

The Hardware/Software Interface

CSE351 Winter 2013

x86 Programming I

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

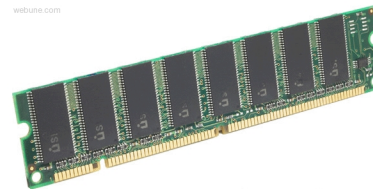
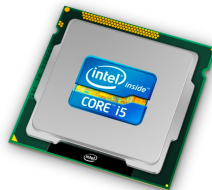
Assembly language:

```
get_mpg:
    pushq    %rbp
    movq    %rsp, %rbp
    ...
    popq    %rbp
    ret
```

Machine code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

Computer system:



Data & addressing
 Integers & floats
 Machine code & C
x86 assembly programming
 Procedures & stacks
 Arrays & structs
 Memory & caches
 Processes
 Virtual memory
 Memory allocation
 Java vs. C

OS:



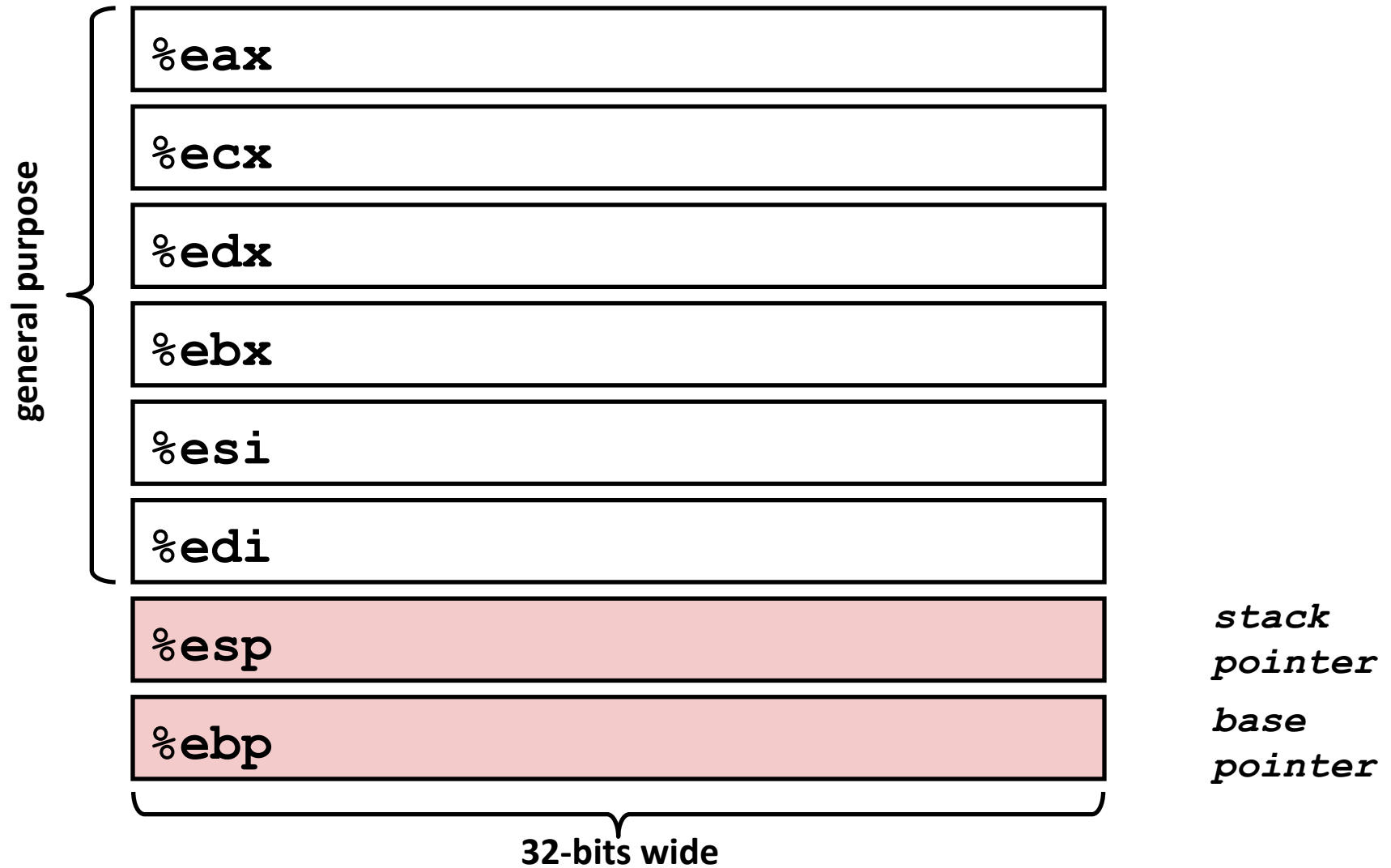
Today

- Move instructions, registers, and operands
- Memory addressing modes
- swap example: 32-bit vs. 64-bit
- Arithmetic operations

Why learn assembly language?

- **Not to be able to write programs directly in assembly**
 - Compilers do that for you
- **But to be able to *understand* the code generated by compilers, so that you can then:**
 - Optimize performance of critical sections of code
 - Investigate unexpected or even buggy behavior
 - Understand how security vulnerabilities arise, and how to protect against them

Integer Registers (IA32)



Three Basic Kinds of Instructions

■ Transfer data between memory and register

- *Load* data from memory into register
 - $\%reg = Mem[address]$
- *Store* register data into memory
 - $Mem[address] = \%reg$

Remember:
memory is indexed
just like an array[]!

