# **CSE 344**

#### **JANUARY 17<sup>TH</sup> – SUBQUERIES**

Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

| SELECT   | <pre>product, Sum(quantity) AS TotalSales</pre> |
|----------|---|
| FROM     | Purchase  |
| WHERE    | price > 1                                       |
| GROUP BY | product   |
|          |   |

How is this query processed?

Purchase(product, price, quantity)

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|          |   |

Do these queries return the same number of rows? Why?

SELECT product, Sum(quantity) AS TotalSales
FROM Purchase
GROUP BY product

Purchase(product, price, quantity)

Find total quantities for all sales over \$1, by product.

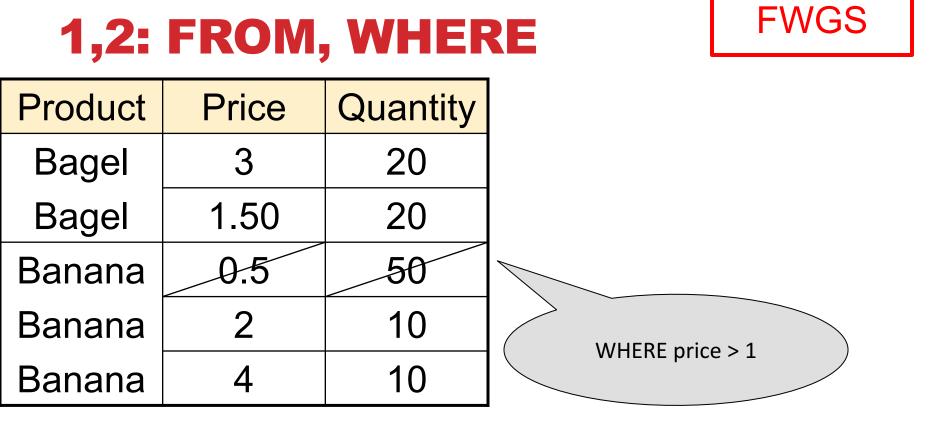
| SELECT   | <pre>product, Sum(quantity) AS TotalSales</pre> |
|----------|---|
| FROM     | Purchase  |
| WHERE    | price > 1                                       |
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|          |   |

Do these queries return the same number of rows? Why?

| SELECTproduct, Sum(quFROMPurchase | antity) AS TotalSales   |  |
|-----------------------------------|---|--|
| GROUP BY product                  | Empty groups are removed, hence first query may return fewer groups |  |

- 1. Compute the FROM and WHERE clauses.
- 2. Group by the attributes in the GROUPBY
- 3. Compute the SELECT clause: grouped attributes and aggregates.





| SELECT   | product, Sum(quantity) <mark>AS</mark> TotalSales |
|----------|---|
| FROM     | Purchase  |
| WHERE    | price > 1   |
| GROUP BY | product   |
|          |   |



| Product | Price | Quantity |  |         |            |
|---------|-------|----------|--|---------|------------|
| Bagel   | 3     | 20       |  | Product | TotalSales |
| Dayci   |       |          |  |         | 10         |
| Bagel   | 1.50  | 20       |  | Bagel   | 40         |
| Banana  | 0.5   | 50       |  | Banana  | 20         |
| Banana  | 2     | 10       |  |         |            |
| Banana  | 4     | 10       |  |         |            |

**FWGS** 

| SELECT   | product, Sum(quantity) AS TotalSales |
|----------|--------------------------------------|
| FROM     | Purchase                             |
| WHERE    | price > 1                            |
| GROUP BY | Y product                            |
|          |                                      |

### **ORDERING RESULTS**

SELECT product, sum(price\*quantity) as rev
FROM Purchase
GROUP BY product
ORDER BY rev desc



Note: some SQL engines want you to say ORDER BY sum(price\*quantity) desc

### **HAVING CLAUSE**

Same query as before, except that we consider only products that had at least 30 sales.

| SELECT   | <pre>product, sum(price*quantity)</pre> |
|----------|---|
| FROM     | Purchase                                |
| WHERE    | price > 1                               |
| GROUP BY | product                                 |
| HAVING   | <pre>sum(quantity) &gt; 30</pre>        |
|          |   |
|          |   |

HAVING clause contains conditions on aggregates.

