

## Goals

---

- Contrast type synonyms with new types
- See pattern-matching for built-in “one of” types (not really a concept, but important for ML programming) and “each of” types
- Investigate why accumulator-style recursion can be more efficient

## CSE 341: Programming Languages

Autumn 2005

Lecture 5 — Type synonyms, more pattern-matching, accumulators

CSE 341 Autumn 2005, Lecture 5

1

CSE 341 Autumn 2005, Lecture 5

2

## Type synonyms

---

You can bind a *type name* to a type. Example:

```
type intpair = int * int
```

(We call something else a *type variable*.)

In ML, this creates a *synonym*, also known as a *transparent* type definition. Recursion not allowed.

So a type name is *equivalent* to its definition.

To contrast, the type a datatype binding introduces is not equivalent to any other type (until possibly a later type binding).

CSE 341 Autumn 2005, Lecture 5

3

CSE 341 Autumn 2005, Lecture 5

4

## Review: datatypes and pattern-matching

---

Evaluation rules for datatype bindings and case expressions:

```
datatype t = C1 of t1 | C2 of t2 | ... | Cn of tn
```

Adds constructors  $C_i$  where  $C_i$   $v$  is a value (and  $C_i$  has type  $t_i \rightarrow t$ ).

```
case e of p1 => e1 | p2 => e2 | ... | pn => en
```

- Evaluate  $e$  to  $v$
- If  $p_i$  is the first pattern to *match*  $v$ , then result is evaluation of  $e_i$  in environment extended by the match.
- If  $C$  is a constructor of type  $t_1 * \dots * t_n \rightarrow t$ , then  $C(x_1, \dots, x_n)$  is a pattern that matches  $C(v_1, \dots, v_n)$  and the match extends the environment with  $x_1$  to  $v_1 \dots x_n$  to  $v_n$ .
- Coming soon: many more pattern forms.

CSE 341 Autumn 2005, Lecture 5

3

CSE 341 Autumn 2005, Lecture 5

4

## Why patterns?

Even without more pattern forms, this design has advantages over functions for “testing and destructing” (e.g., `null`, `hd`, and `tl`):

- easier to check for missing and redundant cases
- more concise syntax by combining “test, destruct, and bind”
- you can easily define testing and destructing in terms of pattern-matching

In fact, case expressions are the preferred way to test variants and extract values from all ML’s “one-of” types, including predefined ones (`[]` and `::`: just funny syntax).

So: Do *not* use functions `hd`, `tl`, `null`, `isSome`, `valOf`

Teaser: These functions are useful for *passing as values*

CSE 341 Autumn 2005, Lecture 5

5

CSE 341 Autumn 2005, Lecture 5

6

## Tuple/record patterns

You can also use patterns to extract fields from tuples and records:

pattern `{f1=x1, ..., fn=xn}` (or `(x1,...,xn)`) matches `{f1=v1, ..., fn=vn}` (or `(v1,...,vn)`).

For record-patterns, field-order does not matter.

This is better style than `#1` and `#foo`, and it means you do not (ever) need to write function-argument types.

Instead of a case with one pattern, better style is a pattern directly in a `val` binding.

Next time: “deep” (i.e., nested) patterns.

## Recursion

You should now have the hang of recursion:

- It’s no harder than using a loop (whatever that is)
- It’s much easier when you have multiple recursive calls (e.g., with functions over ropes or trees)

But there are idioms you should learn for *elegance, efficiency*, and *understandability*.

Today: using an *accumulator*.

## Accumulator lessons

- Accumulators can avoid data-structure copying
- Accumulators can reduce the depth of recursive calls that are not *tail calls*
- Key idioms:
  - Non-accumulator: compute recursive results and combine
  - Accumulator: use recursive result as new accumulator
  - The base case becomes the initial accumulator

You will use recursion in non-functional languages—this lesson still applies.

Let’s investigate the evaluation of `to_list_1` and `to_list_2`.

CSE 341 Autumn 2005, Lecture 5

7

CSE 341 Autumn 2005, Lecture 5

8