■ Perform arithmetic function on register or memory data

- $c = a + b;$

■ Transfer control

- Unconditional jumps to/from procedures
- Conditional branches

Moving Data: IA32

■ Moving Data

- `movx Source, Dest`
- `x` is one of {`b`, `w`, `l`}

- `movl Source, Dest`:
Move 4-byte “long word”
- `movw Source, Dest`:
Move 2-byte “word”
- `movb Source, Dest`:
Move 1-byte “byte”

■ Lots of these in typical code

`%eax`

`%ecx`

`%edx`

`%ebx`

`%esi`

`%edi`

`%esp`

`%ebp`

Moving Data: IA32

■ Moving Data

`movl Source, Dest:`

■ Operand Types

- **Immediate:** Constant integer data
 - Example: `$0x400`, `$-533`
 - Like C constant, but prefixed with ``$'`
 - Encoded with 1, 2, or 4 bytes
- **Register:** One of 8 integer registers
 - Example: `%eax`, `%edx`
 - But `%esp` and `%ebp` reserved for special use
 - Others have special uses for particular instructions
- **Memory:** 4 consecutive bytes of memory at address given by register
 - Simplest example: `(%eax)`
 - Various other “address modes”

<code>%eax</code>
<code>%ecx</code>
<code>%edx</code>
<code>%ebx</code>
<code>%esi</code>
<code>%edi</code>
<code>%esp</code>
<code>%ebp</code>

movl Operand Combinations

	Source	Dest	Src, Dest	C Analog
movl	Imm	Reg	movl \$0x4, %eax	var_a = 0x4;
		Mem	movl \$-147, (%eax)	*p_a = -147;
	Reg	Reg	movl %eax, %edx	var_d = var_a;
		Mem	movl %eax, (%edx)	*p_d = var_a;
	Mem	Reg	movl (%eax), %edx	var_d = *p_a;

Cannot do memory-memory transfer with a single instruction.

Memory vs. registers

- Why both?

Memory Addressing Modes: Basic

■ Indirect (R) Mem[Reg[R]]

- Register R specifies the memory address

```
movl (%ecx), %eax
```

■ Displacement D(R) Mem[Reg[R]+D]

- Register R specifies a memory address
 - (e.g. the start of some memory region)
- Constant displacement D specifies the offset from that address

```
movl 8(%ebp), %edx
```

Using Basic Addressing Modes

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

```
pushl %ebp
movl  %esp, %ebp
pushl %ebx
```

} Set
Up

```
movl 12(%ebp), %ecx
movl 8(%ebp), %edx
movl (%ecx), %eax
movl (%edx), %ebx
movl %eax, (%edx)
movl %ebx, (%ecx)
```

} Body

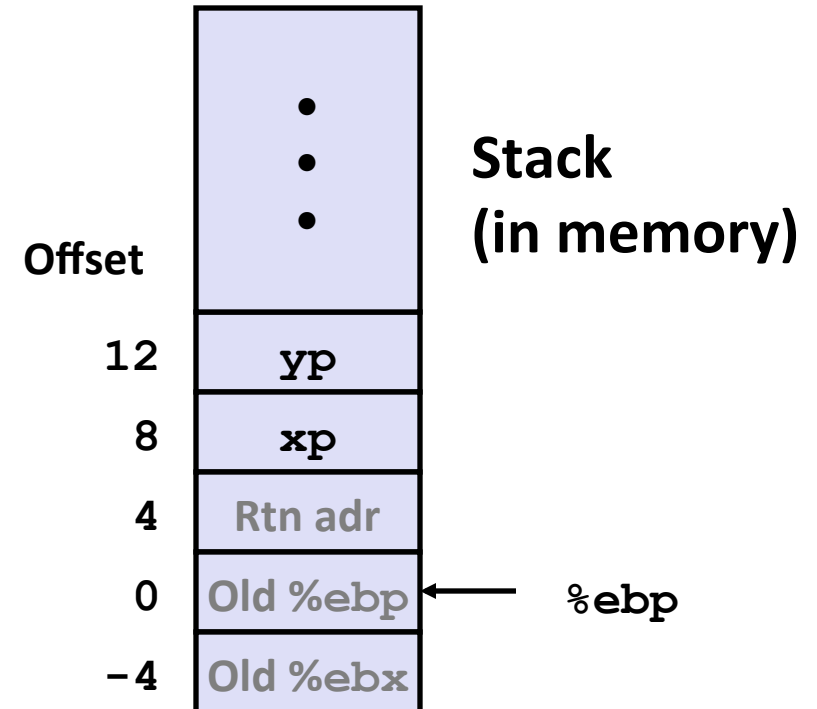
```
movl -4(%ebp), %ebx
movl %ebp, %esp
popl %ebp
ret
```