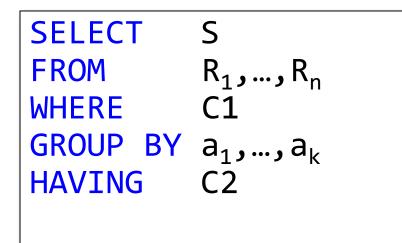
### GENERAL FORM OF GROUPING AND AGGREGATION

SELECTSFROM $R_1, \dots, R_n$ WHEREC1GROUP BY $a_1, \dots, a_k$ HAVINGC2

- S = may contain attributes a<sub>1</sub>,...,a<sub>k</sub> and/or any aggregates but NO OTHER ATTRIBUTES
- C1 = is any condition on the attributes in  $R_1, ..., R_n$
- C2 = is any condition on aggregate expressions and on attributes a<sub>1</sub>,...,a<sub>k</sub>

Why?

# SEMANTICS OF SQL WITH GROUP-BY





#### **Evaluation steps:**

- 1. Evaluate FROM-WHERE using Nested Loop Semantic
- 2. Group by the attributes  $a_1, \dots, a_k$
- 3. Apply condition C2 to each group (may have aggrega
- 4. Compute aggregates in S and return the result





Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

FROM

Purchase



Compute the total income per month Show only months with less than 10 items sold Order by quantity sold and display as "TotalSold"

FROMPurchaseGROUP BYmonth



| FROM     | Purchase           |
|----------|--------------------|
| GROUP BY |                    |
| HAVING   | sum(quantity) < 10 |



| SELECT   | <pre>month, sum(price*quantity), sum(quantity) as TotalSold</pre> |
|----------|---|
| FROM     | Purchase  |
| GROUP BY | month   |
| HAVING   | sum(quantity) < 10  |
|          |   |



| SELECT   | <pre>month, sum(price*quantity), sum(quantity) as TotalSold</pre> |
|----------|---|
| FROM     | Purchase  |
| GROUP BY | month   |
| HAVING   | sum(quantity) < 10  |
| ORDER BY | <pre>sum(quantity)</pre>  |
|          |   |
|          |   |

# WHERE VS HAVING

#### WHERE condition is applied to individual rows

- The rows may or may not contribute to the aggregate
- No aggregates allowed here
- Occasionally, some groups become empty and are removed

#### HAVING condition is applied to the entire group

- Entire group is returned, or removed
- May use aggregate functions on the group



What do they compute?

SELECTmonth, sum(quantity), max(price)FROMPurchaseGROUP BYmonth

SELECTmonth, sum(quantity)FROMPurchaseGROUP BYmonth

SELECTmonthFROMPurchaseGROUP BYmonth



What do they compute?

SELECTmonth, sum(quantity), max(price)FROMPurchaseGROUP BYmonth

SELECTmonth, sum(quantity)FROMPurchaseGROUP BYmonth

SELECTmonthFROMPurchaseGROUP BYmonth

Lesson: DISTINCT is a special case of GROUP BY

# **AGGREGATE + JOIN**

For each manufacturer, compute how many products with price > \$100 they sold

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Problem: manufacturer is in Purchase, price is in Product...

### **AGGREGATE + JOIN**

For each manufacturer, compute how many products with price > \$100 they sold

Problem: manufacturer is in Purchase, price is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
```

| manu<br>facturer | <br>price |  |
|------------------|-----------|--|
| Hitachi          | 150       |  |
| Canon            | 300       |  |
| Hitachi          | 180       |  |

# **AGGREGATE + JOIN**

For each manufacturer, compute how many products with price > \$100 they sold

Problem: manufacturer is in Purchase, price is in Product...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
```

| manu<br>facturer | <br>price |  |
|------------------|-----------|--|
| Hitachi          | 150       |  |
| Canon            | 300       |  |
| Hitachi          | 180       |  |

```
-- step 2: do the group-by on the join
SELECT x.manufacturer, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
GROUP BY x.manufacturer
```

| manu<br>facturer | count(*) |
|------------------|----------|
| Hitachi          | 2        |
| Canon            | 1        |
|                  |          |

# **AGGREGATE + JOIN**

Variant:

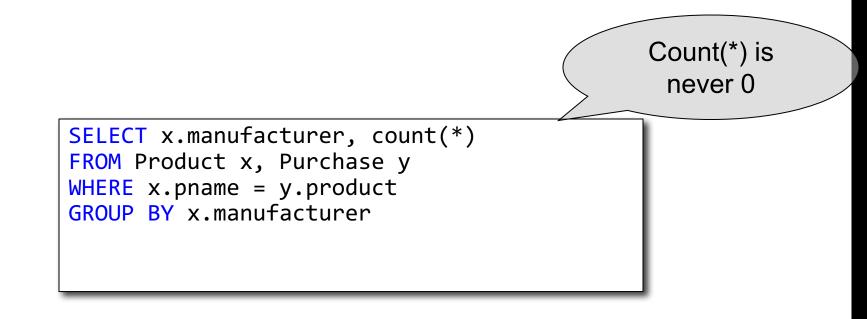
For each manufacturer, compute how many products with price > \$100 they sold in each month

```
SELECT x.manufacturer, y.month, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
GROUP BY x.manufacturer, y.month
```

| manu<br>facturer | month | count(*) |
|------------------|-------|----------|
| Hitachi          | Jan   | 2        |
| Hitachi          | Feb   | 1        |
| Canon            | Jan   | 3        |
|                  |       |          |

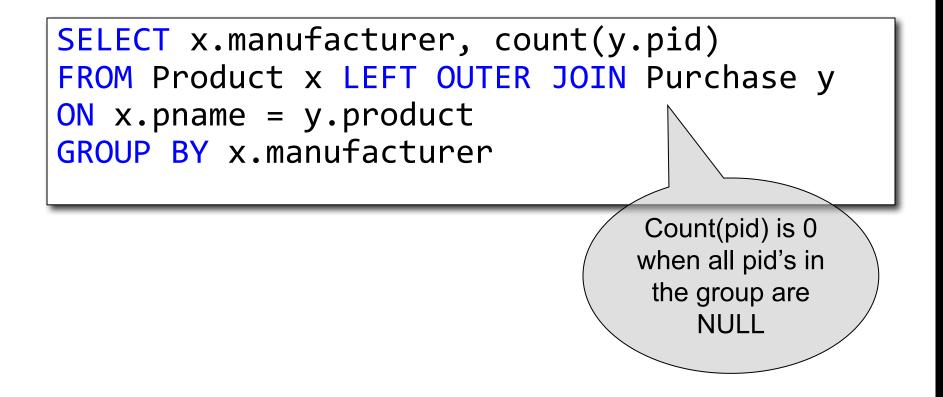
# INCLUDING EMPTY GROUPS

In the result of a group by query, there is one row per group in the result



**FWGHOS** 

# INCLUDING EMPTY GROUPS



### **SUBQUERIES**

A subquery is a SQL query nested inside a larger query

Such inner-outer queries are called nested queries

#### A subquery may occur in:

- A SELECT clause
- A FROM clause
- A WHERE clause

#### Rule of thumb: avoid nested queries when possible

• But sometimes it's impossible, as we will see

### **SUBQUERIES...**

Can return a single value to be included in a SELECT clause

Can return a relation to be included in the FROM clause, aliased using a tuple variable

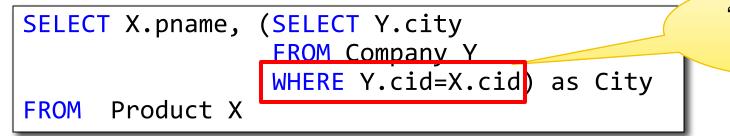
Can return a single value to be compared with another value in a WHERE clause

Can return a relation to be used in the WHERE or HAVING clause under an existential quantifier

# **1. SUBQUERIES IN SELECT**

Product (pname, price, cid)
Company (cid, cname, city)

