CSE 344 Final Examination

March 15, 2016, 2:30pm - 4:20pm

Name: _____

Question	Points	Score
1	47	
2	17	
3	36	
4	54	
5	46	
Total:	200	

- This exam is CLOSED book and CLOSED devices.
- You are allowed TWO letter-size pages with notes (both sides).
- You have 1h:50 minutes; budget time carefully.
- Please read all questions carefully before answering them.
- Some questions are easier, others harder; if a question sounds hard, skip it and return later.
- Good luck!

1 SQL

- 1. (47 points)
 - (a) (7 points) We represent sparse matrices as tables with three attributes: row, colummn, value. Write a SQL query that computes the product of two sparse matrices called A and B. Recall that the product C = AB of two matrices is defined as:

$$C_{ik} = \sum_{j} A_{ij} B_{jk}$$

For example, consider the matrix product below:

	1	0	0]	0	2	1] [0	2	1]
A =	0	3	0	B =	1	0	0	AB =	3	0	0
	2	0	1		0	0	1 _] [0	4	3

Then the inputs and output to your query would be:

А	row	col	val	I	3 [row	col	val
	1	1	1			1	2	2
	2	2	3			1	3	1
	3	1	2			2	1	1
	3	3	1			3	3	1

Answer	row	col	val
	1	2	2
	1	3	1
	2	1	3
	3	2	4
	3	3	3

Write a SQL query:

(b) (10 points) The traces of a matrix A is the sum of the elements on the diagonal: $Tr(A) = \sum_{i} A_{ii}$. Write a SQL query that computes Tr(ABC), where A, B, C are three sparse matrices.

Write a SQL query:

- CSE 344
 - (c) Consider three relations with schemas:

R(A, B), S(C, D), T(E, F)

All attributes are NOT NULL, and no attribute is a key. For each query below, write an equivalent SQL query that is unnested. For example, if the given query were

select distinct x.A from (select y.A as A from R y where y.B = 2) x; then your answer would be:

select x.A from R x where x.B = 2;

If query cannot be unnested, then indicate so.

i. (5 points) Unnest this query: select distinct x.A, (select sum(y.B) from R y where x.A = y.A) as S from R x; With the select sum(y.B) from R y where x.A = y.A) as S

Write an unnested SQL query or say *impossible*:

ii. (5 points) Unnest this query: select x.A, x.B, (select sum(y.D) from S y where x.B = y.C) as S from R x; Write an unnested SQL query or say *impossible*: iii. (5 points) Unnest this query: select distinct x.A as A, y.F as F from R x, (select u.C as C, v.F as F from S u, T v where u.d = v.E) y where x.B = y.C; Write an unnested SQL query or say impossible:

iv. (5 points) Unnest this query: select x.A as A, y.F as F from R x, (select distinct u.C as C, v.F as F from S u, T v where u.D = v.E) y where x.B = y.C; Write an unnested SQL query or say impossible: (d) Consider four relations with the following schemas:

$$R(A, B), R2(A, B), S(C, D), V(G, H, K)$$

For each of the identities in the Relational Algebra below, indicate whether they hold. Assume set semantics for all operators:

i. (2 points) Does this identity hold?

$$(R - R2) - R2 = R$$

Yes/No:

ii. (2 points) Does this identity hold?

$$((R - R2) - R2) - R2 = R - R2$$

ii. _____

Yes/No:

iii. (2 points) Does this identity hold?

$$\Pi_A(R \bowtie_{B=C} S) = \Pi_A(R)$$

iii. _____

Yes/No:

iv. (2 points) Does this identity hold?

$$\Pi_{AB}(R \bowtie_{B=C} S) - R2 = \Pi_{AB}((R - R2) \bowtie_{B=C} S)$$

iv. _____

Yes/No:

v. (2 points) Does this identity hold?

$$\gamma_{G,\operatorname{sum}(K)\to L}(V) = \gamma_{G,\operatorname{sum}(M)\to L}(\gamma_{G,H,\operatorname{sum}(K)\to M}(V))$$

V. _____

Yes/No:

i. _____

2 Semistructured Data and JSon

- 2. (17 points)
 - (a) Answer the multi-choice questions below:
 - i. (3 points) What do we mean when we say that the data is *not in First Normal Form*? Check all that apply.
 - 1. There exists a non-trivial functional dependency $X \to Y$ where X is not a superkey.
 - 2. The data is represented in a human-readable form, like JSon.
 - 3. The value of an attribute of a table is a collection, such as a table or an array.
 - 4. The table has no clustered index.

i. _____

Select from 1,2,3, and/or 4:

- ii. (3 points) Check which of the following statements are true about the semistructured data model:
 - 1. JSon is semistructured data.
 - 2. In semistructured data the value of an attribute can be another collection.
 - 3. Semistructured data means that the data is compressed.
 - 4. There are no query languages for semistructured data.