} Finish

Understanding Swap

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

Register	Value
<code>%ecx</code>	<code>yp</code>
<code>%edx</code>	<code>xp</code>
<code>%eax</code>	<code>t1</code>
<code>%ebx</code>	<code>t0</code>



```
movl 12(%ebp), %ecx # ecx = yp
movl 8(%ebp), %edx  # edx = xp
movl (%ecx), %eax   # eax = *yp (t1)
movl (%edx), %ebx   # ebx = *xp (t0)
movl %eax, (%edx)   # *xp = eax
movl %ebx, (%ecx)   # *yp = ebx
```

Understanding Swap

%eax	
%edx	
%ecx	
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

		Offset	Address
			123
			456
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
	0		0x104
%ebp	→ 0		0x104
	-4		0x100

```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx       # edx = xp
movl (%ecx), %eax        # eax = *yp (t1)
movl (%edx), %ebx        # ebx = *xp (t0)
movl %eax, (%edx)        # *xp = eax
movl %ebx, (%ecx)        # *yp = ebx

```

Understanding Swap

<code>%eax</code>	
<code>%edx</code>	
<code>%ecx</code>	0x120
<code>%ebx</code>	
<code>%esi</code>	
<code>%edi</code>	
<code>%esp</code>	
<code>%ebp</code>	0x104

		Address	
		123	0x124
		456	0x120
			0x11c
			0x118
	Offset		0x114
<code>yp</code>	12	0x120	0x110
<code>xp</code>	8	0x124	0x10c
	4	Rtn adr	0x108
<code>%ebp</code>	→ 0		0x104
	-4		0x100

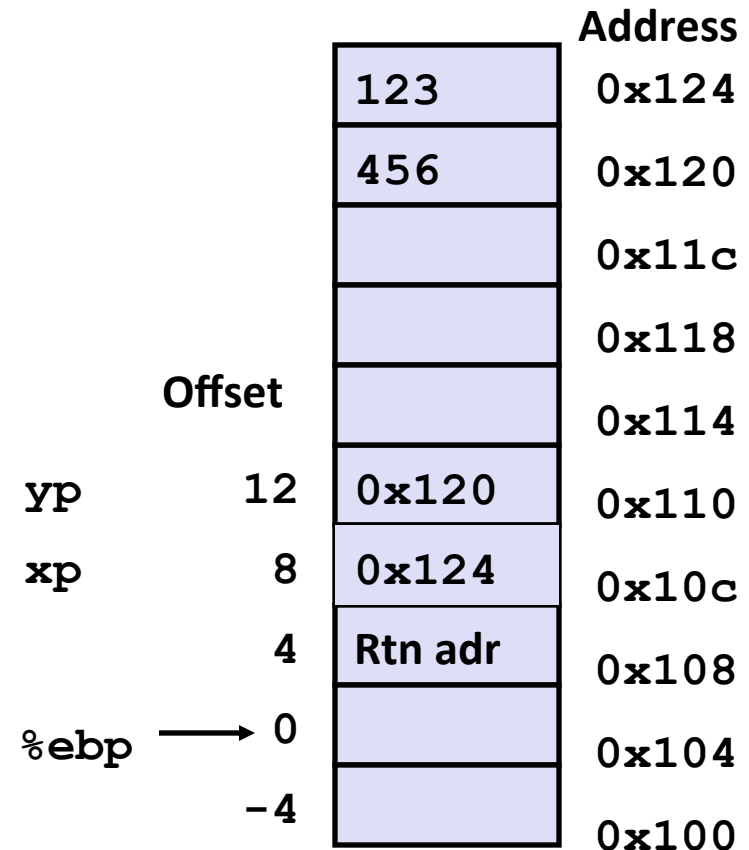
```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx      # edx = xp
movl (%ecx), %eax       # eax = *yp (t1)
movl (%edx), %ebx       # ebx = *xp (t0)
movl %eax, (%edx)       # *xp = eax
movl %ebx, (%ecx)       # *yp = ebx

```

Understanding Swap

<code>%eax</code>	
<code>%edx</code>	0x124
<code>%ecx</code>	0x120
<code>%ebx</code>	
<code>%esi</code>	
<code>%edi</code>	
<code>%esp</code>	
<code>%ebp</code>	0x104



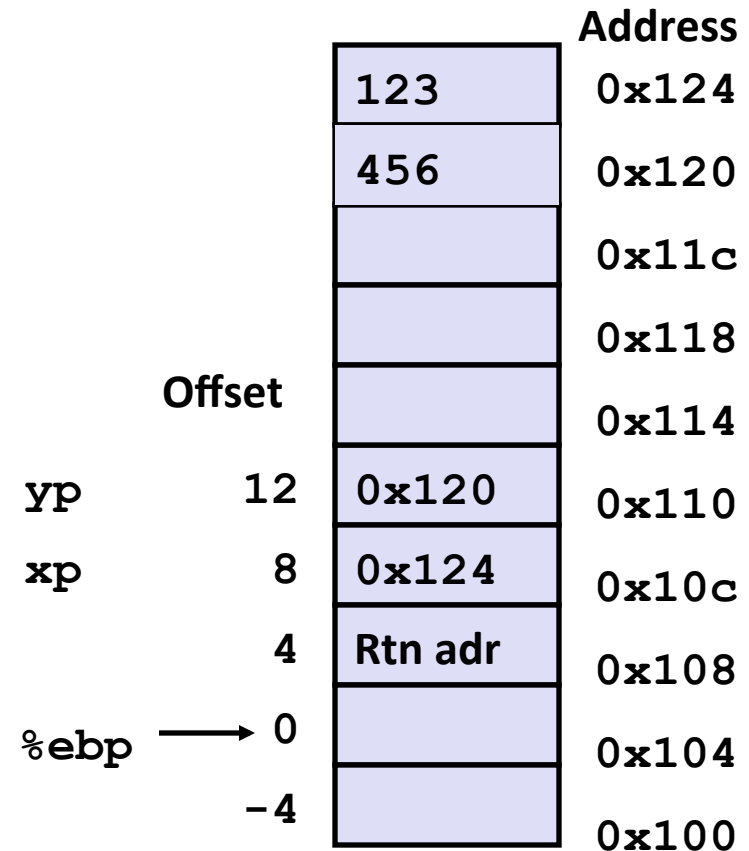
```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax       # eax = *yp (t1)
movl (%edx), %ebx       # ebx = *xp (t0)
movl %eax, (%edx)       # *xp = eax
movl %ebx, (%ecx)       # *yp = ebx

