Introduction to Data Management CSE 344

Lecture 7: Nested Queries in SQL

Announcements

- Homework 2: due Thursday night, 11 pm
- Sections tomorrow:
 - SQL Azure remote database access for HW3
 - Nested queries
- Thursday office hours in 006 lab this week only
- Next webquiz and HW3 posted shortly

Homework 3

- More IMDB queries using SQL Azure
- Login details in the assignment (out later today)
 - Userid = your UW netid; password on board
 - Change password on first login
- Demo in sections tomorrow
- Software (tested on these)
 - Web browsers: Windows IE, Chrome, Firefox; OS X
 Firefix and Safari. Need sliverlite plugin (already on lab windows machines)
 - SQL Management Studio 2014 (windows app)

Lecture Goals

 Today we will learn how to write more powerful SQL queries

They are needed in Homework 3

• Reading: Ch. 6.3

Subqueries

- A subquery is a SQL query nested inside a larger query
- Such inner-outer queries are called nested queries
- A subquery may occur in:
 - A SELECT clause
 - A FROM clause
 - A WHERE clause
- Rule of thumb: avoid writing nested queries when possible; keep in mind that sometimes it's impossible

Subqueries...

- Can return a single constant and this constant can be compared with another value in a WHERE clause
- Can return relations that can be used in various ways in WHERE clauses
- Can appear in FROM clauses, followed by a tuple variable that represents the tuples in the result of the subquery
- Can appear as computed values in a SELECT clause

1. Subqueries in SELECT

Product (<u>pname</u>, price, cid) Company(<u>cid</u>, cname, city)

For each product return the city where it is manufactured

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
FROM Product X
```

"correlated subquery"

What happens if the subquery returns more than one city?
We get a runtime error
(SQLite simply ignores the extra values)

1. Subqueries in SELECT

Whenever possible, don't use a nested queries:

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
```

FROM Product X

SELECT X.pname, Y.city FROM Product X, Company Y WHERE X.cid=Y.cid We have "unnested" the query

1. Subqueries in SELECT

Compute the number of products made by each company

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
```

FROM Company C

Better: we can unnest by using a GROUP BY

SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname

1. Subqueries in SELECT

But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
```

FROM Company C

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

No! Different results if a company has no products

```
SELECT C.cname, count(pname)
FROM Company C LEFT OUTER JOIN Product P
ON C.cid=P.cid
GROUP BY C.cname
```

2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

```
SELECT X.pname
FROM (SELECT * FROM Product AS Y WHERE price > 20) as X
WHERE X.price < 500
```

Unnest this query!

2. Subqueries in FROM

 At the end of the lecture we will see that sometimes we really need a subquery and one option will be to put it in the FROM clause (see "finding witnesses").

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using EXISTS:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE EXISTS (SELECT *
FROM Product P
WHERE C.cid = P.cid and P.price < 200)
```

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using IN

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price < 200)
```

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Using ANY:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

3. Subqueries in WHERE

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

Now let's unnest it:

```
SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid= P.cid and P.price < 200
```

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

Universal quantifiers

Universal quantifiers are hard!

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies: i.e. s.t. <u>some</u> product ≥ 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

2. Find all companies s.t. <u>all</u> their products have price < 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid NOT IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using EXISTS:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE NOT EXISTS (SELECT *
FROM Product P
WHERE P.cid = C.cid and P.price >= 200)
```

3. Subqueries in WHERE

Find all companies s.t. <u>all</u> their products have price < 200

Universal quantifiers

Using ALL:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 >= ALL (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

Question for Database Fans and their Friends

Can we unnest the universal quantifier query?

Monotone Queries

- Definition A query Q is monotone if:
 - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of of the tuples

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c003
Camera	149.99	c001

Company

cid	cname	city
c001	Sunworks	Bonn
c002	DB Inc.	Lyon
c003	Builder	Lodtz



Is the mystery

query monotone?

А	В
149.99	Lodtz
19.99	Lyon

Product

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c003
Camera	149.99	c001
iPad	499.99	c001

Company

cid	cname	city
c001	Sunworks	Bonn
c002	DB Inc.	Lyon
c003	Builder	Lodtz



А	В
149.99	Lodtz
19.99	Lyon
19.99	Bonn
149.99	Bonn

Monotone Queries

- Theorem: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.
- Proof. We use the nested loop semantics: if we insert a tuple in a relation R_i, this will not remove any tuples from the answer

```
SELECT a_1, a_2, ..., a_k
FROM R_1 AS x_1, R_2 AS x_2, ..., R_n AS x_n
WHERE Conditions
```

```
\begin{array}{c} \text{for } x_1 \text{ in } R_1 \text{ do} \\ \text{ for } x_2 \text{ in } R_2 \text{ do} \\ & \dots \\ \text{ for } x_n \text{ in } R_n \text{ do} \\ \text{ if Conditions} \\ \text{ output } (a_1, \dots, a_k) \end{array}
```

Monotone Queries

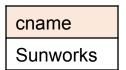
The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

pname	price	cid
Gizmo	19.99	c001

cid	cname	city
c001	Sunworks	Bonn

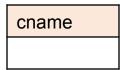




pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c001

cid	cname	city
c001	Sunworks	Bonn





• <u>Consequence</u>: we cannot write it as a SELECT-FROM-WHERE query without nested subqueries

Queries that must be nested

- Queries with universal quantifiers or with negation
- Queries that use aggregates in funny ways
 - Note: sum(..) and count(*) are NOT monotone,
 because they do not satisfy set containment
 - select count(*) from R is not monotone!