For each product return the city where it is manufactured

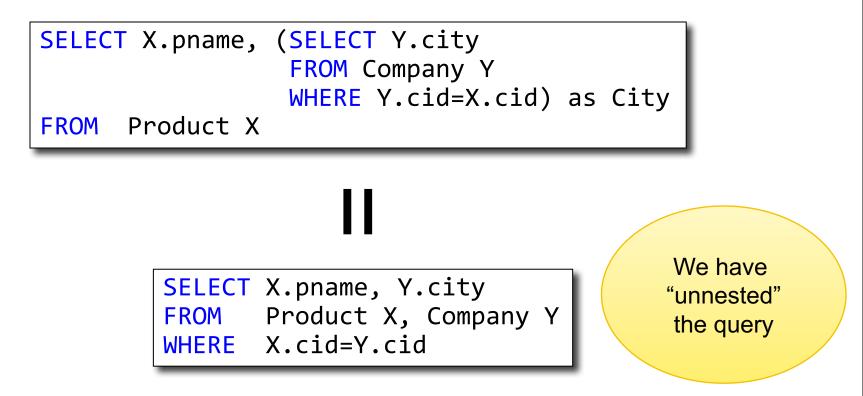


"correlated subquery"

What happens if the subquery returns more than one city?

We get a runtime error (and SQLite simply ignores the extra values...)

Whenever possible, don't use a nested queries:



Compute the number of products made by each company

SELECT DISTINCT C.cname, (SELECT count(\*) FROM Product P WHERE P.cid=C.cid) FROM Company C

Compute the number of products made by each company

| SELEC1 | <b>DISTINCT</b> | C.cname, | (SELECT count(*)   |
|--------|-----------------|----------|--------------------|
|        |                 |          | FROM Product P     |
|        |                 |          | WHERE P.cid=C.cid) |
| FROM   | Company C       |          |                    |

Better: we can unnest using a GROUP BY SELECT C.cname, count(\*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

SELECT C.cname, count(\*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
No! Different results if a company
has no products
```