ii. _____

Select from 1,2,3, and/or 4:

(b) Consider the following database, given in JSon:

```
{"Course":
1
      [ {"title": "Math101",
\mathbf{2}
         "room": "F777",
3
         "instructor": {"Name": "Bob", "Office": "E999"};
4
         "enrollment": [{"name": "David", "year": 2},
5
                          {"name": "Erol", "year": 3}
6
                         ]
7
        },
8
        {"title": "Phys202",
9
         "room": "H909",
10
         "instructor": {"name": "Alice", "office": "C222"};
11
         "enrollment": [{"name": "Carol", "year": 2},
12
                          {"name": "Erol", "year": 3},
13
                          {"name": "Fred", "year": 1}
14
                         ]
15
        },
16
        {"title": "CSE703",
17
         "room": "G080",
18
         "instructor": {"name": "Bob", "office": "E999"}
19
        }
20
      ]
21
    }
22
```

Your task is to convert this data into a relational database.

i. (6 points) Design a schema for the a relational database capable of storing the database above. Turn in relation names, their attributes, and underline the key. For example, you may write:

Room(<u>roomnumber</u>, instructorname, floorid), Floor(floorid, level)

(not a real answer). Write your answer below:

	JSon file repeated from previous page:
1	{"Course":
2	[{"title": "Math101",
3	"room": "F777",
4	"instructor": {"Name": "Bob", "Office": "E999"};
5	<pre>"enrollment": [{"name": "David", "year": 2},</pre>
6	{"name": "Erol", "year": 3}
7]
8	},
9	{"title": "Phys202",
10	"room": "H909",
11	"instructor": {"name": "Alice", "office": "C222"};
12	<pre>"enrollment": [{"name": "Carol", "year": 2},</pre>
13	{"name": "Erol", "year": 3},
14	{"name": "Fred", "year": 1}
15]
16	},
17	{"title": "CSE703",
18	"room": "G080",
19	"instructor": {"name": "Bob", "office": "E999"}
20	}
21]]
22	}

ii. (5 points) Show the content of your tables, representing the same data as the JSon file. For example, you may write:

Room	roomnumber	instructorname	floorid	Fleen	floorid	laval
	E999	Carol	flr0521	Floor	filoofid fr0521	1ever
	C222	Alice	flr0521		1110521	2

(not a real answer) Turn in several table instances:

3 E/R Diagrams, Constraints, Conceptual Design

3. (36 points)

A company selling products both online and in their own stores has a database having a schema described by the E/R diagram below:



- Sales represents individual sales. One sale contains several products bought by one single customer. The discount is a real number representing a percent, for example discount = 33.33 means a discount of 1/3 from the original product price.
- inStoreSales consists of the sales in brick-and-mortar stores; each such sale contains the store name.
- onlineSales consists of online sales; each such sale includes the name of the browser and the operating system used during the purchase.
- cid, sid, pid are int, price and discount are float, the rest are text.
- No attributes may be null, except for inStoreSales.store and onlineSales.os.



(a) (15 points) Create a database schema for the E/R diagram in the figure. You should turn an a set of CREATE TABLE statements. Your schema should include all constraints captured in the E/R diagrams.



(b) (10 points) Write a SQL query that computes, for each customer, the total amount that the customer spent online. Your query does not need to include customers who did not purchase anything oline.

i. _____

ii. _____

(c) Consider three relations:

$$R(A, B, C), S(C, \underline{D}), T(D, A)$$

AB is a key in R, and D is a key in S. For each of the queries below, show the key of the query's answer, and compute D+:

i. (2 points) Query Q1: select R.A, R.B, R.C, S.D from R, S where R.C = S.C and R.A = 20;

Key=? D + =?

ii. (2 points) Query Q2: select T.A, S.C, S.D from S, T where S.D = T.D;

Key=? D + =?

iii. (2 points) Query Q3: select R.A, R.B, R.C, S.D from R, S, T where R.A=T.A and R.C = S.C and S.D = T.D;

iii. _____

Key=? D + =?

(d) (5 points) Consider a relation $R(A_1, A_2, ..., A_n)$ satisfying the following functional dependencies:

$$A_1 \to A_2$$

$$A_2 \to A_3$$

$$A_3 \to A_4$$

$$\dots$$

$$A_{n-1} \to A_n$$

Decompose this relation into BCNF. You need to indicate only your answer by showing the relation names, their attributes, and their key, for example you may write:

$$R_1(\underline{A_2}, A_3, A_4, \dots, A_n), R_2(A_1, \underline{A_3}, A_4, \dots, A_n), R_2(A_1, A_2, \underline{A_4}, \dots, A_n), \dots, R_n(\underline{A_1}, A_2, \dots, A_{n-1})$$

(not the real answer).

