CSE341 Spring 2019, Midterm Examination
May 3, 2019

Please do not turn the page until 12:30.

Rules:

- The exam is closed-book, closed-note, etc. except for one side of one 8.5x11in piece of paper.
- Please stop promptly at 1:20.
- There are 100 points, distributed unevenly among 6 questions (all with multiple parts):
- The exam is printed double-sided.

Advice:

- Read questions carefully. Understand a question before you start writing.
- Write down thoughts and intermediate steps so you can get partial credit. But clearly indicate what is your final answer.
- The questions are not necessarily in order of difficulty. Skip around. Make sure you get to all the questions.
- If you have questions, ask.
- Relax. You are here to learn.
1. (23 points) This problem uses this datatype binding for trees where all the data is at the leaves and each leaf can hold an int or a string.

   datatype tree = I of int | S of string | N of tree * tree

(a) Give an example of a value of type `tree` that contains exactly three ints and zero strings.

(b) Define a function `sum` of type `tree -> int` that computes the sum of all the ints contained in the tree (ignoring any strings).

(c) Define a function `stringify` of type `tree -> tree` that returns a result that is like its argument except everywhere the argument has some `I i`, the result has a `S s` where `s` is the string representation of `i`. For example, `stringify (N(I 3, S "hi")) = N(S "3", S "hi")`. Use library function `Int.toString` of type `int -> string`, which returns the string representation of its argument.

(d) Define a function `intify` of type `tree -> tree` that returns a result that is like its argument except everywhere the argument has some `S s and the s is the string representation of a number, the result has an `I i where s is the string representation of i`. For example, `intify (N(I 3, S "hi")) = N(I 3, S "hi")` and `intify (N(S "3", S "hi")) = N(I 3, S "hi")`. Use library function `Int.fromString` of type `string -> int option`, which returns an option containing the number represented by the string or `NONE` if the string does not represent a number.
2. (17 points) This problem uses the same datatype binding as problem 1 and an incorrect function that is supposed to evaluate to true if and only if its argument contains exactly two ints.

datatype tree = I of int | S of string | N of tree * tree

fun has_exactly_two_ints t = (* this is buggy! *)
  case t of
  (* 1 *) N(I _, I _) => true
  (* 2 *) | N(t1, t2) => has_exactly_two_ints t1 orelse has_exactly_two_ints t2
  (* 3 *) | S _ => false
  (* 4 *) | I _ => false

(a) Give an example tree where has_exactly_two_ints evaluates to true correctly, i.e., the tree has exactly two ints.

(b) Give an example tree where has_exactly_two_ints evaluates to true incorrectly, i.e., the tree does not have exactly two ints.

(c) Give an example tree where has_exactly_two_ints evaluates to false correctly, i.e., the tree does not have exactly two ints.

(d) Give an example tree where has_exactly_two_ints evaluates to false incorrectly, i.e., the tree has exactly two ints.

(e) For each of i–iii below, choose from the options A–D:
   A. The function is still buggy in exactly the same way.
   B. The function is still buggy but is not equivalent to the version above.
   C. The function is now correct.
   D. The function no longer type-checks.
   Do not consider that the first branch of a case expression has no | character and the others do. That is, assume that syntactic detail is fixed.
   i. The branch labeled (* 2 *) is moved to before line (* 1 *) (so the order is 2, 1, 3, 4).
   ii. The branch labeled (* 3 *) is moved to before line (* 1 *) (so the order is 3, 1, 2, 4).
   iii. The branch labeled (* 4 *) is moved to before line (* 1 *) (so the order is 4, 1, 2, 3).
3. (12 points) This problem asks you to give good error messages for why a short ML program does not type-check or has a syntax error. A specific phrase or short sentence is plenty.

For example, for the program,

fun f1 (x,y) = if x then y + 1 else x

a fine answer would be, “the then-branch-expression and the else-branch-expression do not have the same type.”

Give good error messages for each of the following:

(a) fun f1 (x,y) = 
   let
       val y = x
       val x = g 9
       fun g z = x * y
   in
       x + y + (g 7)
   end

(b) fun f2 (x,y,z) = 
   if x=y
   then 0
   else if y=z
   then 2
   else if y > 3
   then 3

(c) fun f3 xs = 
   case xs of
     | [] => 0::[]
     | x::[] => x+1
     | x::xs => (x+1)::(f3 xs)

(d) fun f4 (x,y) = 
   if x > 0.0 orelse x = y
   then y
   else 3.0
4. (18 points) Consider this ML function

```ml
fun foo f g xs =
  case xs of
    [] => []
  | x::xs' =>
    let
      val (y1,i1) = f x
      val (y2,i2) = g x
    in
      (if i1 > i2 then y1 else y2)::(foo f g xs')
    end
```

(a) What does \( \text{foo} \ (\text{fn } x \Rightarrow (0, x \cdot 2)) \ (\text{fn } x \Rightarrow (x + 1, x + 1)) \ [0, 2, 1, 3] \) evaluate to?

(b) What is the type of \( \text{foo} \)?

(c) Complete this alternate definition of \( \text{foo} \) so that it is equivalent to the function above. (Hint: you need several short lines of code.)

```ml
fun foo f g = List.map (fn x =>

```

(d) What is the type of \( \text{foo} \ (\text{fn } x \Rightarrow (x, x)) \ (\text{fn } x \Rightarrow (0 - x, 0 - x)) \)?

