

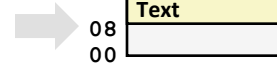
Buffer Overflow

- Basics of memory allocation
- Buffers on stacks
- Overwriting buffers
- Injecting code

IA32 Linux Memory Layout

- **Stack**
 - Runtime stack (8MB limit)
- **Heap**
 - Dynamically allocated storage
 - When call `malloc()`, `calloc()`, `new()`
- **Data**
 - Statically allocated data
 - E.g., arrays & strings declared in code
- **Text**
 - Executable machine instructions
 - Read-only

Upper 2 hex digits
= 8 bits of address



Memory Allocation Example

```

char big_array[1<<24]; /* 16 MB */
char huge_array[1<<28]; /* 256 MB */

int beyond;
char *p1, *p2, *p3, *p4;

int useless() { return 0; }

int main()
{
  p1 = malloc(1 <<28); /* 256 MB */
  p2 = malloc(1 << 8); /* 256 B */
  p3 = malloc(1 <<28); /* 256 MB */
  p4 = malloc(1 << 8); /* 256 B */
  /* Some print statements ... */
}

```

Where does everything go?

not drawn to scale



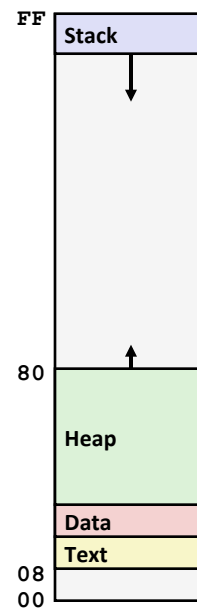
IA32 Example Addresses

address range $\sim 2^{32}$

\$esp	0xffffbcd0
p3	0x65586008
p1	0x55585008
p4	0x1904a110
p2	0x1904a008
&p2	0x18049760
beyond	0x08049744
big_array	0x18049780
huge_array	0x08049760
main()	0x080483c6
useless()	0x08049744
final malloc()	0x006be166

`malloc()` is dynamically linked
address determined at runtime

not drawn to scale



Internet Worm

- **November, 1988**
 - Internet Worm attacks thousands of Internet hosts.
 - How did it happen?

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- **November, 1988**
 - Internet Worm attacks thousands of Internet hosts.
 - How did it happen?

- **The Internet Worm was based on *stack buffer overflow* exploits!**
 - many Unix functions do not check argument sizes
 - allows target buffers to overflow

String Library Code

■ Implementation of Unix function `gets()`

```

/* Get string from stdin */
char *gets(char *dest)
{
    int c = getchar();
    char *p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}

```

- Anything interesting in the above?

String Library Code

■ Implementation of Unix function `gets()`

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char *gets(char *dest)
{
    int c = getchar();
    char *p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}

```

- No way to specify limit on number of characters to read
- **Similar problems with other Unix functions**
 - `strcpy`: Copies string of arbitrary length
 - `scanf`, `fscanf`, `sscanf`, when given `%s` conversion specification

Vulnerable Buffer Code

```

/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}

```

```

int main()
{
    printf("Type a string:");
    echo();
    return 0;
}

```

```

unix> ./bufdemo
Type a string:1234567
1234567

```

```

unix> ./bufdemo
Type a string:12345678
Segmentation Fault

```

```

unix> ./bufdemo
Type a string:123456789ABC
Segmentation Fault

```

Buffer Overflow Disassembly

```

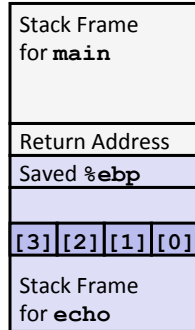
080484f0 <echo>:
80484f0: 55                push   %ebp
80484f1: 89 e5            mov    %esp,%ebp
80484f3: 53                push   %ebx
80484f4: 8d 5d f8        lea   0xffffffff8(%ebp),%ebx
80484f7: 83 ec 14        sub   $0x14,%esp
80484fa: 89 1c 24        mov   %ebx,(%esp)
80484fd: e8 ae ff ff ff  call  80484b0 <gets>
8048502: 89 1c 24        mov   %ebx,(%esp)
8048505: e8 8a fe ff ff  call  8048394 <puts@plt>
804850a: 83 c4 14        add   $0x14,%esp
804850d: 5b                pop   %ebx
804850e: c9                leave
804850f: c3                ret