4 Transactions

- 4. (54 points)
 - (a) Consider a concurrency control manager that uses strict two phase locking that schedules three transactions:
 - $T1: R_1(A), R_1(B), W_1(A), W_1(B), Co_1$
 - $T2: R_2(B), W_2(B), R_2(C), W_2(C), Co_2$
 - $T3: R_3(C), W_3(C), R_3(A), W_3(A), Co_3$

Each transaction begins with its first read operation, and commits with the Co statement. Assume the following questions for each of the schedules below:

- Is the schedule conflict-serializable? If yes, indicate a serialization order.
- Is this schedule possible under a strict 2PL protocol?
- If strict 2PL does not allow this schedule beause it denies a read or a write request, is the system in a deadlock at the time when the request is denied?

i. Schedule 1:

 $R_2(B), W_2(B), R_3(C), W_3(C), R_3(A), W_3(A), Co_3, R_2(C), W_2(C), Co_2, R_1(A), R_1(B), W_1(A), W_1(B), Co_1, W_2(C), W_$

 $\alpha)$ (3 points) Is this schedule conflict-serializable? If yes, indicate a serial zation order.

 β) (2 points) Is it possible under strict 2PL?

ii. Schedule 2:

 $R_{2}(B), W_{2}(B), R_{3}(C), W_{3}(C), R_{1}(A), R_{1}(B), W_{1}(A), W_{1}(B), Co_{1}, R_{2}(C), W_{2}(C), Co_{2}, R_{3}(A), W_{3}(A), Co_{3}(A), W_{3}(A), W_{3}(A),$

 $\alpha)$ (3 points) Is this schedule conflict-serializable? If yes, indicate a serial zation order.

 β) (2 points) Is it possible under strict 2PL?

iii. Schedule 3:

 $R_1(A), R_1(B), R_2(B), W_2(B), R_2(C), W_2(C), Co_2, R_3(C), W_3(C), R_3(A), W_3(A), Co_3, W_1(A), W_1(B), Co_1, W_2(C), W_2(C), W_2(C), W_3(C), W_3(C), W_3(A), W_3(A), W_1(A), W_1(B), Co_1, W_2(C), W_2(C), W_2(C), W_2(C), W_3(C), W_3(C), W_3(A), W_3(A), W_3(A), W_1(A), W_1(B), Co_1, W_2(C), W_2(C)$

 $\alpha)$ (3 points) Is this schedule conflict-serializable? If yes, indicate a serial zation order.

 β) (2 points) Is it possible under strict 2PL?

iv. Schedule 4:

 $R_1(A), R_1(B), W_1(A), R_3(C), W_3(C), R_3(A), W_3(A), Co_3, W_1(B), R_2(B), W_2(B), Co_1, R_2(C), W_2(C), Co_2, W_3(C), W_3(C), W_3(A), W_$

 $\alpha)~(3~{\rm points})$ Is this schedule conflict-serializable? If yes, indicate a serial zation order.

 β) (2 points) Is it possible under strict 2PL?

- (b) (10 points) Consider the following three transactions:
 - $T1: R_1(A), W_1(B), Co_1$
 - $T2: R_2(B), W_2(C), Co_2$
 - $T3: R_3(C), W_3(D), Co_3$

Given an example of a conflict-serializable schedule that has the following properties: transaction T1 commits before transaction T3 starts, and the equivalent serial order is T3, T2, T1.

- (c) A read-only transaction is a transaction that only reads from the database, without writing/inserting deleting. Answer the questions below.
 - i. (2 points) If all transactions are read-only, then every schedule is serializable.

True or False?

ii. (2 points) If no transaction reads the same element twice, then the serialization level READ COMMITTED is equivalent to REPEATABLE READS.

ii. _____

iii.

i. _____

True or False?

iii. (2 points) If no transaction inserts or deletes records to/from the database, then the serialization level REPEATABLE READS is equivalent to SERIAL-IZABLE.

True or False?

iv. (2 points) The reason why some applications use serialization levels other than SERIALIZABLE is because they would not be correct under the SERIALIZ-ABLE isolation level.

True or False?

v. (2 points) In Sqlite phantoms are not possible.

True or False?

vi. (2 points) The difference between Two Phase Locking and Strict Two Phase Locking is that the latter avoids deadlocks, while the former may allow deadlocks.

True or False?

vii. (2 points) Only one transaction can hold a *shared lock* at any time.

True or False?

viii. (2 points) Only one transaction can hold an *exclusive lock* at any time.