```
SELECT C.cname, count(pname)
FROM Company C LEFT OUTER JOIN Product P
ON C.cid=P.cid
GROUP BY C.cname
```

#### Product (<u>pname</u>, price, cid) Comp**2**y **SUBQUERIES IN FROM**

Find all products whose prices is > 20 and < 500

SELECT X.pname FROM (SELECT \* FROM Product AS Y WHERE price > 20) as X WHERE X.price < 500

Find all products whose prices is > 20 and < 500

SELECT X.pname FROM (SELECT \* FROM Product AS Y WHERE price > 20) as X WHERE X.price < 500

Try unnest this query !

Find all products whose prices is > 20 and < 500

**SELECT** X.pname FROM (SELECT \* FROM Product AS Y WHERE price > 20) as X WHERE X.price < 500

Side note: This is not a correlated subquery. (why?)

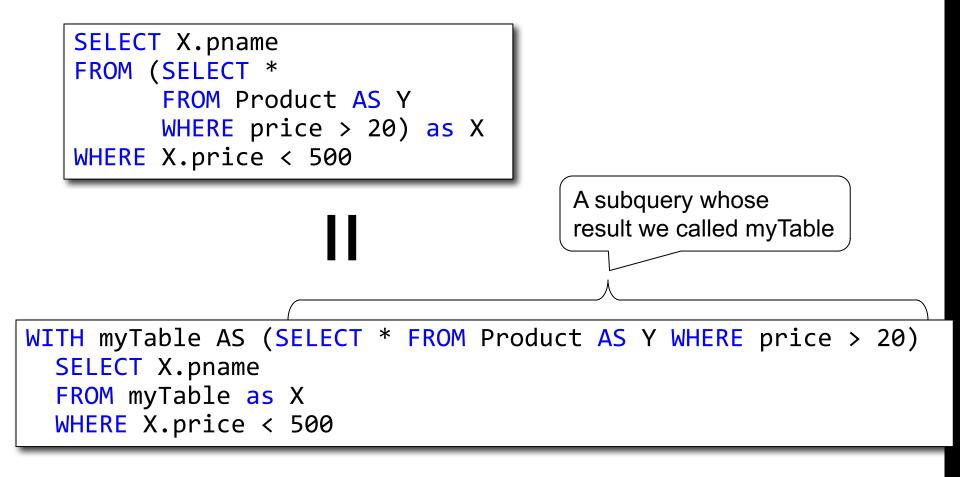
Try unnest this query !

# 2. SUBQUERIES IN FROM

Sometimes we need to compute an intermediate table only to use it later in a SELECT-FROM-WHERE

**Option 1: use a subquery in the FROM clause** 

**Option 2: use the WITH clause** 







**Existential quantifiers** 



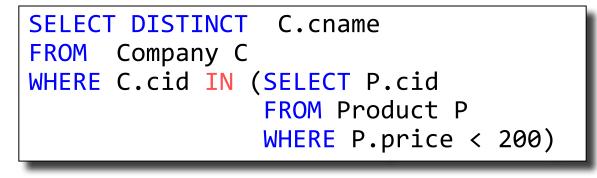
**Existential quantifiers** 

Using **EXISTS**:



**Existential quantifiers** 

Using IN





**Existential quantifiers** 

Using ANY:

SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)



**Existential quantifiers** 

Using ANY:

SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)

Not supported in sqlite



**Existential quantifiers** 

Now let's unnest it:

SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>



**Existential quantifiers** 

Now let's unnest it:

SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>

Existential quantifiers are easy! ©

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make <u>only</u> products with price < 200

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

**Universal quantifiers** 

Find all companies s.t. <u>all</u> their products have price < 200

same as:

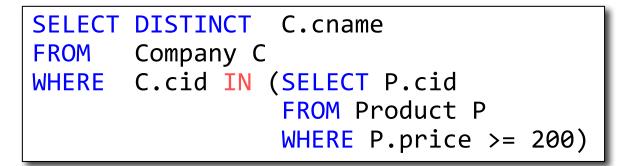
Find all companies that make only products with price < 200

**Universal quantifiers** 

### Universal quantifiers are hard! 🛞

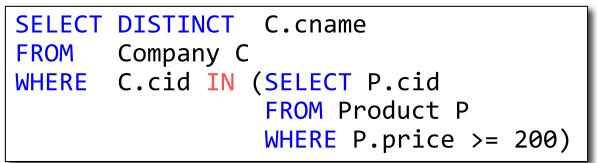
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1. Find *the other* companies that make <u>some</u> product  $\geq$  200



Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies that make <u>some</u> product  $\geq$  200



2. Find all companies s.t. <u>all</u> their products have price < 200



Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using **EXISTS**:

SELECT DISTINCT C.cname FROM Company C WHERE NOT EXISTS (SELECT \* FROM Product P WHERE P.cid = C.cid and P.price >= 200)



Find all companies s.t. <u>all</u> their products have price < 200

**Universal quantifiers** 

Using ALL:

SELECT DISTINCT C.cname
FROM Company C
WHERE 200 >= ALL (SELECT price
FROM Product P
WHERE P.cid = C.cid)



Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using ALL:

Not supported in sqlite

## QUESTION FOR DATABASE THEORY FANS AND THEIR FRIENDS

Can we unnest the *universal quantifier* query?

We need to first discuss the concept of *monotonicity* 

# **MONOTONE QUERIES**

#### **Definition A query Q is monotone if:**

 Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

# **MONOTONE QUERIES**

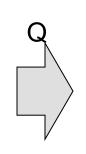
#### **Definition A query Q is monotone if:**

 Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

#### Company

| pname  | price  | cid  |
|--------|--------|------|
| Gizmo  | 19.99  | c001 |
| Gadget | 999.99 | c004 |
| Camera | 149.99 | c003 |

| cid  | cname    | city  |
|------|----------|-------|
| c002 | Sunworks | Bonn  |
| c001 | DB Inc.  | Lyon  |
| c003 | Builder  | Lodtz |



| pname  | city  |
|--------|-------|
| Gizmo  | Lyon  |
| Camera | Lodtz |

# **MONOTONE QUERIES**

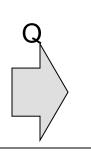
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|------|----------|-------|
| c002 | Sunworks | Bonn  |
| c001 | DB Inc.  | Lyon  |
| c003 | Builder  | Lodtz |
|      | -        |       |



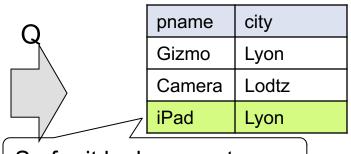
| pname  | city  |
|--------|-------|
| Gizmo  | Lyon  |
| Camera | Lodtz |

#### Product

| pname  | price  | cid  |
|--------|--------|------|
| Gizmo  | 19.99  | c001 |
| Gadget | 999.99 | c004 |
| Camera | 149.99 | c003 |
| iPad   | 499.99 | c001 |

#### Company

| -    | -        |       |
|------|----------|-------|
