



# CSE 341

## Section 2

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# *Today's Agenda*

- Testing
- Lists, Let-Expression (Review)
- Options
- Type synonyms
- Type generality
- Equality types
- Syntactic sugar

**Reminder:** Check out the [CSE341 style guide](#) as you work on HW!

Also check out the style guides in [section 1 slide](#)!

# Testing

- You should still test your code!
- We will assign points to your testing file
- Just do something like this:

```
val test1 = ((4 div 4) = 1);
```

*"Is expected output = actual output"*

# *Section Learning Objectives*

- Review building/accessing new types (e.g. **datatypes**)
- Recognize **type synonyms** as “convenient” feature
- Be able to generalize specific types with **polymorphism** (e.g. int list into ‘a list) and **equality** types
- Practice using pattern-matching with **case expressions**

# *Lists*

- Lots of new types: For any type  $t$ , the type  $t$  list describes lists where all elements have type  $t$ 
  - Examples: `int list`, `bool list`, `int list list`, `(int * int) list`, `(int list * int) list`
- So `[]` can have type  $t$  list for any type  $t$ 
  - SML uses type `'a list` to indicate this (“tick a” or “alpha”)
- For  $e1 :: e2$  to type-check, we need a  $t$  such that  $e1$  has type  $t$  and  $e2$  has type  $t$  list. Then the result type is  $t$  list
  - `null: 'a list -> bool`
  - `hd: 'a list -> 'a`
  - `tl: 'a list -> 'a list`

# *Let-Expression*

- Syntax: `let b1 b2 ... bn in e end`
  - Each **bi** is any binding and **e** is any expression
- Type-checking: Type-check each **bi** and **e** in a static environment that includes the previous bindings.
- Type of whole let-expression is the type of **e**.
- Evaluation: Evaluate each **bi** and **e** in a dynamic environment that includes the previous bindings.

Result of whole let-expression is result of evaluating **e**.

# Options

**t** option is a type for any type **t**

- (much like `t list`, but a different type, not a list)

Building:

- **NONE** has type `'a option` (much like `[]` has type `'a list`)
- **SOME e** has type `t option` if `e` has type `t` (much like `e::[]`)

Accessing:

- **isSome** has type `'a option -> bool`
- **valOf** has type `'a option -> 'a` (exception if given `NONE`)

# Type Synonyms

- What does `int * int * int` represent?
- In HW1 we called it a date
- Wouldn't it be nice to reflect this representation in the source code itself?

```
type date = int * int * int
```



# *Datatypes*

- What if we want something **unique**? A **new** type?
- We can't just use type synonyms because they can only be built from existing types.
- **Datatypes** give us the ability to define **custom types**.

```
datatype foo = bar | baz of int | qux of bool
```

# type VS datatype

- **datatype** introduces a new type name, distinct from all existing types

```
datatype suit = Club | Diamond | Heart | Spade
datatype rank = Jack | Queen | King | Ace
              | Num of int
```

- **type** is just another name

```
type card = suit * rank
```

# *Type Synonyms*

## **Why?**

- For now, just for convenience
- It doesn't let us do anything new

Later in the course we will see another use related to modularity.

# *Type Generality*

Write a function that appends two string lists...

# Type Generality

- We would expect

```
string list * string list -> string list
```

- But the type checker found

```
`a list * `a list -> `a list
```

- `a are called Polymorphic Types
- Why is this OK?

# More General Types

- The type

```
'a list * 'a list -> 'a list
```

is more general than the type

```
string list * string list -> string list
```

and “can be used” as any less general type, such as

```
int list * int list -> int list
```

- But it is not more general than the type

```
int list * string list -> int list
```

# *The Type Generality Rule*

The “more general” rule

A type  $t1$  is more general than the type  $t2$  if you can take  $t1$ , replace its type variables **consistently**, and get  $t2$

What does **consistently** mean?

# *Equality Types*

Write a list “contains” function...



# *Equality Types*

- The double quoted variable arises from use of the = operator
  - We can use = on most types like **int**, **bool**, **string**, tuples (that contain only “equality types”)
  - Functions and **real** are not “equality types”
- Generality rules work the same, except substitution must be some type which can be compared with =
- You can ignore warnings about “calling polyEqual”

# *More Syntactic Sugar*

- Tuples are just records
- If-then-else is implemented as syntactic sugar for a case statement

# *If-then-else*

- We've just covered case statements
- How could we implement if-then-else

```
case x of  
  true => "apple"  
| false => "banana"
```

```
if x then "apple" else "banana"
```

# *val*-Pattern Matching

Remember our unit test?

```
(* Neat trick for creating hard-fail tests: *)  
  
val true = ((4 div 4) = 1);
```

Just a pattern match!

*“Match the left hand side against the value ‘template’ true, binding any variables (there aren’t any!)”*

# *Adventures in pattern matching*

- Shape example
- Function-pattern syntax if we get to it