

CSE 341:

Programming Languages

Section AC with Nate Yazdani

recap

- boolean operators
 - good style to use **andalso**, **orelse**, and **not**
 - syntactic sugar for certain uses of if-then-else

```
(* e1 andalso e2 *)  
if e1  
then e2  
else false
```

```
(* e1 orelse e2 *)  
if e1  
then true  
else e2
```

- style grading will be restrained this quarter

announcement

- sections have typically been like extra lectures
- today, we're trying short coding exercises instead of me live-coding
- hopefully more engaging and useful for you guys
- if not, we'll switch back next time

questions?

agenda

- type synonyms
- type generality
- equality types
- syntactic sugar

type synonyms

- what is the meaning of `int * int * int`?
 - literally, a triple of integers
 - conceptually, it could be a date, a co-ordinate, or some other thing
- it'd sure be nice if our code could reflect the *purpose* of a type in addition to its “literal meaning”

```
type date = int * int * int
```

type vs. datatype

- **datatype** defines a *new* type and a name for it
 - different from all existing types

```
datatype suit = Club | Diamond | Heart | Spade
datatype rank = Jack | Queen | King | Ace
              | Number of int (* 2-10 *)
```

- **type** gives a new name to an “existing” type
 - might be built out of smaller types
 - still just a name

```
type card = suit * rank
```

type synonyms: why bother?

- really really good for documentation
 - for this reason, languages without them often have popular conventions for variable names
- doesn't let us do anything we couldn't do before
- later in the course, we'll see how they help with modularity

coding exercise

please work with the people around you to write an SML function to reverse a string list (without type annotations)

type generality

- what type did SML give your function?
- probably `'a list -> 'a list`
- why not `string list -> string list`?

```
fun rev xs =  
  case xs of  
    x::xs' => rev(xs') @ [x]  
  | []     => []
```

fast list reverse

```
fun rev xs =  
  let fun aux (xs, ys) =  
        case xs of  
          x::xs' => aux(xs', x::ys)  
        | []     => ys  
      in  
        aux(xs, [])  
      end
```

type generality

- the type inferred by SML is *more general* than the one that we had in mind

```
'a list -> 'a list
```

- it works wherever any *less general* type is expected
- so just as good as these other types:

```
string list -> string list
```

```
int list -> int list
```

- but *not* this one:

```
string list -> int list
```

rule for generality

A type t_1 is **more general** than a type t_2
if you can **substitute** the type variables of t_1
consistently to get t_2 .

example of generality

The type

`'a list -> 'a list`

is more general than the type

`int list -> int list`

because you can substitute `int` for `'a`

more coding!

please work together again and write a list-contains function (without type annotations)

...and without the **List.exists** library function :-)

equality types

- a type variable with double quotes (*e.g.*, `''a`) can only be substituted with an *equality type*
- an equality type is a type that supports the `=` operator, such as `int`, `bool`, or `string`
- function types and `real` are *not* equality types
- **you can completely ignore warnings about “calling polyEqual”**

syntactic sugar

- under the hood, the if-then-else syntax form is actually translated into a case statement

```
(* if e1 then e2 else e3 *)  
case e1 of  
  true => e2  
| false => e3
```

- so the **andalso** and **orelse** operators are syntactic sugar for if-then-else, which is syntactic sugar for a case statement!
- SML is pretty “sweet” like that!

...yeah okay that was pretty bad

before you go...

some quick feedback

- did the exercises help at all?
- given what you've learned so far, were they...
 - too small?
 - not enough?
 - annoying and/or confusing?
 - distracting from the main point?
- any other suggestions for how to make section work better for you?

thanks!