The Hardware/Software Interface

CSE351 Winter 2013

x86 Programming I

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
     c.getMPG();
```

Assembly language:

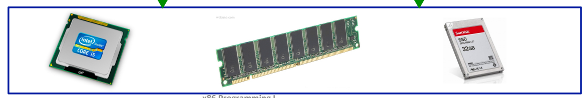
```
get_mpg:
    pushq   %rbp
    movq    %rsp, %rbp
    ...
    popq    %rbp
    ret
```

Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

OS:

Computer system:



- Data & addressing
- Integers & floats
- Machine code & C
- **x86 assembly programming**
- Procedures & stacks
- Arrays & structs
- Memory & caches
- Processes
- Virtual memory
- Memory allocation
- Java vs. C

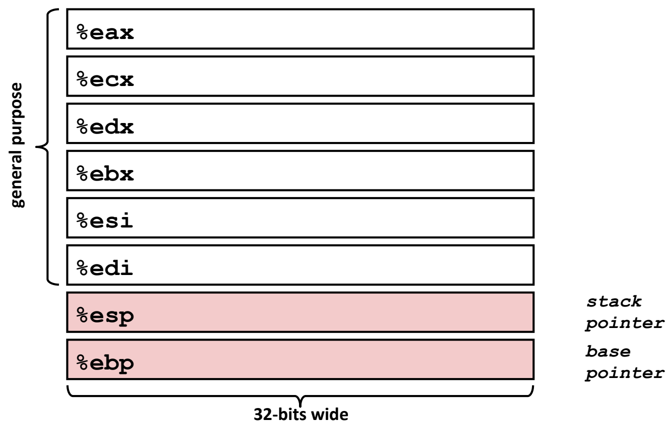
Today

- **Move instructions, registers, and operands**
- **Memory addressing modes**
- **`swap` example: 32-bit vs. 64-bit**
- **Arithmetic operations**

Why learn assembly language?

- **Not to be able to write programs directly in assembly**
  - Compilers do that for you
- **But to be able to *understand* the code generated by compilers, so that you can then:**
  - Optimize performance of critical sections of code
  - Investigate unexpected or even buggy behavior
  - Understand how security vulnerabilities arise, and how to protect against them

Integer Registers (IA32)



Three Basic Kinds of Instructions

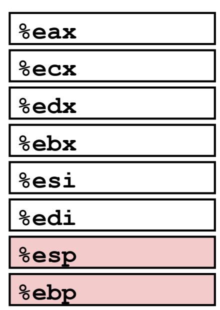
- **Transfer data between memory and register**
  - *Load* data from memory into register
    - %reg = Mem[address]
  - *Store* register data into memory
    - Mem[address] = %reg

Remember: memory is indexed just like an array[]!

- **Perform arithmetic function on register or memory data**
  - c = a + b;
- **Transfer control**
  - Unconditional jumps to/from procedures
  - Conditional branches

Moving Data: IA32

- **Moving Data**
  - **`movx`** ***Source, Dest***
  - **`x`** is one of {**`b, w, l`**}
  - **`movl`** ***Source, Dest***:
    Move 4-byte "long word"
  - **`movw`** ***Source, Dest***:
    Move 2-byte "word"
  - **`movb`** ***Source, Dest***:
    Move 1-byte "byte"
- **Lots of these in typical code**



Moving Data: IA32

- **Moving Data**
  `movl` ***Source, Dest***:
- **Operand Types**
  - ***Immediate:*** Constant integer data
    - Example: **`$0x400, $-533`**
    - Like C constant, but prefixed with **`'$'`**
    - Encoded with 1, 2, or 4 bytes
  - ***Register:*** One of 8 integer registers
    - Example: **`%eax, %edx`**
    - But **`%esp`** and **`%ebp`** reserved for special use
    - Others have special uses for particular instructions
  - ***Memory:*** 4 consecutive bytes of memory at address given by register
    - Simplest example: **`(%eax)`**
    - Various other "address modes"

movl Operand Combinations

	Source	Dest	Src,Dest	C Analog
movl	*Imm*	*Reg*	`movl $0x4,%eax`	`var_a = 0x4;`
		Mem	`movl $-147,(%eax)`	`*p_a = -147;`
	Reg	*Reg*	`movl %eax,%edx`	`var_d = var_a;`
		Mem	`movl %eax,(%edx)`	`*p_d = var_a;`
	Mem	*Reg*	`movl (%eax),%edx`	`var_d = *p_a;`

Cannot do memory-memory transfer with a single instruction.

Memory vs. registers

- **Why both?**

Memory Addressing Modes: Basic

- **Indirect (R) Mem[Reg[R]]**
  - Register R specifies the memory address