```


Understanding Swap

<code>%eax</code>	456
<code>%edx</code>	0x124
<code>%ecx</code>	0x120
<code>%ebx</code>	
<code>%esi</code>	
<code>%edi</code>	
<code>%esp</code>	
<code>%ebp</code>	0x104



```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx       # edx = xp
movl (%ecx), %eax       # eax = *yp (t1)
movl (%edx), %ebx       # ebx = *xp (t0)
movl %eax, (%edx)       # *xp = eax
movl %ebx, (%ecx)       # *yp = ebx

```

Understanding Swap

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

		Offset	Address
			123
			456
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp	→ 0		0x104
	-4		0x100

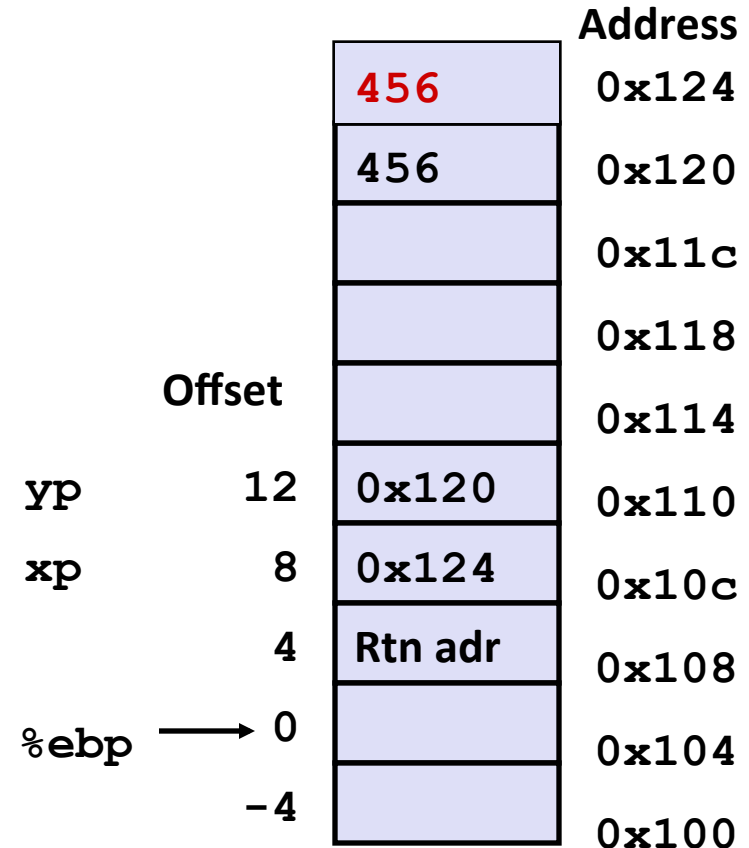
```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx       # edx = xp
movl (%ecx), %eax        # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)        # *xp = eax
movl %ebx, (%ecx)        # *yp = ebx

```

Understanding Swap

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104



```
movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx       # edx = xp
movl (%ecx), %eax        # eax = *yp (t1)
movl (%edx), %ebx        # ebx = *xp (t0)
movl %eax, (%edx)       # *xp = eax
movl %ebx, (%ecx)
```