True or False?

iv. _____

V. _____

vi. _____

vii. _____

viii. _____

$\mathrm{CSE}~344$

5 Parallel Data Processing

5. (46 points)

(a) Consider a social network database with two relations shown below:

User(<u>uid</u> ,name)	1	Million	tuples
Follows(uid1,uid2)	10	Million	tuples

The table User contains user information, while Follows tells us that user1 follows user2. Suppose we are computing the following query:

```
select x.uid1,x.uid2,y.name
from Follows x, User y
where x.uid2 = y.uid
```

We use a distributed system with p servers, and compute the join using partitioned hash-join. In other words:

- The system partitions User(<u>uid</u>, name) by applying a hash function to uid,
- partitions Follows by applying a hash function to uid2,
- then each server computes a join of its local data.

On p = 10 servers, the query runs in 1000 seconds. Estimate the runtime of the system in each of the cases below, assuming the number of servers is increased as shown. Your numbers are only *estimates*: try to estimate within a factor of 2. For example, if the question were *what is the runtime on one server?* then you would answer 10,000 seconds, since one server must do the work of all the 10 servers that took 1000 seconds, although one server could run much faster than 10×1000 second.

i. (5 points) Assume that every user follows at most 5 users, and is followed by at most 5 users:

p =	10	100	1000	100000
Time=	1000s			

ii. (5 points) As in item i, every user follows and is followed by at most 5 users, except for user 'JB' who is followed by 10,000 users.

p =	10	100	1000	100000
Time=	1000s			

iii. (5 points) As in item i, every user follows and is followed by at most 5 users, except for user 'JB' who is followed by 100,000 users.

p =	10	100	1000	100000
Time=	1000s			

i. _____

- (b) We are running a MapReduce job over HDFS. Our input file has 10^{10} records, its size is $1TB = 10^{12}B$. Hadoop's block size is configured at 100KB. Answer each of the questions below.
 - i. (2 points) How many map tasks will the MapReduce system create by default? If there is no default, then then indicate so.

Number of map tasks:

ii. (2 points) How many reduce tasks will the MapReduce system create by default? If there is no default, then then indicate so.

Number of reduce tasks:

ii. _____

For the next few questions, recall the steps of a MapReduce job from the lecture notes:



iii. (2 points) The *Copy* phase of the reduce tasks may start immediately after the first map tasks finish without having to wait for all map tasks to finish.

iii. _____

Yes or no?

iv. (2 points) The *Sort* phase of the reduce tasks may start immediately after the first map tasks finish without having to wait for all map tasks to finish.

iv. _____

v. (2 points) The *Reduce* phase of the reduce tasks may start immediately after the first map tasks finish without having to wait for all map tasks to finish.

V. _____

Yes or no?

Yes or no?

- (c) (5 points) A MapReduce job runs on 100 workers and has 500 reduce tasks. At some point in time, all map tasks have finished, and 150 reduce tasks have finished too: the system is executing 100 reduce tasks, while another 250 reduce tasks are still waiting to be scheduled. At this point worker number 44 fails: the worker and its local disk are lost and not recoverable. Indicate which of the following will happen:
 - 1. The system continues executing the 99 active tasks, then will schedule the remaining 250 tasks on the 99 remaining workers.
 - 2. The system continues executing the 99 active tasks, then will schedule the remaining 251 tasks on the 99 remaining workers (including the reduce task that was running on worker 44).
 - 3. The system reruns all 500 reduce tasks.
 - 4. The system continues executing the 99 active tasks, then reruns the map tasks that had been ran on worker 44, and after that continues executing the remaining reduce tasks.
 - 5. The system needs to restart the entire job (all map tasks and all reduce tasks) on 99 workers.

(c) _____

Your answer:

Page 26

1.	execution of several separate MapReduce jobs.	raine) requires the
		i
	True or False?	
ii.	(2 points) A program that involves iteration (such as page execution of several separate Spark program.	e rank) requires the
		ii
	True or False?	
iii.	(2 points) In a MapReduce program, all intermediate result	s are stored on disk.
		iii
	True or False?	
iv.	(2 points) In a Spark program, all intermediate results are	stored on disk.
		iv
	True or False?	
v.	(2 points) If a worker fails during the execution of a MapRe the entire program needs to be restarted.	educe program, then
		V
	True or False?	
vi.	(2 points) If a worker fails during the execution of a Spark entire program needs to be restarted.	x program, then the
		vi
	True or False?	
vii.	(2 points) MapReduce is ideally suited for OLTP application	lons.
		vii
	True or False?	
viii.	(2 points) Spark is ideally suited for OLTP applications.	
		viii
	True or False?	

(d) The following questions compare MapReduce to Spark. For each statement indicate whether it is true or false.

Final Examination

i. (2 points) A program that involves iteration (such as page rank) requires the