Introduction to Data Management CSE 344

Lectures 10: System's Architecture and Relational Algebra

Announcements

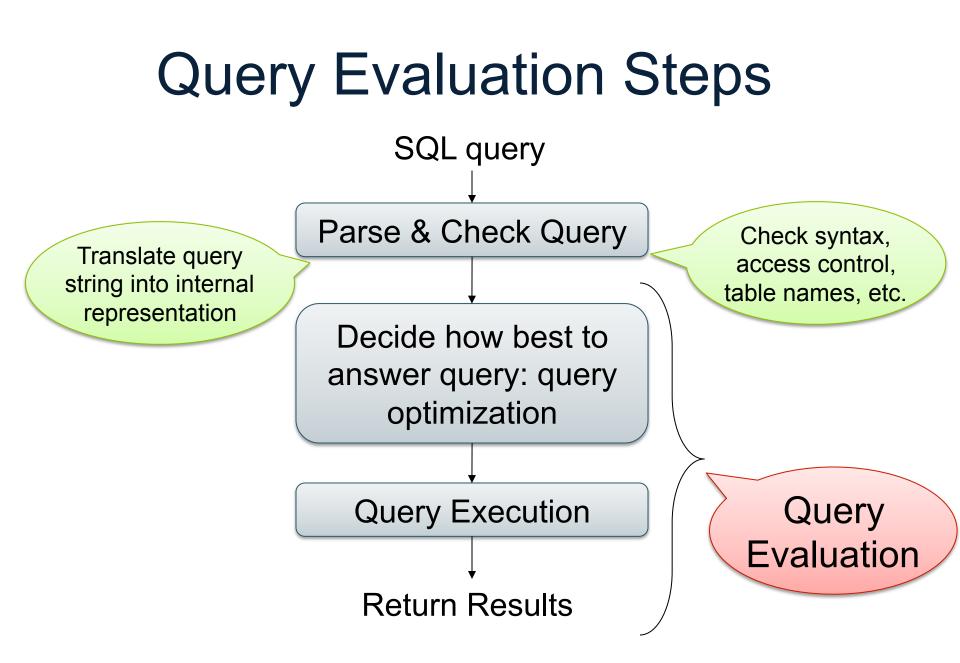
• Webquiz 3 due tonight!

• Today's lecture: 2.4 and 5.1

Where We Are

- Motivation for using a DBMS for managing data
- SQL, SQL, SQL
 - Declaring the schema for our data (CREATE TABLE)
 - Inserting data one row at a time or in bulk (INSERT/.import)
 - Modifying the schema and updating the data (ALTER/UPDATE)
 - Querying the data (SELECT)
 - Tuning queries (CREATE INDEX)
- Next step: More knowledge of how DBMSs work
 - Client-server architecture
 - Relational algebra and query execution

CSE 344 - Fall 2013



The WHAT and the HOW

- SQL = WHAT we want to get form the data
- Relational Algebra = HOW to get the data we want
- The passage from WHAT to HOW is called query optimization