Understanding Swap

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

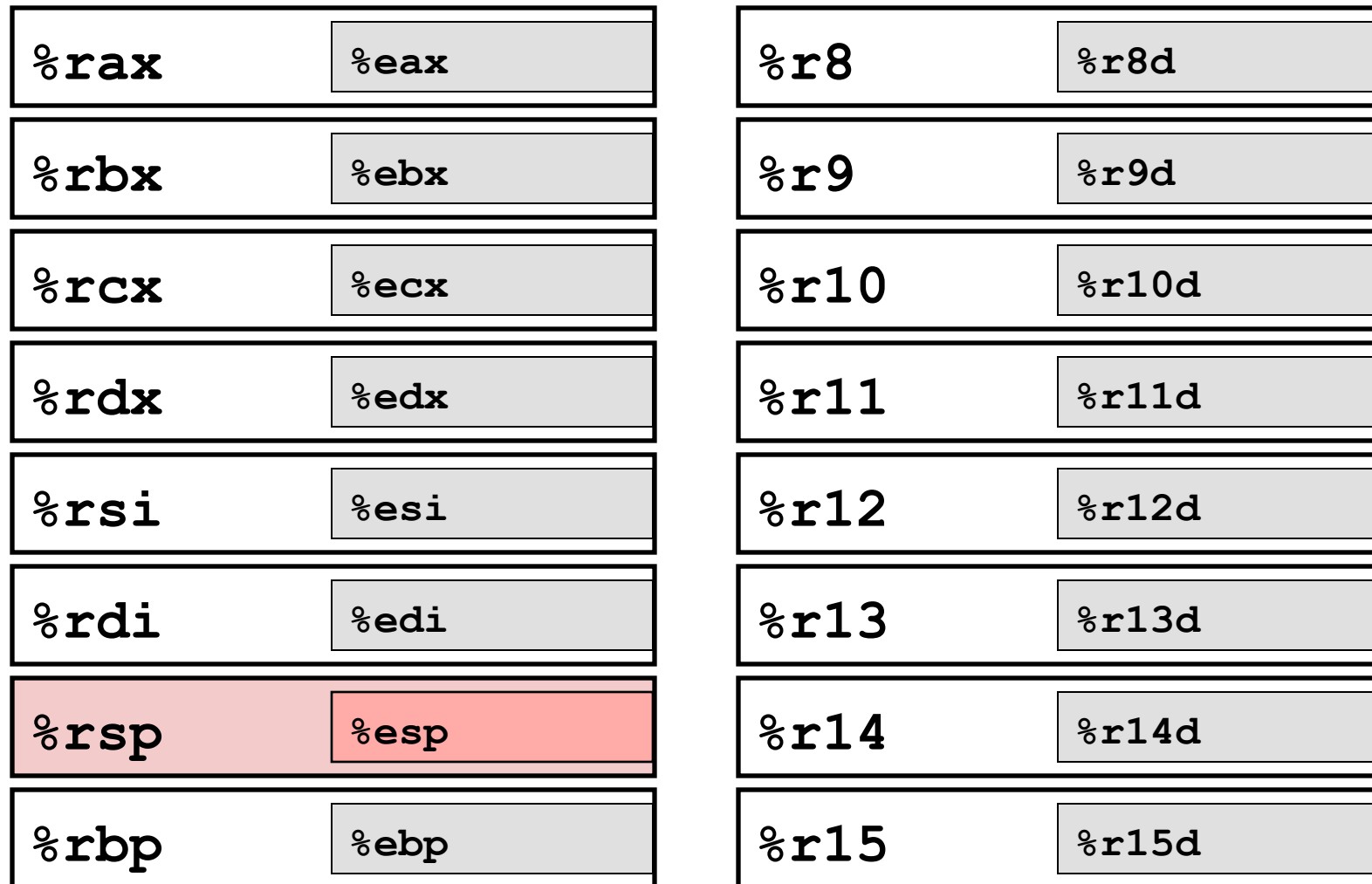
Offset		Address
		456 0x124
		123 0x120
		0x11c
		0x118
		0x114
yp	12	0x120 0x110
xp	8	0x124 0x10c
	4	Rtn adr 0x108
%ebp	→ 0	0x104
	-4	0x100

```

movl 12(%ebp), %ecx      # ecx = yp
movl 8(%ebp), %edx       # edx = xp
movl (%ecx), %eax        # eax = *yp (t1)
movl (%edx), %ebx        # ebx = *xp (t0)
movl %eax, (%edx)        # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx

```

x86-64 Integer Registers



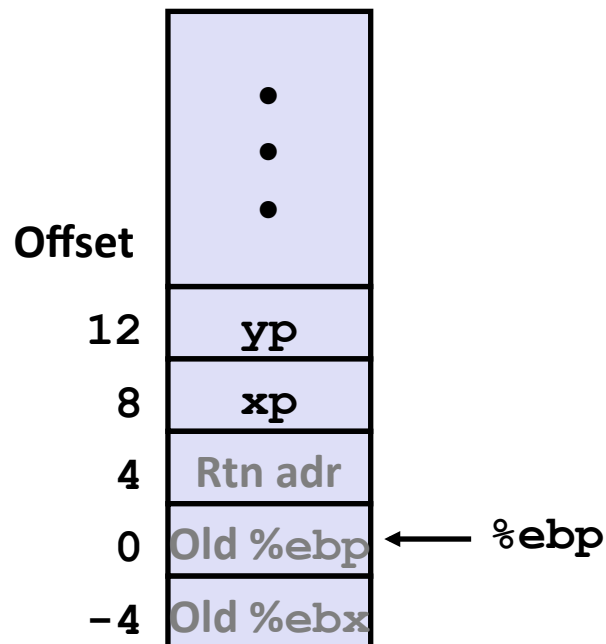
- Extend existing registers, and add 8 new ones; *all* accessible as 8, 16, 32, 64 bits.