| cid  | cname    | city  |
| c002 | Sunworks | Bonn  |
| c001 | DB Inc.  | Lyon  |
| c003 | Builder  | Lodtz |
|      |          |       |



So far it looks monotone...

# **MONOTONE QUERIES**

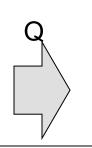
#### **Definition A query Q is monotone if:**

 Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

#### Company

| pname  | price  | cid  |
|--------|--------|------|
| Gizmo  | 19.99  | c001 |
| Gadget | 999.99 | c004 |
| Camera | 149.99 | c003 |
| Camera | 140.00 | 0000 |

| cid  | cname    | city  |
|------|----------|-------|
| c002 | Sunworks | Bonn  |
| c001 | DB Inc.  | Lyon  |
| c003 | Builder  | Lodtz |



| pname  | city  |
|--------|-------|
| Gizmo  | Lyon  |
| Camera | Lodtz |

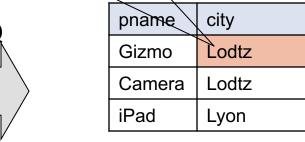
#### Product

| pname  | price  | cid  |
|--------|--------|------|
| Gizmo  | 19.99  | c001 |
| Gadget | 999.99 | c004 |
| Camera | 149.99 | c003 |
| iPad   | 499.99 | c001 |

#### Company

|      | J        |       |
|------|----------|-------|
| cid  | cname    | city  |
| c002 | Sunworks | Bonn  |
| c001 | DB Inc.  | Lyon  |
| c003 | Builder  | Lodtz |
| c004 | Crafter  | Lodtz |

# Q is not monotone!



# **MONOTONE QUERIES**

<u>Theorem</u>: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.

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<u>Theorem</u>: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.

**Proof.** We use the nested loop semantics: if we insert a tuple in a relation  $R_i$ , this will not remove any tuples from the answer

SELECT 
$$a_1$$
,  $a_2$ , ...,  $a_k$   
FROM  $R_1$  AS  $x_1$ ,  $R_2$  AS  $x_2$ , ...,  $R_n$  AS  $x_n$   
WHERE Conditions

for  $x_1$  in  $R_1$  do for  $x_2$  in  $R_2$  do for  $x_n$  in  $R_n$  do if Conditions output (a<sub>1</sub>,...,a<sub>k</sub>)

## **MONOTONE QUERIES** The query:

Find all companies s.t. <u>all</u> their products have price < 200 **is not monotone** 

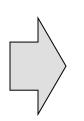
## **MONOTONE QUERIES** The query:

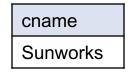
Find all companies s.t. <u>all</u> their products have price < 200

#### is not monotone

| pname | price | cid  |
|-------|-------|------|
| Gizmo | 19.99 | c001 |

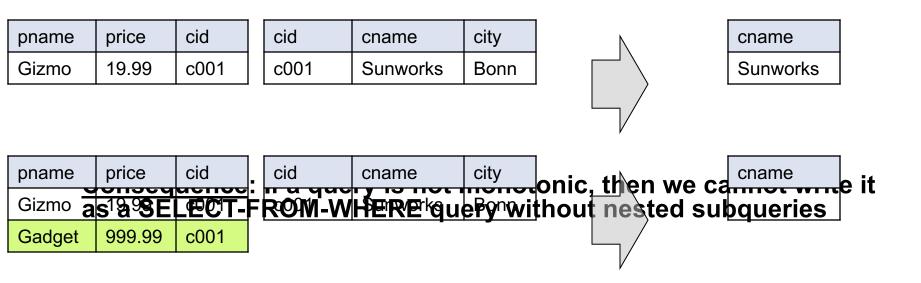
| cid  | cname    | city |
|------|----------|------|
| c001 | Sunworks | Bonn |





#### MONOTONE QUERIES The query:

Find all companies s.t. <u>all</u> their products have price < 200



# QUERIES THAT MUST BE NESTED

Queries with universal quantifiers or with negation

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Queries with universal quantifiers or with negation

# Queries that use aggregates in certain ways

- sum(..) and count(\*) are NOT monotone, because they do not satisfy set containment
- select count(\*) from R is not monotone!

# ON TO DATA MANAGEMEN

# **RELATIONAL ALGEBRA**

CSE 344 - 2017au

# ANNOUNCEMENTS

#### You received invitation email to @cs

#### You will be prompted to choose passwd

- Problems with existing account?
- In the worst case we will ask you to create a new @outlook account just for this class

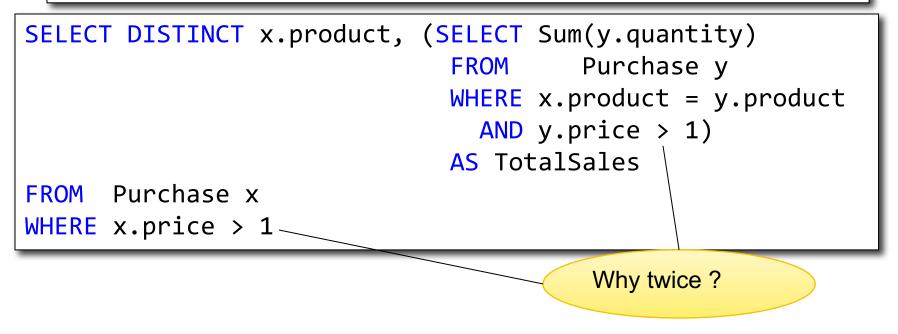
#### If OK, create the database server

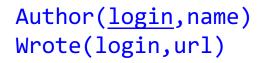
• Choose cheapest pricing tier!

Remember: WQ2 due on Friday

### Purchase(<u>pid</u>, product, quantity, price) GROUP BY V.S. NESTED QUERIES

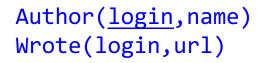
- SELECT product, Sum(quantity) AS TotalSales
- FROM Purchase
- WHERE price > 1
- **GROUP BY** product





## **MORE UNNESTING**

Find authors who wrote  $\geq$  10 documents:



## **MORE UNNESTING**

Find authors who wrote  $\geq$  10 documents:

Attempt 1: with nested queries

SELECT DISTINCT Author.name FROM Author WHERE (SELECT count(Wrote.url) FROM Wrote WHERE Author.login=Wrote.login) >= 10 This is SQL by a novice

## **MORE UNNESTING**

Find authors who wrote  $\geq$  10 documents:

Attempt 1: with nested queries

Attempt 2: using GROUP BY and HAVING

SELECTAuthor.nameFROMAuthor, WroteWHEREAuthor.login=Wrote.loginGROUP BY Author.nameHAVINGcount(wrote.url) >= 10

This is SQL by an expert

## **FINDING WITNESSES**

For each city, find the most expensive product made in that city

## **FINDING WITNESSES**

For each city, find the most expensive product made in that city

Finding the maximum price is easy...