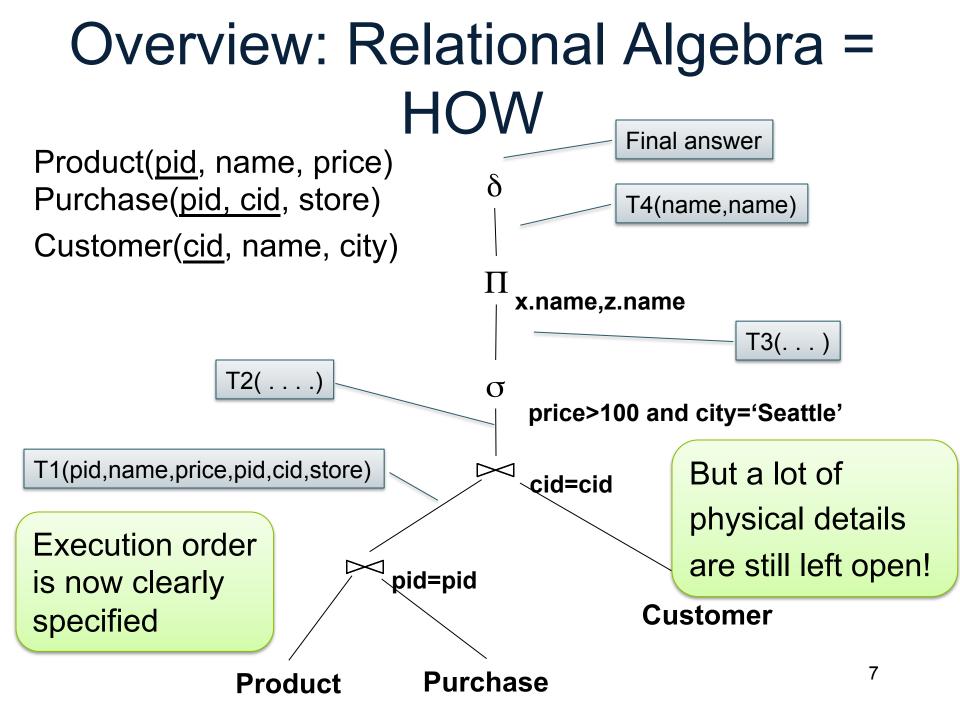
Overview: SQL = WHAT

Product(<u>pid</u>, name, price) Purchase(<u>pid, cid</u>, store)

Customer(<u>cid</u>, name, city)

SELECT DISTINCT x.name, z.name FROM Product x, Purchase y, Customer z WHERE x.pid = y.pid and y.cid = y.cid and x.price > 100 and z.city = 'Seattle'

It's clear WHAT we want, unclear HOW to get it



Relational Algebra

Relational Algebra Operators

- Union U, intersection ∩, difference -
- Selection o
- Projection
- Cartesian product ×, join ⋈
- Rename p
- Duplicate elimination δ
- Grouping and aggregation γ
- Sorting τ

Extended RA

RA

From Logical Plans to Physical Plans

Supplier(<u>sid</u>, sname, scity, sstate) Supply(<u>sid, pno</u>, quantity)

Example

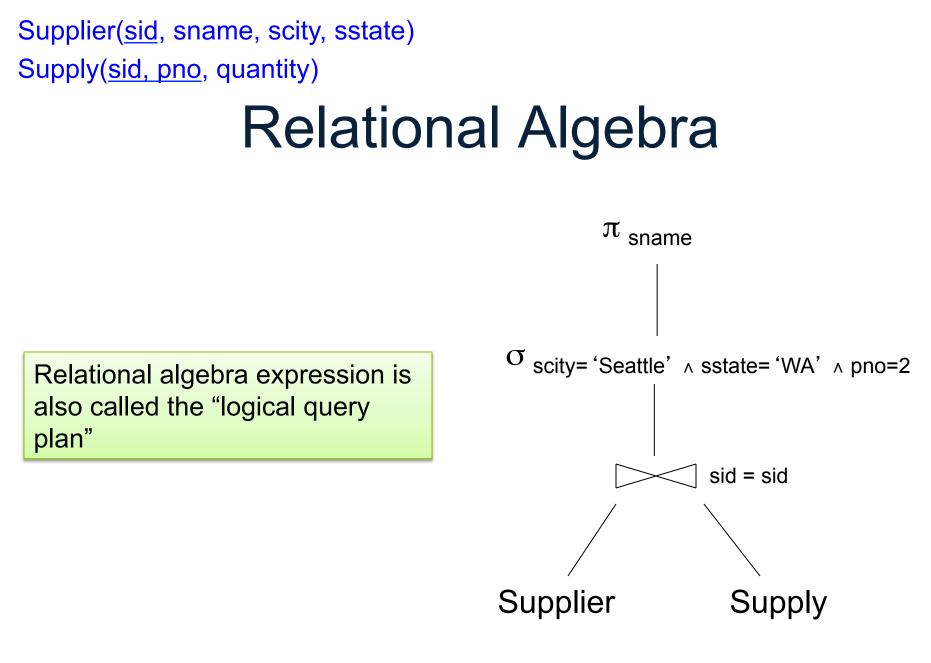
SELECT sname FROM Supplier x, Supply y WHERE x.sid = y.sid and y.pno = 2 and x.scity = 'Seattle' and x.sstate = 'WA'

