

# Introduction to Data Management CSE 344

## Lecture 17: E/R Diagrams and Constraints

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## Announcements

- HW5 due this Friday
  - Please note minor update to the instructions
- WQ6 due next Wednesday

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## Database Design

What it is:

- Starting from scratch, design the database schema: relation, attributes, keys, foreign keys, constraints etc

Why it's hard

- The database will be in operation for a very long time (years). Updating the schema while in production is very expensive (why?)

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## Database Design

- Consider issues such as:
  - What entities to model
  - How entities are related
  - What constraints exist in the domain
- Several formalisms exists
  - We discuss E/R diagrams
- Reading: Sec. 4.1-4.6

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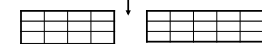
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## Database Design Process

Conceptual Model:

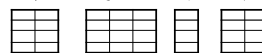


Relational Model:  
Tables + constraints  
And also functional dep.



Normalization:  
Eliminates anomalies

Conceptual Schema



Physical storage details

Physical Schema



## Entity / Relationship Diagrams

- Entity set = a class
  - An entity = an object
- Attribute
- Relationship

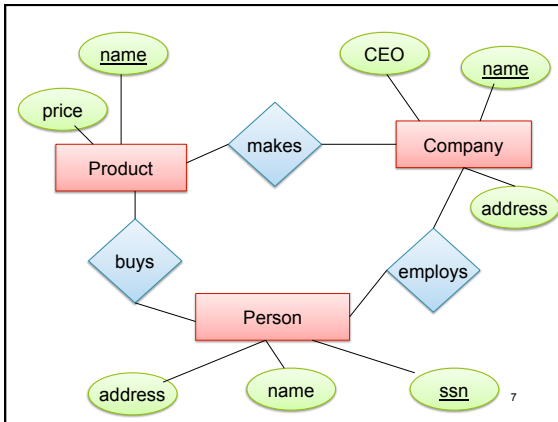
Product

city

makes

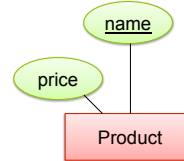
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## Keys in E/R Diagrams

- Every entity set must have a key

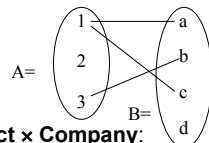


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## What is a Relation ?

- A mathematical definition:
  - if A, B are sets, then a relation R is a subset of  $A \times B$
- $A = \{1, 2, 3\}$ ,  $B = \{a, b, c, d\}$ ,  
 $A \times B = \{(1, a), (1, b), \dots, (3, d)\}$   
 $R = \{(1, a), (1, c), (3, b)\}$



- **makes** is a subset of **Product**  $\times$  **Company**:



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## Multiplicity of E/R Relations

- one-one:



- many-one

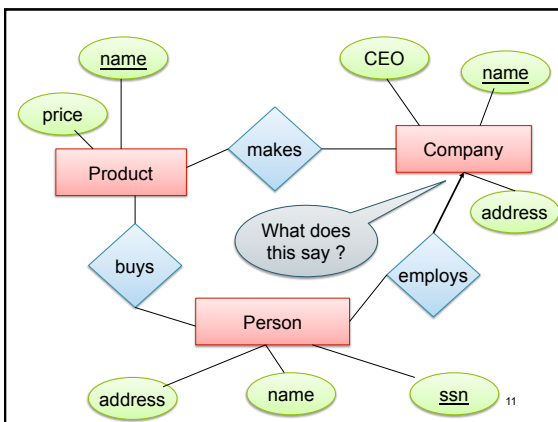


- many-many



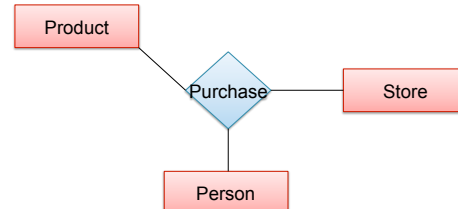
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## Multi-way Relationships

How do we model a purchase relationship between buyers, products and stores?



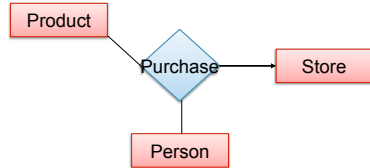
Can still model as a mathematical set (Q. how ?)