80485f2: e8 f9 fe ff ff  call  80484f0 <echo>
80485f7: 8b 5d fc        mov   0xffffffffc(%ebp),%ebx
80485fa: c9                leave
80485fb: 31 c0            xor   %eax,%eax
80485fd: c3                ret

```

Buffer Overflow Stack

Before call to gets



```

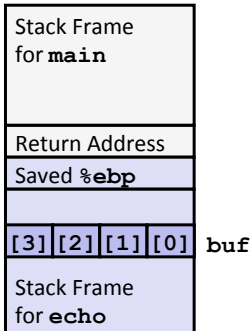
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
    
```

```

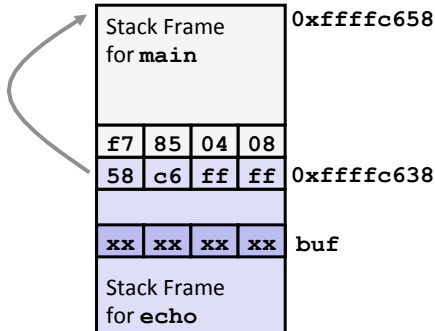
echo:
    pushl %ebp           # Save %ebp on stack
    movl  %esp, %ebp
    pushl %ebx          # Save %ebx
    leal -8(%ebp),%ebx  # Compute buf as %ebp-8
    subl $20, %esp      # Allocate stack space
    movl  %ebx, (%esp)  # Push buf addr on stack
    call  gets          # Call gets
    . . .
    
```

Buffer Overflow Stack Example

Before call to gets



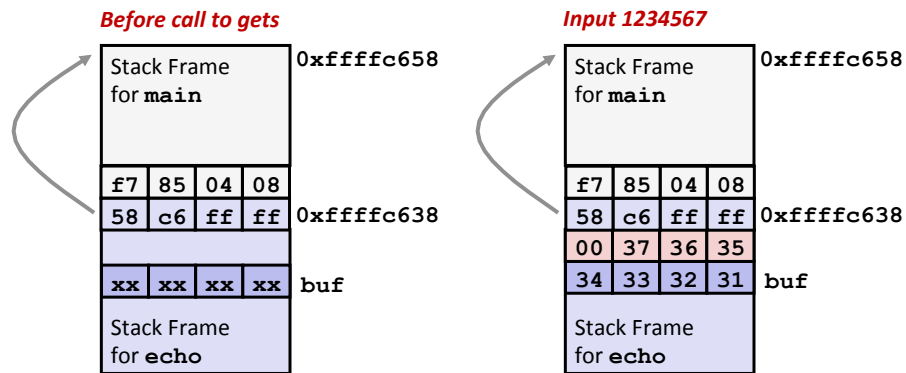
Before call to gets



```

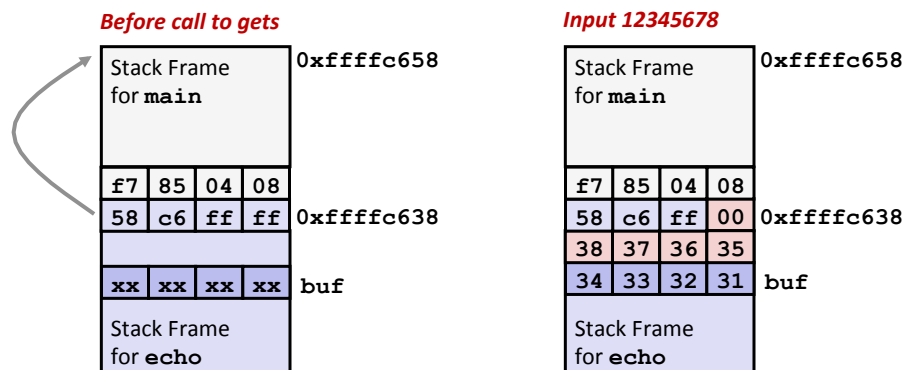
80485f2: call 80484f0 <echo>
80485f7: mov  0xffffffc(%ebp),%ebx # Return Point
    