Give a relational algebra expression for this query

Supplier(<u>sid</u>, sname, scity, sstate) Supply(<u>sid</u>, pno, quantity)

Relational Algebra

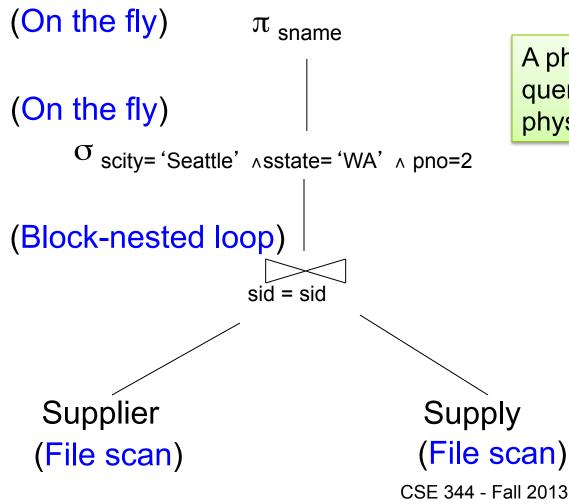
 $\pi_{\text{sname}}(\sigma_{\text{scity='Seattle'} \land \text{sstate='WA'} \land \text{pno=2}} (\text{Supplier} \Join))$



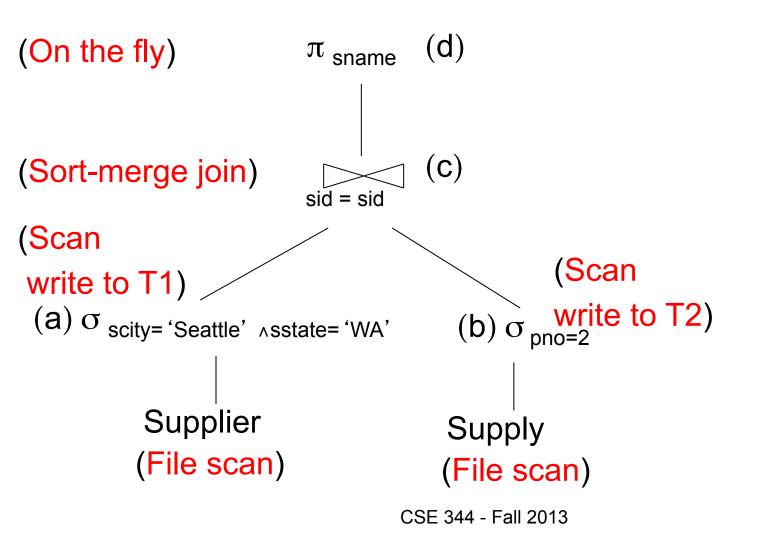
Supplier(sid, sname, scity, sstate)

