CSE 344 Introduction to Data Management

Section 4: Relational Algebra

Outline

- HW3 Check-in
- Relational Algebra Review
- Translate nested SQL Queries to RA
- Translate from RA to SQL

Relational Algebra

- SQL = WHAT we want to get from the data
- Relational Algebra = HOW to get the data we want
- SQL → Relational Algebra →Physical Plan
- Relational Algebra = Logical Plan (usually written as a tree)

Relational Algebra Operators

Standard:

- Selection: σ
  Projection: π
  Rename: ρ

Sets:

- Union: ∪
  Intersection: ∩
  Difference: -

Joins:

- Cartesian Product: X
  Join: ⋈

Extended:

- Duplicate Elimination: δ
  Grouping and Aggregation: ɣ
  Sorting: 

SQL to RA Review

Dancer(did, name, birthyear, country)
Show(sid, title, choreographer, composer, year)
Role(did, sid, role, company)

- Write a Relational Algebra plan for the following query:

```
SELECT d.did, d.name, count(*)
FROM Dancer d, Show s, Role r
WHERE d.did=r.did
AND r.sid=s.sid
AND s.composer='Tchaikovsky'
GROUP BY d.did, d.name
ORDER BY d.name;
```

SQL to RA Solution

```
SELECT d.did, d.name, count(*)
FROM Dancer d, Show s, Role r