A. As a set of triples  $\subseteq \text{Person} \times \text{Product} \times \text{Store}$

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## Arrows in Multiway Relationships

Q: What does the arrow mean ?



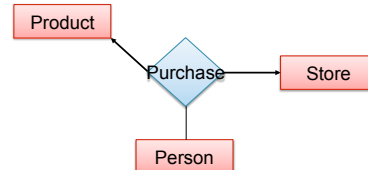
A: A given person buys a given product from at most one store  
 [Arrow pointing to E means that if we select one entity from each of the other entity sets in the relationship, those entities are related to at most one entity in E]

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## Arrows in Multiway Relationships

Q: What does the arrow mean ?

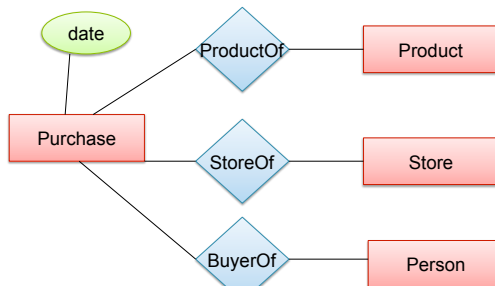


A: A given person buys a given product from at most one store  
 AND every store sells to every person at most one product

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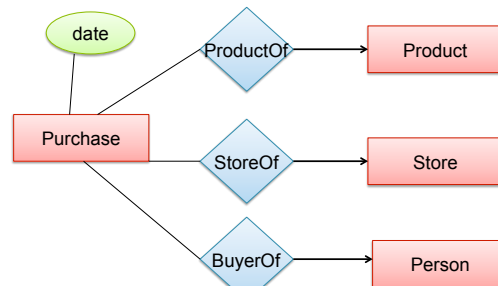
## Converting Multi-way Relationships to Binary



Arrows go in which direction?

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## Converting Multi-way Relationships to Binary

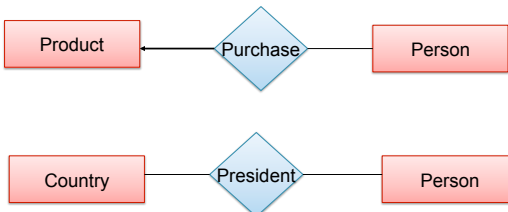


Make sure you understand why!

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## 3. Design Principles

What's wrong?

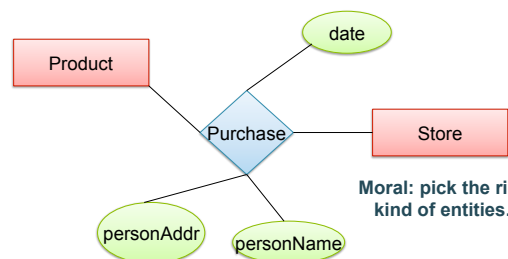


Moral: be faithful to the specifications of the app!

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## Design Principles: What's Wrong?

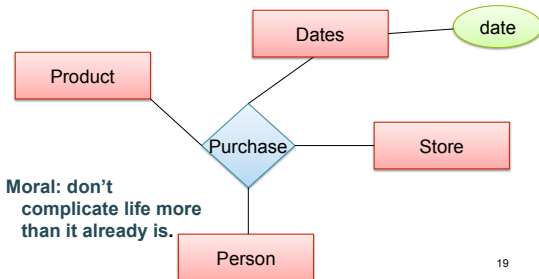


Moral: pick the right kind of entities.

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## Design Principles: What's Wrong?



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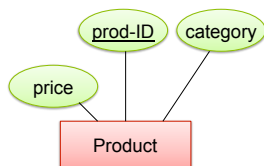
## From E/R Diagrams to Relational Schema

- Entity set  $\rightarrow$  relation
- Relationship  $\rightarrow$  relation

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## Entity Set to Relation

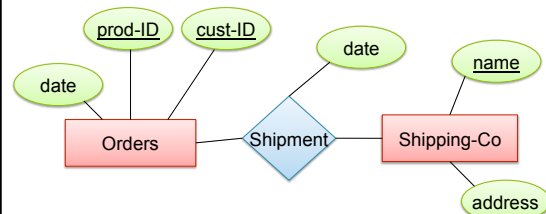


**Product**(prod-ID, category, price)

<u>prod-ID</u>	category	price
Gizmo55	Camera	99.99
Pokemn19	Toy	29.99

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## N-N Relationships to Relations

