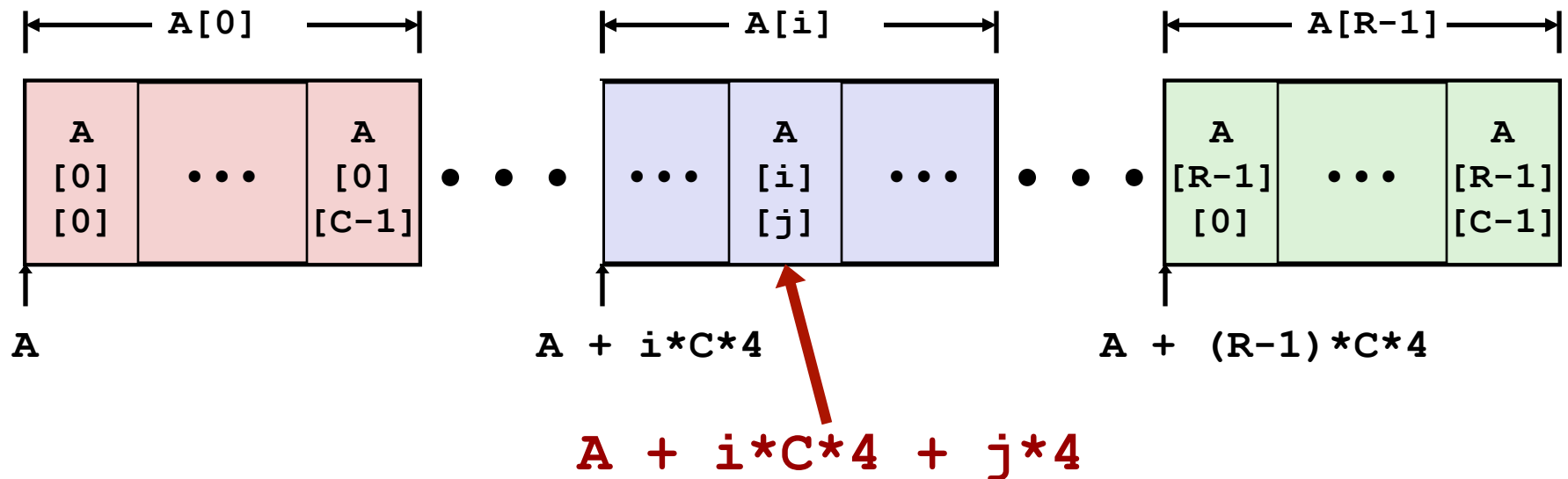


Nested Array Row Access

■ Array Elements

- $A[i][j]$ is element of type T , which requires K bytes
- Address $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



Nested Array Element Access Code

```
int get_sea_digit
  (int index, int dig)
{
  return sea[index][dig];
}
```

```
# %ecx = dig
# %eax = index
leal 0(,%ecx,4),%edx      # 4*dig
leal (%eax,%eax,4),%eax   # 5*index
movl sea(%edx,%eax,4),%eax # *(sea + 4*dig + 20*index)
```

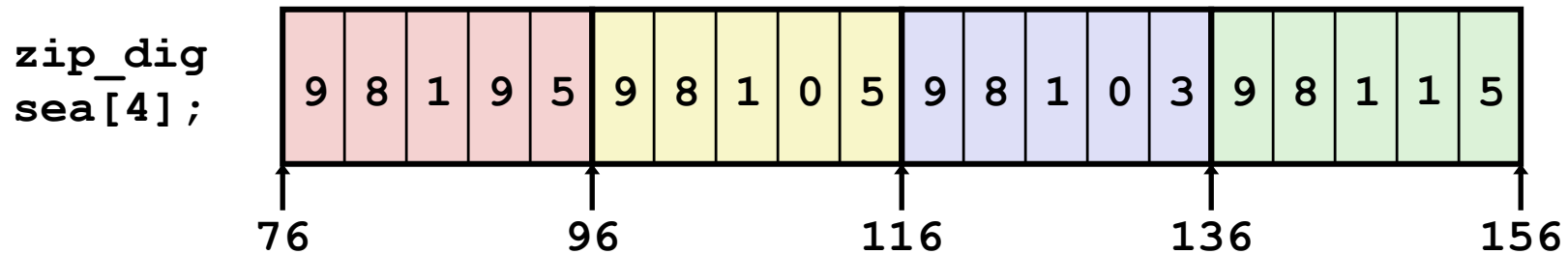
■ Array Elements

- sea[index][dig] is int
- Address: sea + 20*index + 4*dig

■ IA32 Code

- Computes address sea + 4*dig + 4*(index+4*index)
- movl performs memory reference

Strange Referencing Examples



Reference	Address	Value	Guaranteed?
sea[3][3]	$76+20*3+4*3 = 148$	1	Yes
sea[2][5]	$76+20*2+4*5 = 136$	9	Yes
sea[2][-1]	$76+20*2+4*-1 = 112$	5	Yes
sea[4][-1]	$76+20*4+4*-1 = 152$	5	Yes
sea[0][19]	$76+20*0+4*19 = 152$	5	Yes
sea[0][-1]	$76+20*0+4*-1 = 72$??	No