```
movl (%ecx),%eax
```

- **Displacement D(R) Mem[Reg[R]+D]**
  - Register R specifies a memory address
    - (e.g. the start of some memory region)
  - Constant displacement D specifies the offset from that address

```
movl 8(%ebp),%edx
```


Using Basic Addressing Modes

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
   pushl %ebp
   movl  %esp,%ebp
   pushl %ebx

   movl 12(%ebp),%ecx
   movl 8(%ebp),%edx
   movl (%ecx),%eax
   movl (%edx),%ebx
   movl %eax,(%edx)
   movl %ebx,(%ecx)

   movl -4(%ebp),%ebx
   movl %ebp,%esp
   popl %ebp
   ret
```

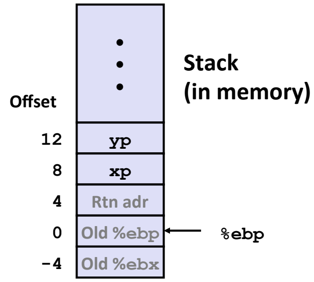
Set Up

Body

Finish

Understanding Swap

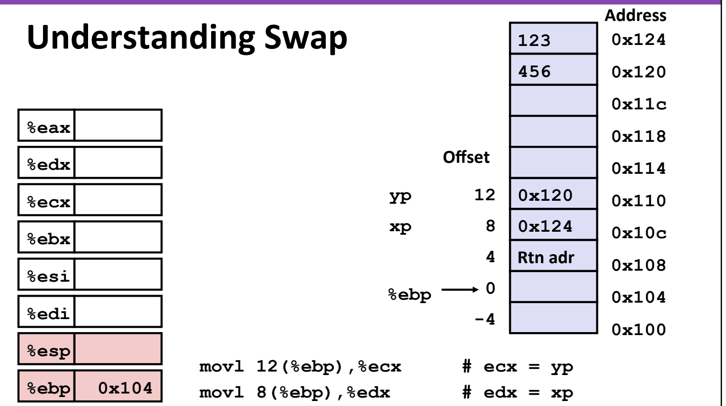
```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```



Register	Value
%ecx	yp
%edx	xp
%eax	t1
%ebx	t0

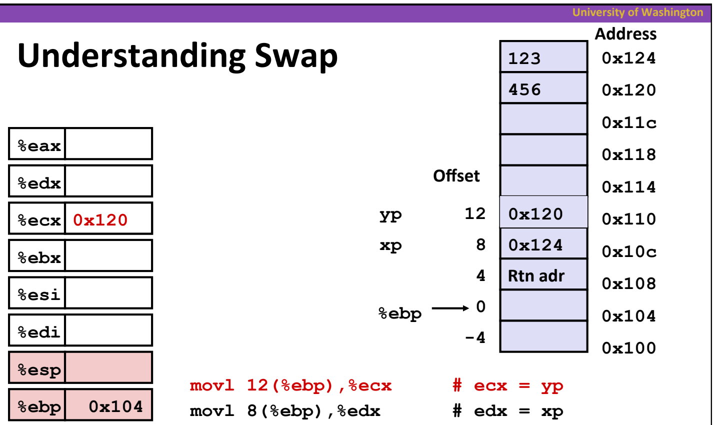
```
movl 12(%ebp),%ecx   # ecx = yp
movl 8(%ebp),%edx    # edx = xp
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```