```

Buffer Overflow Example #1



Overflow buf, but no problem

Buffer Overflow Example #2



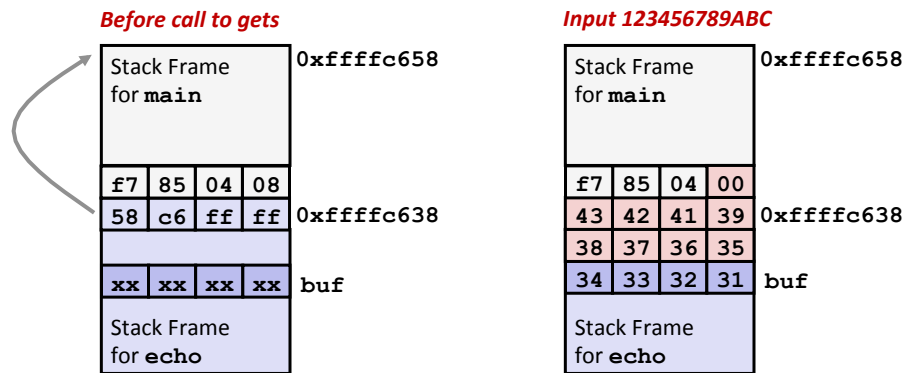
Base pointer corrupted

```

. . .
804850a: 83 c4 14 add    $0x14,%esp # deallocate space
804850d: 5b      pop    %ebx     # restore %ebx
804850e: c9      leave          # movl %ebp, %esp; popl %ebp
804850f: c3      ret           # Return

```

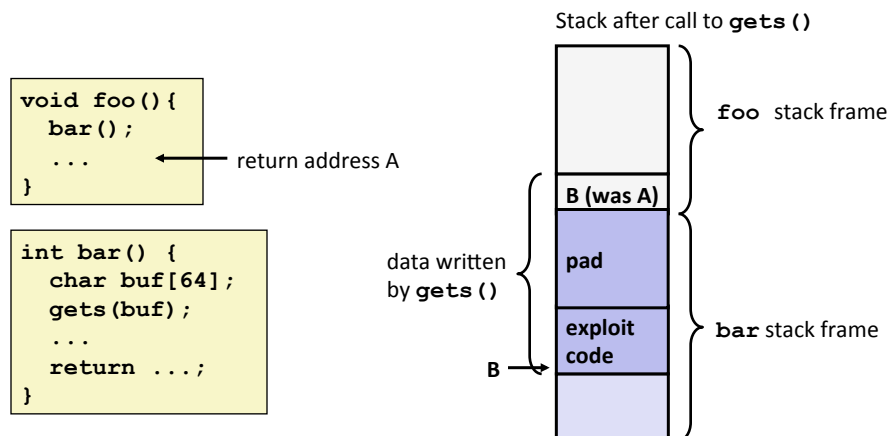
Buffer Overflow Example #3



Return address corrupted

```
80485f2: call 80484f0 <echo>
80485f7: mov 0xffffffff(%ebp),%ebx # Return Point
```

Malicious Use of Buffer Overflow



- Input string contains byte representation of executable code
- Stack frame must be big enough to hold exploit code
- Overwrite return address with address of buffer (need to know B)
- When `bar()` executes `ret`, will jump to exploit code (instead of A)

Exploits Based on Buffer Overflows

- **Buffer overflow bugs allow remote machines to execute arbitrary code on victim machines**
- **Internet worm**
 - Early versions of the finger server (fingerd) used `gets()` to read the argument sent by the client:
 - `finger droh@cs.cmu.edu`
 - Worm attacked fingerd server by sending phony argument:
 - `finger "exploit-code padding new-return-address"`
 - exploit code: executed a root shell on the victim machine with a direct TCP connection to the attacker

Avoiding Overflow Vulnerability

```

/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    fgets(buf, 4, stdin);
    puts(buf);
}

```

- **Use library routines that limit string lengths**
 - `fgets` instead of `gets` (second argument to `fgets` sets limit)
 - `strncpy` instead of `strcpy`
 - Don't use `scanf` with `%s` conversion specification
 - Use `fgets` to read the string
 - Or use `%ns` where `n` is a suitable integer

System-Level Protections

- **Randomized stack offsets**
 - At start of program, allocate random amount of space on stack
 - Makes it difficult for hacker to predict beginning of inserted code
- **Nonexecutable code segments**
 - Only allow code to execute from “text” sections of memory
 - Do NOT execute code in stack, data, or heap regions
 - Hardware support