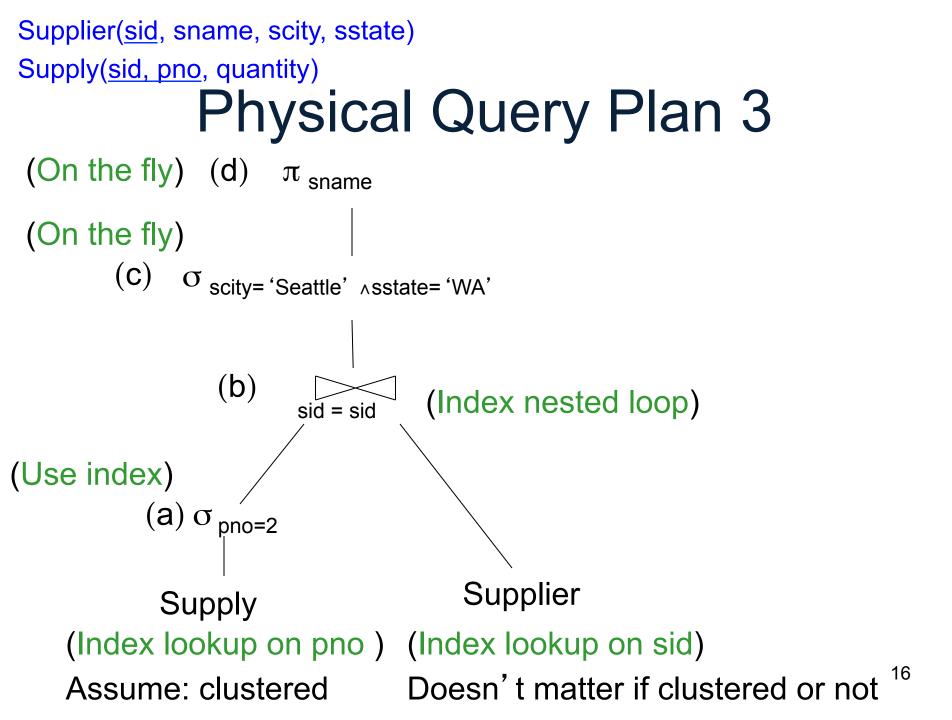
Supply(sid, pno, quantity)

Physical Query Plan 1



A physical query plan is a logical query plan annotated with physical implementation details Supplier(<u>sid</u>, sname, scity, sstate) Supply(<u>sid, pno</u>, quantity) Physical Query Plan 2





Physical Data Independence

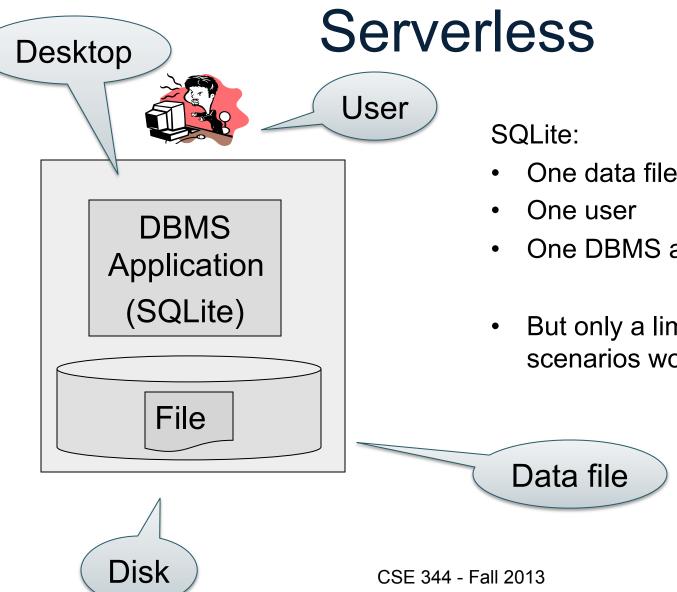
- Means that applications are insulated from changes in physical storage details
 - E.g., can add/remove indexes without changing apps
 - Can do other physical tunings for performance
- SQL and relational algebra facilitate physical data independence because both languages are "set-at-a-time": Relations as input and output