Understanding Swap




```
movl 12(%ebp),%ecx   # ecx = yp
movl 8(%ebp),%edx    # edx = xp
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```


Understanding Swap




```
movl 12(%ebp),%ecx   # ecx = yp
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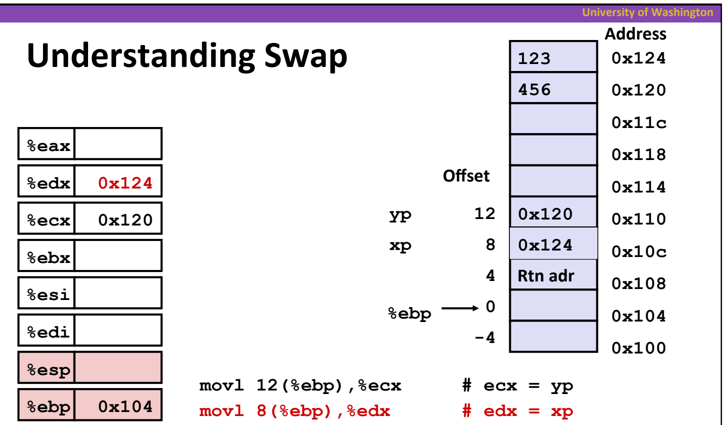
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```


Understanding Swap




```
movl 12(%ebp),%ecx   # ecx = yp
movl 8(%ebp),%edx    # edx = xp
movl (%ecx),%eax     # eax = *yp (t1)
movl (%edx),%ebx     # ebx = *xp (t0)
movl %eax,(%edx)     # *xp = eax
movl %ebx,(%ecx)     # *yp = ebx
```

Understanding Swap

Register	Value
%eax	456
%edx	0x124
%ecx	0x120
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

	Offset		Address
		123	0x124
		456	0x120
			0x11c
			0x118
			0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp →	0		0x104
	-4		0x100

```
movl 12(%ebp),%ecx     # ecx = yp
movl 8(%ebp),%edx      # edx = xp
movl (%ecx),%eax       # eax = *yp (t1)
movl (%edx),%ebx       # ebx = *xp (t0)
movl %eax,(%edx)       # *xp = eax
movl %ebx,(%ecx)       # *yp = ebx
```


Understanding Swap

Register	Value
%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

	Offset		Address
		123	0x124
		456	0x120
			0x11c
			0x118
			0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp →	0		0x104
	-4		0x100

```
movl 12(%ebp),%ecx     # ecx = yp
movl 8(%ebp),%edx      # edx = xp
movl (%ecx),%eax       # eax = *yp (t1)
movl (%edx),%ebx       # ebx = *xp (t0)
movl %eax,(%edx)       # *xp = eax
movl %ebx,(%ecx)       # *yp = ebx
```


Understanding Swap

Register	Value
%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

	Offset		Address
		456	0x124
		456	0x120
			0x11c
			0x118
			0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp →	0		0x104
	-4		0x100

```
movl 12(%ebp),%ecx     # ecx = yp
movl 8(%ebp),%edx      # edx = xp
movl (%ecx),%eax       # eax = *yp (t1)
movl (%edx),%ebx       # ebx = *xp (t0)
movl %eax,(%edx)       # *xp = eax
movl %ebx,(%ecx)       # *yp = ebx
```