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Dancer d

Role r

Show s

SQL to RA Solution

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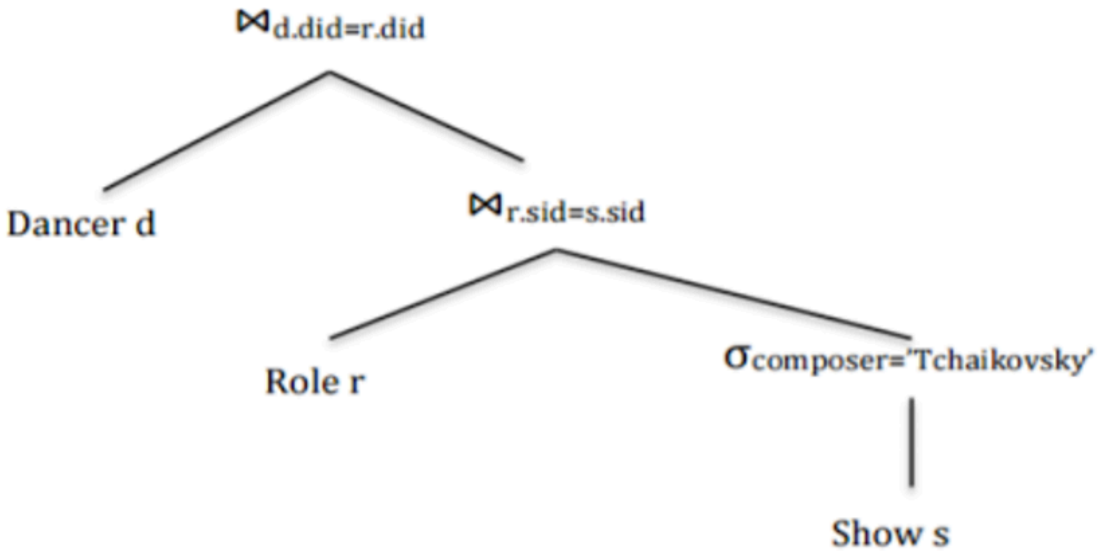

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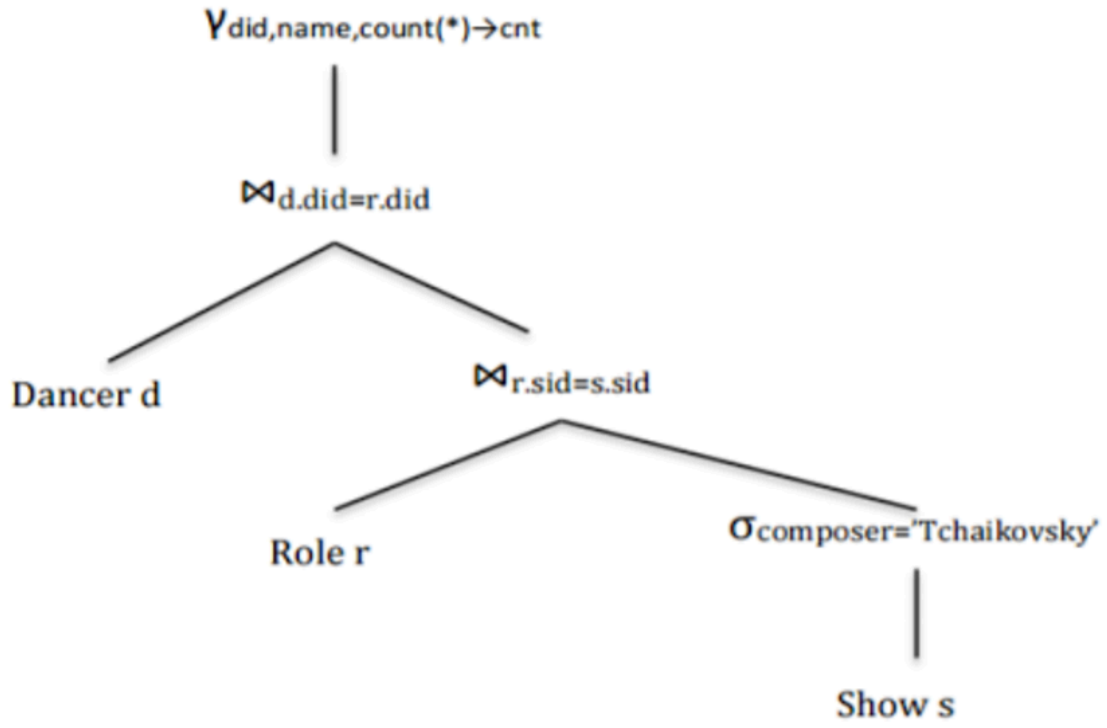

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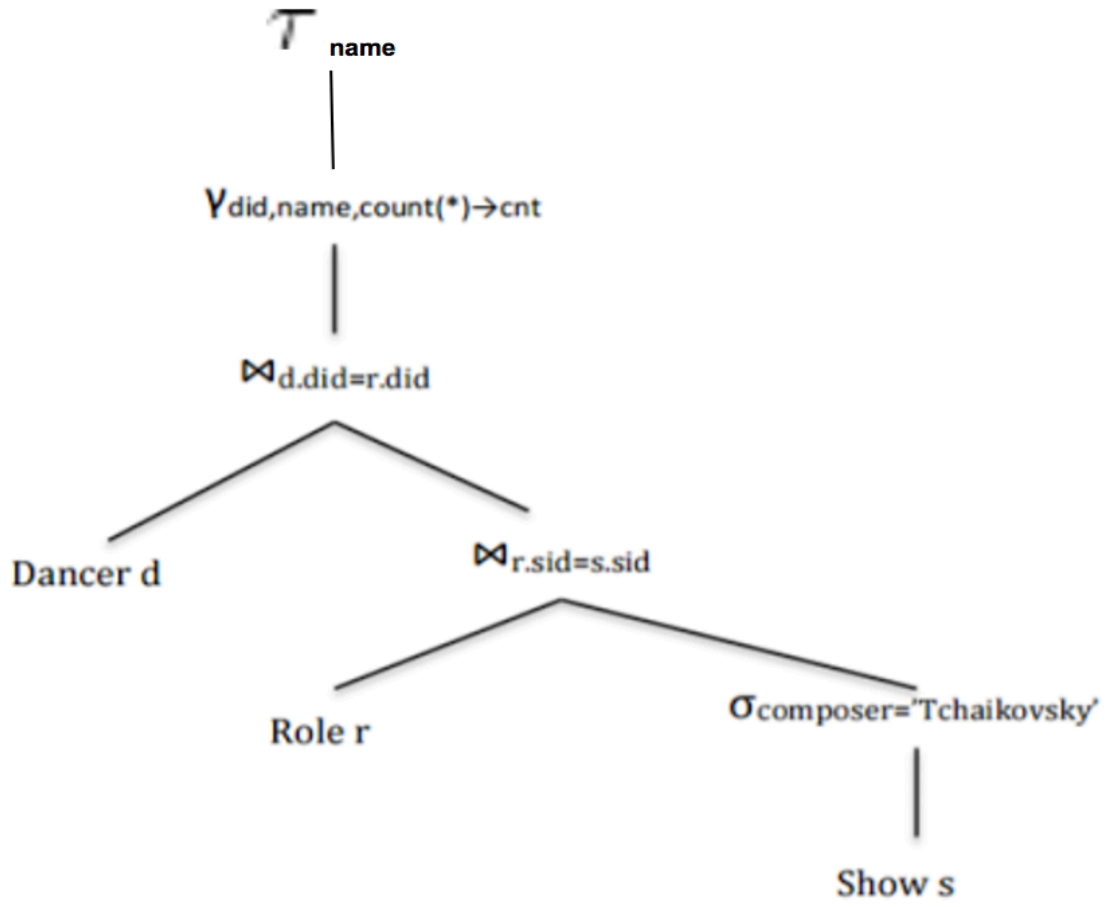

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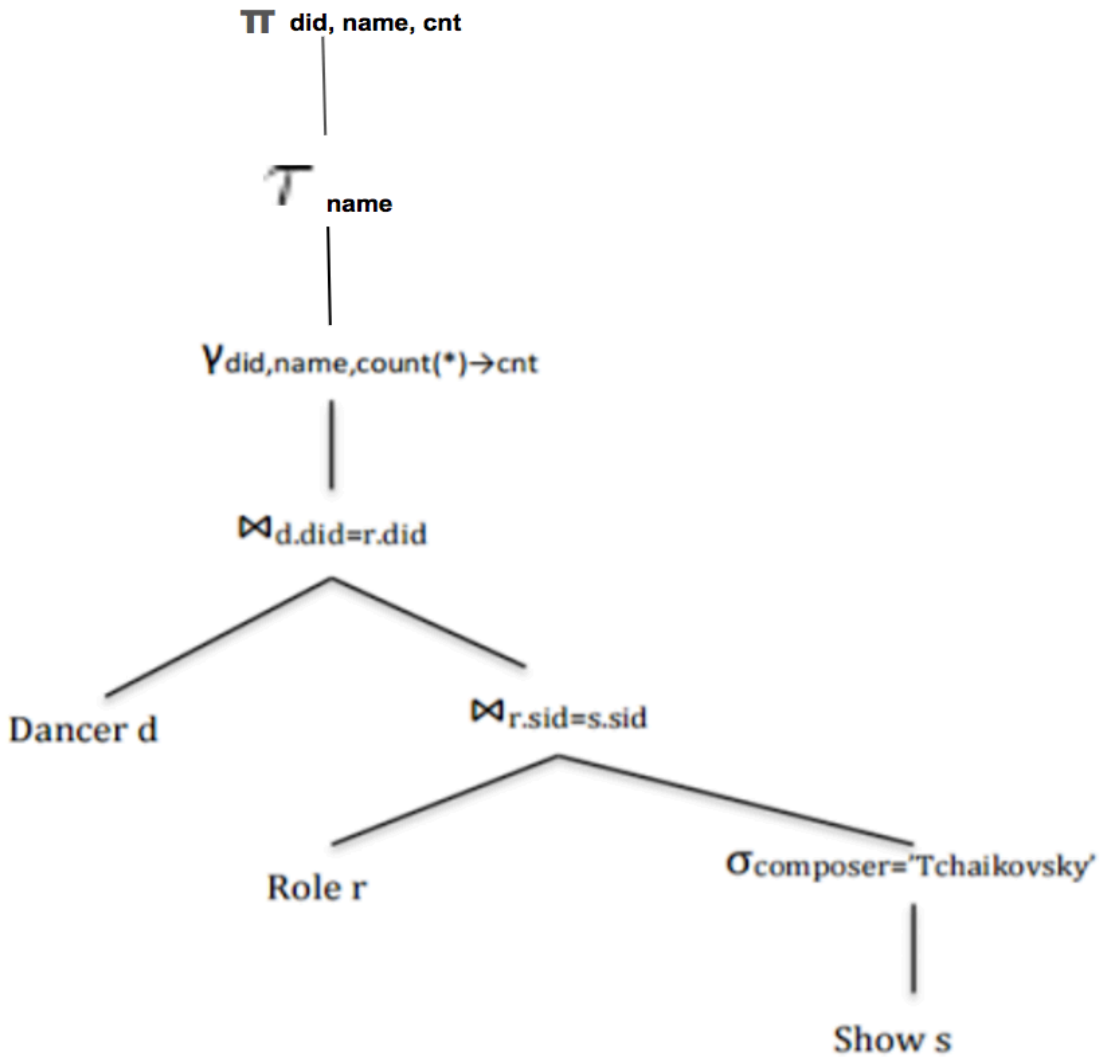

Translate nested SQL Queries to RA

Member(mid, name, age)

Picture(pid, year)

Tagged(mid, pid)