- Code does not do any bounds checking
- Ordering of elements within array guaranteed

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig uw  = { 9, 8, 1, 9, 5 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3  
int *univ[UCOUNT] = {uw, cmu, ucb};
```

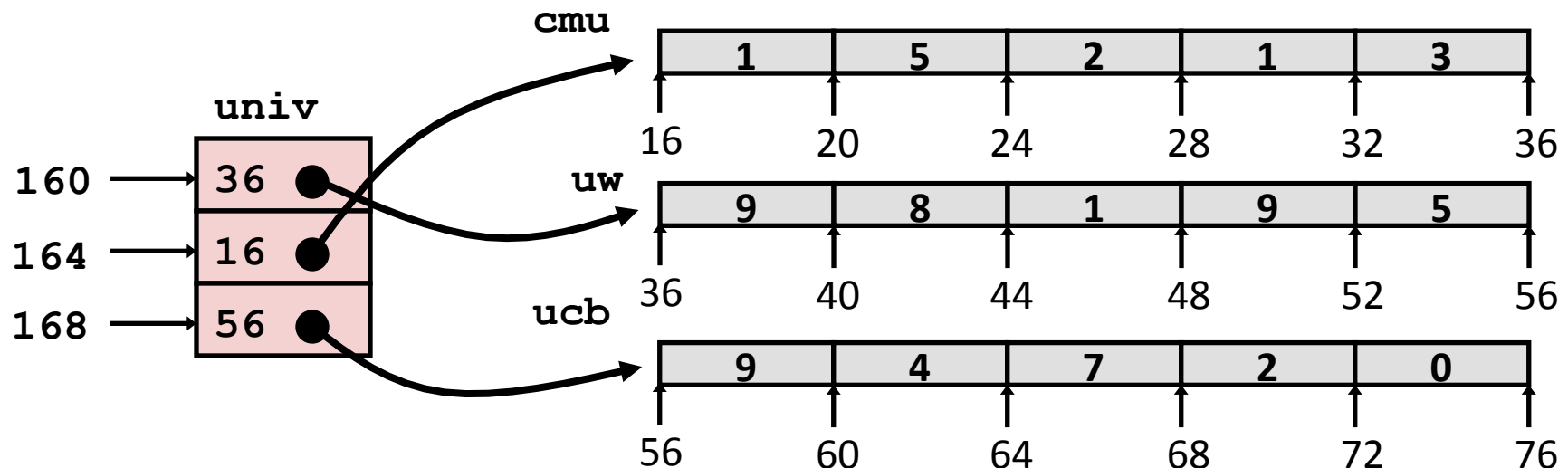
Same thing as a 2D array?

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig uw  = { 9, 8, 1, 9, 5 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3
int *univ[UCOUNT] = {uw, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
 - 4 bytes
- Each pointer points to array of ints



Note: this is how Java represents multi-dimensional arrays.

Element Access in Multi-Level Array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

```
# %ecx = index
# %eax = dig
leal 0(,%ecx,4),%edx    # 4*index
movl univ(%edx),%edx    # Mem[univ+4*index]
movl (%edx,%eax,4),%eax # Mem[...+4*dig]
```

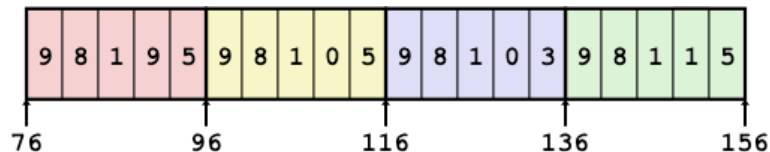
■ Computation (IA32)

- Element access `Mem[Mem[univ+4*index]+4*dig]`
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

Array Element Accesses

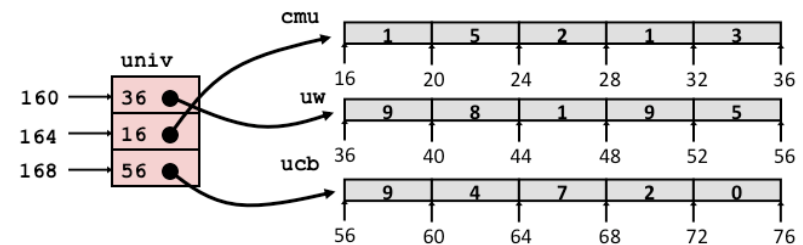
Nested array