```
SELECT x.city, max(y.price)
FROM Company x, Product y
WHERE x.cid = y.cid
GROUP BY x.city;
```

But we need the witnesses, i.e., the products with max price

## **FINDING WITNESSES**

To find the witnesses, compute the maximum price in a subquery (in FROM or in WITH)

```
WITH CityMax AS
 (SELECT x.city, max(y.price) as maxprice
 FROM Company x, Product y
 WHERE x.cid = y.cid
 GROUP BY x.city)
SELECT DISTINCT u.city, v.pname, v.price
FROM Company u, Product v, CityMax w
WHERE u.cid = v.cid
 and u.city = w.city
 and v.price = w.maxprice;
```

## **FINDING WITNESSES**

To find the witnesses, compute the maximum price in a subquery (in FROM or in WITH)

```
SELECT DISTINCT u.city, v.pname, v.price
FROM Company u, Product v,
    (SELECT x.city, max(y.price) as maxprice
    FROM Company x, Product y
    WHERE x.cid = y.cid
    GROUP BY x.city) w
WHERE u.cid = v.cid
    and u.city = w.city
    and v.price = w.maxprice;
```

## **FINDING WITNESSES**

Or we can use a subquery in where clause

## **FINDING WITNESSES**

There is a more concise solution here:

```
SELECT u.city, v.pname, v.price
FROM Company u, Product v, Company x, Product y
WHERE u.cid = v.cid and u.city = x.city
and x.cid = y.cid
GROUP BY u.city, v.pname, v.price
HAVING v.price = max(y.price)
```