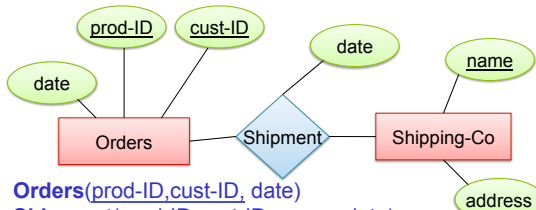


Represent this in relations

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## N-N Relationships to Relations



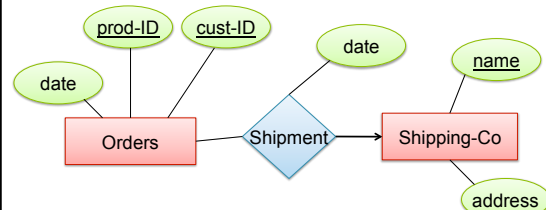
**Orders**(prod-ID, cust-ID, date)

**Shipment**(prod-ID, cust-ID, name, date)

**Shipping-Co**(name, address)

<u>prod-ID</u>	<u>cust-ID</u>	<u>name</u>	date
Gizmo55	Joe12	UPS	4/10/2011
Gizmo55	Joe12	FEDEX	4/9/2011

## N-1 Relationships to Relations

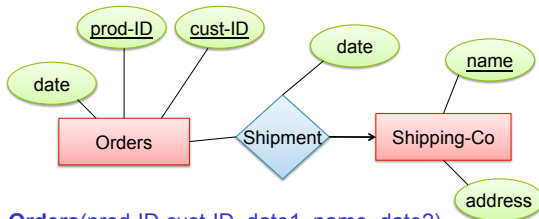


Represent this in relations

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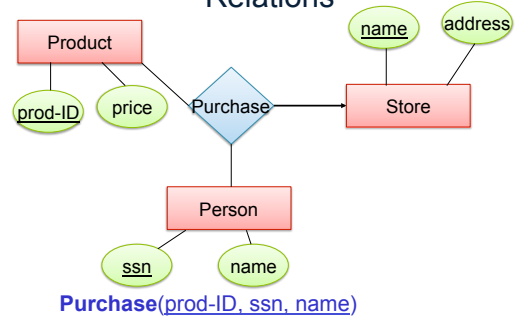
## N-1 Relationships to Relations



**Orders**(prod-ID, cust-ID, date1, name, date2)  
**Shipping-Co**(name, address)

Remember: no separate relations for many-one relationship

## Multi-way Relationships to Relations



**Purchase**(prod-ID, ssn, name)

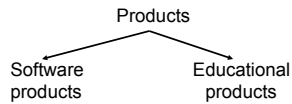
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## Modeling Subclasses

Some objects in a class may be special

- define a new class
- better: define a *subclass*

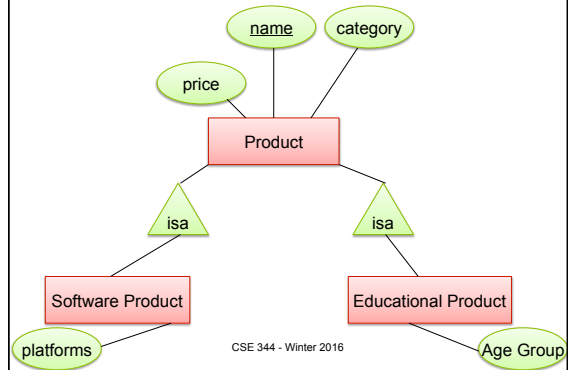


So --- we define subclasses in E/R

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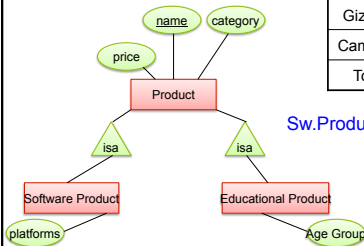
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## Subclasses



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## Subclasses to Relations



Other ways to convert are possible

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**Product**

<u>Name</u>	<u>Price</u>	<u>Category</u>
Gizmo	99	gadget
Camera	49	photo
Toy	39	gadget

**Sw.Product**

<u>Name</u>	<u>platforms</u>
Gizmo	unix

**Ed.Product**

<u>Name</u>	<u>Age Group</u>
Gizmo	toddler
Toy	retired

## Modeling Union Types with Subclasses

FurniturePiece

Person

Company

Say: each piece of furniture is owned either by a person or by a company

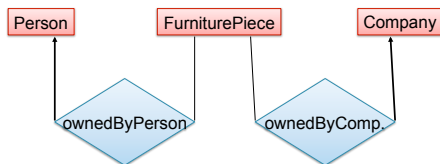
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## Modeling Union Types with Subclasses

Say: each piece of furniture is owned either by a person or by a company

Solution 1. Acceptable but imperfect (What's wrong ?)