Understanding Swap

Register	Value
%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

	Offset		Address
		456	0x124
		123	0x120
			0x11c
			0x118
			0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp →	0		0x104
	-4		0x100

```
movl 12(%ebp),%ecx     # ecx = yp
movl 8(%ebp),%edx      # edx = xp
movl (%ecx),%eax       # eax = *yp (t1)
movl (%edx),%ebx       # ebx = *xp (t0)
movl %eax,(%edx)       # *xp = eax
movl %ebx,(%ecx)       # *yp = ebx
```

x86-64 Integer Registers

64-bits wide

%rax	%eax	%r8	%r8d
%rbx	%ebx	%r9	%r9d
%rcx	%ecx	%r10	%r10d
%rdx	%edx	%r11	%r11d
%rsi	%esi	%r12	%r12d
%rdi	%edi	%r13	%r13d
%rsp	%esp	%r14	%r14d
%rbp	%ebp	%r15	%r15d

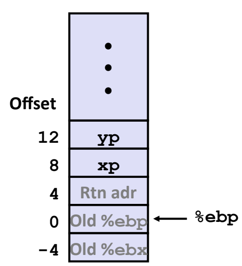
- Extend existing registers, and add 8 new ones; *all* accessible as 8, 16, 32, 64 bits.

32-bit vs. 64-bit operands

- **Long word `l` (4 Bytes) ↔ Quad word `q` (8 Bytes)**
- **New instruction forms:**
  - `movl` → `movq`
  - `addl` → `addq`
  - `sall` → `salq`
  - etc.
- **x86-64 can still use 32-bit instructions that generate 32-bit results**
  - Higher-order bits of destination register are just set to 0
  - Example: `addl`

Swap Ints in 32-bit Mode

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```




```
swap:
   pushl %ebp
   movl  %esp,%ebp       } Setup
   pushl %ebx

   movl 12(%ebp),%ecx
   movl 8(%ebp),%edx
   movl (%ecx),%eax      } Body
   movl (%edx),%ebx
   movl %eax,(%edx)
   movl %ebx,(%ecx)

   movl -4(%ebp),%ebx
   movl %ebp,%esp        } Finish
   popl %ebp
   ret
```


Swap Ints in 64-bit Mode

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
   movl    (%rdi), %edx
   movl    (%rsi), %eax
   movl    %eax, (%rdi)
   movl    %edx, (%rsi)
   retq
```

- **Arguments passed in registers (why useful?)**
  - First (**xp**) in `%rdi`, second (**yp**) in `%rsi`
  - 64-bit pointers
- **No stack operations required**
- **32-bit data**
  - Data held in registers `%eax` and `%edx`
  - `movl` operation (the `l` refers to data width, not address width)

Swap Long Ints in 64-bit Mode