32-bit vs. 64-bit operands

- Long word `l` (4 Bytes) ↔ Quad word `q` (8 Bytes)
- New instruction forms:
 - `movl` → `movq`
 - `addl` → `addq`
 - `sall` → `salq`
 - etc.
- **x86-64 can still use 32-bit instructions that generate 32-bit results**
 - Higher-order bits of destination register are just set to 0
 - Example: `addl`

Swap Ints in 32-bit Mode

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```



swap:

```
pushl %ebp
movl  %esp, %ebp
pushl %ebx
```

} Setup

```
movl 12(%ebp), %ecx
movl 8(%ebp), %edx
movl (%ecx), %eax
movl (%edx), %ebx
movl %eax, (%edx)
movl %ebx, (%ecx)
```

} Body

```
movl -4(%ebp), %ebx
movl %ebp, %esp
popl %ebp
ret
```

} Finish

Swap Ints in 64-bit Mode

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movl    (%rdi), %edx
    movl    (%rsi), %eax
    movl    %eax, (%rdi)
    movl    %edx, (%rsi)
    retq
```

- **Arguments passed in registers (why useful?)**
 - First (**xp**) in `%rdi`, second (**yp**) in `%rsi`
 - 64-bit pointers
- **No stack operations required**
- **32-bit data**
 - Data held in registers `%eax` and `%edx`
 - `movl` operation (the `l` refers to data width, not address width)

Swap Long Ints in 64-bit Mode

```
void swap_l
(long int *xp, long int *yp)
{
    long int t0 = *xp;
    long int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap_l:
    movq    (%rdi), %rdx
    movq    (%rsi), %rax
    movq    %rax, (%rdi)
    movq    %rdx, (%rsi)
    retq
```

■ 64-bit data

- Data held in registers **%rax** and **%rdx**
- **movq** operation
- “q” stands for quad-word

Complete Memory Addressing Modes

- Remember, the addresses used for accessing memory in `mov` (and other) instructions can be computed in several different ways

- Most General Form:**

$$D(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri] + D]$$

- D: Constant “displacement” 1, 2, or 4 bytes
- Rb: Base register: Any of the 8/16 integer registers
- Ri: Index register: Any, except for `%esp` or `%rsp`
 - Unlikely you’d use `%ebp`, either
- S: Scale: 1, 2, 4, or 8 (*why these numbers?*)

- Special Cases: can use any combination of D, Rb, Ri and S**

$$(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]]$$

$$D(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D]$$

$$(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri]]$$

Address Computation Examples

<code>%edx</code>	<code>0xf000</code>
<code>%ecx</code>	<code>0x100</code>

(Rb,Ri)

Mem[Reg[Rb]+Reg[Ri]]

D(,Ri,S)

Mem[S*Reg[Ri]+D]

(Rb,Ri,S)

Mem[Reg[Rb]+S*Reg[Ri]]

D(Rb)

Mem[Reg[Rb] +D]

Expression	Address Computation	Address
<code>0x8(%edx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%edx,%ecx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%edx,%ecx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%edx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

Address Computation Instruction

■ `leal Src, Dest`

- `Src` is address mode expression
- Set `Dest` to address computed by expression
 - (lea stands for *load effective address*)
- Example: `leal (%edx, %ecx, 4), %eax`

■ Uses

- Computing addresses without a memory reference
 - E.g., translation of `p = &x[i];`
- Computing arithmetic expressions of the form $x + k*i$
 - $k = 1, 2, 4, \text{ or } 8$

Some Arithmetic Operations

■ Two Operand (Binary) Instructions:

Format

Computation

addl *Src, Dest*

$Dest = Dest + Src$

subl *Src, Dest*

$Dest = Dest - Src$

imull *Src, Dest*

$Dest = Dest * Src$

sall *Src, Dest*

$Dest = Dest \ll Src$

Also called shll

sarl *Src, Dest*

$Dest = Dest \gg Src$

Arithmetic

shrl *Src, Dest*

$Dest = Dest \gg Src$

Logical

xorl *Src, Dest*

$Dest = Dest \wedge Src$

andl *Src, Dest*

$Dest = Dest \& Src$

orl *Src, Dest*

$Dest = Dest | Src$

■ Watch out for argument order! (especially `subl`)

■ No distinction between signed and unsigned int (why?)

Some Arithmetic Operations

■ One Operand (Unary) Instructions

`incl Dest` $Dest = Dest + 1$

`decl Dest` $Dest = Dest - 1$

`negl Dest` $Dest = -Dest$

`notl Dest` $Dest = \sim Dest$

- See textbook section 3.5.5 for more instructions: `mull`, `cld`, `idivl`, `divl`

Using `leal` for Arithmetic Expressions