```
int get_sea_digit
(int index, int dig)
{
    return sea[index][dig];
}
```



Multi-level array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

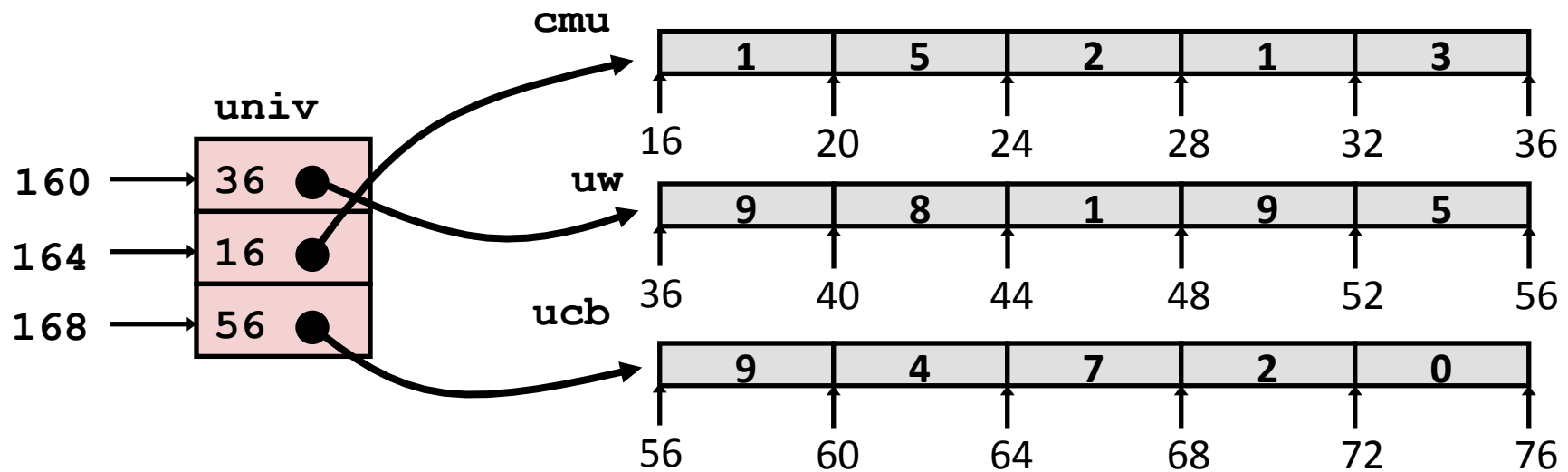


Access looks similar, but it isn't:

$\text{Mem}[\text{sea} + 20 * \text{index} + 4 * \text{dig}]$

$\text{Mem}[\text{Mem}[\text{univ} + 4 * \text{index}] + 4 * \text{dig}]$

Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>univ[2][3]</code>	$56+4*3 = 68$	2	Yes
<code>univ[1][5]</code>	$16+4*5 = 36$	9	No
<code>univ[2][-1]</code>	$56+4*-1 = 52$	5	No
<code>univ[3][-1]</code>	??	??	No
<code>univ[1][12]</code>	$16+4*12 = 64$	7	No

- Code does not do any bounds checking
- Location of each lower-level array in memory is not guaranteed

Using Nested Arrays

```
#define N 16
typedef int fix_matrix[N][N];

/* Compute element i,k of
   fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
 int i, int k)
{
    int j;
    int result = 0;
    for (j = 0; j < N; j++)
        result += a[i][j]*b[j][k];
    return result;
}
```

Using Nested Arrays

■ Strengths

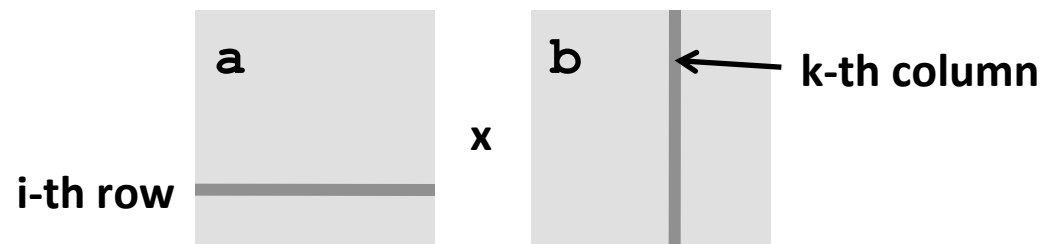
- Generates very efficient assembly code
- Avoids multiply in index computation

■ Limitation

- Only works for fixed array size

```
#define N 16
typedef int fix_matrix[N][N];
```

```
/* Compute element i,k of
   fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
 int i, int k)
{
    int j;
    int result = 0;
    for (j = 0; j < N; j++)
        result += a[i][j]*b[j][k];
    return result;
}
```



Dynamic Nested Arrays

■ Strength

- Can create matrix of any size

■ Programming

- Must do index computation explicitly

■ Performance

- Accessing single element costly
- Must do multiplication

```
int * new_var_matrix(int n)
{
    return (int *)
        calloc(sizeof(int), n*n);
}
```

```
int var_ele
(int *a, int i, int j, int n)
{
    return a[i*n+j];
}
```

```
movl 12(%ebp),%eax      # i
movl 8(%ebp),%edx       # a
imull 20(%ebp),%eax     # n*i
addl 16(%ebp),%eax     # n*i+j
movl (%edx,%eax,4),%eax # Mem[a+4*(i*n+j)]
```

Arrays in C

- **Contiguous allocations of memory**
- **No bounds checking**
- **Can usually be treated like a pointer to first element**