(e) In approximately one English sentence, what does the function produced by \( \text{foo} \ (\text{fn } x \Rightarrow (x, x)) \ (\text{fn } x \Rightarrow (0 - x, 0 - x)) \) compute?
5. (9 points) Recall `List.filter` has type `(a -> bool) -> 'a list -> 'a list`. Consider this function of type `(a -> bool) -> (a -> bool) -> 'a list -> 'a list`:

```haskell
fun filter2 f g xs =
  case xs of
    [] => []
    | x::xs => if f x andalso g x
      then x :: (filter2 f g xs)
      else filter2 f g xs
```

(a) Is the recursive call in the then-branch above a tail call?

(b) Is the recursive call in the else-branch above a tail call?

(c) Is the call to `f` in the body of `filter2` a tail call?

(d) Is the call to `g` in the body of `filter2` a tail call?

(e) Reimplement `filter2` by filling in the blank below:

```haskell
fun filter2 f g = List.filter __________________________
```
6. (21 points) This problem considers two ML modules \texttt{Direction1} and \texttt{Direction2}, and a signature \texttt{DIRECTION}. They are on the next page. Separate that page from your exam and do not turn it in.

Provided information: Under the signature \texttt{DIRECTION}, the two modules are equivalent to each other.

(a) If we deleted the \texttt{turn} function in \texttt{Direction2} and replaced it with a copy of the \texttt{turn} function in \texttt{Direction1}, would the program (A) no longer type-check, (B) still type-check and have equivalent modules, or (C) type-check but with modules not equivalent?

(b) If we deleted the \texttt{turn} function in \texttt{Direction1} and replaced it with a copy of the \texttt{turn} function in \texttt{Direction2}, would the program (A) no longer type-check, (B) still type-check and have equivalent modules, or (C) type-check but with modules not equivalent?

(c) Given signature \texttt{DIRECTION}, consider client code outside of module \texttt{Direction2} (for the rest of Problem 6, \texttt{Direction1} is irrelevant).

i. If possible, fill in the blank so that \texttt{nope} evaluates to \texttt{false}. If impossible, write “impossible.”

\begin{verbatim}
val x = ______________
val nope = Direction2.isNS x orelse Direction2.isEW x
\end{verbatim}

ii. If possible, fill in the blank so \texttt{isWest} correctly implements a function that evaluates to \texttt{true} if and only if its argument is \texttt{Direction2}'s representation of west. If impossible, write “impossible.”

\begin{verbatim}
fun isWest x = ______________
\end{verbatim}

iii. If possible, fill in the blank so \texttt{yep} evaluates to \texttt{true}. If impossible, write “impossible.”

\begin{verbatim}
val west = ______________
val yep = Direction2.isEW west
\end{verbatim}

(d) Repeat part (c) but assuming the first line of \texttt{DIRECTION} is type \texttt{t = int}:

i. If possible, fill in the blank so that \texttt{nope} evaluates to \texttt{false}. If impossible, write “impossible.”

\begin{verbatim}
val x = ______________
val nope = Direction2.isNS x orelse Direction2.isEW x
\end{verbatim}

ii. If possible, fill in the blank so \texttt{isWest} correctly implements a function that evaluates to \texttt{true} if and only if its argument is \texttt{Direction2}'s representation of west. If impossible, write “impossible.”

\begin{verbatim}
fun isWest x = ______________
\end{verbatim}

iii. If possible, fill in the blank so \texttt{yep} evaluates to \texttt{true}. If impossible, write “impossible.”

\begin{verbatim}
val west = ______________
val yep = Direction2.isEW west
\end{verbatim}

(e) Repeat part (c) with \texttt{t} again abstract but assuming the line \texttt{val north : t} is removed.

i. If possible, fill in the blank so that \texttt{nope} evaluates to \texttt{false}. If impossible, write “impossible.”

\begin{verbatim}
val x = ______________
val nope = Direction2.isNS x orelse Direction2.isEW x
\end{verbatim}

ii. If possible, fill in the blank so \texttt{isWest} correctly implements a function that evaluates to \texttt{true} if and only if its argument is \texttt{Direction2}'s representation of west. If impossible, write “impossible.”

\begin{verbatim}
fun isWest x = ______________
\end{verbatim}

iii. If possible, fill in the blank so \texttt{yep} evaluates to \texttt{true}. If impossible, write “impossible.”

\begin{verbatim}
val west = ______________
val yep = Direction2.isEW west
\end{verbatim}
These definitions are used in Problem 6. Rip this page out from the rest of the exam. Do not put answers on this page and do not turn it in.

signature DIRECTION =
  sig
    type t
    val turn : t * int -> t
    val north : t
    val isNS : t -> bool
    val isEW : t -> bool
  end

structure Direction1 :> DIRECTION =
  struct
    datatype t = North | East | South | West
    fun turnClockwise x = case x of North => East | East => South | South => West | West => North
    fun turnCounterClockwise x = case x of North => West | West => South | South => East | East => North
    fun turn (x,n) =
      if n = 0
      then x
      else if n > 0
      then turn(turnClockwise x, n-1)
      else turn(turnCounterClockwise x, n+1)
    val north = North
    fun isNS x =
      case x of
        North => true
        | South => true
        | _ => false
    fun isEW x =
      case x of
        East => true
        | West => true
        | _ => false
  end

structure Direction2 :> DIRECTION =
  struct
    type t = int (* 0 = North, 1 = East, 2 = South, 3 = West *)
    fun turnClockwise x = if x=3 then 0 else x+1
    fun turnCounterClockwise x = if x=0 then 3 else x-1
    fun turn (x,n) = (x+n) mod 4
    val north = 0
    fun isNS x = (x = 0) orelse (x=2)
    fun isEW x = (x = 1) orelse (x=3)
  end