```
int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

arith:

```
pushl %ebp
movl %esp,%ebp
```

} Set
Up

```
movl 8(%ebp),%eax
movl 12(%ebp),%edx
leal (%edx,%eax),%ecx
leal (%edx,%edx,2),%edx
sall $4,%edx
addl 16(%ebp),%ecx
leal 4(%edx,%eax),%eax
imull %ecx,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

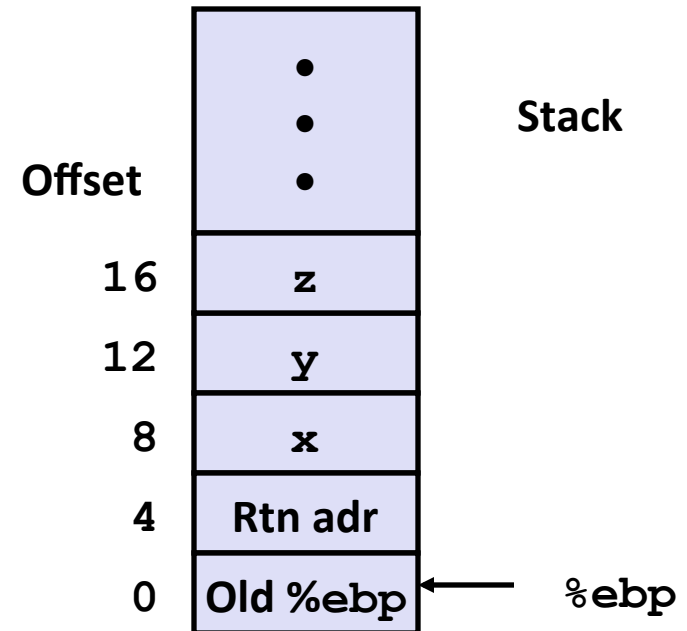
} Finish

Understanding arith

```

int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}

```



```

movl 8(%ebp), %eax      # eax = x
movl 12(%ebp), %edx     # edx = y
leal (%edx, %eax), %ecx # ecx = x+y (t1)
leal (%edx, %edx, 2), %edx # edx = y + 2*y = 3*y
sall $4, %edx          # edx = 48*y (t4)
addl 16(%ebp), %ecx    # ecx = z+t1 (t2)
leal 4(%edx, %eax), %eax # eax = 4+t4+x (t5)
imull %ecx, %eax      # eax = t5*t2 (rval)

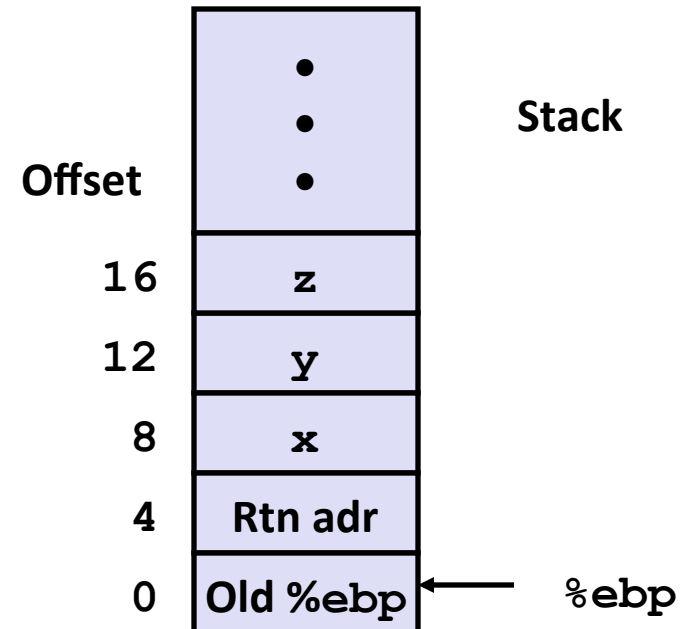
```


Understanding arith

```

int arith
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  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}

```



```

movl 8(%ebp), %eax      # eax = x
movl 12(%ebp), %edx     # edx = y
leal (%edx, %eax), %ecx # ecx = x+y (t1)
leal (%edx, %edx, 2), %edx # edx = y + 2*y = 3*y
sall $4, %edx          # edx = 48*y (t4)
addl 16(%ebp), %ecx    # ecx = z+t1 (t2)
leal 4(%edx, %eax), %eax # eax = 4+t4+x (t5)
imull %ecx, %eax       # eax = t5*t2 (rval)

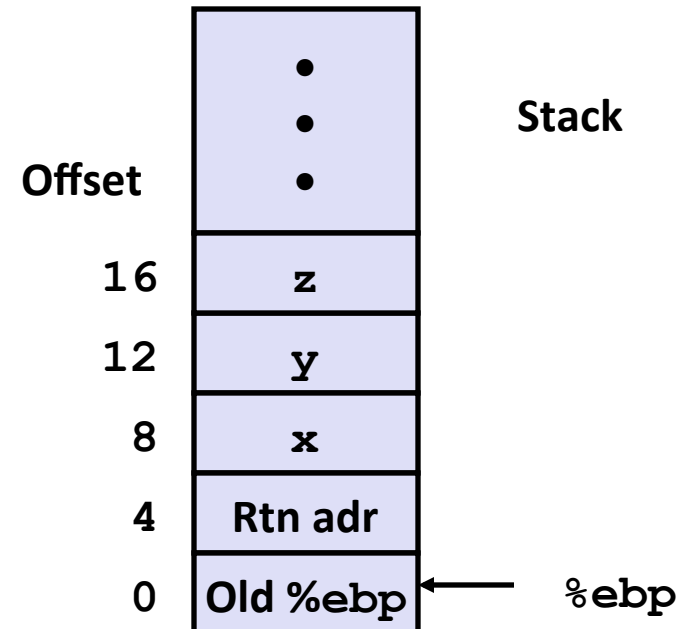
```

Understanding arith

```

int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}

```



```

movl 8(%ebp), %eax      # eax = x
movl 12(%ebp), %edx     # edx = y
leal (%edx, %eax), %ecx # ecx = x+y (t1)
leal (%edx, %edx, 2), %edx # edx = y + 2*y = 3*y
sall $4, %edx          # edx = 48*y (t4)
addl 16(%ebp), %ecx    # ecx = z+t1 (t2)
leal 4(%edx, %eax), %eax # eax = 4+t4+x (t5)
imull %ecx, %eax      # eax = t5*t2 (rval)

