

# Introduction to Data Management

## CSE 344

### Lecture 13: Relational Calculus

# Announcements

- WQ 4 is out
- HW 4 is out
  
- Midterm review session in class next Fri (11/4)
  
- OHs changes next few weeks
  - Check website for most updated info

# Big Picture

- Relational data model
  - Instance
  - Schema
  - Query language
    - SQL
    - Relational algebra
    - Relational calculus
    - Datalog
- Query processing
  - Logical & physical plans
  - Indexes
  - Cost estimation
  - Query optimization

# Why bother with another QL?

- SQL and RA are good for query planning
  - They are not good for *formal reasoning*
  - How do you show that two SQL queries are equivalent / non-equivalent?
  - Two RA plans?
- RC was the first language proposed with the relational model (Codd)
- Influenced the design of datalog as we will see

# Relational Calculus

- Aka predicate calculus or first order logic
  - 311 anyone?
- TRC = Tuple Relational Calculus
  - See book
- DRC = Domain Relational Calculus
  - We study only this one
  - Also see *Query Language Primer* on course website

# Relational Calculus

Query Q:

$$Q(x_1, \dots, x_k) = P$$

Relational predicate  $P$  is a formula given by this grammar:

$$P ::= \text{atom} \mid P \wedge P \mid P \vee P \mid P \Rightarrow P \mid \text{not}(P) \mid \forall x.P \mid \exists x.P$$

Atomic predicate is either a relational or interpreted predicate:

$$\text{atom} ::= R(x_1, \dots, x_k) \mid x = y \mid x > k \mid \dots$$

$R(x,y)$  means  $(x,y)$  is in  $R$

Actor(pid,fName,lName)

Casts(pid,mid)

Movie(mid,title,year)

# Relational Calculus

Query Q:

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Relational predicate P is a formula given by this grammar:

$$P ::= \text{atom} \mid P \wedge P \mid P \vee P \mid P \Rightarrow P \mid \text{not}(P) \mid \forall x.P \mid \exists x.P$$

Atomic predicate is either a relational or interpreted predicate:

$$\text{atom} ::= R(x_1, \dots, x_k) \mid x = y \mid x > k \mid \dots \quad R(x,y) \text{ means } (x,y) \text{ is in } R$$

---

Example: find the first/last names of actors who acted in 1940

$$Q(f,l) = \exists x. \exists y. \exists z. (\text{Actor}(z,f,l) \wedge \text{Casts}(z,x) \wedge \text{Movie}(x,y,1940))$$

What does this query return ?

$$Q(f,l) = \exists z. (\text{Actor}(z,f,l) \wedge \forall x. (\text{Casts}(z,x) \Rightarrow \exists y. \text{Movie}(x,y,1940)))$$

Likes(drinker, beer)

Frequents(drinker, bar)

Serves(bar, beer)

# Important Observation

Find all bars that serve all beers that Fred likes

$$A(x) = \forall y. \text{Likes}(\text{"Fred"}, y) \Rightarrow \text{Serves}(x, y)$$

- Note:  $P \Rightarrow Q$  (read P implies Q) is the same as  $(\text{not } P) \vee Q$

In this query: If Fred likes a beer the bar must serve it ( $P \Rightarrow Q$ )

In other words: Either Fred does not like the beer ( $\text{not } P$ ) OR the bar serves that beer ( $Q$ ).

$$A(x) = \forall y. \text{not}(\text{Likes}(\text{"Fred"}, y)) \vee \text{Serves}(x, y)$$



Likes(drinker, beer)  
Frequents(drinker, bar)  
Serves(bar, beer)

# More Examples

Average Joe

Find drinkers that frequent some bar that serves some beer they like.

Likes(drinker, beer)  
Frequents(drinker, bar)  
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# More Examples

Average Joe

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$$Q(x) = \exists y. \exists z. \text{Frequents}(x, y) \wedge \text{Serves}(y, z) \wedge \text{Likes}(x, z)$$

Likes(drinker, beer)  
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$$Q(x) = \exists y. \exists z. \text{Frequents}(x, y) \wedge \text{Serves}(y, z) \wedge \text{Likes}(x, z)$$

Prudent Peter

Find drinkers that frequent only bars that serves some beer they like.

Likes(drinker, beer)  
Frequents(drinker, bar)  
Serves(bar, beer)

# More Examples

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Cautious Carl

Find drinkers that frequent some bar that serves only beers they like.

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# Remember your logical equivalences!

- $A \Rightarrow B = \text{not}(A) \vee B$
- $\text{not}(A \wedge B) = \text{not}(A) \vee \text{not}(B)$
- $\text{not}(A \vee B) = \text{not}(A) \wedge \text{not}(B)$
- $\forall x. P(x) = \text{not}(\exists x. \text{not}(P(x)))$
  
- Example:
  - $\forall z. \text{Serves}(y,z) \Rightarrow \text{Likes}(x,z)$
  - $\forall z. \text{not}(\text{Serves}(y,z)) \vee \text{Likes}(x,z)$
  - $\text{not}(\exists z. \text{Serves}(y,z) \wedge \text{not}(\text{Likes}(x,z)))$

Likes(drinker, beer)

Frequents(drinker, bar)

Serves(bar, beer)

# Domain Independent Relational Calculus

- An unsafe RC query, aka domain dependent, returns an answer that does not depend just on the relations, but on the entire domain of possible values

$A1(x) = \text{not Likes}(\text{"Fred"}, x)$

$A1(x) = \exists y \text{ Serves}(y, x) \wedge \text{not Likes}(\text{"Fred"}, x)$

Make sure x is a beer

Likes(drinker, beer)

Frequents(drinker, bar)

Serves(bar, beer)

# Domain Independent Relational Calculus

- An unsafe RC query, aka domain dependent, returns an answer that does not depend just on the relations, but on the entire domain of possible values

Make sure x is a beer

$A1(x) = \text{not Likes}(\text{"Fred"}, x)$

$A1(x) = \exists y \text{ Serves}(y, x) \wedge \text{not Likes}(\text{"Fred"}, x)$

$A2(x, y) = \text{Likes}(\text{"Fred"}, x) \vee \text{Serves}(\text{"Bar"}, y)$

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$A2(x, y) = \text{Likes}(\text{"Fred"}, x) \vee \text{Serves}(\text{"Bar"}, y)$

Same here

$A2(x, y) = \exists u \text{ Serves}(u, x) \wedge \exists w \text{ Serves}(w, y) \wedge [\text{Likes}(\text{"Fred"}, x) \vee \text{Serves}(\text{"Bar"}, y)]$

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$$A3(x) = \forall y. \text{Serves}(x, y)$$

Likes(drinker, beer)

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$$A3(x) = \forall y. \text{Serves}(x,y)$$

$$A3(x) = \exists u. \text{Serves}(x,u) \wedge \forall y. \exists z. \text{Serves}(z,y) \rightarrow \text{Serves}(x,y)$$

Likewise

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Likewise

**Lesson: make sure your RC queries are domain independent**