```
void swap_l
  (long int *xp, long int *yp)
{
  long int t0 = *xp;
  long int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap_l:
    movq    (%rdi), %rdx
    movq    (%rsi), %rax
    movq    %rax, (%rdi)
    movq    %rdx, (%rsi)
    retq
```

- 64-bit data
  - Data held in registers `%rax` and `%rdx`
  - `movq` operation
  - "q" stands for quad-word

Complete Memory Addressing Modes

- **Remember, the addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways**
- **Most General Form:**

  D(Rb,Ri,S) Mem[Reg[Rb] + S*Reg[Ri] + D]
  - D: Constant "displacement" 1, 2, or 4 bytes
  - Rb: Base register: Any of the 8/16 integer registers
  - Ri: Index register: Any, except for `%esp` or `%rsp`
    - Unlikely you'd use `%ebp`, either
  - S: Scale: 1, 2, 4, or 8 (*why these numbers?*)
- **Special Cases: can use any combination of D, Rb, Ri and S**

  (Rb,Ri) Mem[Reg[Rb]+Reg[Ri]]

  D(Rb,Ri) Mem[Reg[Rb]+Reg[Ri]+D]

  (Rb,Ri,S) Mem[Reg[Rb]+S*Reg[Ri]]

Address Computation Examples

%edx	0xf000
%ecx	0x100

(Rb,Ri) Mem[Reg[Rb]+Reg[Ri]]
D(,Ri,S) Mem[S*Reg[Ri]+D]
(Rb,Ri,S) Mem[Reg[Rb]+S*Reg[Ri]]
D(Rb) Mem[Reg[Rb] +D]

Expression	Address Computation	Address
0x8(%edx)	0xf000 + 0x8	0xf008
(%edx,%ecx)	0xf000 + 0x100	0xf100
(%edx,%ecx,4)	0xf000 + 4*0x100	0xf400
0x80(,%edx,2)	2*0xf000 + 0x80	0x1e080

Address Computation Instruction

- `leal` ***Src,Dest***
  - *Src* is address mode expression
  - Set *Dest* to address computed by expression
    - (lea stands for *load effective address)*
  - Example: `leal (%edx,%ecx,4), %eax`
- **Uses**
  - Computing addresses without a memory reference
    - E.g., translation of `p = &x[i];`
  - Computing arithmetic expressions of the form x + k*i
    - k = 1, 2, 4, or 8

Some Arithmetic Operations

- **Two Operand (Binary) Instructions:**

Format		*Computation*	
`addl`	*Src,Dest*	*Dest = Dest + Src*	
`subl`	*Src,Dest*	*Dest = Dest − Src*	
`imull`	*Src,Dest*	*Dest = Dest * Src*	
`sall`	*Src,Dest*	*Dest = Dest << Src*	*Also called shll*
`sarl`	*Src,Dest*	*Dest = Dest >> Src*	*Arithmetic*
`shrl`	*Src,Dest*	*Dest = Dest >> Src*	*Logical*
`xorl`	*Src,Dest*	*Dest = Dest ^ Src*	
`andl`	*Src,Dest*	*Dest = Dest & Src*	
`orl`	*Src,Dest*	*Dest = Dest | Src*	

- **Watch out for argument order! (especially `subl`)**
- **No distinction between signed and unsigned int (why?)**

Some Arithmetic Operations

- **One Operand (Unary) Instructions**

`incl`	*Dest*	*Dest = Dest + 1*
`decl`	*Dest*	*Dest = Dest − 1*
`negl`	*Dest*	*Dest = −Dest*
`notl`	*Dest*	*Dest = ~Dest*

- **See textbook section 3.5.5 for more instructions: `mull`, `cltd`, `idivl`, `divl`**

Using `leal` for Arithmetic Expressions

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

```
arith:
   pushl %ebp
   movl %esp,%ebp

   movl 8(%ebp),%eax
   movl 12(%ebp),%edx
   leal (%edx,%eax),%ecx
   leal (%edx,%edx,2),%edx
   sall $4,%edx
   addl 16(%ebp),%ecx
   leal 4(%edx,%eax),%eax
   imull %ecx,%eax

   movl %ebp,%esp
   popl %ebp
   ret
```

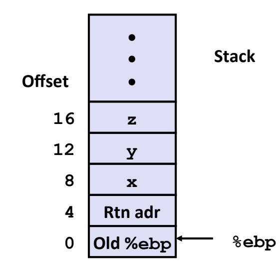
Set Up

Body

Finish

Understanding `arith`

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```




```
movl 8(%ebp),%eax          # eax = x
movl 12(%ebp),%edx         # edx = y
leal (%edx,%eax),%ecx      # ecx = x+y  (t1)
leal (%edx,%edx,2),%edx    # edx = y + 2*y = 3*y
sall $4,%edx               # edx = 48*y (t4)
addl 16(%ebp),%ecx         # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax     # eax = 4+t4+x (t5)
imull %ecx,%eax            # eax = t5*t2 (rval)
```

Understanding `arith`

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
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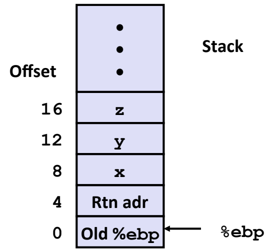
  return rval;
}
```




```
movl 8(%ebp),%eax          # eax = x
movl 12(%ebp),%edx         # edx = y
leal (%edx,%eax),%ecx      # ecx = x+y   (t1)
leal (%edx,%edx,2),%edx    # edx = y + 2*y = 3*y
sall $4,%edx               # edx = 48*y (t4)
addl 16(%ebp),%ecx         # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax     # eax = 4+t4+x (t5)
imull %ecx,%eax            # eax = t5*t2 (rval)
```


