CSE 341: Programming Languages

Winter 2006

Lecture 4— "One-of" types; user-defined types

Where are we

- Done features: functions, tuples, lists, local bindings
- Done concepts: syntax vs. semantics, environments, mutation-free
- Today features: record types, datatypes, type synonyms, pattern-matching
- Today concepts: "One-of" types, constructors/destructors, case-coverage

Base types and compound types

Languages typically provide a small number of "built-in" types and ways to build compound types out of simpler ones:

- Base types examples: int, bool
- Type builder examples: tuples, lists, records

Base types *clutter* a language definition; better to make them *libraries* when possible.

 ML does this to a remarkable extent (e.g., we will soon define away bool and conditionals)

Good to let programmers bind types to type names, just like we bind values to variables.

Compound-type flavors

Conceptually, just a few ways to build compound types:

- 1. "Each-of": A t contains a t1 and a t2
- 2. "One-of": A t contains a t1 or a t2
- 3. "Self-reference": The definition of t may refer to t

Examples:

- int * bool
- int option
- int list

Remarkable: A lot of data can be described this way.

Convenient to think of as trees.

(optional) jargon: Product types, sum types, recursive types

User-defined types

There are many reasons to define your own types:

- 1. Using a tuple with 12 fields is incomprehensible
- 2. Writing down large types is unpleasant; we have computers for that
- 3. Large programs can use abstract types to be robust to change
 - A couple weeks ahead
- 4. So the language doesn't have to "bake in" lists and options and ...

Datatype

One-of types are less similar across languages

• We'll discuss OO's approach to one-of in a few weeks

In ML, we use make a new type with a datatype binding, e.g.:

Semantics: Extend the environment with three *constructors* (in part, functions/constants that produce values of type mytype)

So we have a way to build them... what's missing?

The old way

For lists, we had a way to:

- Test which variant a value was
- Extract the values from *value-carrying* variants
 - Makes no sense if you have the wrong variant

What would this look like for mytype?

The new way

Rather than add *variant-tests* and *variant-destructors* (non-standard jargon and nothing to do with C++ destructors), ML has a *case* expression that uses *pattern-matching*.

In its simplest form, case has one pattern for each constructor in a dataype and binds one variable for each value carried. Example:

```
case e of
TwoInts(i1,i2) => e1
| Str s => e2
| Pizza => e3
```

What are the typing rules? What are the evaluation rules?

Type-checking case

In addition to binding local variables and requiring branches to have the same type, the typing rules for case prevent some run-time errors:

- Exhaustiveness: No test can "fail" (a warning)
- Redundancy: No test can be "impossible" (an error)

So far, case gives us what we *need* to use datatypes:

- A (combined) way to test variants and extract values
- Powerful enough to define our own tests and destructors

In fact, pattern-matching is far more general and elegant:

- Can use it for datatypes already in the top-level environment
- Can use it for any type (later)
- Can have deep patterns (later)