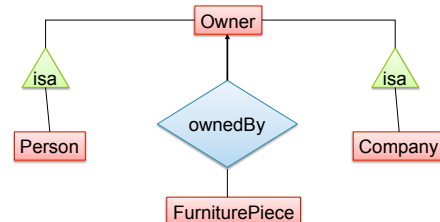


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## Modeling Union Types with Subclasses

Solution 2: better, more laborious

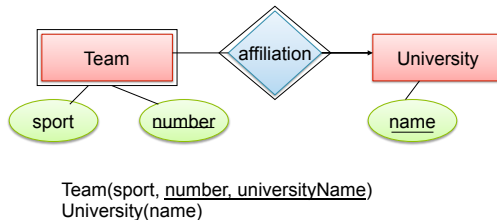


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## Weak Entity Sets

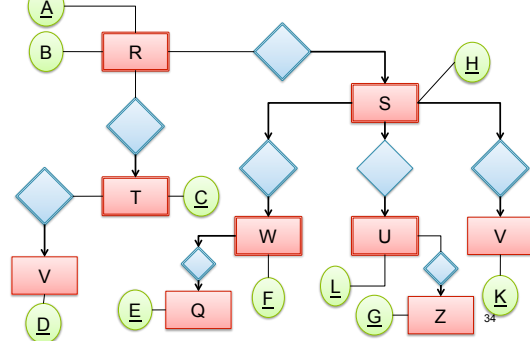
Entity sets are weak when their key comes from other classes to which they are related.



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## What Are the Keys of R ?



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## What makes good schemas?



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## Integrity Constraints Motivation

An integrity constraint is a condition specified on a database schema that restricts the data that can be stored in an instance of the database.

- ICs help prevent entry of incorrect information
- How? DBMS enforces integrity constraints
  - Allows only legal database instances (i.e., those that satisfy all constraints) to exist
  - Ensures that all necessary checks are always performed and avoids duplicating the verification logic in each application

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## Constraints in E/R Diagrams

Finding constraints is part of the modeling process.  
Commonly used constraints:

**Keys:** social security number uniquely identifies a person.

**Single-value constraints:** a person can have only one father.

**Referential integrity constraints:** if you work for a company, it must exist in the database.

**Other constraints:** peoples' ages are between 0 and 150.

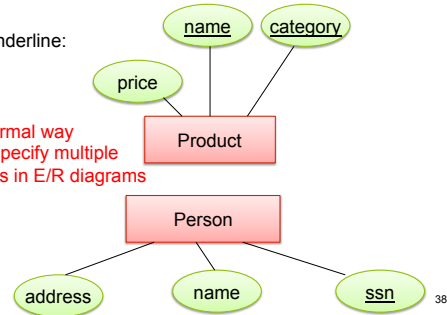
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## Keys in E/R Diagrams

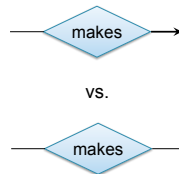
Underline:

No formal way  
to specify multiple  
keys in E/R diagrams



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## Single Value Constraints



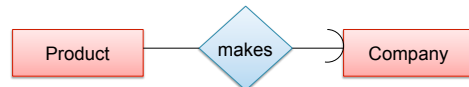
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## Referential Integrity Constraints



Each product made by at most one company.  
Some products made by no company

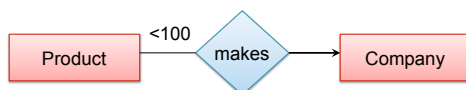


Each product made by exactly one company.

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## Other Constraints



Q: What does this mean ?

