

# CSE 344 Final Examination

December 16, 2015, 8:30am - 10:20am

Name: \_\_\_\_\_

Question	Points	Score
1	20	
2	30	
3	20	
4	30	
Total:	100	

- This exam is CLOSED book and CLOSED devices.
- You are allowed TWO letter-size pages with notes (both sides).
- You have 1h:50 minutes; budget time carefully.
- Please read all questions carefully before answering them.
- Some questions are easier, others harder; if a question sounds hard, skip it and return later.
- Good luck!

# 1 XML, XPath, and XQuery

1. (20 points)

(a) (10 points) Consider the following XML document stored in a file called data.xml:

```
<cluster>
  <machine id='M1'>
    <software>
      <dbms>MySQL</dbms>
      <dbms>PostgreSQL</dbms>
    </software>
    <hardware>
      <memory>16</memory>
      <cores>4</cores>
    </hardware>
  </machine>

  <machine id='M2'>
    <software>
      <dbms>MySQL</dbms>
    </software>
    <hardware>
      <memory>8</memory>
      <cores>2</cores>
    </hardware>
  </machine>

  <machine id='M3'>
    <software>
      <dbms>PostgreSQL</dbms>
    </software>
    <hardware>
      <memory>64</memory>
      <cores>16</cores>
    </hardware>
  </machine>

  <machine id='M4'>
    <software>
      <dbms>MongoDB</dbms>
    </software>
    <hardware>
      <memory>64</memory>
      <cores>16</cores>
    </hardware>
  </machine>
</cluster>
```



- (b) (10 points) Write an XQuery expression that will transform data.xml into the following document:

```
<result>
  <dbms>
    <name>MySQL</name>
    <nb_servers>2</nb_servers>
    <cores>4</cores>
    <cores>2</cores>
  </dbms>
  <dbms>
    <name>PostgreSQL</name>
    <nb_servers>2</nb_servers>
    <cores>4</cores>
    <cores>16</cores>
  </dbms>
  <dbms>
    <name>MongoDB</name>
    <nb_servers>1</nb_servers>
    <cores>16</cores>
  </dbms>
</result>
```

## 2 E/R Diagrams, Constraints, Conceptual Design

2. (30 points)

(a) (10 points) Design an E/R diagram describing the following domain:

- A **person** has attributes **id** (key), **fname**, and **lname**.
- A **company** has attributes **id** (key) and **name**. Each company name is unique.
- A **party** has attributes **pid**, **date**, and **duration**. The date attribute cannot be NULL. A party cannot last longer than 4 hours.
- A party is **booked** by either a person or a company. We need to capture information about both **when the party takes place** and **when the party was booked**. The booking date must be earlier than the party date. A party is booked by zero or one organizers. A person or company can book many parties.
- A party **occurs** in either a **house**, an **apartment**, or a **suite**. All these locations are uniquely identified with a **location id (lid)**. A house has an attribute **street address**. The address is an atomic string attribute. An apartment additionally has an **apartment number**. A suite has a street address and a **suite number**. A party occurs in zero or one place. A place can host multiple parties.
- A company **occupies** zero or one suite and a suite hosts zero or one company (i.e., the application does not store historical data only current data).
- A person **lives** in either a house or an apartment. A person lives in zero or one place. A given place can host many people.

Answer (Draw an E/R Diagram on the next page):

Answer (Draw an E/R Diagram):

- (b) (10 points) Write the CREATE TABLE statements necessary to capture the **subset of the above ER diagram** that corresponds to **party, person, and company** entities as well as the **books** relationship. Include all primary key, foreign key, and other constraints. Avoid creating unnecessary tables when possible. Dates can be represented with the `datetime` type.

Answer (Write CREATE TABLE statements for **SUBSET** of ER diagram):

- (c) (10 points) Consider the following relational schema and set of functional dependencies.

$R(A,B,C,D,E,F,G)$  with functional dependencies:

$E \rightarrow C$

$G \rightarrow AD$

$B \rightarrow E$

$C \rightarrow BF$

- Give one example of non-trivial functional dependency implied by the ones above:

Answer (Example FD):

- Compute  $E^+$ , the closure of  $E$ .

Answer ( $E^+$ ):

- Why do we bother to decompose relations into BCNF? What does this normal form ensure?

