CSE 341 | Section 3

Q1: datatypes and pattern-matching

Colors are often represented based on red, green, and blue values (RGB), where the value for each of these three components is an integer between 0 and 255 (inclusive). The RGB value for Red is (255, 0, 0), the RGB value for Green is (0, 255, 0) and the RGB value for Blue is (0, 0, 255). We can represent other colors with a mix of RGB values. Suppose we have the following datatype defined to represent a color conveniently as one of the three primary colors, or a custom color:

datatype color = Red | Green | Blue | RGB of (int * int * int);

Write a function invert_color that takes a color argument and returns its inverse as a new RGB representation (as a color datatype), resulting from subtracting each original r, g, b value from 255. For example, the call invert_color (RED) should return the RGB value (0, 255, 255) and invert_color (RGB(0, 25, 155)) should return the RGB value (255, 220, 100). Assume that if the color argument was constructed with the RGB constructor, the three integers are between 0 and 255 (inclusive). *Hint: use a case expression in your solution.*

Q2:

The color datatype is an example of a "one-of" datatype as discussed in lecture/readings. One alternative is to represent colors with a record:

{ color_name : string, r : int, g : int, b : int }

a. What is one advantage of using a datatype for our color representation?

b. What is an example where it would be more appropriate to represent a type as a record instead of a "one-of" datatype?

Q3: More datatype and Pattern-matching examples

a. Consider the following type and datatype:

Write a function area which takes a shape as an argument and returns its area (as a real value).

b. Now recall this datatype to represent expression trees from lecture:

Write a function const_not_under_add of type exp -> bool that returns true if and only if there
exists a Constant in the expression that is not a child of an Add expression. For example,
const_not_under_add(Constant 341) should return true, as should
const_not_under_add(Multiply(Constant 341, Add(Constant 0, Constant 1))).

Q4: Consider the following code:

```
fun length l =
    case l of
    _::xs => 1 + length xs
    | [] => 0
```

Is it tail-recursive? Why/why not?

```
fun all_positive (accum, l) =
   case l of
```

```
x::xs => all_positive (accum andalso x > 0, xs)
| [] => accum
```

Is it tail-recursive? Why/why not?

Datatypes and Pattern-Matching (17au Midterm Question)

1. (20 points) This problem uses this datatype binding, where an exp is a simple arithmetic expression like we studied in class except instead of negations and multiplications, we have doubling and (integer) division.

datatype exp = Constant of int
 | Double of exp
 | Add of exp * exp
 | Divide of exp * exp

- (a) Write a function eval_exp of type exp -> int that returns the "answer" for "executing" the arithmetic expression. Some notes on division:
 - Use integer division, which in ML is done with the infix operator div. For example, in ML, 6 div 4 is 1.
 - Division by zero will raise an exception, which is fine.
- (b) Give an example of a value of type exp where:
 - Calling eval_exp with your expression causes a division-by-zero exception, but ...
 - ... no use of the Divide constructor has Constant 0 as its second argument.
- (c) Write a function no_literal_zero_divide of type exp -> bool that returns true if and only if no use of the Divide constructor has Constant 0 as its second argument. Notes:
 - So, no_literal_zero_divide applied to your answer to the previous question would evaluate to true.
 - You should *not* use eval_exp this question has nothing to do with evaluating expressions.