A: A Company entity cannot be connected by relationship to more than 99 Product entities

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## Constraints in SQL

Constraints in SQL:

- Keys, foreign keys
- Attribute-level constraints
- Tuple-level constraints
- Global constraints: assertions

simplest

Most complex

- The more complex the constraint, the harder it is to check and to enforce

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## Key Constraints

Product(name, category)

```
CREATE TABLE Product (
  name CHAR(30) PRIMARY KEY,
  category VARCHAR(20))
```

OR: 

```
CREATE TABLE Product (
  name CHAR(30),
  category VARCHAR(20),
  PRIMARY KEY (name))
```

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## Keys with Multiple Attributes

Product(name, category, price)

```
CREATE TABLE Product (
  name CHAR(30),
  category VARCHAR(20),
  price INT,
  PRIMARY KEY (name, category))
```

Name	Category	Price
Gizmo	Gadget	10
Camera	Photo	20
Gizmo	Photo	30
<del>Gizmo</del>	<del>Gadget</del>	<del>40</del>

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## Other Keys

```
CREATE TABLE Product (
  productID CHAR(10),
  name CHAR(30),
  category VARCHAR(20),
  price INT,
  PRIMARY KEY (productID),
  UNIQUE (name, category))
```

There is at most one **PRIMARY KEY**;  
there can be many **UNIQUE**

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## Foreign Key Constraints

```
CREATE TABLE Purchase (
  prodName CHAR(30)
  REFERENCES Product(name),
  date DATETIME)
```

Referential  
integrity  
constraints

prodName is a **foreign key** to Product(name)  
name must be a **key** in Product

May write  
just Product  
if name is PK

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## Foreign Key Constraints

- Example with multi-attribute primary key

```
CREATE TABLE Purchase (
  prodName CHAR(30),
  category VARCHAR(20),
  date DATETIME,
  FOREIGN KEY (prodName, category)
  REFERENCES Product(name, category))
```

- (name, category) must be a **KEY** in Product

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## What happens when data changes?

Types of updates:

- In Purchase: insert/update
- In Product: delete/update

Product		Purchase	
Name	Category	ProdName	Store
Gizmo	gadget	Gizmo	Wiz
Camera	Photo	Camera	Ritz
OneClick	Photo	Camera	Wiz

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## What happens when data changes?

- SQL has three policies for maintaining referential integrity:
- NO ACTION** reject violating modifications (default)
- CASCADE** after delete/update do delete/update
- SET NULL** set foreign-key field to NULL
- SET DEFAULT** set foreign-key field to default value
  - need to be declared with column, e.g.,  
CREATE TABLE Product (pid INT DEFAULT 42)

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## Maintaining Referential Integrity

```
CREATE TABLE Purchase (
  prodName CHAR(30),
  category VARCHAR(20),
  date DATETIME,
  FOREIGN KEY (prodName, category)
  REFERENCES Product(name, category)
  ON UPDATE CASCADE
  ON DELETE SET NULL )
```

Product		Purchase	
Name	Category	ProdName	Category
Gizmo	gadget	Gizmo	Gizmo
Camera	Photo	Snap	Camera
OneClick	Photo	EasyShoot	Camera

## Constraints on Attributes and Tuples

- Constraints on attributes:
  - NOT NULL** -- obvious meaning...
  - CHECK** condition -- any condition !
- Constraints on tuples
  - CHECK** condition

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## Constraints on Attributes and Tuples

```
CREATE TABLE R (
  A int NOT NULL,
  B int CHECK (B > 50 and B < 100),
  C varchar(20),
  D int,
  CHECK (C >= 'd' or D > 0))
```

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## Constraints on Attributes and Tuples

```
CREATE TABLE Product (
  productID CHAR(10),
  name CHAR(30),
  category VARCHAR(20),
  price INT CHECK (price > 0),
  PRIMARY KEY (productID),
  UNIQUE (name, category))
```

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## Constraints on Attributes and Tuples

What does this constraint do?

```
CREATE TABLE Purchase (
  prodName CHAR(30)
  CHECK (prodName IN
  (SELECT Product.name
  FROM Product),
  date DATETIME NOT NULL)
```

What is the difference from Foreign-Key ?

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## General Assertions

```
CREATE ASSERTION myAssert CHECK  
(NOT EXISTS(  
  SELECT Product.name  
  FROM Product, Purchase  
  WHERE Product.name = Purchase.prodName  
  GROUP BY Product.name  
  HAVING count(*) > 200) )
```

But most DBMSs do not implement assertions  
Because it is hard to support them efficiently  
Instead, they provide triggers

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