Understanding `arith`

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
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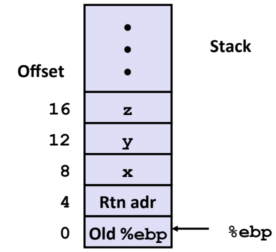
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```




```
movl 8(%ebp),%eax          # eax = x
movl 12(%ebp),%edx         # edx = y
leal (%edx,%eax),%ecx      # ecx = x+y   (t1)
leal (%edx,%edx,2),%edx    # edx = y + 2*y = 3*y
sall $4,%edx               # edx = 48*y (t4)
addl 16(%ebp),%ecx         # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax     # eax = 4+t4+x (t5)
imull %ecx,%eax            # eax = t5*t2 (rval)
```


Understanding `arith`

```
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  (int x, int y, int z)
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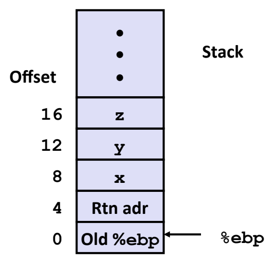
  int rval = t2 * t5;
  return rval;
}
```




```
movl 8(%ebp),%eax          # eax = x
movl 12(%ebp),%edx         # edx = y
leal (%edx,%eax),%ecx      # ecx = x+y   (t1)
leal (%edx,%edx,2),%edx    # edx = y + 2*y = 3*y
sall $4,%edx               # edx = 48*y (t4)
addl 16(%ebp),%ecx         # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax     # eax = 4+t4+x (t5)
imull %ecx,%eax            # eax = t5*t2 (rval)
```


Observations about `arith`

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

- Instructions in different order from C code
- Some expressions require multiple instructions
- Some instructions cover multiple expressions
- Get exact same code when compile:
- `(x+y+z)*(x+4+48*y)`

```
movl 8(%ebp),%eax          # eax = x
movl 12(%ebp),%edx         # edx = y
leal (%edx,%eax),%ecx      # ecx = x+y   (t1)
leal (%edx,%edx,2),%edx    # edx = y + 2*y = 3*y
sall $4,%edx               # edx = 48*y (t4)
addl 16(%ebp),%ecx         # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax     # eax = 4+t4+x (t5)
imull %ecx,%eax            # eax = t5*t2 (rval)
```

Another Example

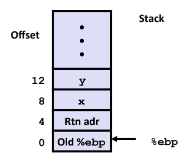
```
int logical(int x, int y)
{
  int t1 = x^y;
  int t2 = t1 >> 17;
  int mask = (1<<13) - 7;
  int rval = t2 & mask;
  return rval;
}
```

```
logical:
    pushl %ebp              Set
    movl %esp,%ebp          Up

    movl 8(%ebp),%eax
    xorl 12(%ebp),%eax
    sarl $17,%eax
    andl $8185,%eax         Body

    movl %ebp,%esp
    popl %ebp               Finish
    ret
```

```
movl 8(%ebp),%eax      # eax = x
xorl 12(%ebp),%eax     # eax = x^y
sarl $17,%eax          # eax = t1>>17
andl $8185,%eax        # eax = t2 & 8185
```



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    popl %ebp               Finish
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```

```
movl 8(%ebp),%eax      eax = x
xorl 12(%ebp),%eax     eax = x^y      (t1)
sarl $17,%eax          eax = t1>>17   (t2)
andl $8185,%eax        eax = t2 & 8185
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}
```

$2^{13} = 8192$, $2^{13} - 7 = 8185$

...0010000000000000, ...0001111111111001

```
logical:
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    movl 8(%ebp),%eax
    xorl 12(%ebp),%eax
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```
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