Answer (Explanation):



- Decompose R into BCNF. Show your work for partial credit. Your answer should consist of a list of table names and attributes and an indication of the keys in each table (underlined attributes).

Answer (Decompose R into BCNF):

### 3 Transactions

3. (20 points)

- (a) (10 points) Consider transactions executing concurrently on the same instance of SQL Server. For each of the execution traces below, circle the isolation levels (if any) that will allow the given interleaved execution to occur.

```
BEGIN TRANSACTION
```

```
SELECT * from Customers
```

```
SELECT * from Products
```

```
COMMIT
```

```
BEGIN TRANSACTION
```

```
SELECT * from Customers
```

```
SELECT * from Products
```

```
COMMIT
```

Serializable   Repeatabe Read   Read Committed

```
BEGIN TRANSACTION
```

```
UPDATE Customers SET id = 20  
WHERE id = 10
```

```
SELECT * from Customers
```

```
COMMIT
```

```
BEGIN TRANSACTION
```

```
UPDATE Products SET price = 0  
WHERE price is NULL
```

```
SELECT * from Products
```

```
COMMIT
```

Serializable Repeatable Read Read Committed

```
BEGIN TRANSACTION
```

```
UPDATE Customers SET id = 10  
WHERE id = 20
```

```
SELECT * from Products
```

```
COMMIT
```

```
BEGIN TRANSACTION
```

```
UPDATE Products SET price = 10  
WHERE price = 0
```

```
SELECT * from Customers
```

```
COMMIT
```

Serializable Repeatable Read Read Committed

```
BEGIN TRANSACTION
```

```
SELECT * from Customers
```

```
COMMIT
```

```
BEGIN TRANSACTION
```

```
INSERT INTO Customers VALUES(1)
```

```
SELECT * from Customers
```

```
COMMIT
```

Serializable Repeatable Read Read Committed

- (b) (5 points) Consider the following transaction schedules. For each schedule, draw the precedence graph and indicate if it is **conflict-serializable** or not.

r1(A); w1(B); r2(B); w2(C); r3(C); w3(A);

Answer (YES/NO):

r1(A); r2(B); r3(B); w3(A); w2(C); r3(D); r3(C); w1(B);

Answer (YES/NO):

(c) (5 points) What is strict 2PL?

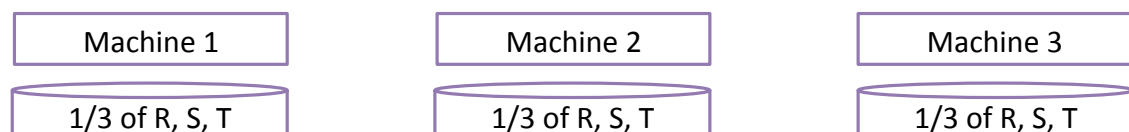
## 4 Parallel Data Processing

4. (30 points)

- (a) (10 points) Consider relations  $R(a,b)$ ,  $S(c,d)$ , and  $T(e,f)$ . All three are horizontally partitioned across  $N = 3$  machines as shown in the diagram below. Each machine locally stores approximately  $\frac{1}{N}$  of the tuples in  $R$ ,  $S$ , and  $T$ . The tuples are randomly organized across machines (i.e.,  $R$  is block partitioned across machines). Show a relational algebra plan for the following query and how it will be executed across the  $N = 3$  machines. **Use hash-join (a.k.a shuffle-join) operators.** Include operators that need to re-shuffle data and add a note explaining how these operators will re-shuffle that data.

```
SELECT *
FROM R, S, T
WHERE R.b = S.c
AND S.d = T.e
AND (R.a - T.f) > 100
```

Answer (Draw the parallel query plan):





- (b) (10 points) Now consider the case where the **R** relation is very large and both **S** and **T** are very small. Show a plan that uses **broadcast joins** to compute the result of the query.

- (c) (10 points) Explain how the query would be executed in **MapReduce** (not Pig). Make sure to specify the computations performed in the map and the reduce functions. Use **hash-joins** (shuffle joins). No need to give the pseudocode. Just state what each function does and what it outputs in plain text format.

**Answer** (Describe the map and reduce functions of the executed MapReduce job(s)):