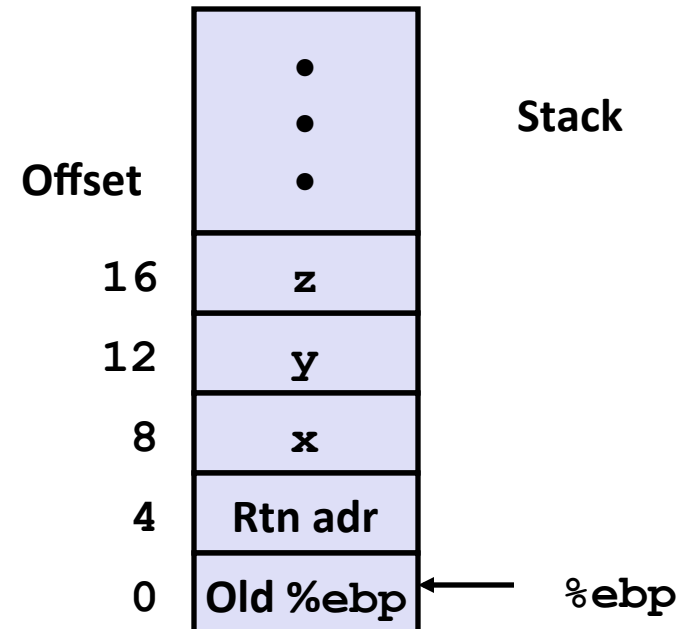
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Understanding arith

```

int arith
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  int rval = t2 * t5;
  return rval;
}

```



```

movl 8(%ebp), %eax      # eax = x
movl 12(%ebp), %edx     # edx = y
leal (%edx, %eax), %ecx # ecx = x+y (t1)
leal (%edx, %edx, 2), %edx # edx = y + 2*y = 3*y
sall $4, %edx          # edx = 48*y (t4)
addl 16(%ebp), %ecx     # ecx = z+t1 (t2)
leal 4(%edx, %eax), %eax # eax = 4+t4+x (t5)
imull %ecx, %eax       # eax = t5*t2 (rval)

```

Observations about arith

```

int arith
  (int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}

```

- Instructions in different order from C code
- Some expressions require multiple instructions
- Some instructions cover multiple expressions
- Get exact same code when compile:
 - $(x+y+z) * (x+4+48*y)$

<code>movl 8(%ebp), %eax</code>	<code># eax = x</code>
<code>movl 12(%ebp), %edx</code>	<code># edx = y</code>
<code>leal (%edx, %eax), %ecx</code>	<code># ecx = x+y (t1)</code>
<code>leal (%edx, %edx, 2), %edx</code>	<code># edx = y + 2*y = 3*y</code>
<code>sall \$4, %edx</code>	<code># edx = 48*y (t4)</code>
<code>addl 16(%ebp), %ecx</code>	<code># ecx = z+t1 (t2)</code>
<code>leal 4(%edx, %eax), %eax</code>	<code># eax = 4+t4+x (t5)</code>
<code>imull %ecx, %eax</code>	<code># eax = t5*t2 (rval)</code>

Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set Up

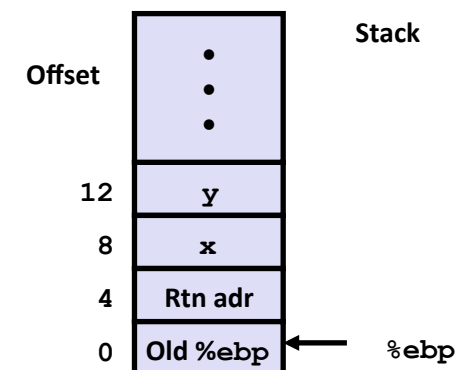
```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax    # eax = x
xorl 12(%ebp),%eax  # eax = x^y
sarl $17,%eax       # eax = t1>>17
andl $8185,%eax     # eax = t2 & 8185
```



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xorl 12(%ebp),%eax
sarl $17,%eax
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```

}
Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
eax = x
eax = x^y      (t1)
eax = t1>>17  (t2)
eax = t2 & 8185
```

Another Example

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{
    int t1 = x^y;
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    int t1 = x^y;
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}
```

$2^{13} = 8192,$ $2^{13} - 7 = 8185$
 $\dots00100000000000000000,$ $\dots000111111111111001$

logical:

```
    pushl %ebp
    movl  %esp,%ebp
```

} Set
Up

```
    movl  8(%ebp),%eax
    xorl  12(%ebp),%eax
    sarl  $17,%eax
    andl  $8185,%eax
```

}
Body

```
    movl  %ebp,%esp
    popl  %ebp
    ret
```

} Finish

```
movl  8(%ebp),%eax
xorl  12(%ebp),%eax
sarl  $17,%eax
andl  $8185,%eax
```

```
eax = x
eax = x^y      (t1)
eax = t1>>17  (t2)
eax = t2 & 8185
```