```
SELECT w.year, max(w.c) AS m
FROM(SELECT x.name, z.year, count(*) AS c
        FROMMember x, Tagged y, Picture z
        WHERE x.mid = y.mid
        AND y.pid = z.pid
        AND age < 20
        GROUP BY x.name, z.year) w
GROUP BY w.year
HAVING sum(w.c) > 100;
```

Nested SQL Queries to RA Solution

```
SELECT w.year, max(w.c) AS m
FROM(SELECT x.name, z.year, count(*) AS c
        FROM Member x, Tagged y, Picture z
        WHERE x.mid = y.mid
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        GROUP BY x.name, z.year) w
GROUP BY w.year
HAVING sum(w.c) > 100;
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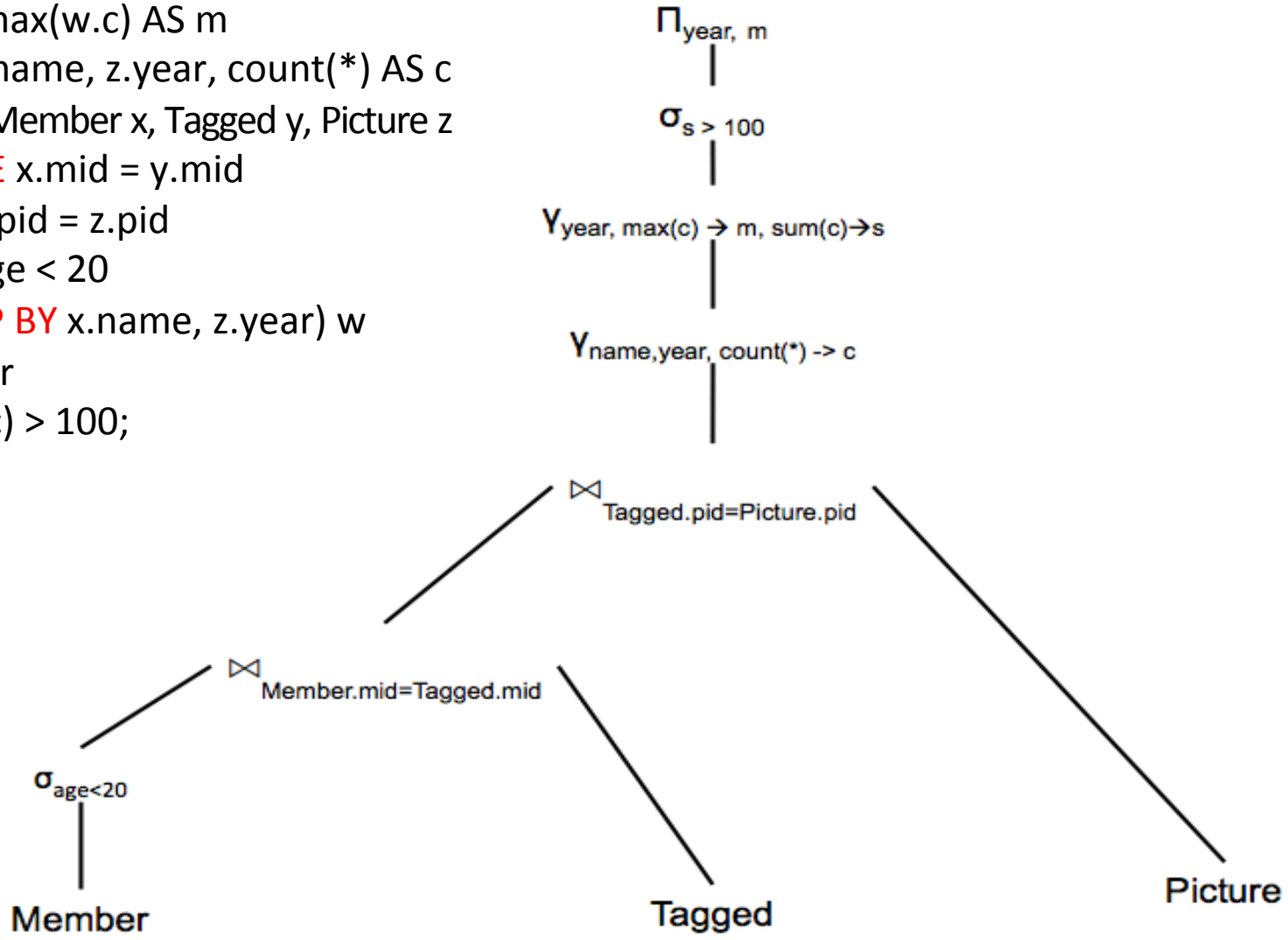