Architectures

1. Serverless

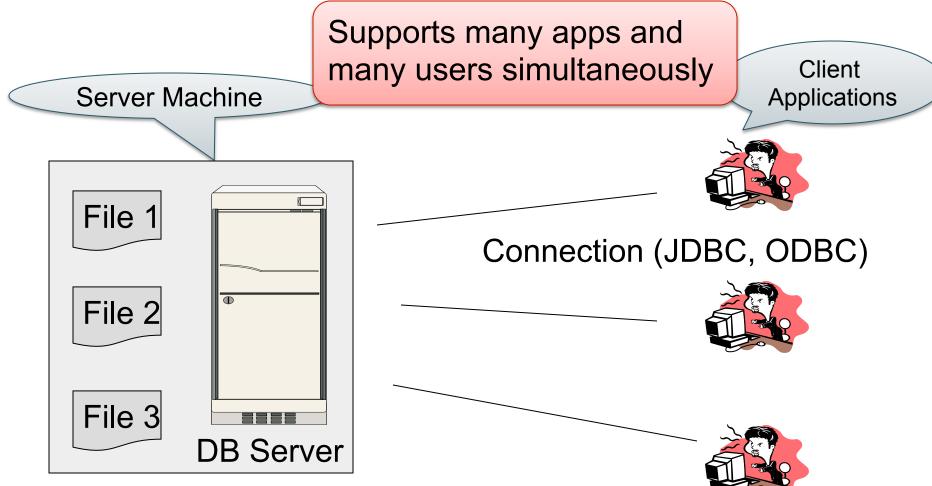
2. Two tier: client/server

3. Three tier: client/app-server/db-server



- One data file
- **One DBMS application**
- But only a limited number of scenarios work with such model