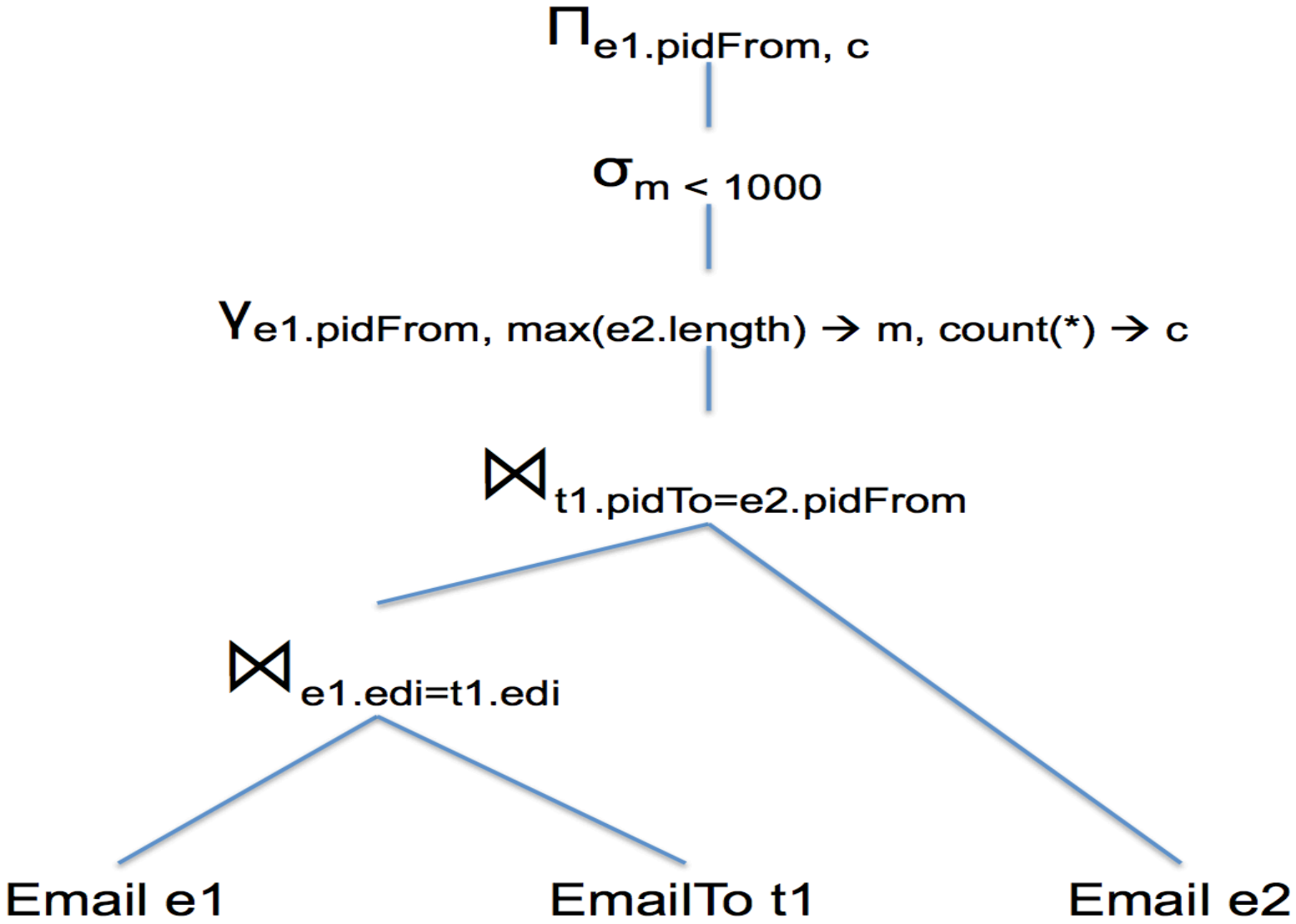


Translate from RA to SQL

- Put tables in FROM clause
- Put join predicates in WHERE clause
- Put selection predicates in WHERE clause
- Translate extended RA symbols to SQL equivalent
- Put selection of aggregates in HAVING clause
- Put projection predicates in SELECT clause

RA to SQL Example

Person(pid,name)
Email(eid, pidFrom, tid, body, length)
EmailTo(eid,pidTo)



RA to SQL Solution

```
SELECT e1.pidFrom, count(*)
FROM Email e1, EmailTo t1, Email e2
WHERE e1.eid = t1.eid
AND t1.pidTo = e2.pidFrom
GROUP BY e1.pidFrom
HAVING max(e2.length) < 1000;
```