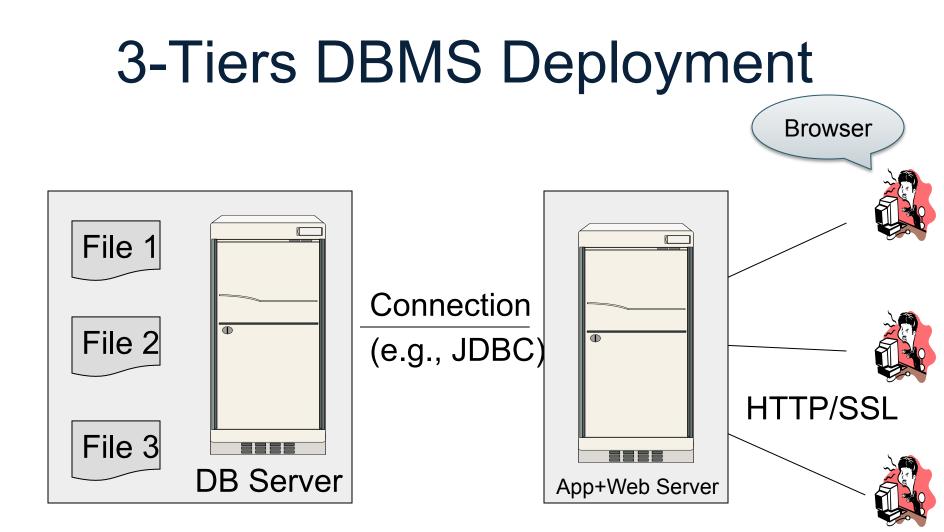
Client-Server



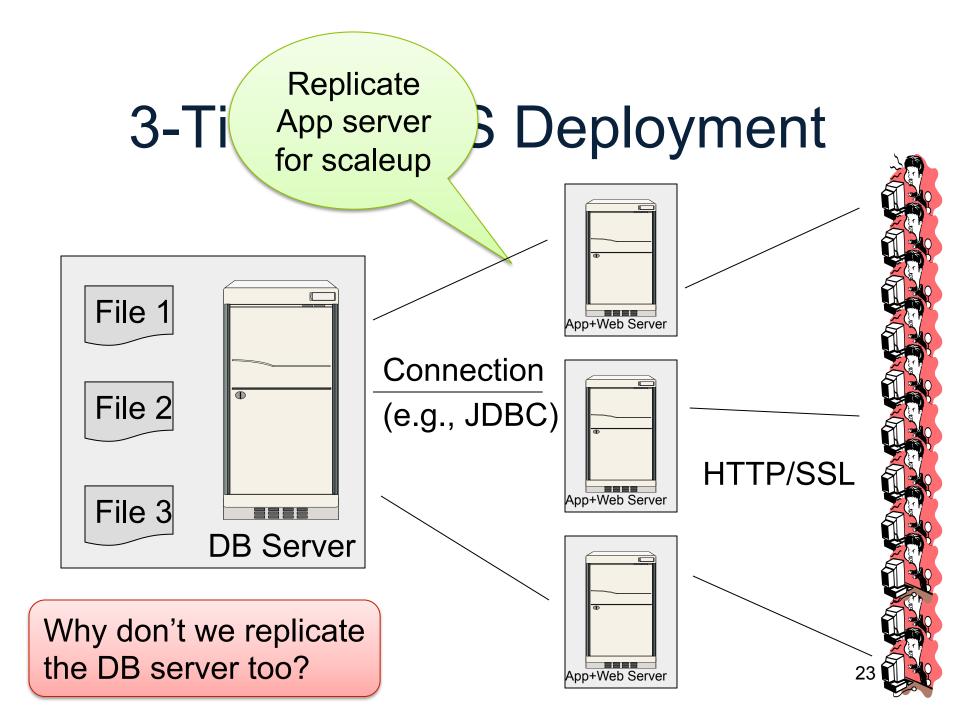
- One server running the database
- Many clients, connecting via the ODBC or JDBC (Java Database Connectivity) protocol

Client-Server

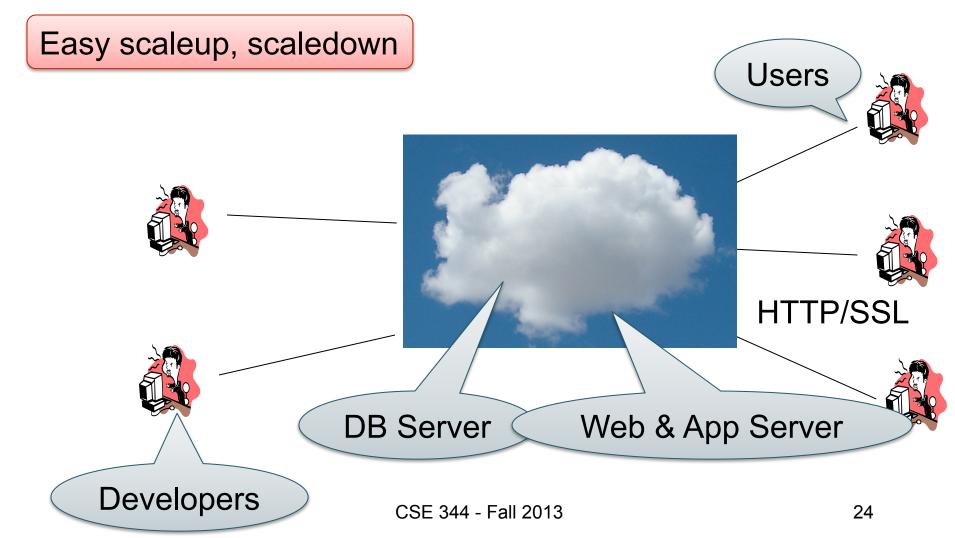
- One *server* that runs the DBMS (or RDBMS):
 - Your own desktop, or
 - Some beefy system, or
 - A cloud service (SQL Azure)
- Many *clients* run apps and connect to DBMS
 - Microsoft's Management Studio (for SQL Server), or
 - psql (for postgres)
 - Some Java program (HW5) or some C++ program
- Clients "talk" to server using JDBC/ODBC protocol



Web-based applications



DBMS Deployment: Cloud



Using a DBMS Server

- 1. Client application establishes connection to server
- 2. Client must authenticate self
- 3. Client submits SQL commands to server
- 4. Server executes commands and returns results

